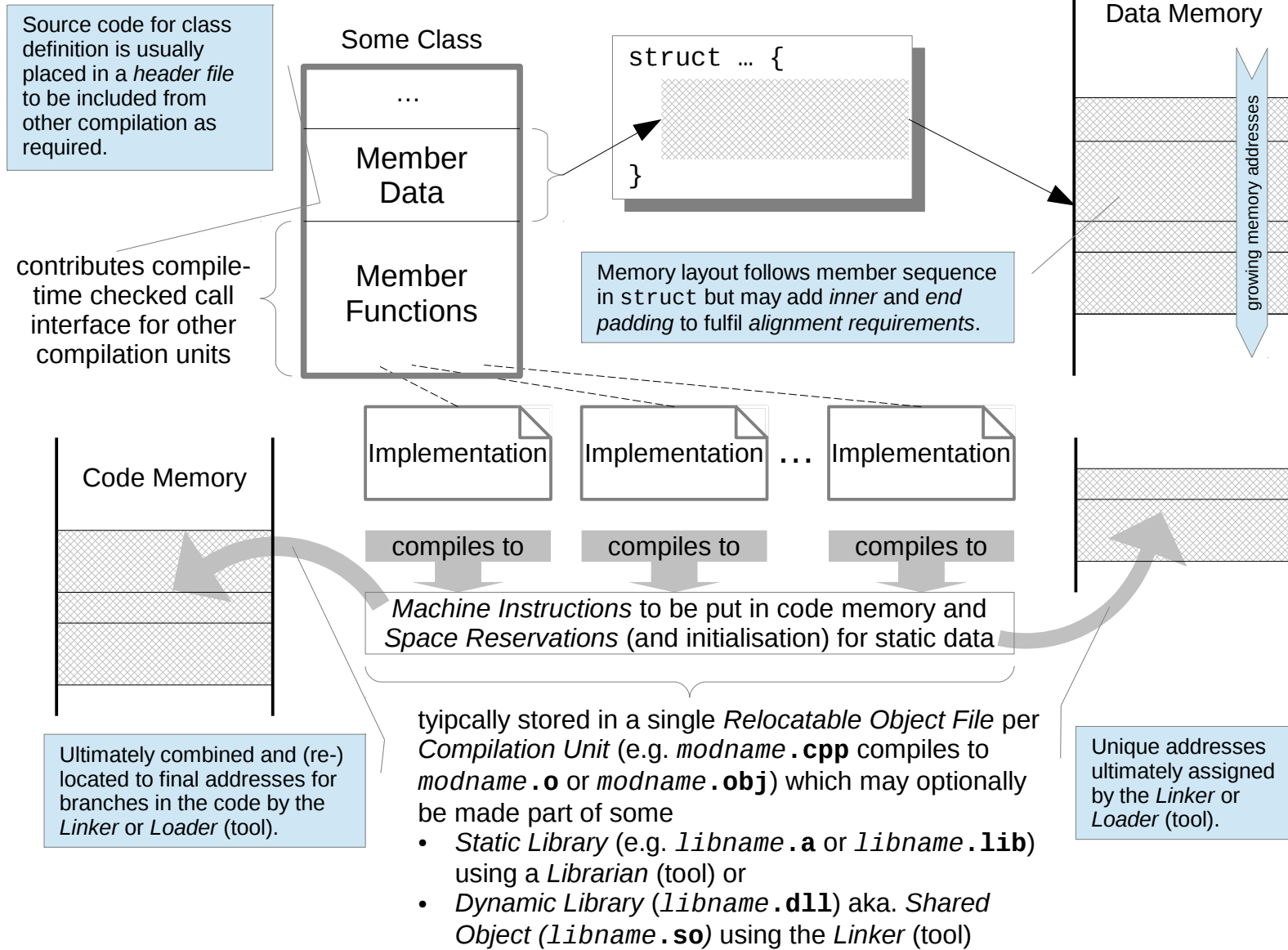
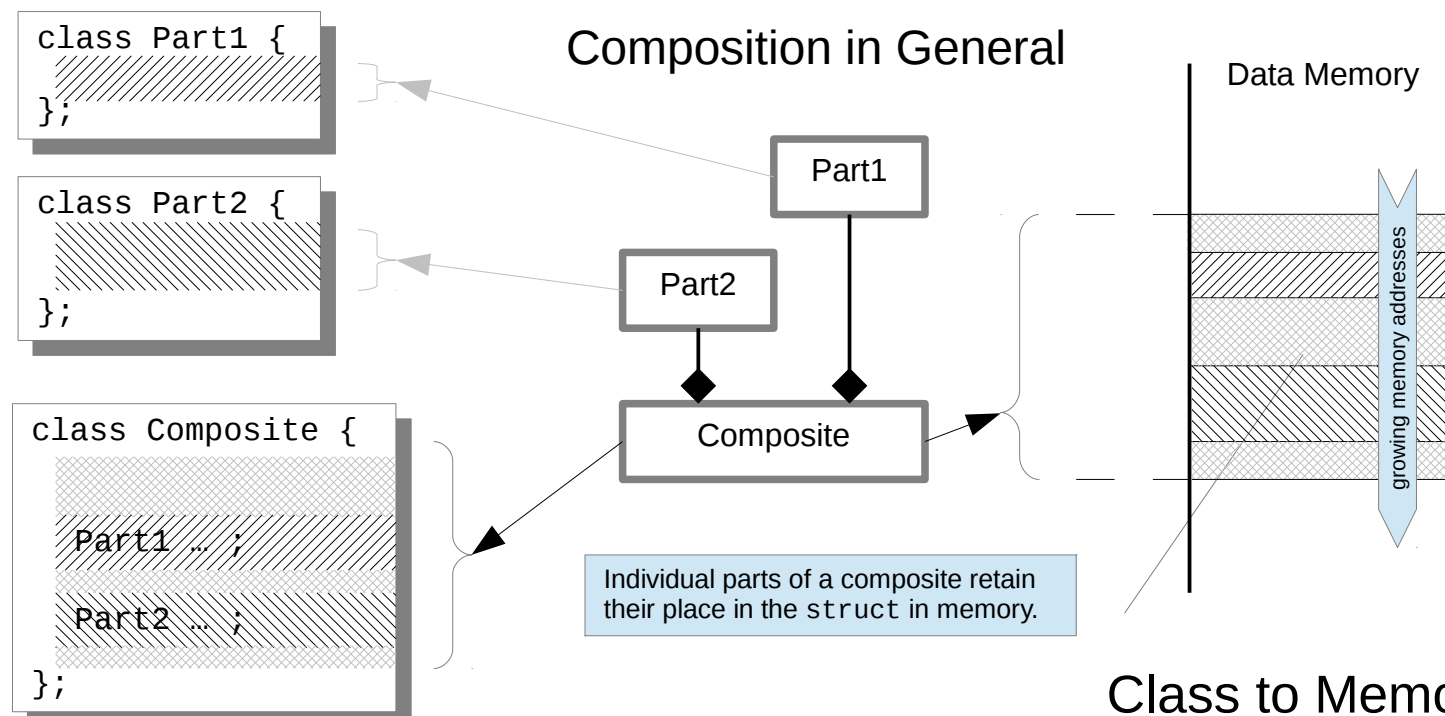


Mapping Classes to Code And Data



Composition in General



Class to Memory Mapping

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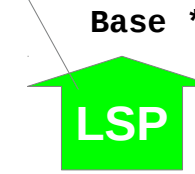
The LSP is part of the type conversion rules built-in to the compiler (and a “no-op” at run-time for single inheritance): Any *Derived* can trivially be used as a *Base* because the memory layout of the latter is part of the former.

```

class Base {
};
    
```

```

class Derived
: public Base {
};
    
```



Base *

Derived *

Base

Derived

Public Base Class (Inheritance)

Data Memory

growing memory addresses

... versus ...

In absence of an explicit type conversion the compiler declines to accept a *Derived* where a *Base* is expected (despite the former contains the memory layout of the latter).

```

class Base {
};
    
```

```

class Derived
: private Base {
};
    
```



Base *

Derived *

Base

Derived

Privat Base Class (Composition)

Data Memory

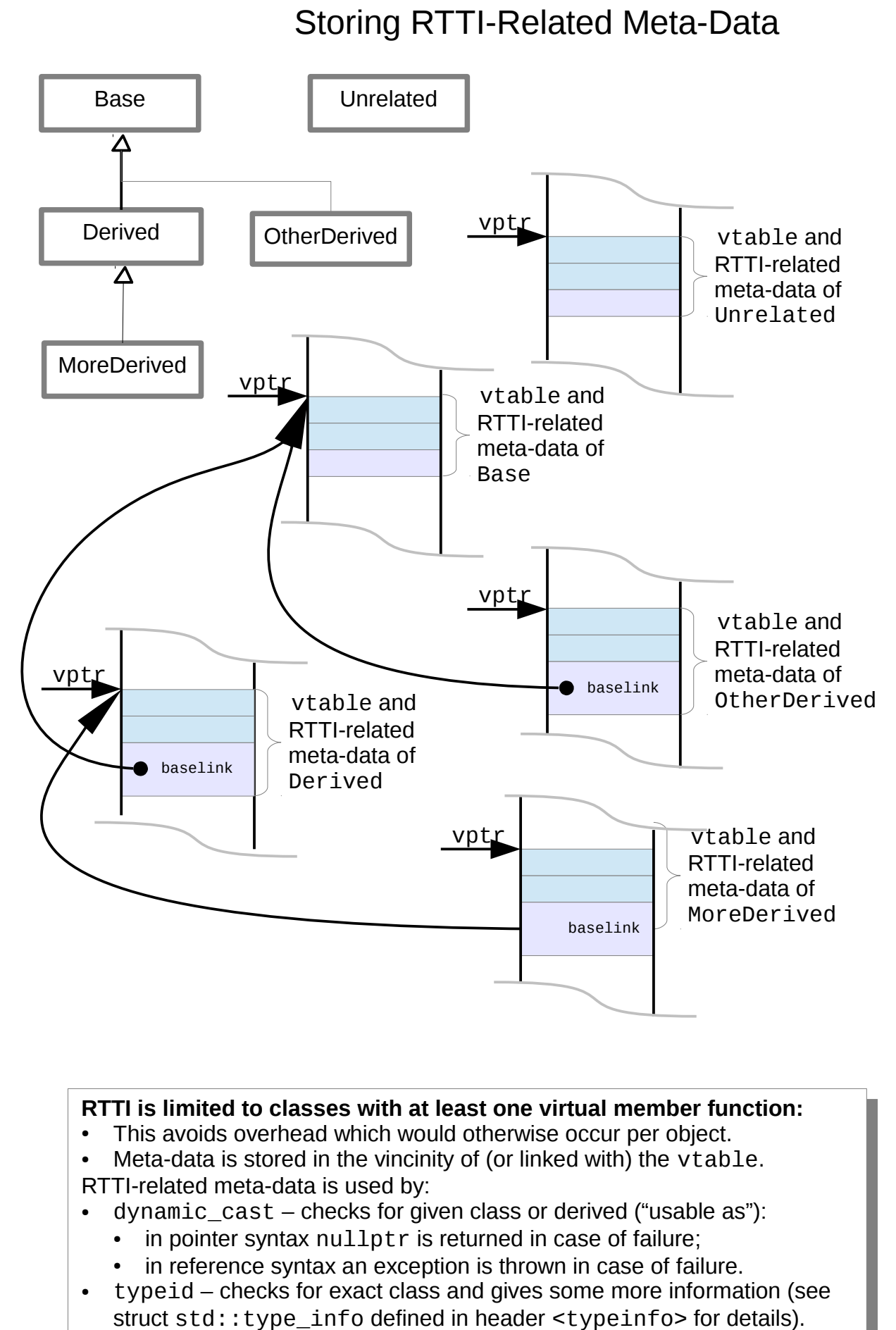
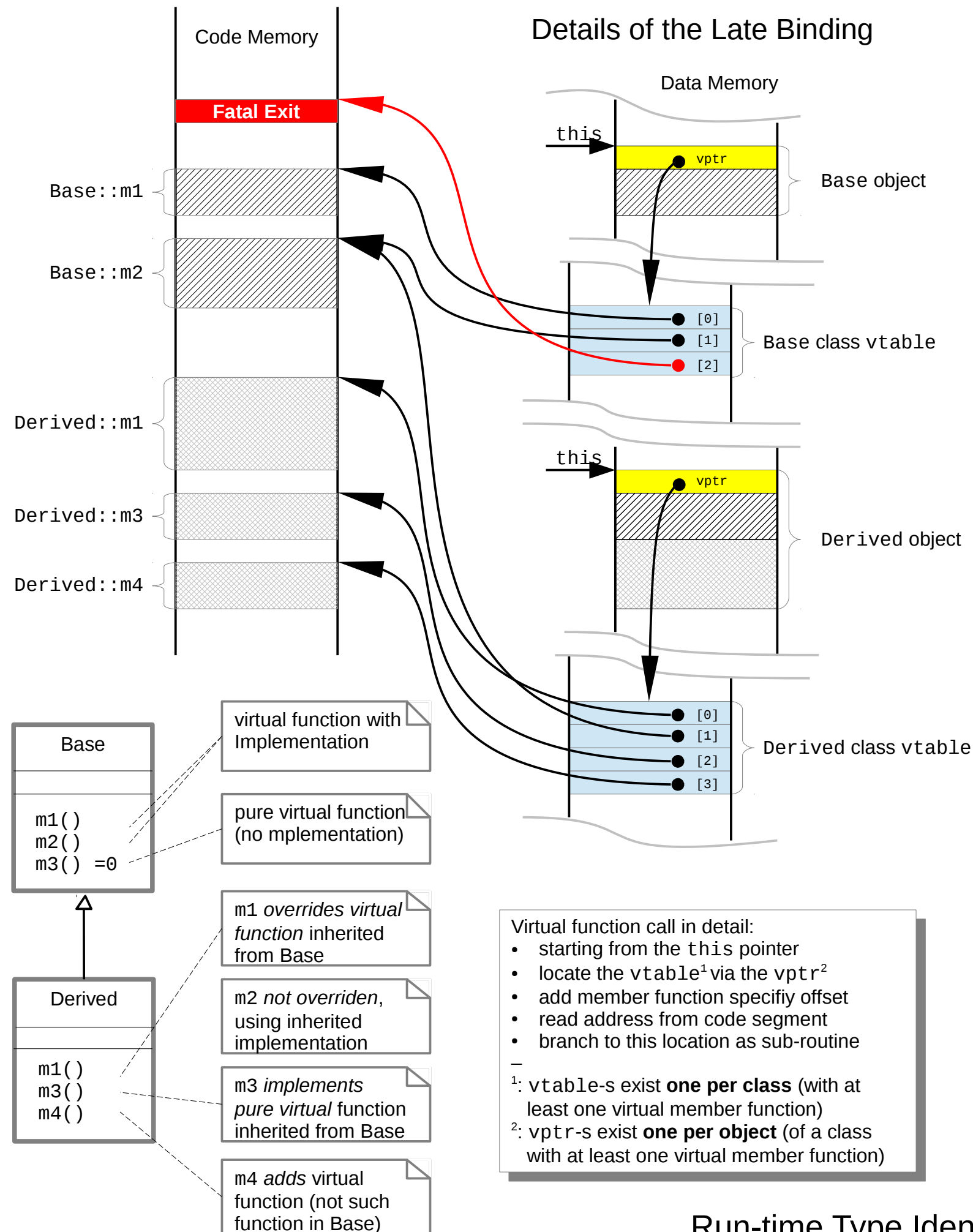
growing memory addresses

The LSP – short for “Liskov Substitution Principle” - was formulated by *Barbara Liskov* and demands:

- Any object of a derived class should be a valid substitute for an object of its – direct or indirect – base classes.

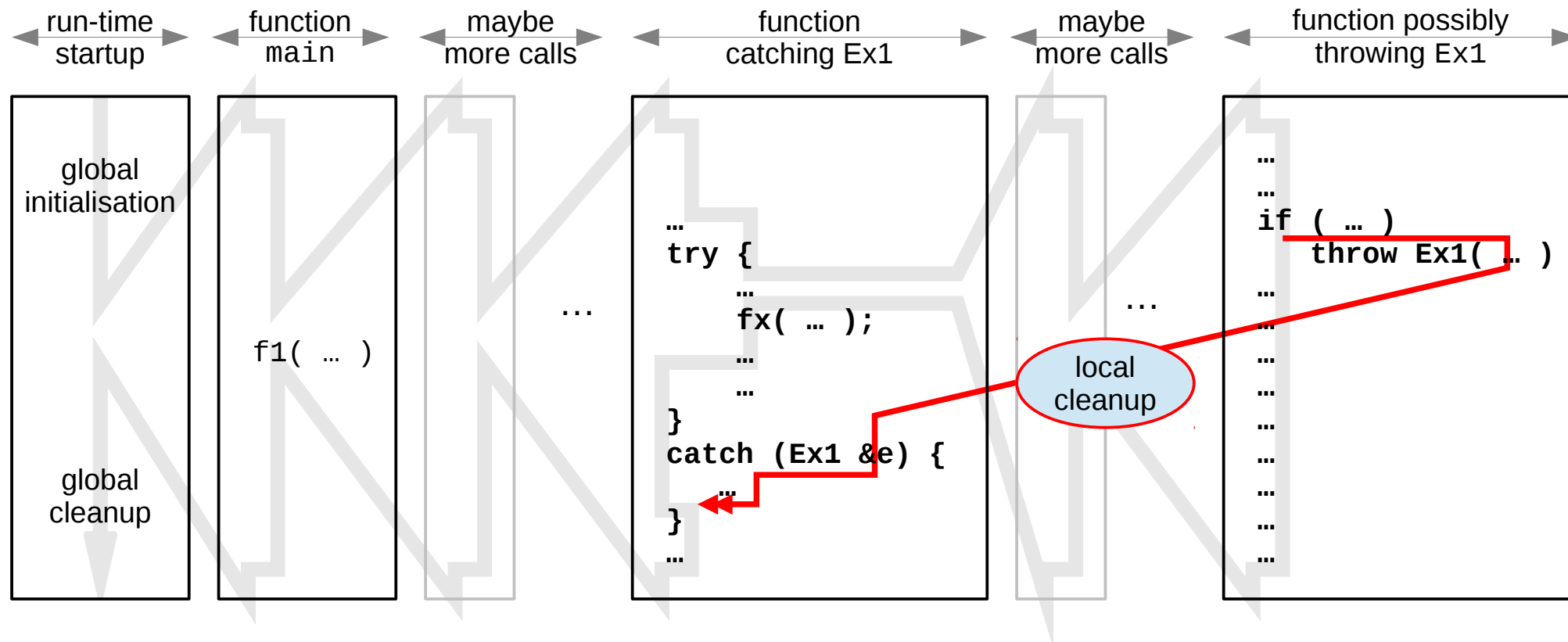
As long as only single inheritance is used the LSP is effectively a “no-op” in C++ since base class objects start at the same memory address as their derived classes.

For private base classes there is no LSP in C++, hence they should be viewed as *Composition*, not *Inheritance*!

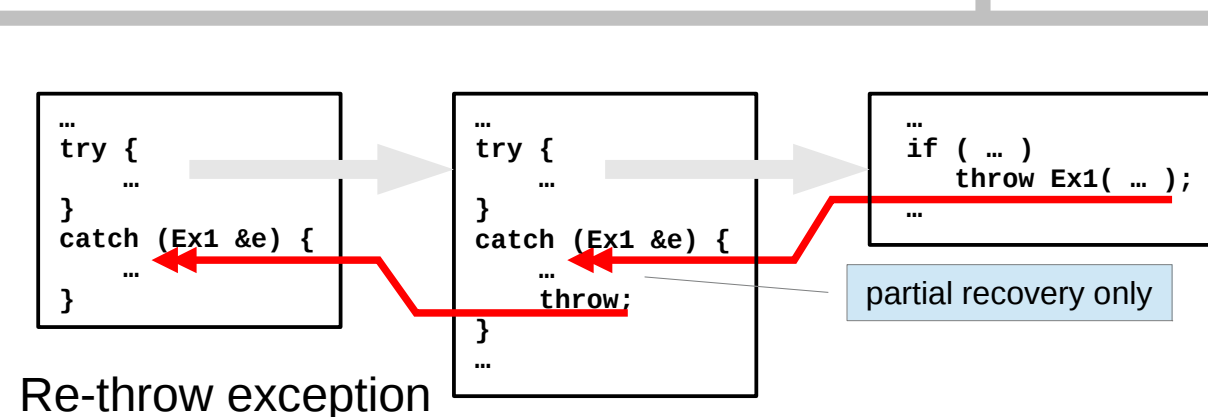
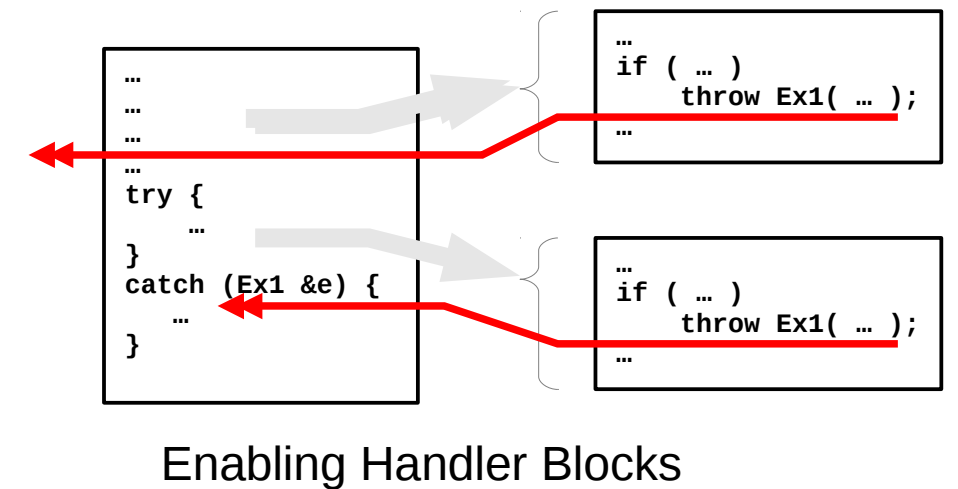
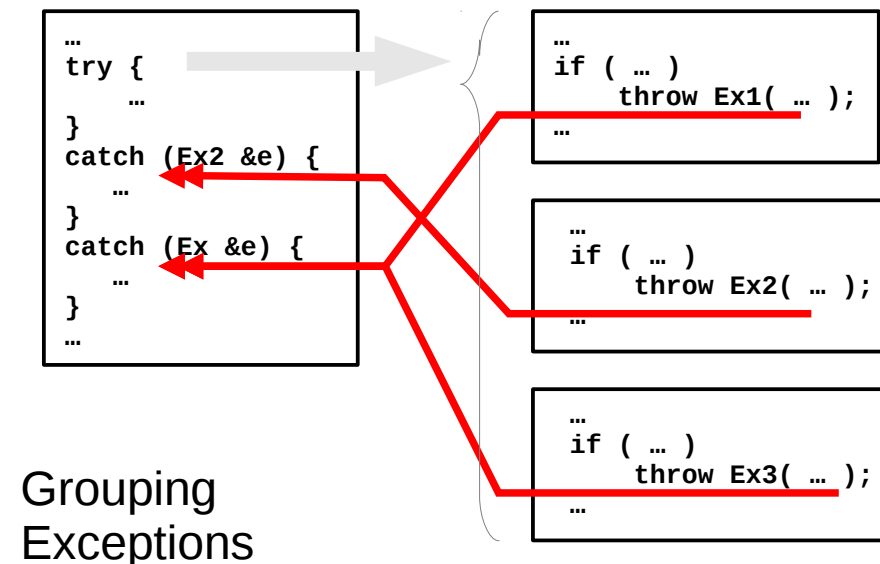
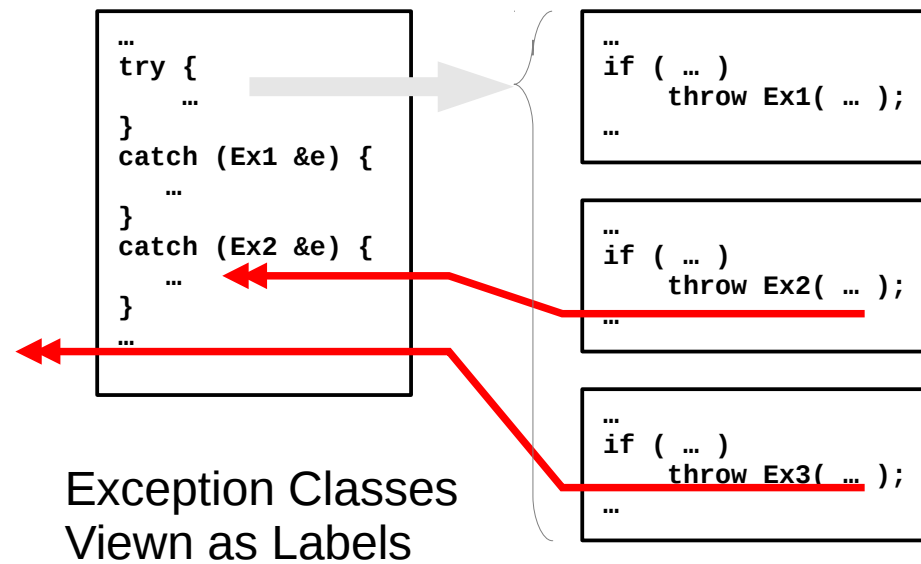
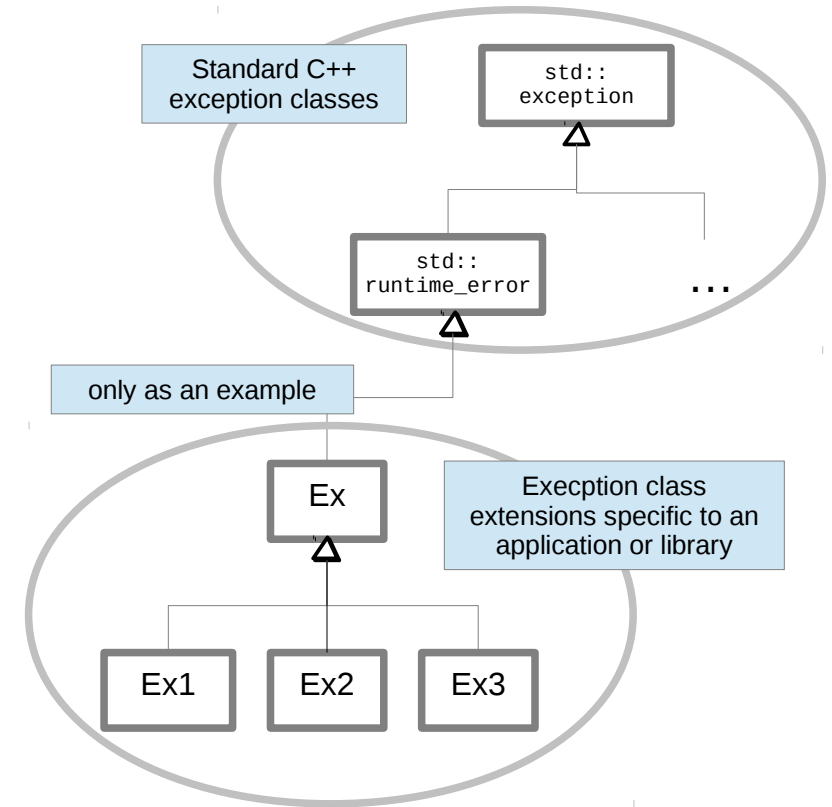


Run-time Type Identification

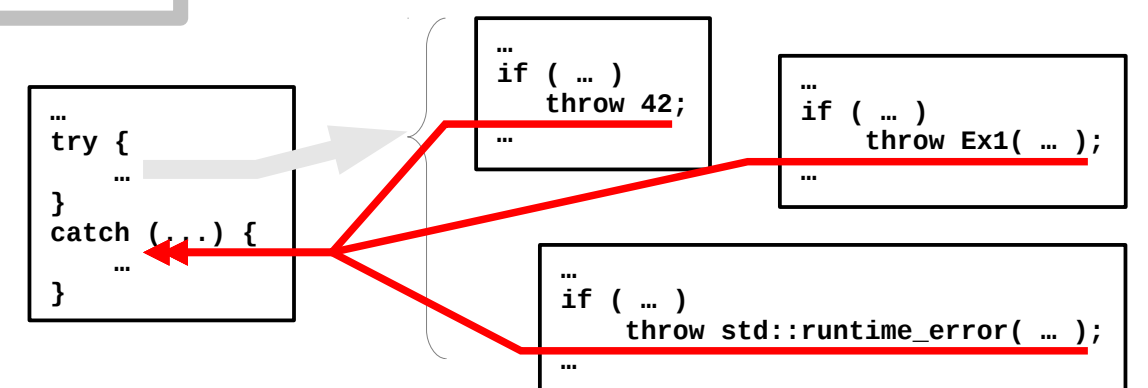
Execution Path taken for Exception



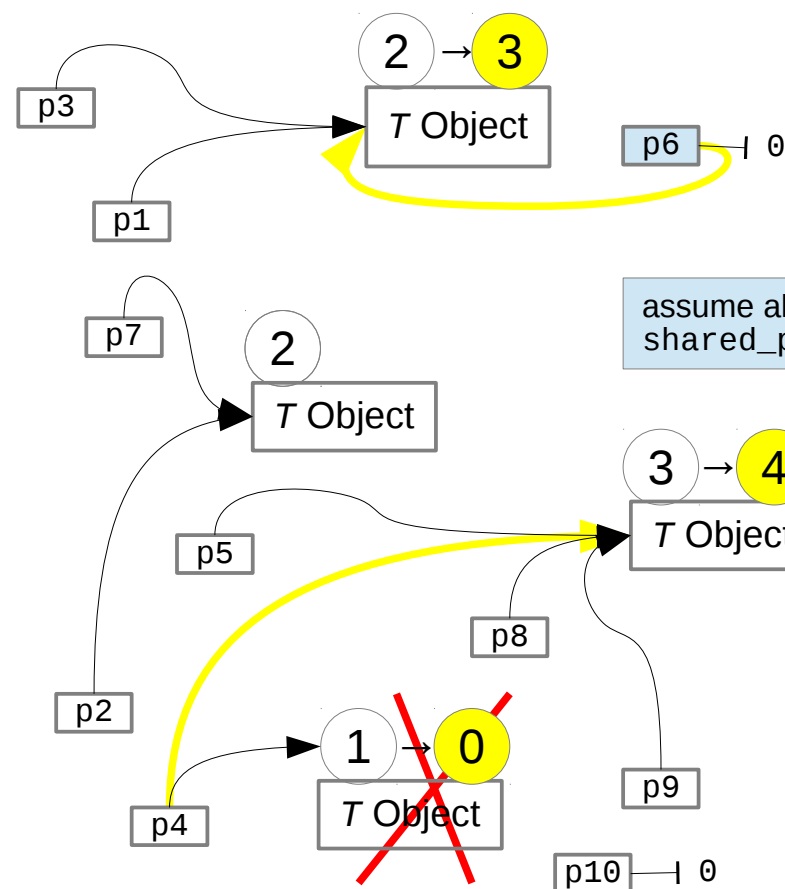
Exception Class Hierarchies



Catch Any Exception



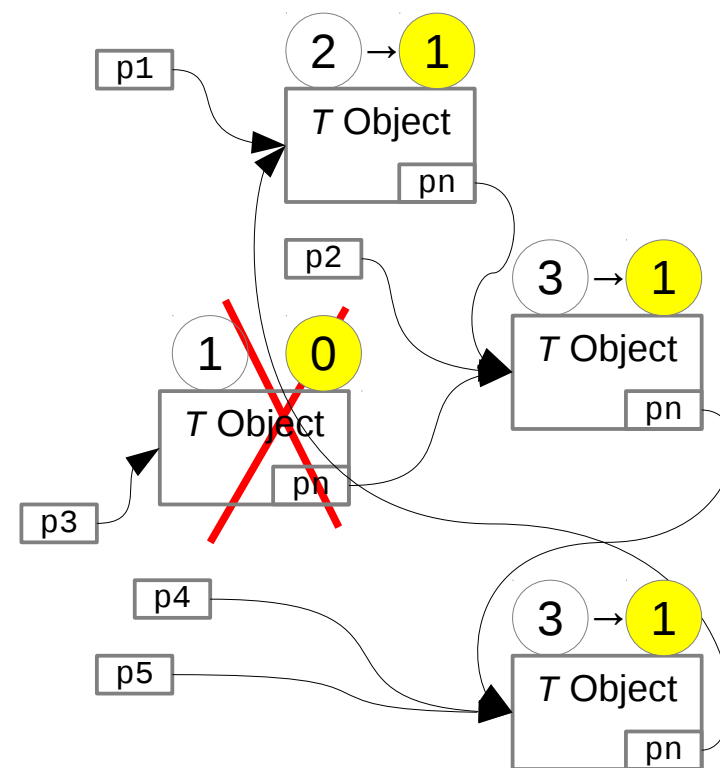
Exception Basics



assume all pointers
shared_ptr<T>

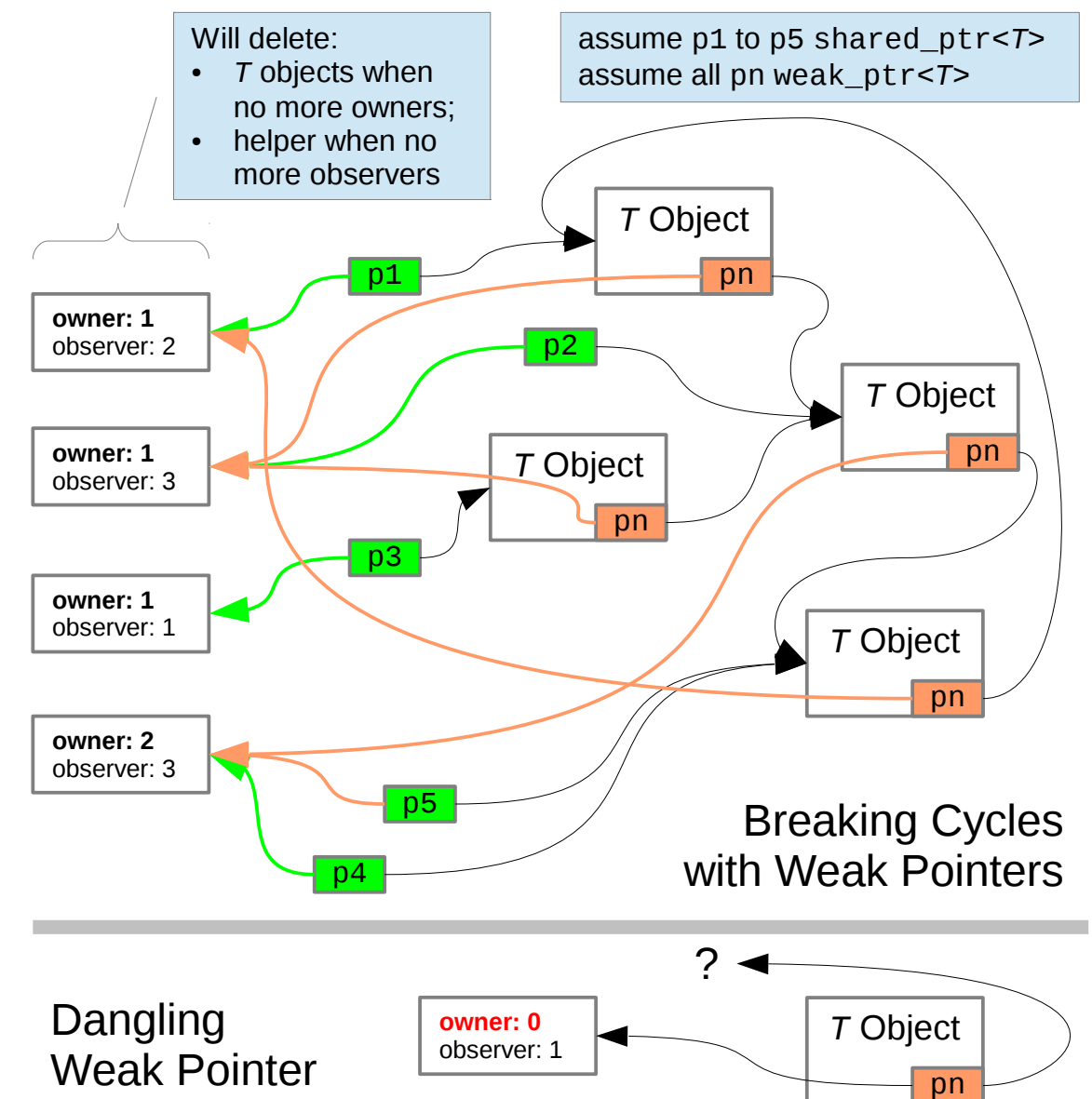
```
// assume assignments
p6 = p3;
p4 = p5;
```

Reference Counting Principle



```
// assume life-time
// of p1 to p5 ends
```

Problem of Cyclic References



Will delete:
• T objects when
no more owners;
• helper when no
more observers

assume p1 to p5 shared_ptr<T>
assume all pn weak_ptr<T>

Breaking Cycles
with Weak Pointers

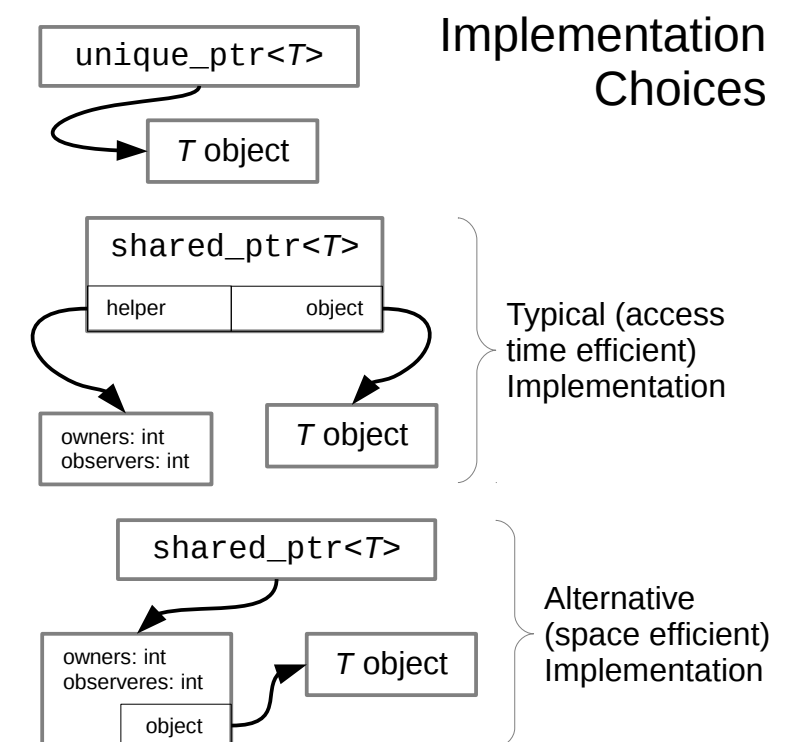
Dangling
Weak Pointer

Comparing ...	std::unique_ptr<T>	std::shared_ptr<T>	Remarks
Characteristic	refers to a single object of type <i>T</i> , uniquely owned	refers to a single object of type <i>T</i> , possibly shared with other referrers	may also refer to "no object" (like a nullptr)
Data Size	same as plain pointer	same as a plain pointer plus some extra space per referred-to object	
Copy Constructor	no*	yes	particularly efficient as only pointers are involved
Move Constructor	yes		
Copy Assignment	no*		a <i>T</i> destructor must also be called in an assignment if the current referrer is the only one referring to the object
Move Assignment	yes		
Destructor (when referrer life-time ends)	always called for referred-to object	called for referred-to object when referrer is the last (and only) one	

*: explicit use of std::move for argument is possible

Smart Pointer Comparison

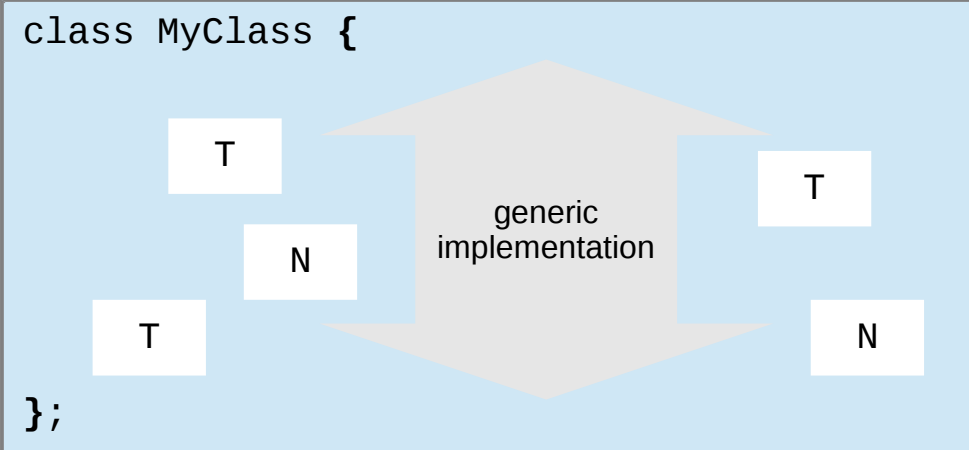
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keywords class and typename have the same meaning in template parameter lists

Template Class

```
template<typename T, int N>
```

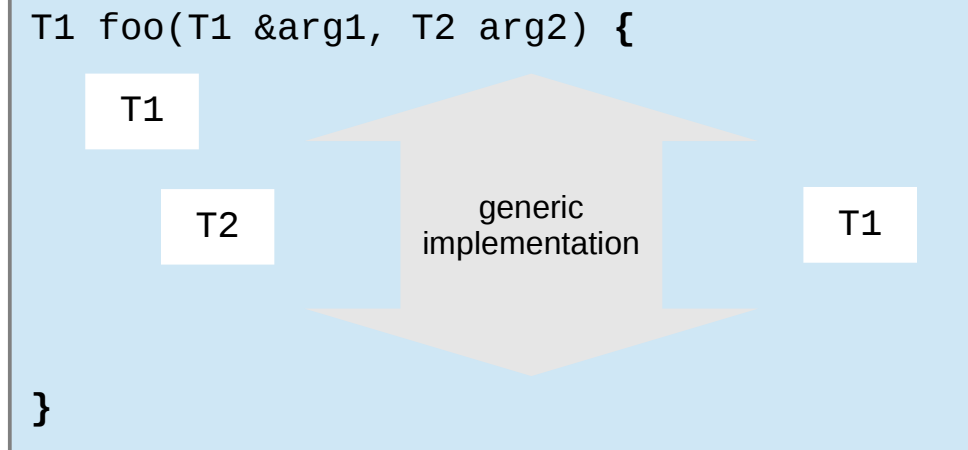


types and constants may be parameterized

preliminary
syntax
checking

Template Function

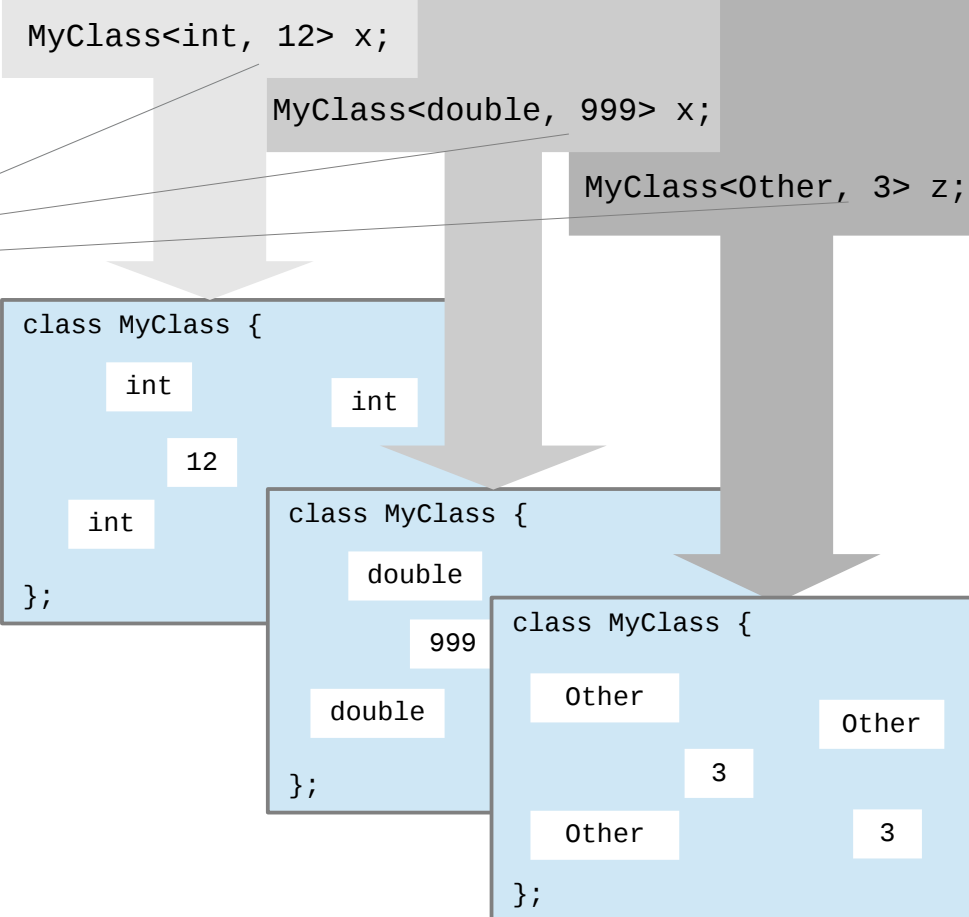
```
template<typename T1, typename T2>
```



typically only types are parameterized

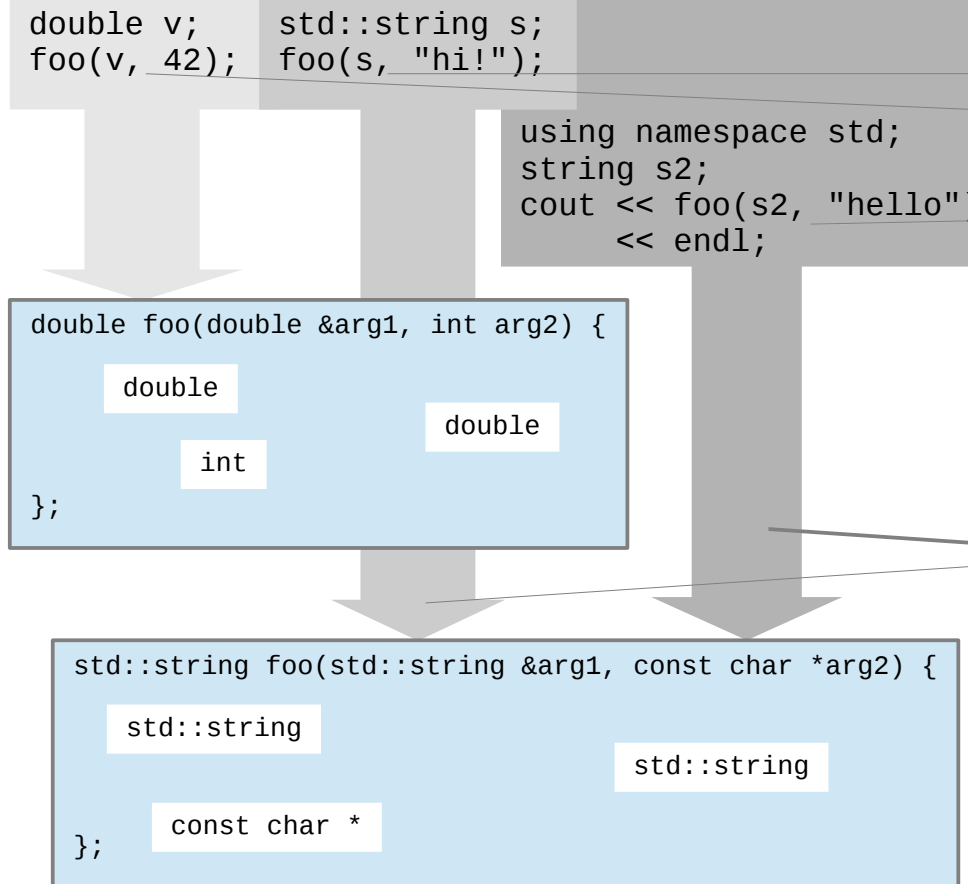
Template definition extends to end of block (i.e. class or function body)

Compiler-Dependant Intermediate Representation



for template classes type and value arguments **must always** be supplied

template-aware
code
generation



for template functions

- types are typically deduced at the call site;
- may optionally be supplied, or
- must** be supplied if **not** present in parameter list (syntax then as for classes: concrete types in angle bracket list following identifier)

duplicated non-inline versions of functions (with identical set of instantiation types) are usually "optimized out" at link time

Code Compiled and Optimised for Specific Template Arguments

Template Basics

Parametrizing *Type* (*double* → *T*) and *Size* (11 → *N+1*)

```
class RingBuffer {
    double data[11];
protected:
    std::size_t iput;
    std::size_t iget;
    static std::size_t wrap(std::size_t idx) {
        return idx % 11;
    }
public:
    RingBuffer()
        : iput(0), iget(0)
    {}
    bool empty() const {
        return (iput == iget);
    }
    bool full() const {
        return (wrap(iput+1) == iget);
    }
    std::size_t size() const {
        return (iput >= iget)
            ? iput - iget
            : iput + 11 - iget;
    }
    void put(const double &);
    void get(double &);
    double peek(std::size_t) const;

    void RingBuffer::put(const double &e) {
        if (full())
            iget = wrap(iget+1);
        assert(!full());
        data[iput] = e;
        iput = wrap(iput+1);
    }

    void RingBuffer::get(double &e) {
        assert(!empty());
        e = data[iget];
        iget = wrap(iget+1);
    }

    double RingBuffer::peek(std::size_t offset = 0) const {
        assert(offset < size());
        return data[wrap(idx + offset)];
    }
};
```

Parametrizing *Type*

```
template<typename Type>
class RingBuffer {
    Type data[11];
    ...
    void put(const Type &);
    void get(Type &);
    Type peek(std::size_t) const;
};

template<typename Type>
void RingBuffer<Type>::put(const Type &e) {
    ...
}

template<typename Type>
void RingBuffer<Type>::get(Type &e) {
    ...
}

template<typename Type>
Type RingBuffer<Type>::peek(std::size_t offset = 0) const {
    ...
}
```

RingBuffer<double> b;
RingBuffer<MyClass> z;

It makes sense to use the net-size here as leaving the last slot empty to differ between an empty and a full buffer can be considered to be an implementation detail.

```
template<std::size_t Size>
class RingBuffer {
    double data[Size+1];
    ...
    static std::size_t wrap(std::size_t idx) {
        return Size+1;
    }
    ...
    std::size_t size() const {
        return (iput >= iget)
            ? iput - iget
            : iput + (Size+1) - iget;
    }
    ...
};

template<std::size_t Size>
void RingBuffer<Size>::put(const double &e) {
    ...
}

template<std::size_t Size>
void RingBuffer<Size>::get(double &e) {
    ...
}

template<std::size_t Size>
double RingBuffer<Size>::peek(std::size_t offset = 0) const {
    ...
}
```

RingBuffer<100> b;
RingBuffer<30> b2;

RingBuffer b;

```
template<typename T, std::size_t N>
class RingBuffer {
    T data[N+1];
protected:
    std::size_t iput;
    std::size_t iget;
    static std::size_t wrap(std::size_t idx) {
        return idx % (N+1);
    }
public:
    RingBuffer()
        : iput(0), iget(0)
    {}
    bool empty() const {
        return (iput == iget);
    }
    bool full() const {
        return (wrap(iput+1) == iget);
    }
    std::size_t size() const {
        return (iput >= iget)
            ? iput - iget
            : iput + (N+1) - iget;
    }
    void put(const T &);
    void get(T &);
    T peek(std::size_t) const;
};

template<typename T, std::size_t N>
void RingBuffer<T, N>::put(const T &e) {
    if (full())
        iget = wrap(iget+1);
    assert(!full());
    data[iput] = e;
    iput = wrap(iput+1);
}

template<typename T, std::size_t N>
void RingBuffer<T, N>::get(T &e) {
    assert(!empty());
    e = data[iget];
    iget = wrap(iget+1);
}

template<typename T, std::size_t N>
T RingBuffer<T, N>::peek(std::size_t offset = 0) const {
    assert(offset < size());
    return data[wrap(idx + offset)];
}
```

RingBuffer<double, 10> b;
RingBuffer<int, 10000> x;
RingBuffer<string, 42> y;
RingBuffer<MyClass, 9> z;

Parametrizing Types and Sizes

C-Compatibility

```
std::string filename;
...
FILE *fp = fopen(filename.c_str(), "r")
```

Where a C-Style string (`const char *`) is expected an `std::string` must be explicitly converted ...

`CharType`
`std::basic_string`

Lookup in reference documentation [here](#) ...

`char`
`std::basic_string`

... but prefer these typedef-s for readability!

`wchar_t`
`std::basic_string` } `std::wstring`

```
void foo(const std::string &);
...
int main() {
    foo("hello, world");
}
```

... the other way round is automatically

```
char c;
std::string s;
while (get_next(c)) {
    ...
    s.append(&c, 1);
}
```

When accepting an `std::string` as read-only argument a const-reference should be used ...

```
void bar(std::string in);
```

... as for value arguments an (avoidable) copy would be created.

Algorithmically filling an `std::string` by always adding to its end can be considered efficient as reservations internally care for extra space.

Efficiency

hi!

a

Hello, world

```
std::string a("hi!");
std::string b("Hello, world");
```

a = b;

becomes (silently) **shared**

a

Hello, world

b

content copied on write ...

a[7] = 'W';

... and again **unshared**

a

Hello, World

May or May Not Have COW-Implementation

Classes

Input and Output

For input

- use `std::getline` to read until given separator ('`\n`' by default);
- use `operator>>` to skip leading white-space first, then read characters up to next white-space;

For output

- `operator<<` has the usual behaviour.

```
int n = 0;
...
cout << ++n
      << line
      << endl;
```

```
using namespace std;
...
// read standard input by line
string line;
while (getline(cin, line))
    ...
```

```
// to given separator
string w;
istringstream s(line);
... getline(s, w, ':') ...
```

```
// by word
... cin >> w ...
```

convert to every builtin **integral** type

convert to every builtin **floating point** type

overloaded for every builtin **integral** and **floating point** type

```
std::string std::to_string( ... );
```

Standard Strings may – more or less – be used like builtin types.

C++11

```
std::string s;
...
std::stoi(s) ...
std::stol(s) ...
std::stoul(s) ...
std::stoll(s) ...
std::stoull(s) ...
std::stof(s) ...
std::stod(s) ...
std::stold(s) ...
...
```

```
std::string str;
...
for (char c : str) {
    // process str
    // char-by-char
}
...
```

Advanced operations:

- too many to list (→ RTFM).

Process like an STL container:

- the usual iterator subclasses are defined as for containers;
- therefore all STL algorithms may be used on `std::string` too.

C++11 has added some “missing” ones:

- e.g. get rid of over-allocation with `shrink_to_fit`

Basic String Operations

Similar to builtin types:

- construction, assignment, including comparison etc. ... (as can be expected);

Similar to builtin arrays:

- single character access via `operator[]`;
- consider to use `at()` for index-checking at run-time;
- both return a reference, so assignment works too.

As in some other programming languages:

- concatenation with `operator+`;
- `operator+=` for combined assignment.

Basic searching:

- `find`, `rfind`, `find_first_of`, ...

Partition:

- `copy` (partially), `substr`, ...

Numeric Conversions

Standard Strings 101

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Advanced String Operations

Container Dimension														
Library Kind of Container Data Structure Class Name Iterator Category Dereferenced Iterator	STL									Standard Strings	Iterator Interface to I/O-Streams		e.g. Boost <ul style="list-style-type: none">Special Containers<ul style="list-style-type: none">ptr_vectorptr_set...More Maps<ul style="list-style-type: none">bimapmulti_index...	Others
	Sequential Containers					Associative Containers					I/O operations for some type T			
	Random Access			Sequential Access		Tree	Hash	Tree	Hash					
	array	vector	deque	list	forward_list	set	unordered_set	map	unordered_map	string wstring ...	istream_iterator	ostream_iterator		
						multiset	unordered_multiset	multimap	unordered_multimap					
	Random Access Iterators			Bidirectional Iterators	Unidirectional Iterators	Bidirectional Iterators				Random Access Iterators	Input Iterators	Output Iterators		
accesses element							accesses key-value-pair		single character	single item of type T				
operations available via iterators														

Algorithm Dimension	STL	
	Access:	• find • search • ...
	Modify:	• remove • sort • ...
	Miscellaneous:	• count • mimmax_seq • ...
e.g. Boost	Algorithm:	• join • ...
	String_algo:	• trimleft • trimright • ...
Others		

operations expected from iterators



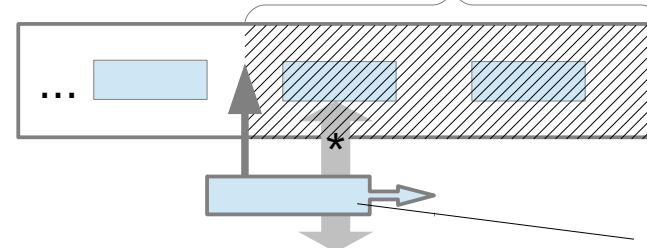
Failure to comply will cause a compile-time error, typically with respect to the header file that defines the algorithm.

Iterators as "Glue" to connect Containers with Algorithms



Failure to comply will either cause a compile-time error or show at runtime and may depend on the kind of container.

elements still physically present though no longer logically part of the container



"Removing" Elements ... returns "New End"

Use of iterators to specify container elements to process:

- starting point is the first element to process
- ending point is the first element **not** to process
- whole container is specified via its `begin()` and `end()`

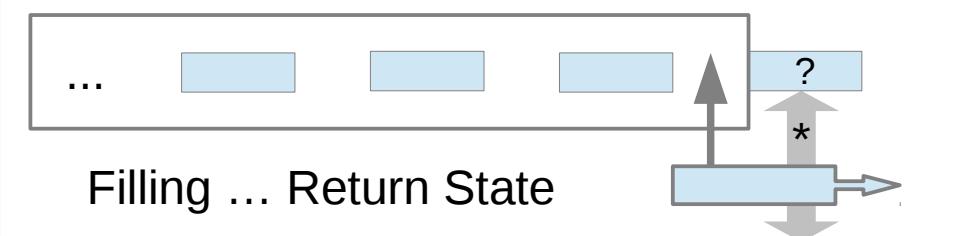


always valid for dereferencing

... Return Success ...

not necessarily valid for dereferencing!

... or Failure



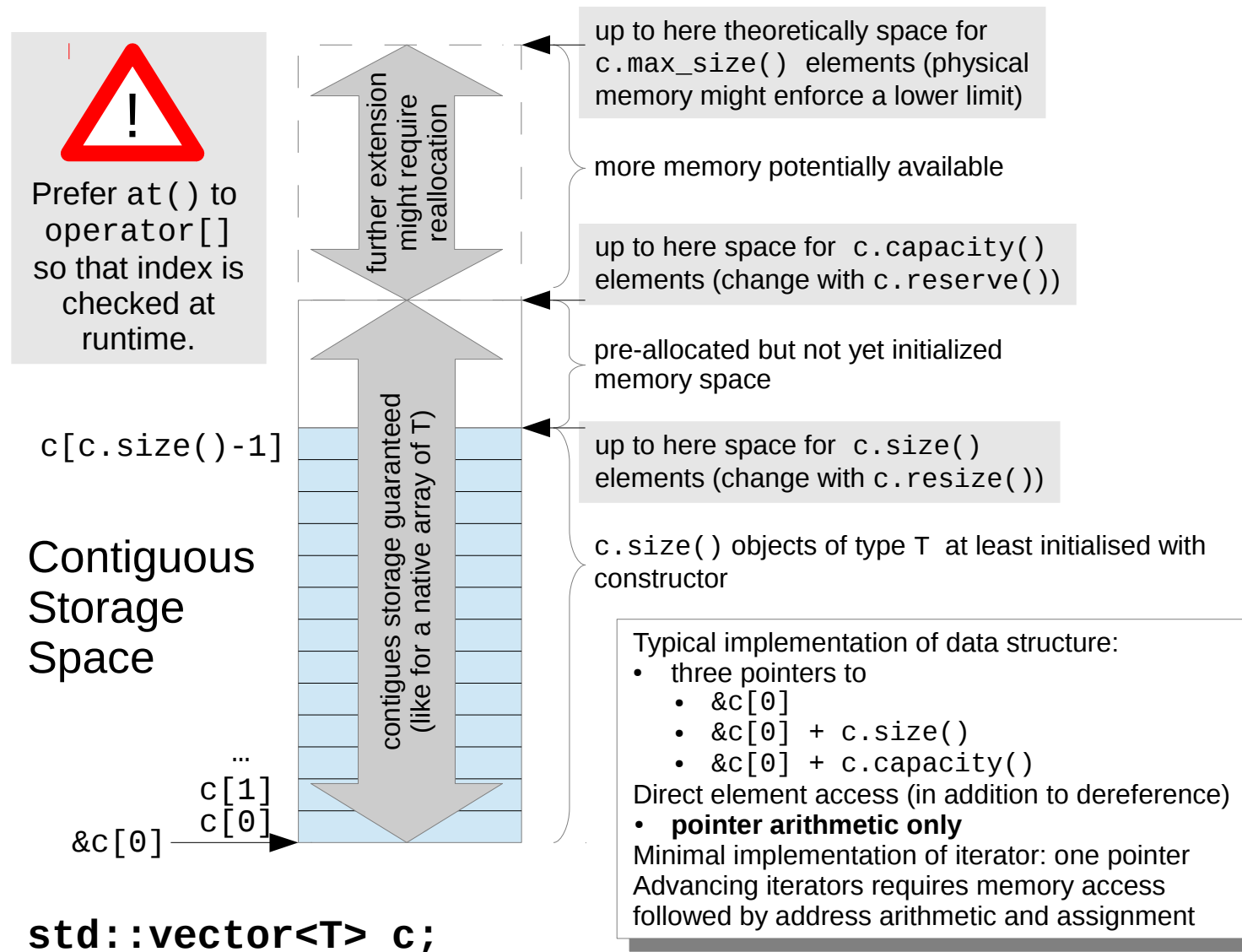
Input Iterators Semantic Restrictions

- * must only be used for read access
- ++ must follow each read exactly once

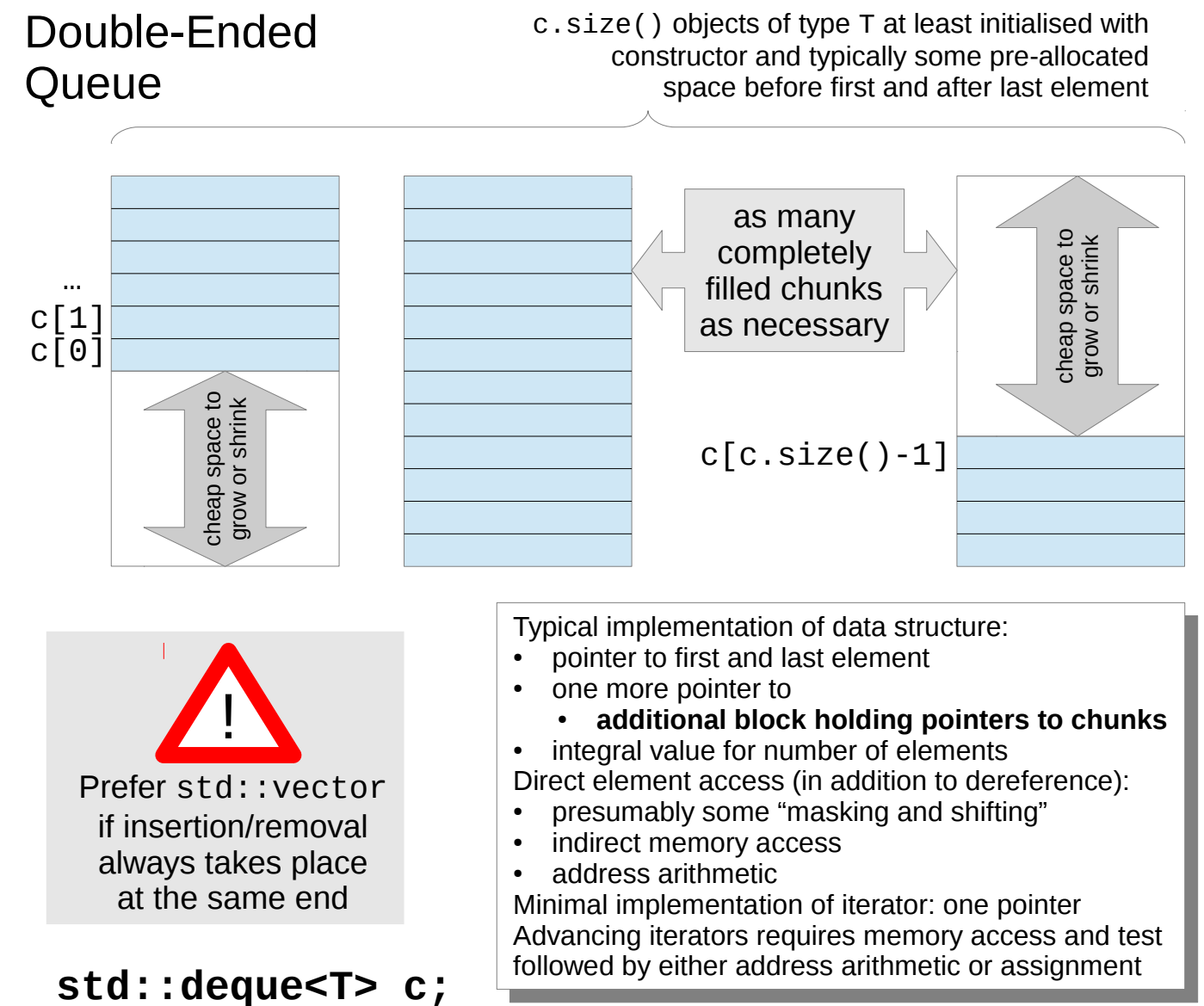
Output Iterators Semantic Restrictions

- * must only be used for write access
- ++ must follow each write exactly once

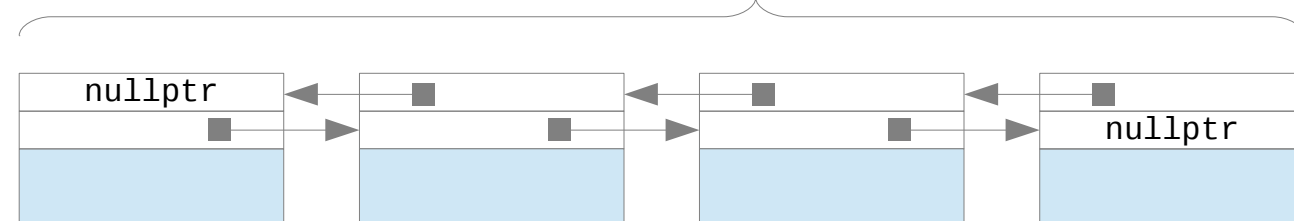
STL – Iterator Usages



Double-Ended Queue



`std::list<T> c;` `c.size()` objects of type T at least initialised with constructor



Typical implementation:

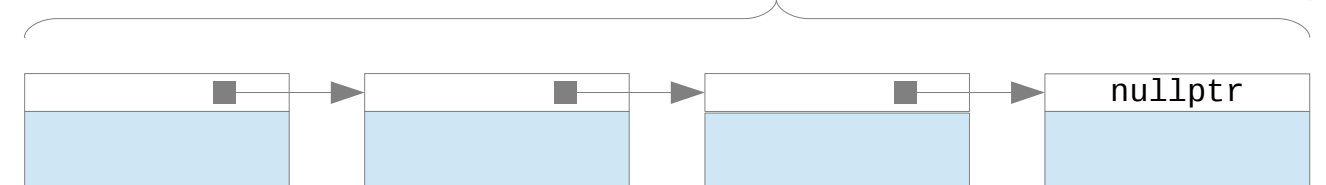
- **two pointers per element**
- pointer to first and last element
- integral value for number of elements

Direct element access not supported!
Minimal implementation of iterator: one pointer
Advancing iterators requires memory access followed by assignment

Double Linked List

Substantial overhead if `sizeof(T)` is small.

`std::forward_list<T> c;` objects of type T initialised with constructor



Singly Linked List

Use `c.empty()` to check whether elements exist.

Typical implementation:

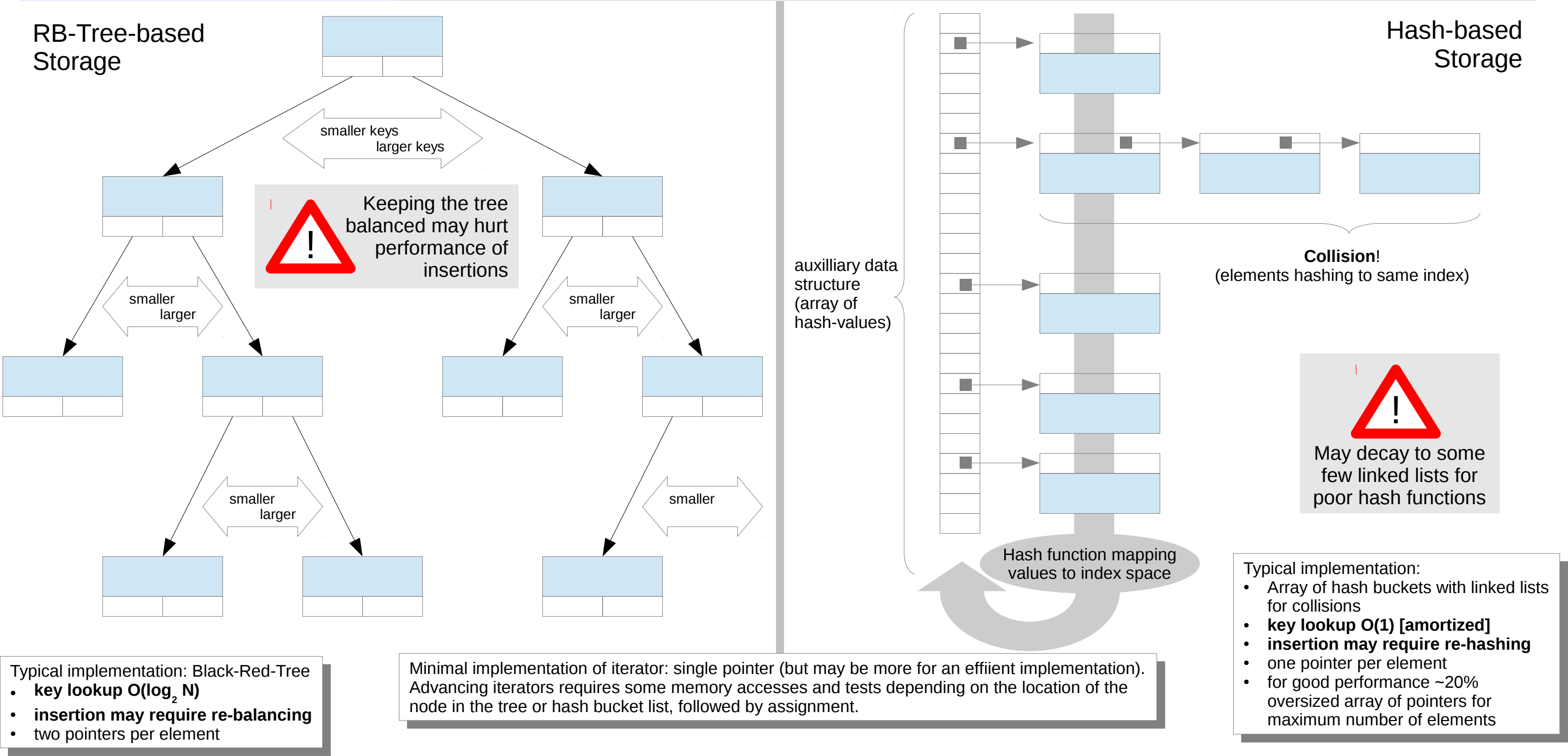
- **one pointer per element**
- only pointer to first element
- **number of elements not stored!**

Direct element access not supported!
Minimal implementation of iterator: one pointer
Advancing iterators requires memory access followed by assignment

STL – Sequence Container Classes

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Contained elements	STL Class Name		Restrictions
objects of type T	<code>std::set</code>	<code>std::unordered_set</code>	unique elements guaranteed
	<code>std::multiset</code>	<code>std::unordered_multiset</code>	multiple elements possible (comparing equal to each other)
pairs of objects of type T_1 (key) and type T_2 (associated value)	<code>std::map</code>	<code>std::unordered_map</code>	unique keys guaranteed
	<code>std::multimap</code>	<code>std::unordered_multimap</code>	multiple keys possible (comparing equal to each other)



Basics



Substantial overhead if
sizeof(T) is small.

```
#include <iostream>
#include <string>
#include <regex>
using namespace std;

...
int main() {
    regex re("he(1+)(o*)");
    ...
    string line;
    while (getline(cin, line)) {
        ... re ... // execute FSM
    }
}
```

Regular Expression
represented in Text Form

constructor

Regex-Object
(FSM representing the RE)

Regular
Expression
Object

Substitution Format:

- may contain any text
- plus the placeholders
\$0, \$1, ... for parts of
the compared string
matching parts of the
regular expression
put in parentheses.

```
...
const char fmt[] = R("
-----
complete match: $0
matching el-s:  $1
matching o-s:   $2
-----
");
...
regex_replace(line, re, fmt) ...
...
```

```
-----
complete match: helloo
matching el-s:  ll
matching o-s:   oo
-----
```

Match-Object:

- allows to access the parts of a string
matching the parts of a regular
expression put into round parentheses;
- has also a size() member function.

```
...
smatch m;
if (regex_search(line, m, re)) {
    // access matching parts
    ... m[0] ...
    ... m[1] ...
    ... m[2] ...
}
...
```

s	a	y		h	e	l	l	o	o	\n
---	---	---	--	---	---	---	---	---	---	----

Eingabezeile

Flexible Comparisons

```
...
string line;
while (getline(cin, line)) {
    if (regex_match(line, re)) {
        // line FULLY matches RE
        // ... whatever
        // ...
    }
    else {
        // no full match
        // ...
    }
}
...
```

```
...
string line;
while (getline(cin, line)) {
    if (regex_search(line, re)) {
        // some PART of line matches RE
        // ... whatever
        // ...
    }
    else {
        // no partial match
        // ...
    }
}
...
```

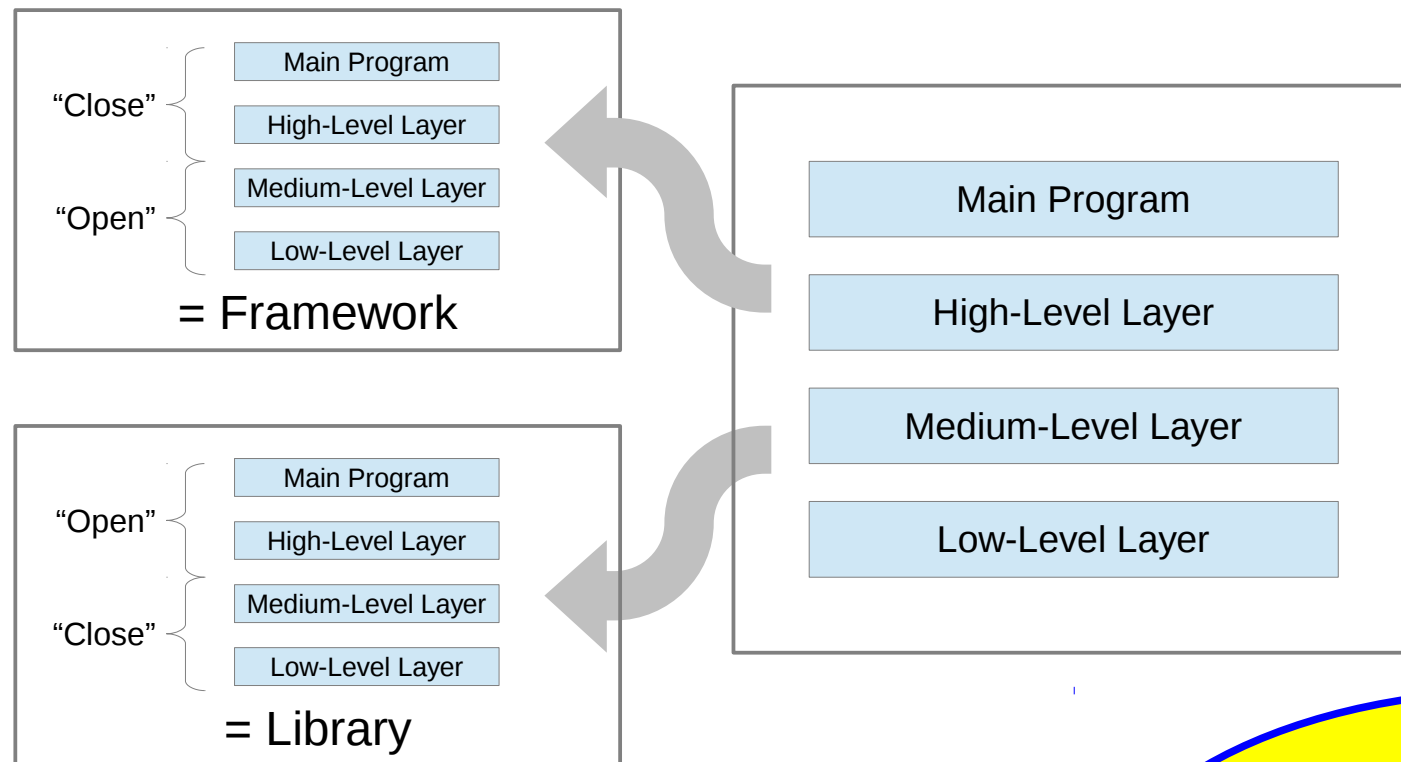
hel hell hellll helllll helo heloo
helloo
say hello
hello?
heo
say hello!
say hello
say hell, oh?!
Othello
say Hello
Goodbye

Powerful Substitutions

Regular Expressions

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Easy Parsings



Design for Reusability

- *Libraries* or *Frameworks* for common components
- Classes for common services or abstractions
- C++-Templates for genericity in types

Use Available Tools and Libraries, e.g.

- *Doxygen* (or similar) to create good-looking documentation from embedded comments
- The *Boost Platform* for a extremely rich choice of "what seems to be missing or forgotten" in the C/C++ Standard Library

Pick the Best from Agility, at least

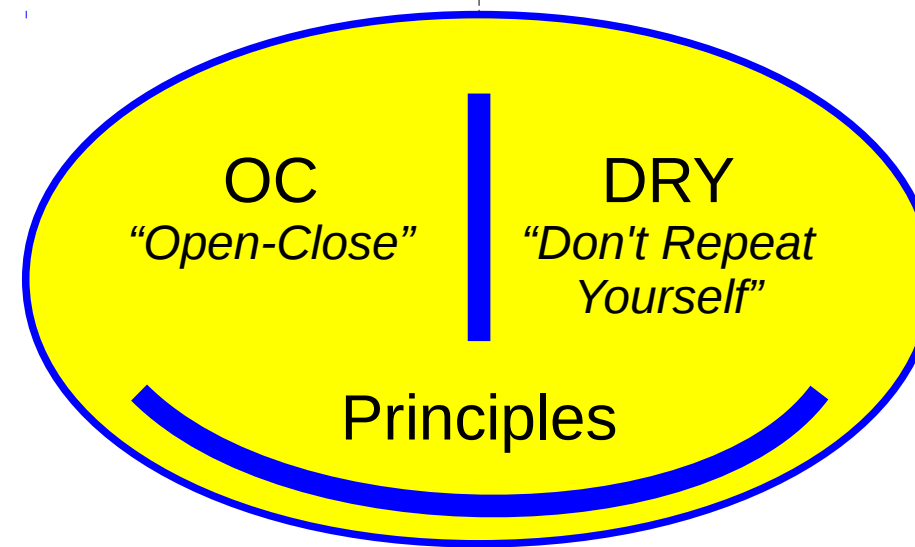
- integrate continuously
- automate boring tests
- (maybe try "pair-programming"?)

Parametrize for Flexibility with

- Run-Time arguments for functions and subroutines
- Compile-Time arguments for templates

Apply Best Practices, e.g.

- Standard Design Patterns (from GoF) like
 - Composite
 - Template Methode
 - ...
- Well-known C++ Idioms like
 - PIMPL (Pointer to Implementation)
 - RAI (Resource Acquisition is Initialisation)
 - CRTP (Curiously Recurring Template Pattern)
 - ...
- Handy Little Techniques where useful
 - "Named Argument" (from C++ FAQ)
 - "Safe delete" (from Boost)
 - ...



Consider to Write Your Own Tools, e.g. to

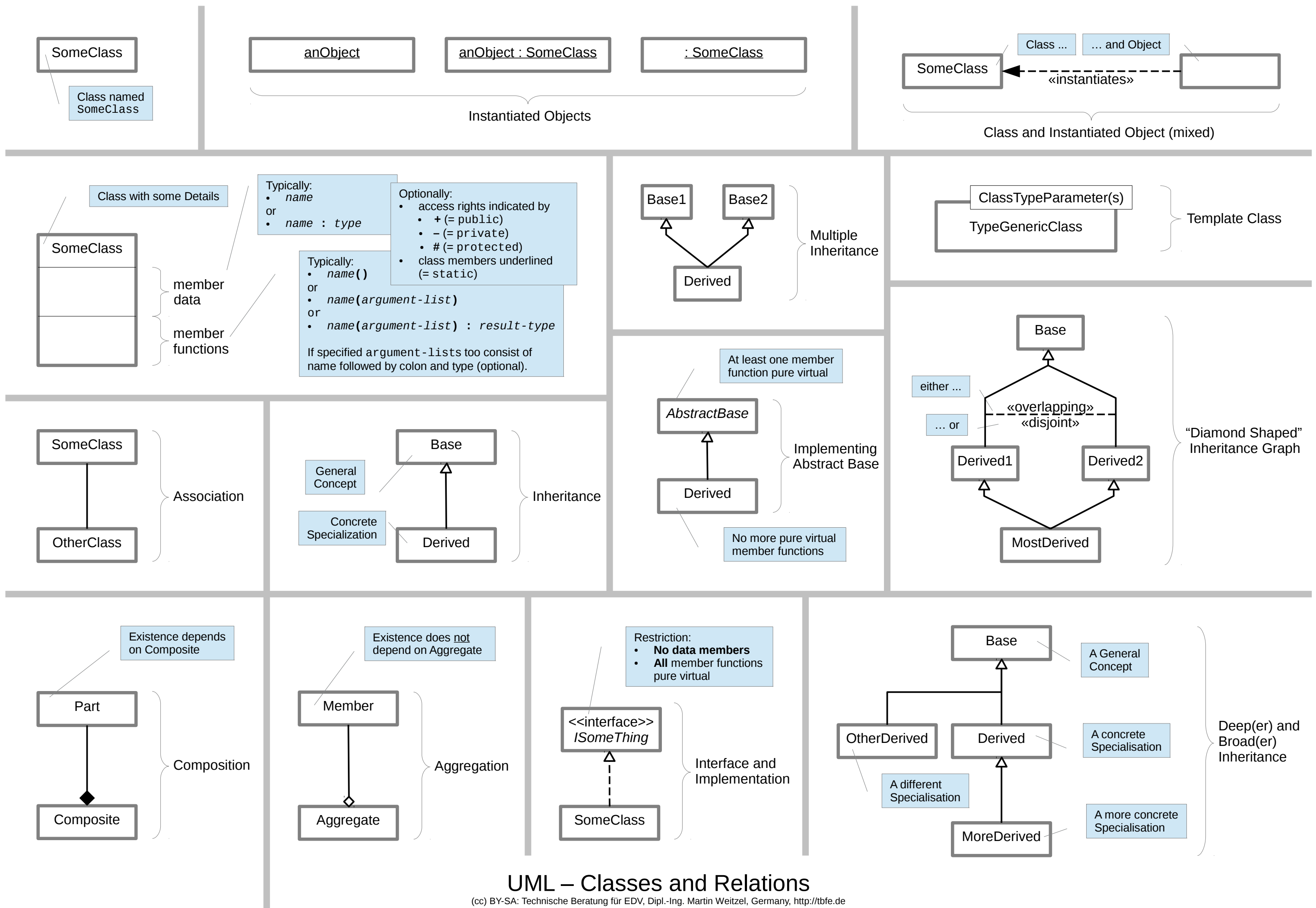
- create a C/C++ header file from a spreadsheet or vice versa
- create a CSV- or XML-document from a source file, or even
- create both, source code and auxilliary documents from a DSL (domain specific language)

But always judiciously decide ... and Don't Overdo!

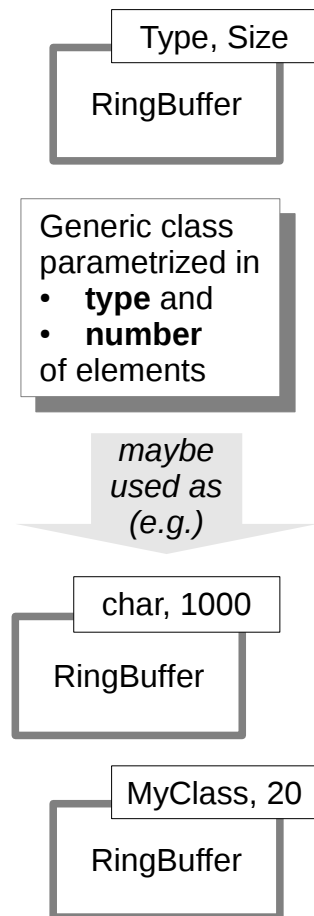
- Not each and every global variable needs to be turned into a Singleton.
- Not each and every little config file needs to be parsed as full XML.
- Not each and every small class needs type genericity.
- ...

If you can't avoid a complex design in the end, at least provide some easy to use defaults for the most common use cases!

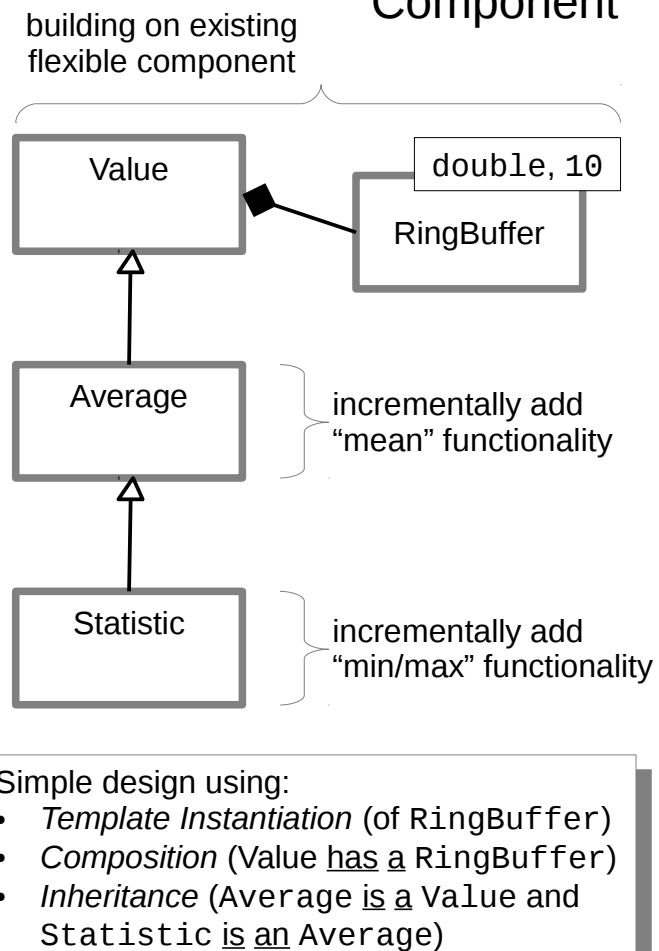
Guiding Principles



Reusable Component

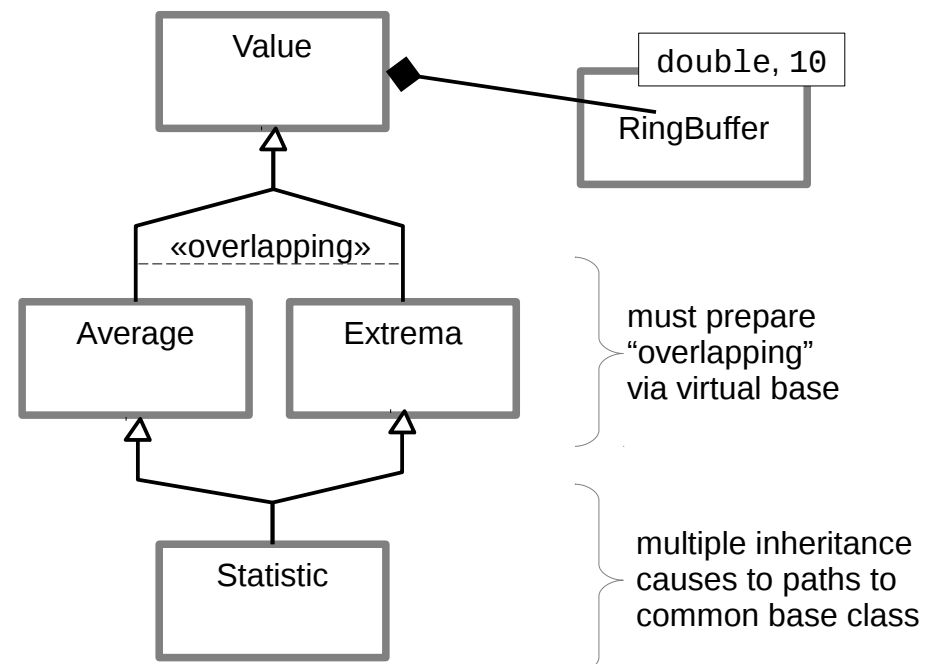


Reusing Adapted Component



offers flexibility in combinations

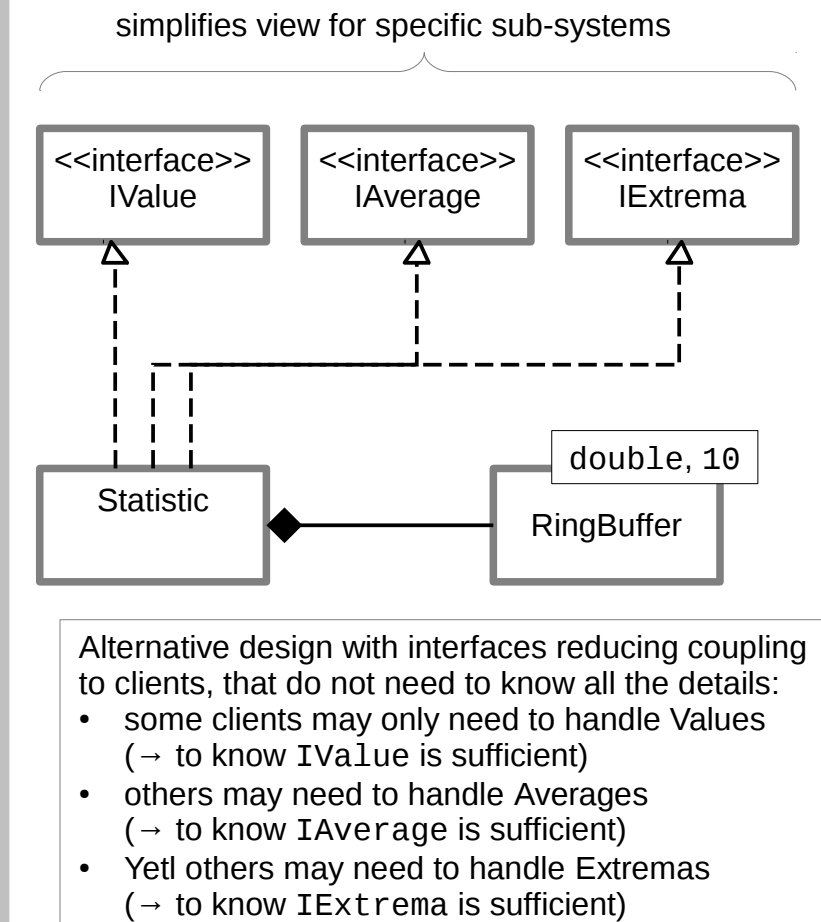
Diamond-Shaped Inheritance



More flexible design with "diamond shaped" inheritance:

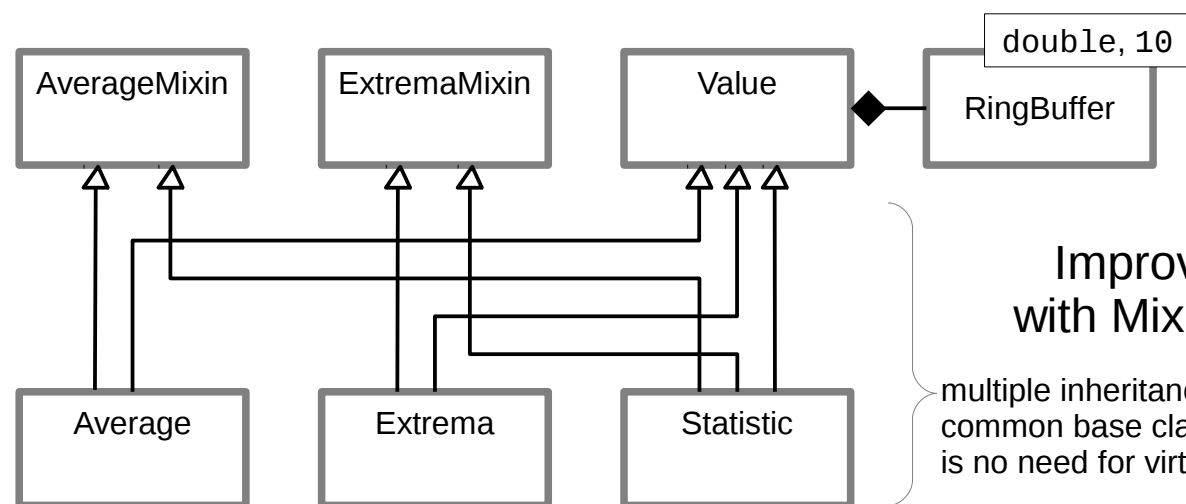
- each of the classes (*Value*, *Average*, *Extrema*, *Statistic*) may be used on its own
- intermediate classes (*Average*, *Value*) must pay the "price" ...
- ... for simple re-use in the most derived class (*Statistic*)

Three Interfaces



incrementally add functionality

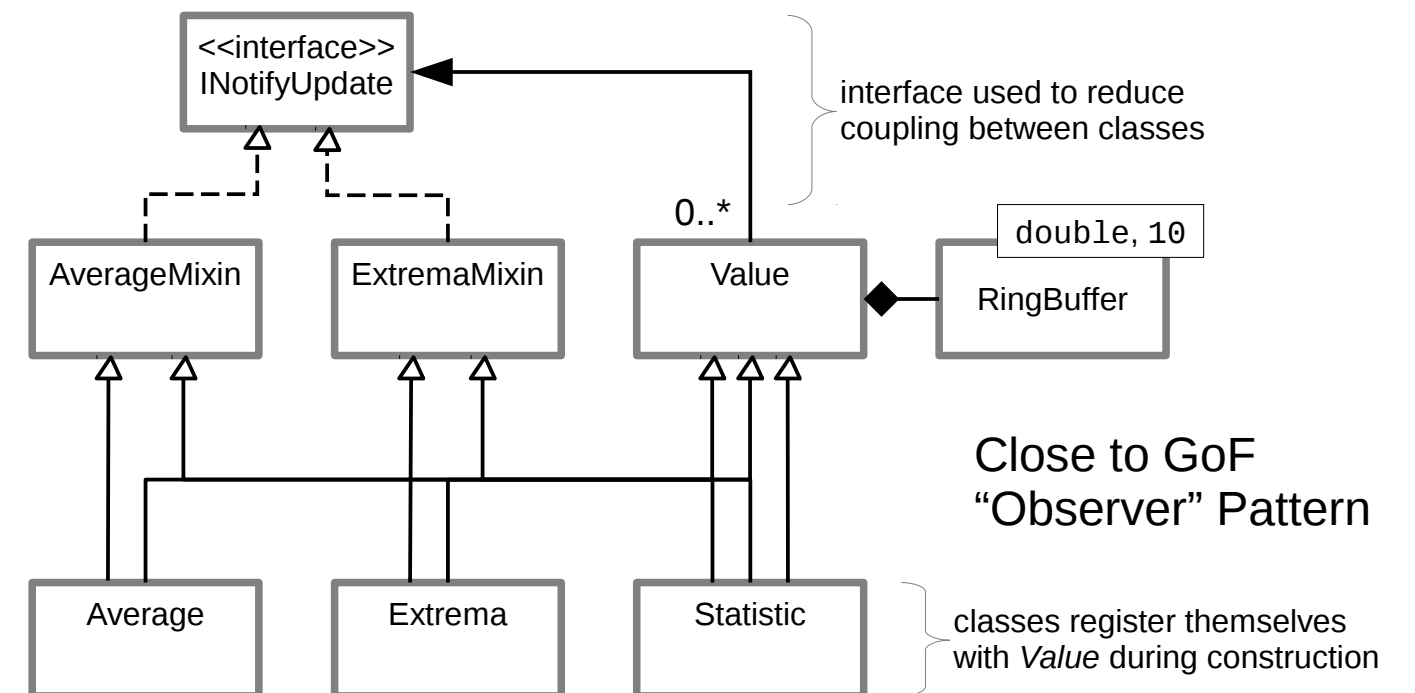
building on existing components



Improved Reuse with Mixin Classes

More elaborate design:

- flexibility achieved with "mixin" classes
- multiple inheritance but not "diamond shaped"



Still more elaborate design:

- Mixins notified via generic interface
- Value only handles INotifyUpdate

Examples – Classes and Relations

Type-dependent Flow of Control

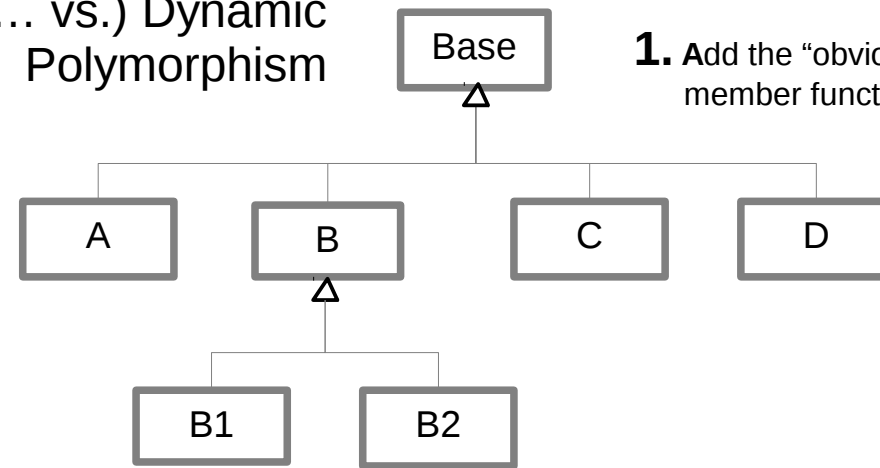
```
void foo(Base &r) {  
    ...  
    if (auto p = dynamic_cast<A*>(&r)) {  
        ...  
    }  
    if (auto p = dynamic_cast<B1*>(&r)) {  
        ...  
    }  
    if (auto p = dynamic_cast<B2*>(&r)) {  
        ...  
    }  
    if (auto p = dynamic_cast<C*>(&r)) {  
        ...  
    }  
    if (auto p = dynamic_cast<D*>(&r)) {  
        ...  
    }  
    ...  
}
```

...
// combining B1 and B2
if (auto p = dynamic_cast<B*>(&r)) {
 ...
}

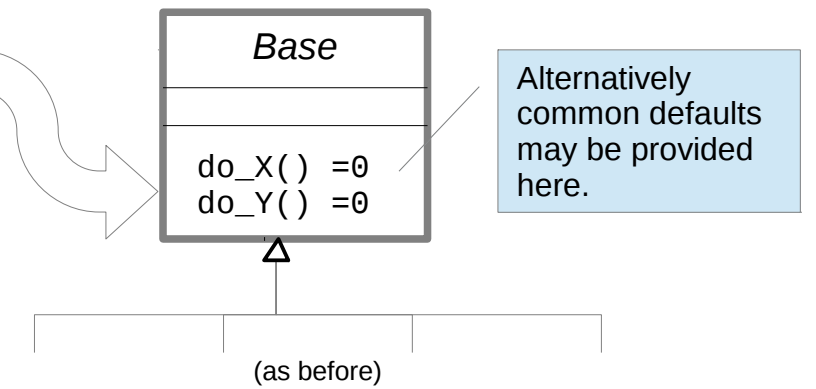
```
void foo(Base &r) {  
    ...  
    if (typeid(r) == typeid(A)) {  
        ...  
    }  
    if (typeid(r) == typeid(B1)) {  
        ...  
    }  
    if (typeid(r) == typeid(B2)) {  
        ...  
    }  
    if (typeid(r) == typeid(C)) {  
        ...  
    }  
    if (typeid(r) == typeid(D)) {  
        ...  
    }  
    ...  
}
```

...
// combining B1 and B2
if (typeid(r) == typeid(B1)
 || typeid(r) == typeid(B2)) {
 ...
}

(... vs.) Dynamic Polymorphism



1. Add the “obviously missing” member functions to Base:



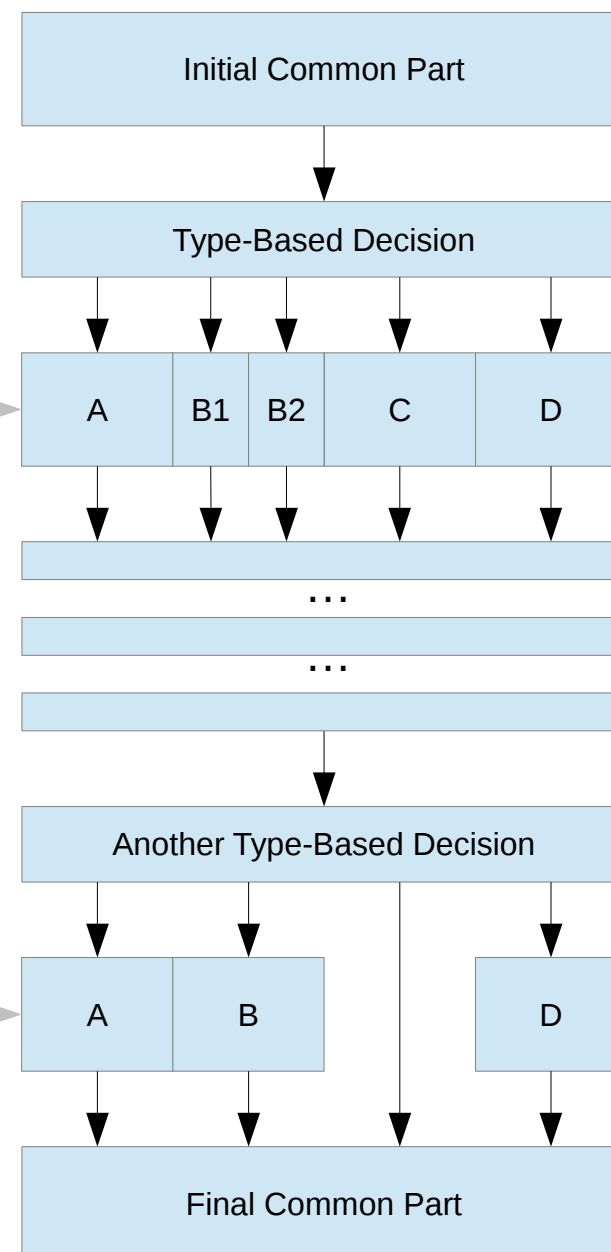
2. Move actions from multiway branches to member function implementations:

```
void A::do_X() {  
    ...  
}  
void B1::do_X() {  
    ...  
}  
void B2::do_X() {  
    ...  
}  
void C::do_X() {  
    ...  
}  
void D::do_X() {  
    ...  
}
```

```
void A::do_Y() {  
    ...  
}  
void B::do_Y() {  
    ...  
}  
void C::do_Y() {  
    /*empty*/  
}  
void D::do_Y() {  
    ...  
}
```

Data can be shared easily in privacy.

The need for sharing data may weaken information hiding.



do X

do Y

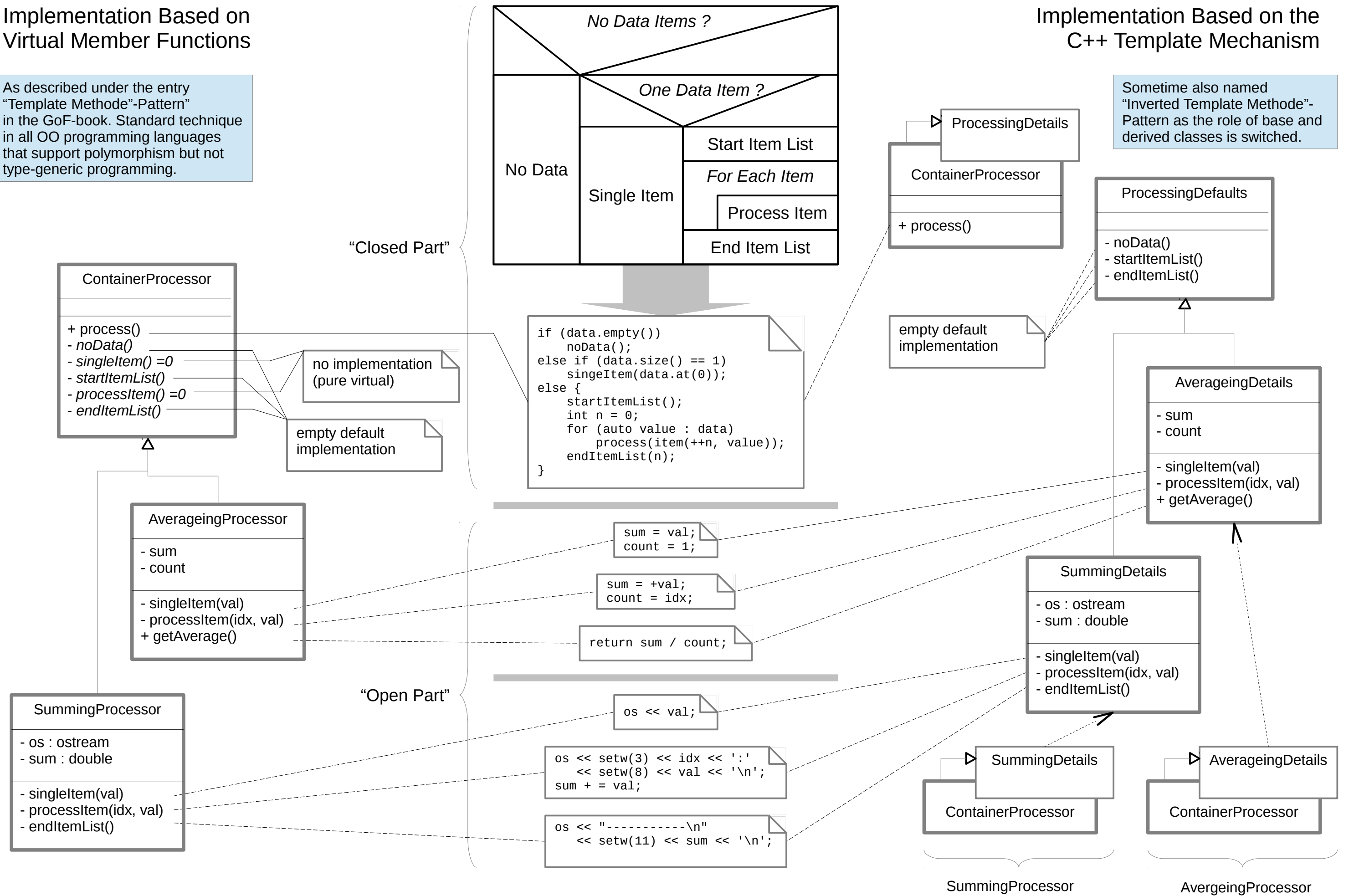
3. Replace multiway branches with member function calls:

```
...  
r.do_X();  
...  
r.do_Y();  
...
```

Type-Based Multiway Branching

Implementation Based on Virtual Member Functions

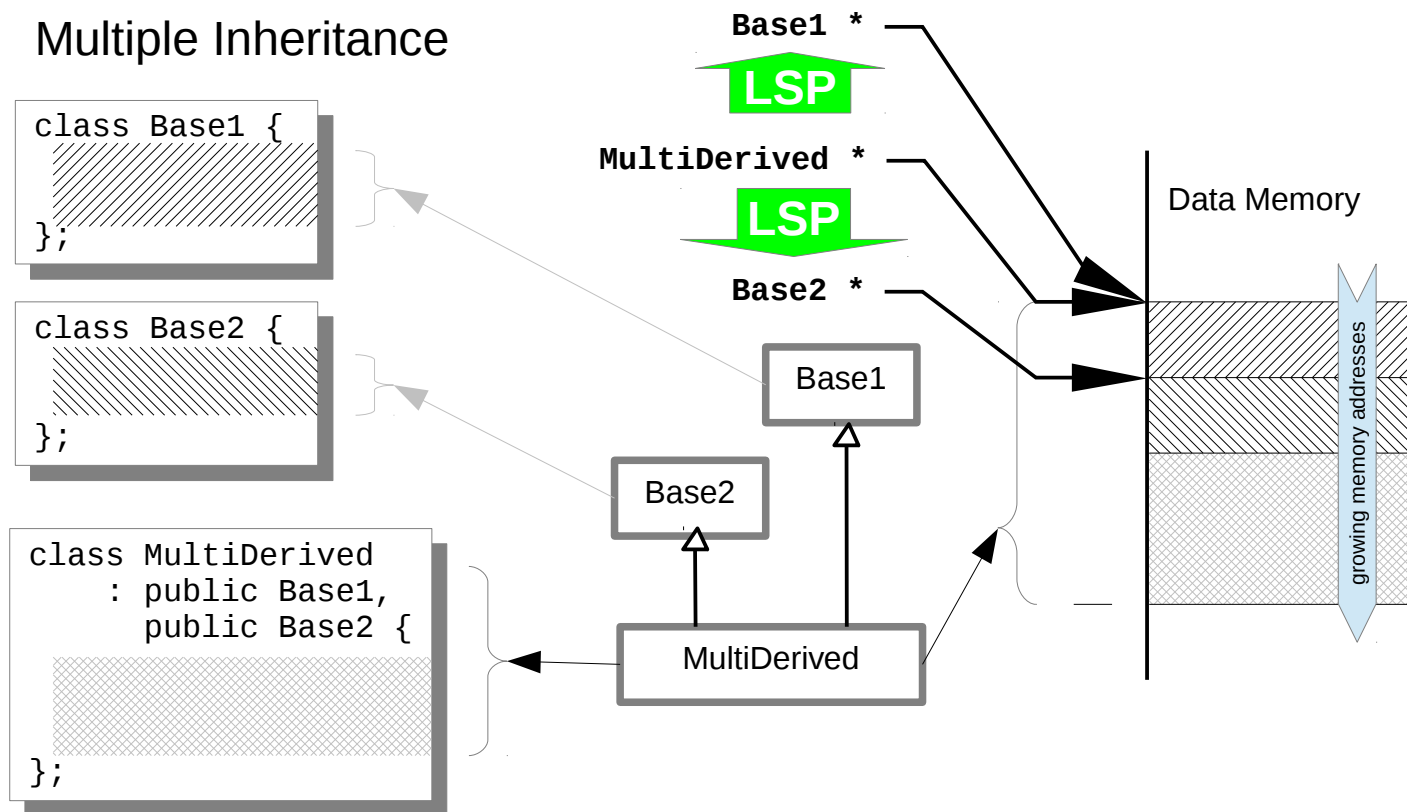
As described under the entry “Template Methode”-Pattern in the GoF-book. Standard technique in all OO programming languages that support polymorphism but not type-generic programming.



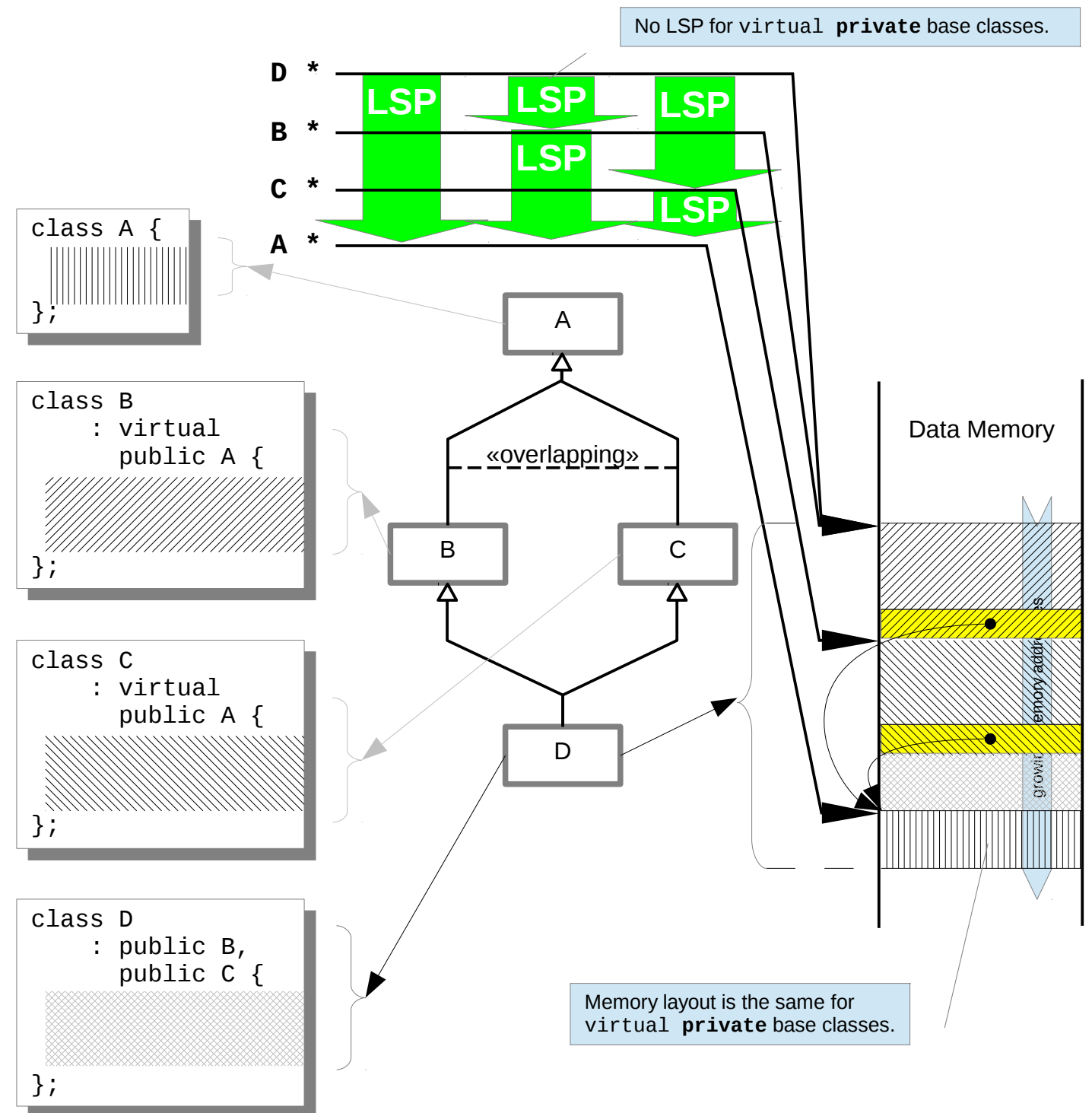
Example – “Open Close”-Principle

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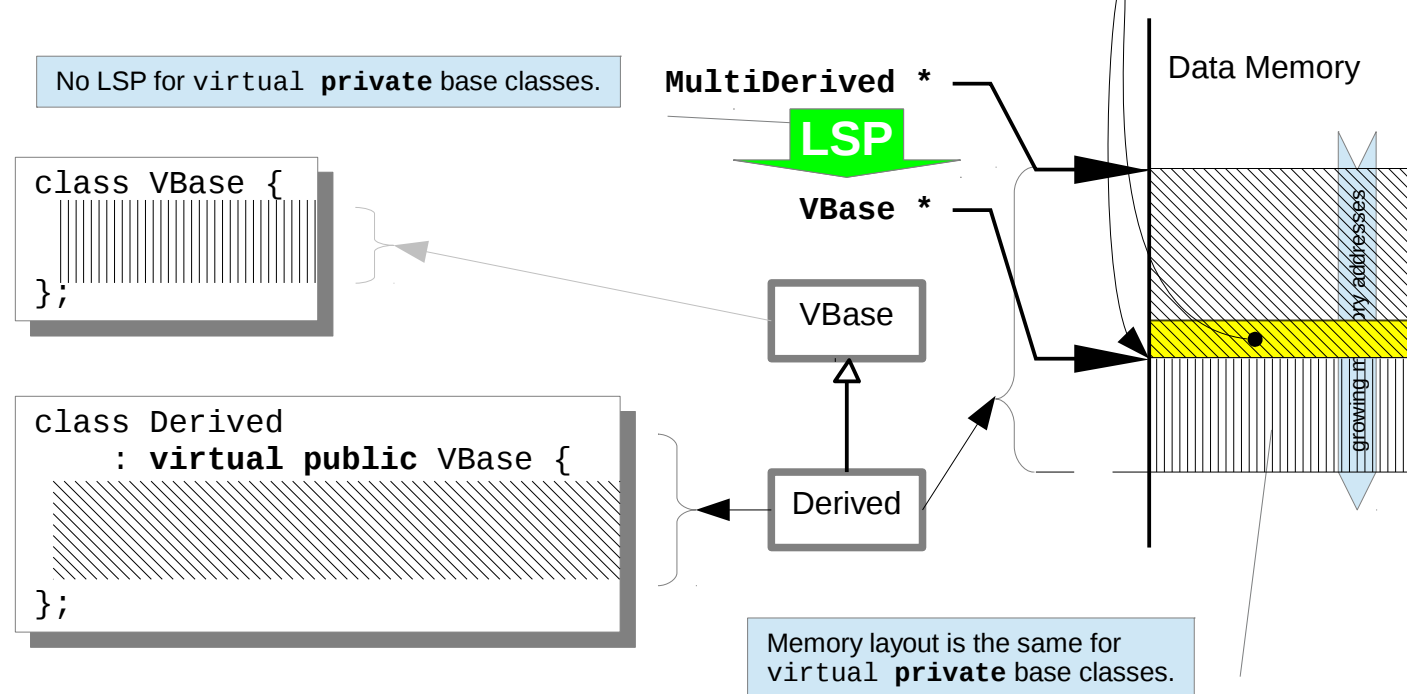
Multiple Inheritance



Overlapping Common Base Class



Virtual Base Class



A virtual base class introduces additional overhead in the derived class:

- space is allocated for an pointer which points to the base class part;
- all access to the base class part is indirect using this pointer.

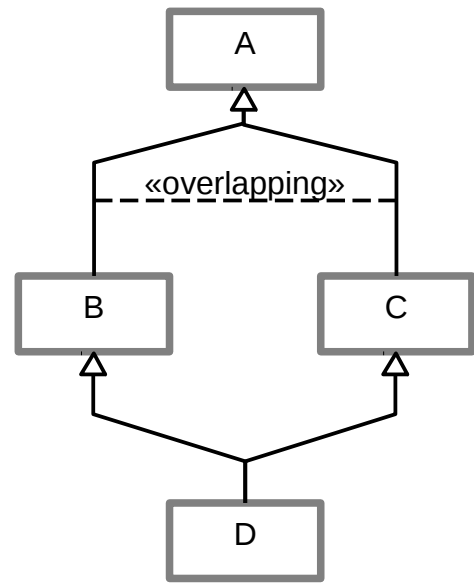
As far as is shown virtual base classes have no advantage.

Virtual base classes and resulting memory layout only shows one of several possible implementations

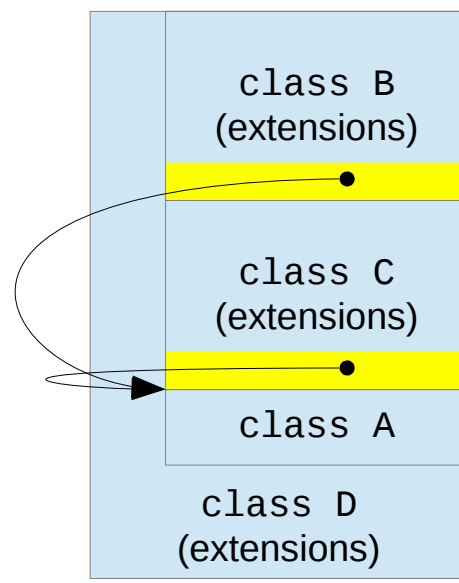
Virtual base classes are the mechanism to make a common base in a “diamond-shaped” inheritance relationship overlapping (see A above).

- This has to be prepared by the classes at the intermediate level (B and C above).
- The most derived class (D above) does not use virtual bases – it finds its direct bases at fixed offsets.
- These bases refer to their base via the embedded pointer (see left side).
- Both pointers are set to point to the same (embedded) base object.

Multiple Inheritance and Virtual Base Classes



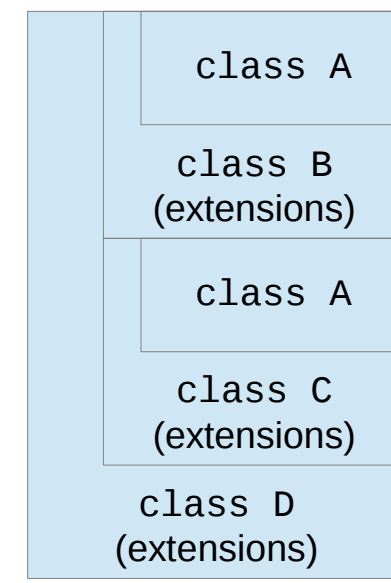
UML Class Graph



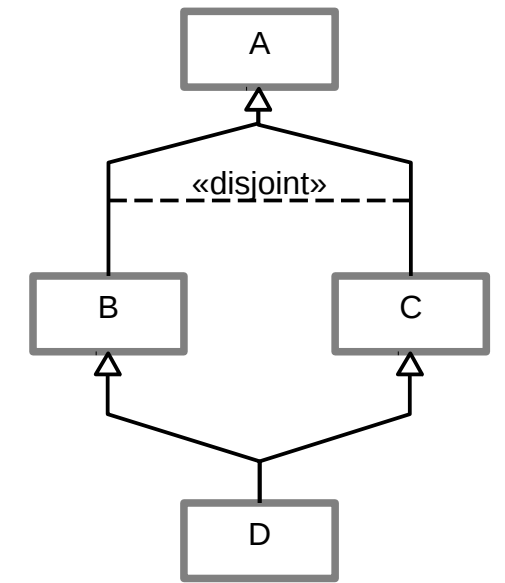
Member Data to Memory Mapping
(showing one of several possible solutions)

Automatic Type Conversions		
<i>to</i> ←	<i>from</i>	→ <i>to</i>
A	A	A
A	B	A
A	C	A
A, B, C	D	B, C

Up-Casts by LSP



Member Data to Memory Mapping
(showing the straight forward solution)



UML Class Graph

```
class A {
    ...
};
class B : virtual public A {
    ...
};
class C : virtual public A {
    ...
};
class D : public B, public C {
    ...
};
```

C++ Source

Creation and Destruction of D objects

C++ Source

```
class A {
    ...
};
class B : public A {
    ...
};
class C : public A {
    ...
};
class D : public B, public C {
    ...
};
```

```
A::A( ... ) { ... };
```

Virtual base
constructed from most derived class
(trying default construction if no explicit constructor)

```
B::B( ... )
    : A( ... )
{ ... };
```

```
C::C( ... )
    : A( ... )
{ ... };
```

```
D::D( ... )
    : A( ... ), B( ... ), C( ... )
{ ... };
```

Order of Constructor Calls

A::A(...)	MI-List, then Body
B::B(...)	(remaining) MI-List except A::A(...), then Body
C::C(...)	(remaining) MI-List except A::A(...), then Body
D::D(...)	MI-list, then Body

Order of Destructor Calls

D::~~D()	Body, chaining to
C::~~C()	Body, chaining to
B::~~B()	Body, chaining to
A::~~A()	

Special rule for calling virtual base class constructors:

- executed when a B or C object is created stand-alone;
- ignored when a B or C base of class of D is created.

Order of Constructor Calls

A::A(...)	base of B	MI-List, then Body
B::B(...)		(remaining) MI-List, then Body
A::A(...)	base of C	MI-List, then Body
C::C(...)		(remaining) MI-List, then Body
D::D(...)		(remaining) MI-List, then Body

Order of Destructor Calls

D::~~D()		Body, chaining to
C::~~C()		Body, chaining to
A::~~A()	base of C	Body, chaining to
B::~~B()		Body, chaining to
A::~~A()	base of B	Body

No special rule for calling (non-virtual) base class constructors:

- each class cares for its direct base(s);
- no knowledge wrt. indirect bases.**

```
A::A( ... ) { ... };
```

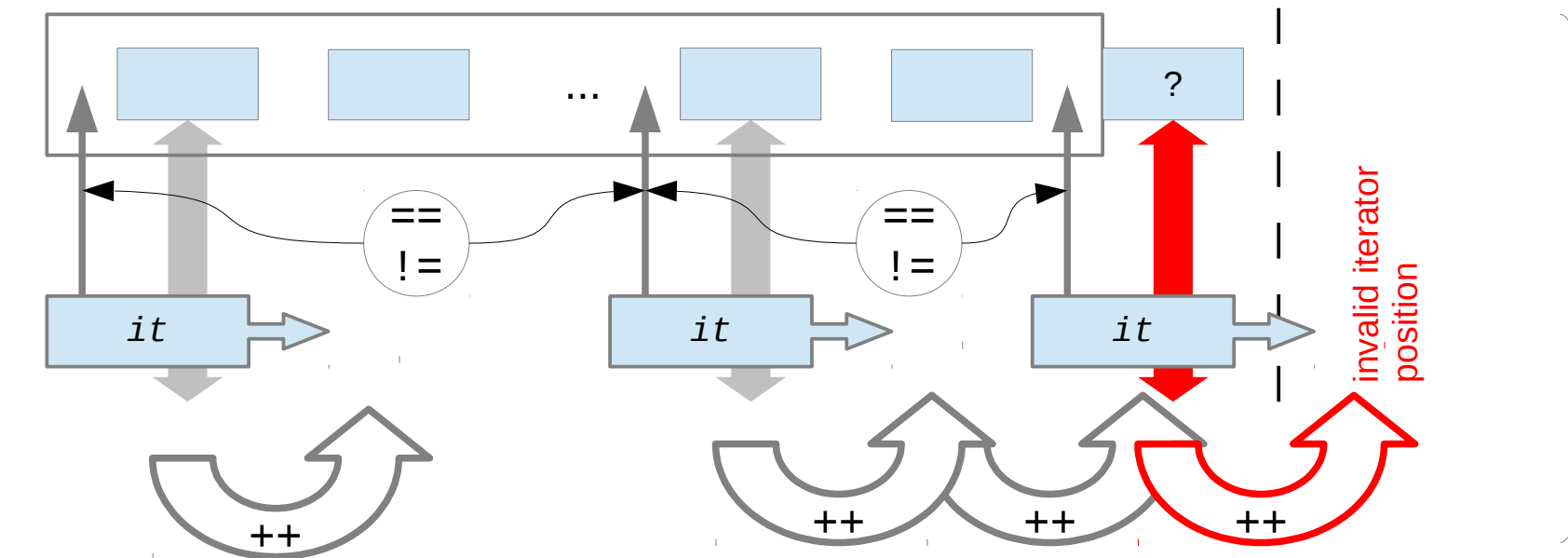
```
B::B( ... )
    : A( ... )
{ ... };
```

```
A::A( ... ) { ... };
```

```
C::C( ... )
    : A( ... )
{ ... };
```

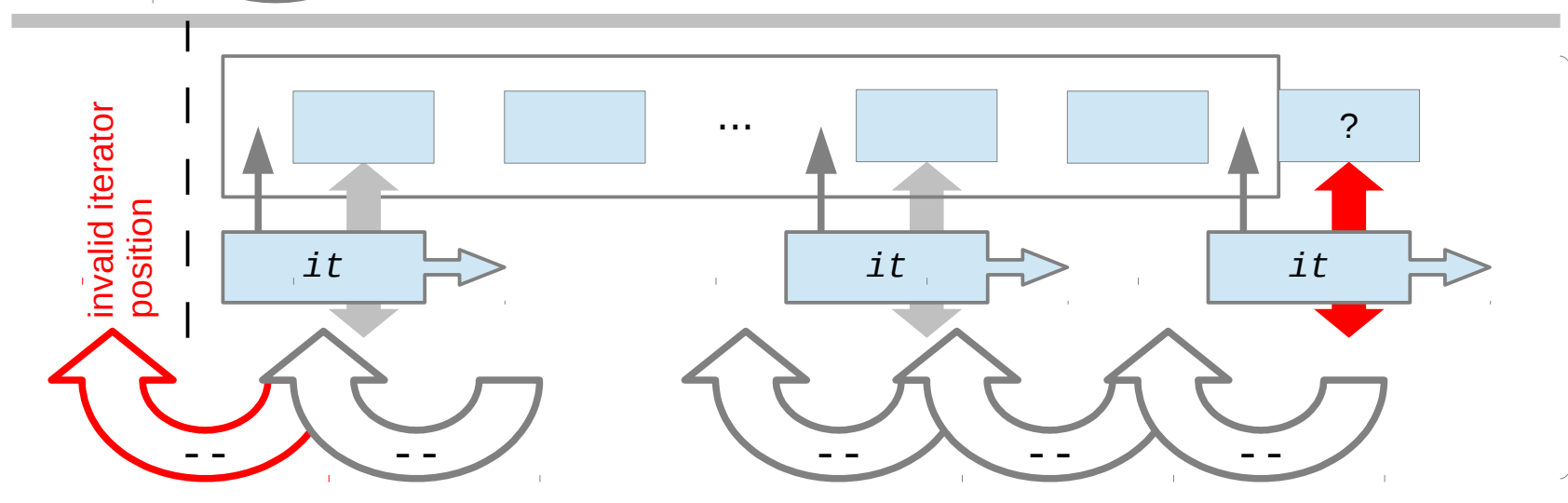
```
D::D( ... )
    : B( ... ), C( ... )
{ ... };
```

Diamond Shaped Inheritance



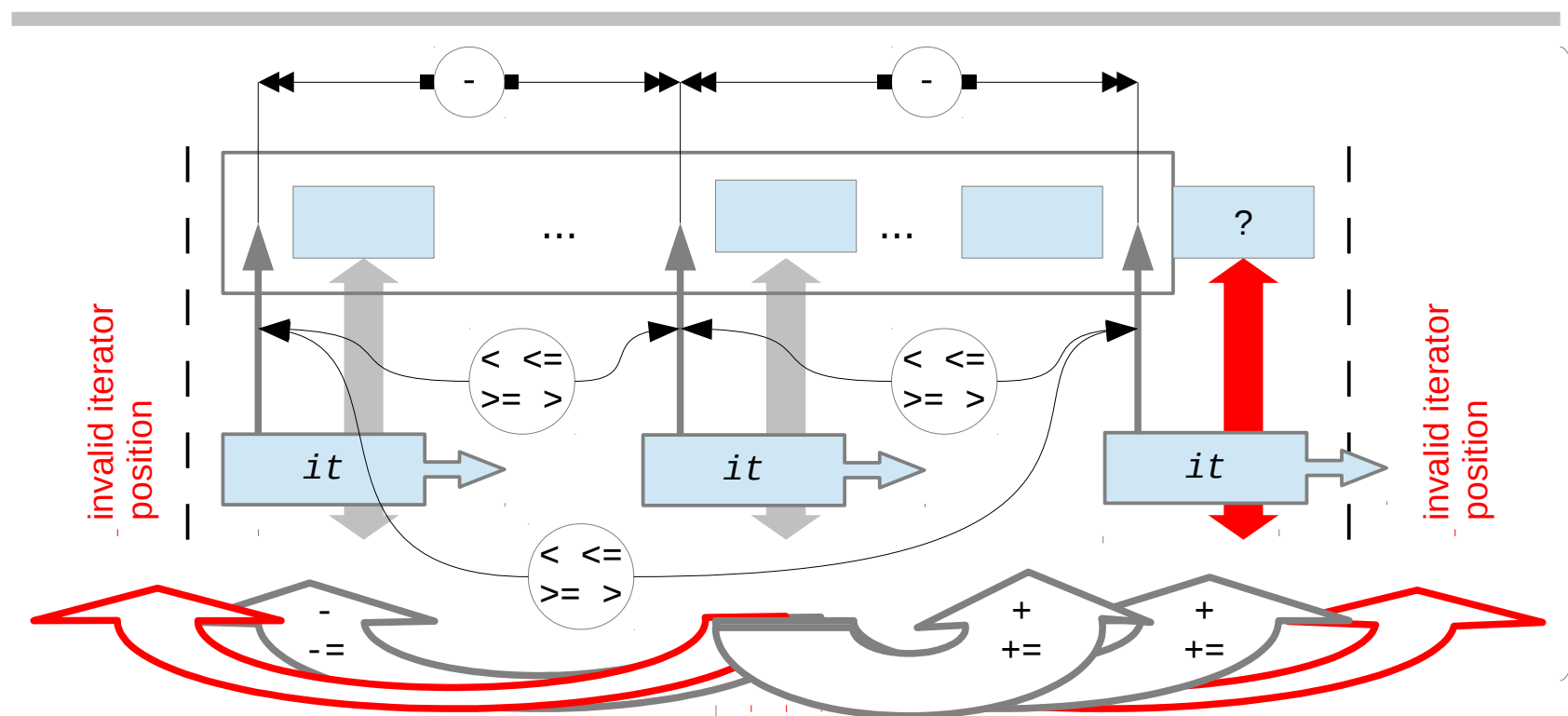
Operations of Unidirectional Iterators

	Effect	Remarks
<code>*it</code>	access referenced element	undefined at container end
<code>++it</code> <code>it++</code>	advance to next element (usual semantic for pre-/postfix version)	
<code>it == it</code>	compare for identical position	operands must denote existing element or end of same container
<code>it != it</code>	compare for different position	



Additional Operations of Bidirectional Iterators

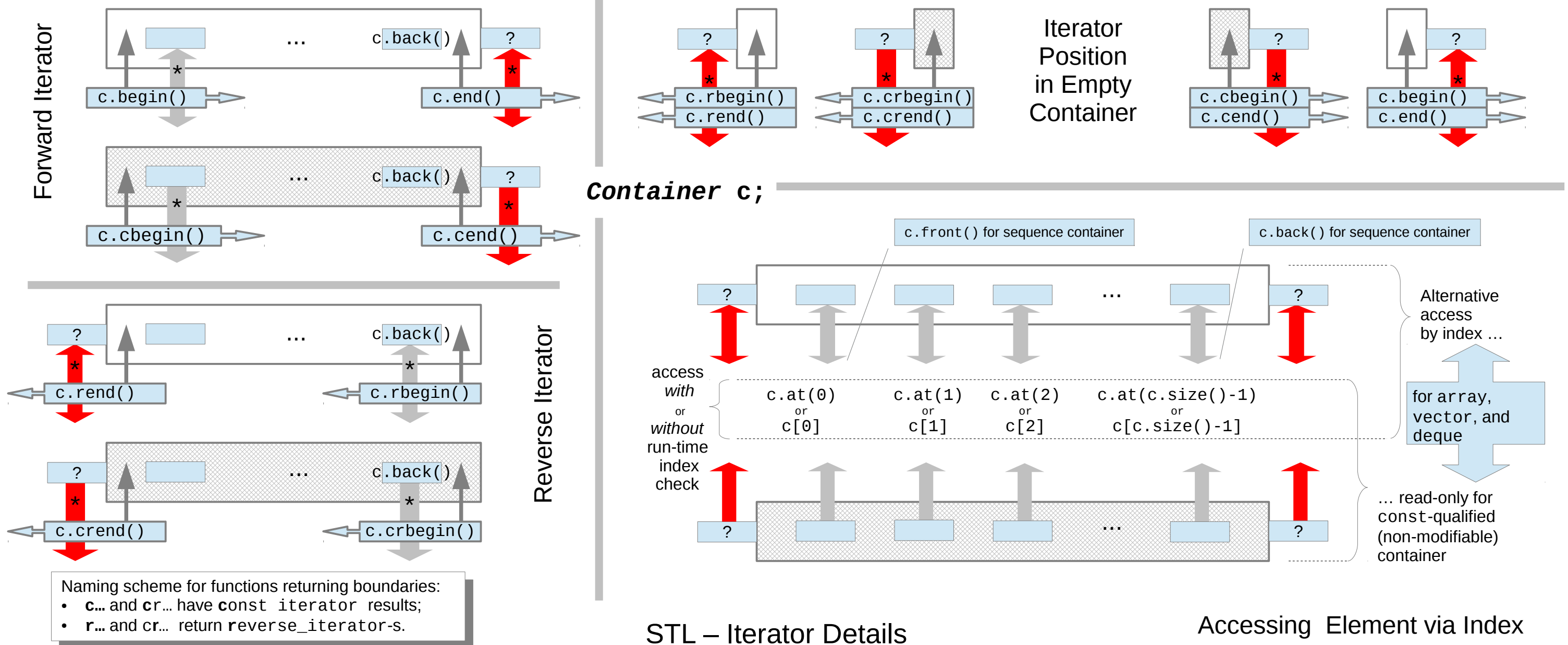
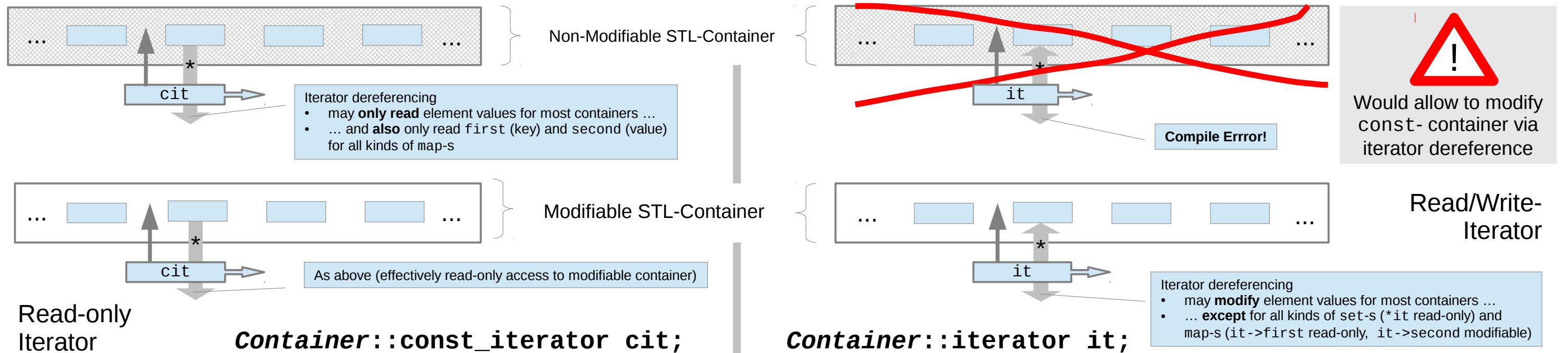
	Effect	Remarks
<code>--it</code> <code>it--</code>	advance to previous element (usual semantic for pre-/postfix version)	undefined at container begin



Additional Operations of Random Access Iterators

	Effect	Remarks
<code>it + n</code> <code>it += n</code>	<code>it</code> advanced to n -th next element (previous if $n < 0$)	resulting iterator position must be inside container (denote existing element or end)
<code>it - n</code> <code>it -= n</code>	<code>it</code> advanced to n -th previous element (next if $n < 0$)	
<code>it - it</code>	number of increments to reach rhs <code>it</code> from lhs <code>it</code>	operands must denote existing element or end of same container
<code>it < it</code>	true lhs <code>it</code> <u>before</u> rhs <code>it</code>	
<code>it <= it</code>	true if lhs <code>it</code> <u>not after</u> rhs <code>it</code>	
<code>it >= it</code>	true if lhs <code>it</code> <u>not before</u> rhs <code>it</code>	
<code>it > it</code>	true if lhs <code>it</code> <u>after</u> rhs <code>it</code>	

STL – Iterator Categories

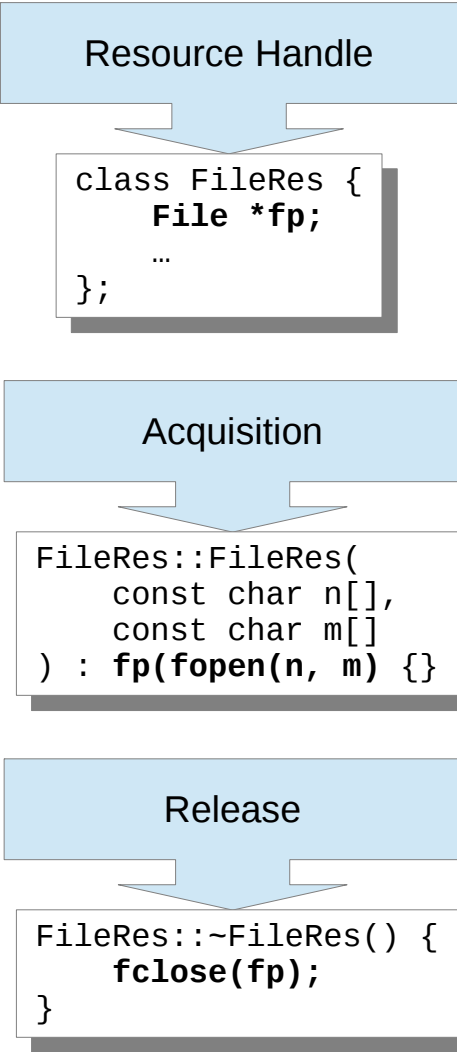


STL – Iterator Details

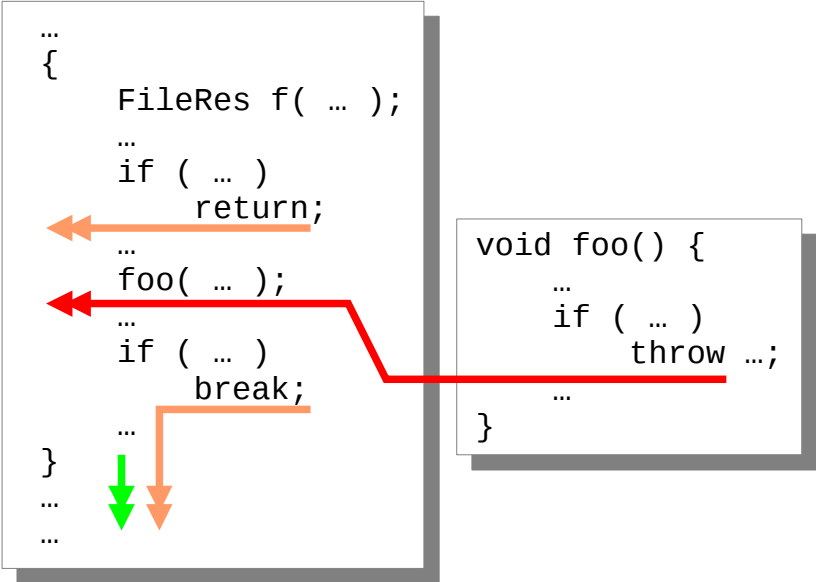
Classic Resource Management APIs

Turn into RAII

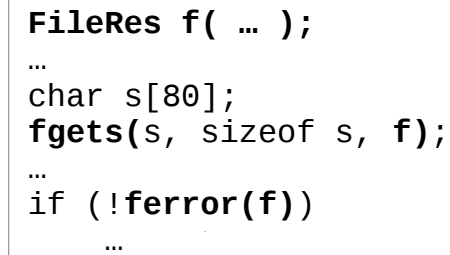
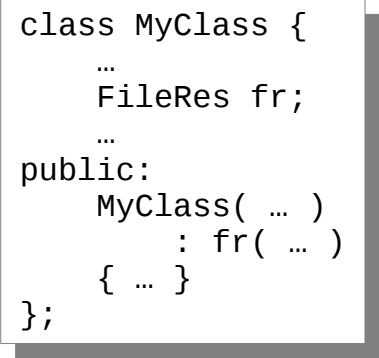
Principles	Examples					
	Unix/Linux		C	C Free Memory (Heap)		C++11
	Processes	Files	Files	C++ Free Memory (Heap)		
Operation to acquire returns ...	fork()	creat(), open()	fopen(), freopen()	malloc(), calloc(), realloc() new T new T[N]		std::mutex m; m.lock(), m.try_lock()
... some handle to identify resource ...	pid_t (some integer)	int	FILE * (pointer to some struct with opaque content)	generic pointer (void*) to otherwise unused storage for (at least) as many bytes as requested T* denoting a pointer to otherwise unused storage for (at least) one object of type T T* denoting a pointer to otherwise unused storage for (at least) N objects of type T at adjacent memory locations like in a builtin array		no special return value (instead state of object is changed)
... in subsequent operations like ...	kill(), ptrace(), ...	read(), write(), seek(), poll(), ...	fread(), fwrite(), fseek(), ftell(), fflush(), ...	after conversion to the target type all builtin pointer operations all builtin pointer operations		m.native_handle()
... until final release (eventually returning resource to a pool)	wait(), waitpid()	close()	fclose()	free() delete ... delete[] ...		m.unlock()
Standard Wrapper	none	none	none	std::unique_ptr<T>	std::unique_ptr<T[]>	std::lock_guard



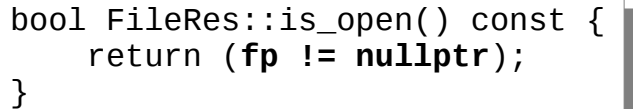
Acquire Resource for Code Segment



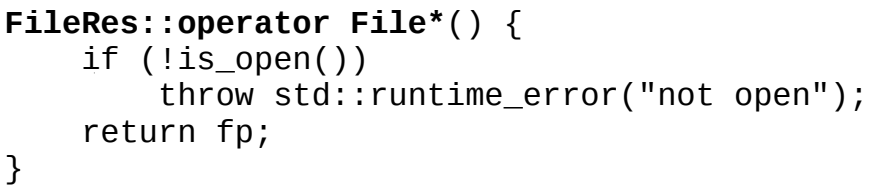
Acquire Resource for Object Lifetime



Optionally add Convenience Operations

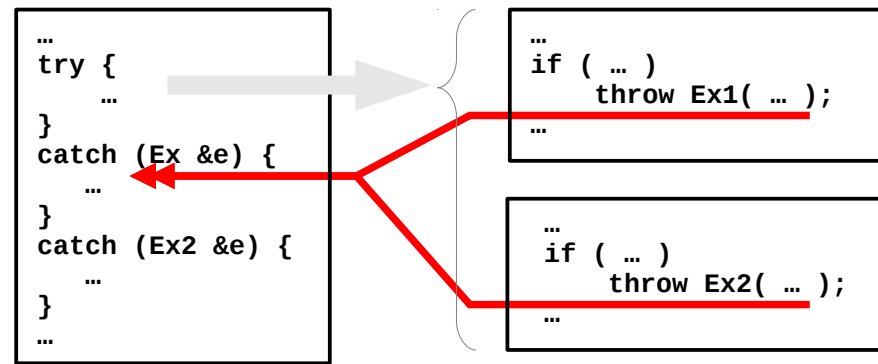


Easy and Secure Use via Automatic Conversion



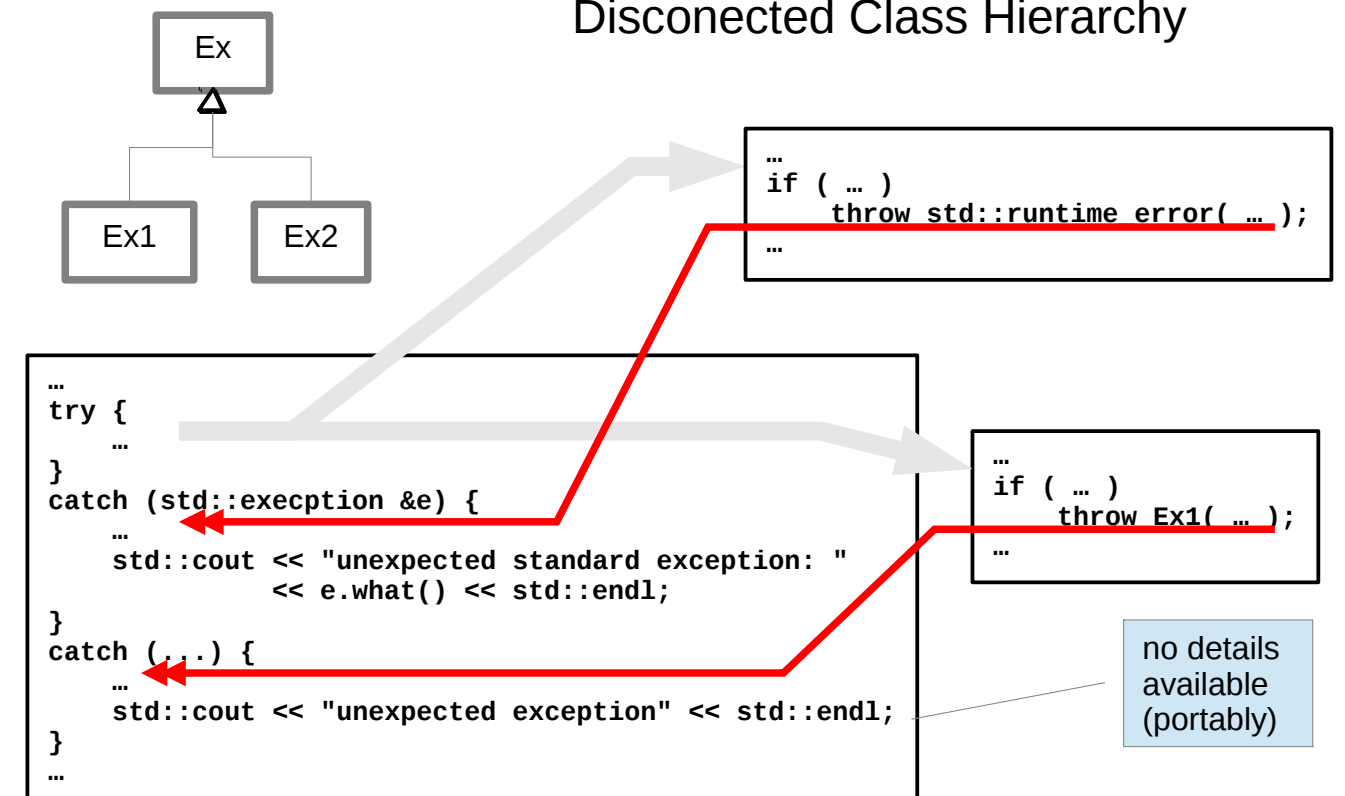
Wrapped Resource

Bad Order of Handlers

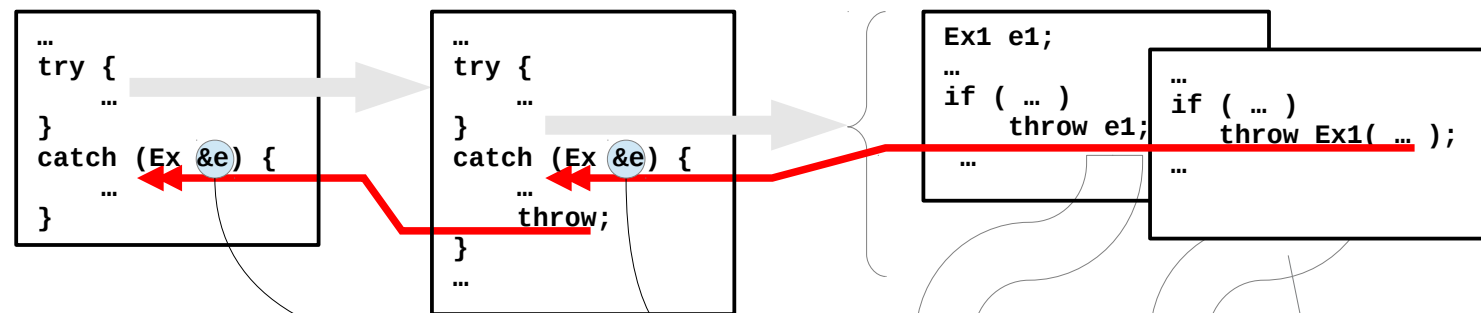


The compiler may issue a warning that the second catch-clause is shadowed by the first but this is not mandatory.

Disconnected Class Hierarchy



Optimal Re-throwing



by reference
Memory location guaranteed to exist until last catch-block accessing exception object

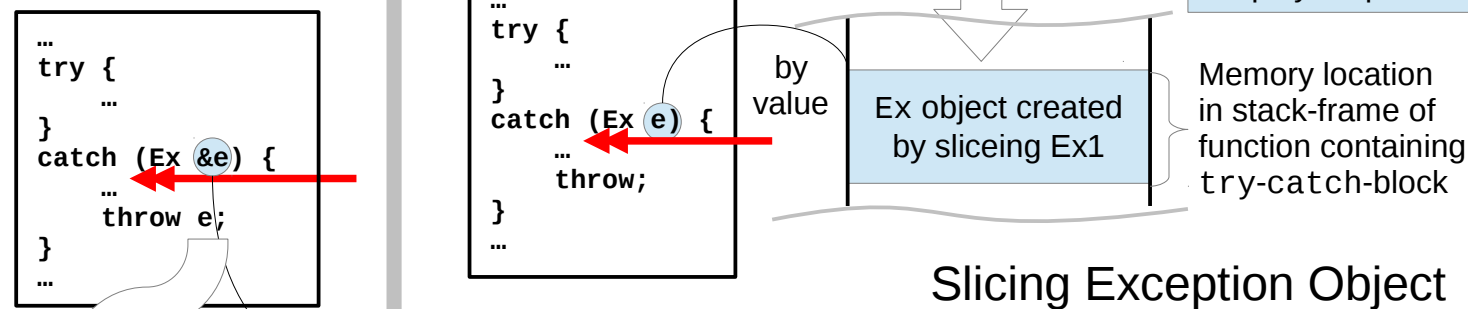
Ex1 object formally created with copy-constructor

may use RVO

physical copy, no polymorphism

Memory location in stack-frame of function containing try-catch-block

Slicing Exception Object



by reference

Ex object created as copy, thereby possibly sliced

Memory location from rethrowing, guaranteed to exist until last catch-block

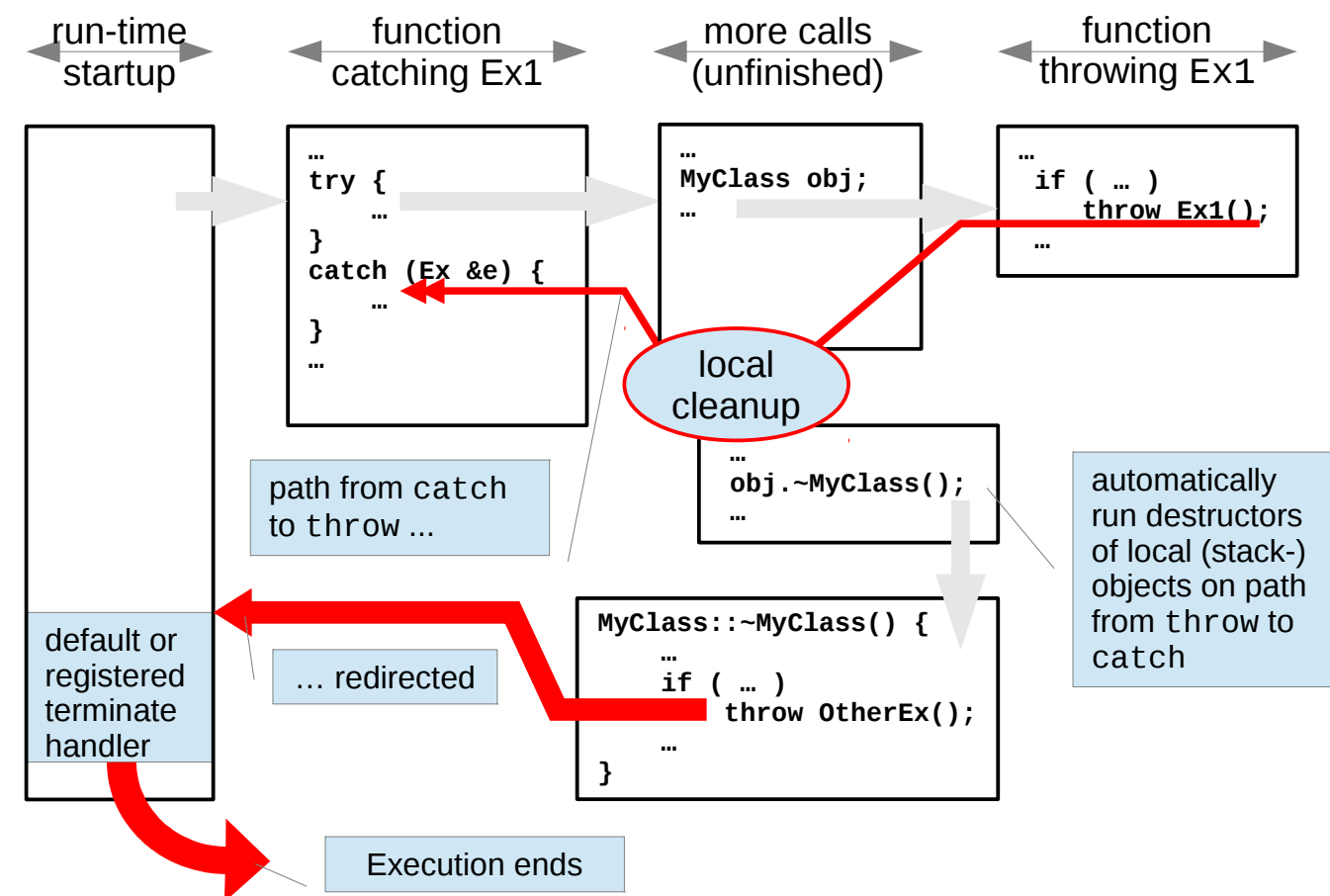
Object of derived class Ex1 or Ex2

Memory location from initial throw, guaranteed to exist until current catch-block ends

Sub-optimal Re-throwing

Exception Details

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path from catch to throw ...

local cleanup

automatically run destructors of local (stack-) objects on path from throw to catch

Execution ends

Throwing from Destructors

