

VE203 Assignment 6

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Exercise 6.1

i)

$$\begin{aligned}
 f(n) &= af\left(\frac{n}{b}\right) + cn^d \\
 &= a\left[af\left(\frac{n}{b^2}\right) + c\left(\frac{n}{b}\right)^d\right] + cn^d \\
 &= a^2f\left(\frac{n}{b^2}\right) + ac\left(\frac{n}{b}\right)^d + cn^d \\
 &= \dots \\
 &= a^{\log_b n} f(1) + \sum_{i=1}^{\log_b n} cn^d \left(\frac{a}{b^d}\right)^{i-1} \\
 &= n^{\log_b a} f(1) + cn^d \sum_{i=1}^{\log_b n} \\
 &= f(1)n^d + cn^d \log_b n
 \end{aligned}$$

ii)

$$\lim_{n \rightarrow \infty} \frac{f(1)n^d + cn^d \log_b n}{n^d \log n} = \lim_{n \rightarrow \infty} c \frac{\log_b n}{\log_b n / \log_b 10} = c \log_b 10 = C$$

iii)

$$\begin{aligned}
 f(n) &= a^{\log_b n} f(1) + \sum_{i=1}^{\log_b n} cn^d \left(\frac{a}{b^d}\right)^{i-1} \\
 &= n^{\log_b a} f(1) + cn^d \frac{1 - \left(\frac{a}{b^d}\right)^{\log_b n}}{1 - \frac{a}{b^d}} \\
 &= n^{\log_b a} f(1) + cn^d \frac{b^d - b^d \left(\frac{a}{b^d}\right)^{\log_b n}}{b^d - a} \\
 &= n^{\log_b a} f(1) + \frac{b^d c}{b^d - a} n^d + cn^d \frac{b^d \frac{a^{\log_b n}}{n^d}}{a - b^d} \\
 &= \frac{b^d c}{b^d - a} n^d + \left[f(1) + \frac{b^d c}{a - b^d}\right] n^{\log_b a}
 \end{aligned}$$

iv)

$$a < b^d \implies \log_b a < d$$

$$\lim_{n \rightarrow \infty} \frac{C_1 n^d + C_2 n^{\log_b a}}{n^d} = C_1 = C$$

v)

$$a > b^d \implies \log_b a > d$$

$$\lim_{n \rightarrow \infty} \frac{C_1 n^d + C_2 n^{\log_b a}}{n^{\log_b a}} = C_2 = C$$

Exercise 6.2

i)

$$f(n) = f(n/2) + 1$$

ii)

$$O(\log n)$$

Exercise 6.3

We can simply find that in a bit string, whenever a series of "0" is switched to "1", we get a "01" string, and whenever a series of "1" is switched to "0", we get a "10" string. Then we only need to consider the first and last bit of the string.

When the first and last bit is the same (both "0" or both "1"), the number of switching from "0" to "1" and "1" to "0" is the same, so "10" and "01" occurs the same time.

When the first bit is "0" and the last bit is "1", the number of switching from "0" to "1" is one more than that of "1" to "0", so "01" occurs one more time than "10".

Similarly, when the first bit is "1" and the last bit is "0", "10" occurs one more time than "01".

In conclusion, in a bit string, the string 01 occurs at most one more time than the string 10.

Exercise 6.4

i) When $q_{n-1} - 1 \neq 0$,

$$(p_n, q_n) = (p_{n-1} + 1, q_{n-1} - 1)$$

When $q_{n-1} - 1 = 0$,

$$(p_n, q_n) = (1, p_{n-1} + 1)$$

Forming them together, we get

$$(p_n, q_n) = ((p_{n-1} + 1) \bmod p_{n-1}, q_{n-1} - 1 + [1 - u(q_{n-1} - 1)](p_{n-1} + 1))$$

ii) For $p = 1$ and $q = 1$, $(p, q) = (p_1, q_1)$

For $p = 1$ and $q = k, k \in N^+$, suppose we can find $(p, q) = (p_m, q_m)$

Then for $p = 1$ and $q = k + 1$, we can also find $(p, q) = (p_{m+k}, q_{m+k})$ according to the recurrence relation $(p_n, q_n) = (1, p_{n-1} + 1)$

So for $p = 1$ and $q \in N^+$, we can find $(p, q) = (p_n, q_n)$

For $p = k, k < q, k \in N^+$ and $q \in N^+$, suppose we can find $(p, q) = (p_m, q_m)$ for all of q
Then for $p = k+1$, we can also find $(p, q) = (p_{m+k+q}, q_{m+k+q})$ according to the recurrence relation
 $(p_n, q_n) = (p_{n-1} + 1, q_{n-1} - 1)$ and $(p_n, q_n) = (1, p_{n-1} + 1)$
So for $p < q, p \in N^+$ and $q \in N^+$, we can find $(p, q) = (p_n, q_n)$

- iii) For $n = 1$ we can find $\phi(p_1/q_1) = 1$, it is true.
For $n = k$, suppose that $\phi(p_n/q_n) = n$
For $n = k + 1$, we should discuss it in two situations
Firstly, for $q_n = 1, p_{n+1} = 1, q_{n+1} = p_n + 1$

$$\begin{aligned}\phi\left(\frac{p_{n+1}}{q_{n+1}}\right) &= \frac{(p_{n+1} + q_{n+1} - 1)(p_{n+1} + q_{n+1} - 2)}{2} + p_{n+1} \\ &= \frac{(q_n + p_n)(q_n + p_n - 1)}{2} + 1 \\ &= \frac{(q_n + p_n - 2)(q_n + p_n - 1)}{2} + 1 + q_n + p_n - 1 \\ &= \frac{(p_n + q_n - 1)(p_n + q_n - 2)}{2} + p_n + 1 \\ &= n + 1\end{aligned}$$

Secondly, for $q_n > 1, p_{n+1} = p_n + 1, q_{n+1} = q_n - 1$

$$\begin{aligned}\phi\left(\frac{p_{n+1}}{q_{n+1}}\right) &= \frac{(p_{n+1} + q_{n+1} - 1)(p_{n+1} + q_{n+1} - 2)}{2} + p_{n+1} \\ &= \frac{(p_n + q_n - 1)(p_n + q_n - 2)}{2} + p_n + 1 \\ &= n + 1\end{aligned}$$

So it is proved.

- iv) Since ϕ is a function from a set p/q to a set n , it is surjective.
Since it can be reversed as is proved in part v), it is injective.
So it is bijective.
- v) When $k \in N^+$, suppose

$$\frac{k(k+1)}{2} < n \leq \frac{(k+1)(k+2)}{2}$$

where

$$\frac{(k+1)(k+2)}{2} - \frac{k(k+1)}{2} = k + 1$$

and

$$\begin{aligned}k^2 + k - 2n &< 0 \\ k &\in \left(\frac{-1 - \sqrt{1 + 8n}}{2}, \frac{-1 + \sqrt{1 + 8n}}{2} \right)\end{aligned}$$

The maximum integer k is what we need here, so

$$k = \left\lceil \frac{-1 + \sqrt{1 + 8n}}{2} \right\rceil$$

According to the recurrence we find

$$p = n - \frac{k(k+1)}{2} \quad q = k + 2 - p$$

So we find the inverse ϕ^{-1}

$$\phi^{-1}(n) = \frac{n - \frac{k(k+1)}{2}}{k + 2 - n + \frac{k(k+1)}{2}} \text{ where } k = \left\lceil \frac{-1 + \sqrt{1 + 8n}}{2} \right\rceil$$

vi)

$$\phi(p, q) = \frac{(p+q)(p+q-1)}{2} + p$$

vii)

$$\phi(p, q) = (|p| + |q|)(|p| + |q| - 1) + 2p - u(pq)$$

Exercise 6.5

Exercise 6.6

Since $\text{card } M = \text{card } N$, the number of elements in set M and N is the same.

Since $M \in N$, if there is an element which is in N but is not in M , then the number of elements in M is more than that in N , which leads to a contradiction.

So there isn't any element which is in N but is not in M , which means $M = N$

Exercise 6.7

If $f : M \rightarrow N$ is surjective, according to theorem 2.4.21, if $\text{card } M \neq \text{card } N$, f is injective.

If $f : M \rightarrow N$ is not surjective, then we can find a $f' : M \rightarrow N'$ which is surjective so that $\text{card } N' < \text{card } N$, so we can use the theorem again. If $\text{card } M \neq \text{card } N'$, f' is injective. So f is also injective.

So it is proved.