

# VE475 Homework 5

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## Ex. 1 — RSA setup

1. In the RSA encryption and decryption, we use

$$ed \equiv 1 \pmod{\varphi(n)}$$

$$m^{ed} \equiv m \pmod{\varphi(n)}$$

This is based on the Euler's theorem, which has a condition that  $m$  and  $n$  be two coprime integers. So it is likely for  $n$  to be coprime with  $m$ .

2. Suppose  $k = a\varphi(n)$ ,  $a \in N^*$ , and  $m < n$ .

(a)

$$\begin{aligned} m^k &\equiv (m^{\varphi(n)})^a \pmod{n} \\ &\equiv 1^a \pmod{n} \\ &\equiv 1 \pmod{n} \end{aligned}$$

So

$$m^k \equiv 1 \pmod{p} \quad \text{and} \quad m^k \equiv 1 \pmod{q}$$

- (b) First, if  $\gcd(m, n) = 1$ , according to (a), it's obvious that

$$m^{k+1} \equiv m \pmod{p} \quad \text{and} \quad m^{k+1} \equiv m \pmod{q}$$

Second, if  $\gcd(m, n) = p$ , so  $\gcd(m/p, q) = 1$

$$\begin{aligned} m^{k+1} &\equiv p \left[ \left( \frac{m}{p} \right)^{k+1} \pmod{q} \right] \pmod{n} \\ &\equiv p \left[ \left( \frac{m}{p} \right)^{a(p-1)\varphi(q)+1} \pmod{q} \right] \pmod{n} \\ &\equiv p \cdot \frac{m}{p} \pmod{n} \\ &\equiv m \pmod{n} \end{aligned}$$

So

$$m^{k+1} \equiv m \pmod{p} \quad \text{and} \quad m^{k+1} \equiv m \pmod{q}$$

Third, if  $\gcd(m, n) = q$ , it is similar to the second case.

We can conclude that for any arbitrary  $m$ ,  $m^{k+1} \equiv m \pmod{p}$  and  $\pmod{q}$ .

3. (a) We know that  $ed \equiv 1 \pmod{\varphi(n)}$ , which means that  $ed = k + 1$  where  $k$  is a multiple of  $\varphi(n)$ . According to part 2(b), we know that for any arbitrary  $m$ ,  $m^{k+1} \equiv m \pmod{p}$  and  $\pmod{q}$ , or we can say  $m^{k+1} \equiv m \pmod{n}$ , so  $m^{ed} \equiv m \pmod{n}$ ,
- (b) From the previous calculation, we can find that for all  $m < n$ , no matter  $m$  and  $n$  are coprime or not, we can both find that  $m^{ed} \equiv m \pmod{n}$ , so that the RSA encryption and decryption can be performed. So we can conclude that it is not necessary that  $\gcd(m, n) = 1$ .

## Ex. 2 — RSA decryption

$$n = 11413 = 101 \times 113$$

So we can find that  $p = 101$  and  $q = 113$ , and we should calculate  $d$  so that  $ed \equiv 1 \pmod{n}$ .

$$m \equiv c^d \pmod{n}$$

## Ex. 3 — Breaking RSA

## Ex. 4 — Programming

## Ex. 5 — Simple Questions

- 1.
- 2.
- 3.
- 4.
- 5.

$$(97 - 1) = 96 = 2^5 \times 3$$

So the generator  $x$  should satisfy that

$$x^{32} \not\equiv 1 \pmod{97} \quad \text{and} \quad x^{48} \not\equiv 1 \pmod{97}$$

$$x^{16} \not\equiv \pm 1, 35, 61 \pmod{97}$$

We can find that

$$2^{16} \equiv 61 \pmod{97}$$

$$3^{16} \equiv 61 \pmod{97}$$

$$4^{16} \equiv 1 \pmod{97}$$

$$5^{16} \equiv 36 \pmod{97}$$

So the smallest generator of  $U(\mathbb{Z}/97\mathbb{Z})$  is 5.