## VE475 Homework 4

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#### Ex. 1 — Euler's totient

#### 1. Suppose

$$\varphi(p^k) = p^{k-1}(p-1) = p^k - p^{k-1}$$

which means, there are  $p^{k-1}$  integers of  $n \in [1, p^k]$  so that

$$\gcd(n, p^k) > 1$$

What's more, if an integer and  $p^k$  is not coprime, it can be divided by p since all of prime factors of  $p^k$  are p.

When k = 1, we know  $\varphi(p) = p - 1$  since p is a prime.

When k = i, suppose  $\varphi(p^i) = p^i - p^{i-1}$ .

When k=i+1, we know that there are  $p^{i-1}$  integers in  $[1,p^i]$  which are not coprime with  $p^i$ , so they are also not coprime with  $p^{i+1}$ . Then consider the integers  $n \in [p^i+1,p^{i+1}]$  which are not coprime with  $p^{i+1}$ , we know that they all have a prime factor p, and  $n/p \in [p^{i-1}+1,p^i]$ , so there are  $(p-1)p^{i-1}$  integers that satisfy this condition. In total, there are  $p^{i-1}+(p-1)p^{i-1}=p^i$  integers which are not coprime with  $p^{i+1}$ , so  $\varphi(p^{i+1})=p^{i+1}-p^i$ .

According to the mathematical induction above, we can concluded that

$$\varphi(p^k) = p^{k-1}(p-1)$$

- 2. According to the Chinese Reminder Theorem, since m and n are coprime, there exists a ring isomorphism between Z/mnZ and  $Z/mZ \times Z/nZ$ , and here  $\varphi(mn)$  is the order of Z/mnZ,  $\varphi(m)$  is the order of Z/mZ and  $\varphi(n)$  is the order of Z/nZ. Suppose MN is the set of counted integers in  $\varphi(mn)$ , M is that in  $\varphi(M)$  and N is that in  $\varphi(N)$ , there is a bijection between MN and  $M \times N$ . So  $\varphi(mn) = \varphi(m)\varphi(n)$ .
- 3. Suppose

$$n = p_1^{k_1} p_2^{k_2} \cdots p_n^{k_n}$$

where  $p_1, p_2, \dots, p_n$  are primes and  $k_1, k_2, \dots, k_n \ge 1$ , it is obvious that  $p_1^{k_1}, p_2^{k_2}, \dots, p_n^{k_n}$  are pairwise coprime, so

$$\begin{split} \varphi(n) &= \varphi(p_1^{k_1}) \varphi(p_n^{k_n}) \cdots \varphi(p_n^{k_n}) \\ &= p_1^{k_1 - 1} (p_1 - 1) p_2^{k_2 - 1} (p_2 - 1) \cdots p_n^{k_n - 1} (p_n - 1) \\ &= p_1^{k_1} \left( 1 - \frac{1}{p_1} \right) p_2^{k_2} \left( 1 - \frac{1}{p_2} \right) \cdots p_n^{k_n} \left( 1 - \frac{1}{p_n} \right) \\ &= n \prod_{p \mid n} \left( 1 - \frac{1}{p} \right) \end{split}$$

4.

$$\varphi(1000) = 1000 \left(1 - \frac{1}{2}\right) \left(1 - \frac{1}{5}\right) = 400$$

According to Euler's Theorem, since 7 is coprime with 1000,

$$7^{400} \equiv 1 \bmod 1000$$

$$7^{803} \equiv 7^3 \mod 1000$$
  
 $\equiv 343 \mod 1000$ 

#### Ex. 2 - AES

1. 128 bits of 1 is used as the key for round 1.

2.

$$K(5) = K(4) \oplus K(1)$$

3. We know for a 4 bit number X,

$$X \oplus 1111 = \overline{X}$$

We also know

$$K(0) = K(1) = K(2) = K(3) = 1111$$

So it's easy to find

$$K(10) = K(9) \oplus K(6)$$

$$= [K(8) \oplus K(5)] \oplus [K(5) \oplus K(2)]$$

$$= K(8) \oplus K(2)$$

$$= \overline{K(8)}$$

$$K(11) = K(10) \oplus K(7)$$

$$= [K(9) \oplus K(6)] \oplus [K(6) \oplus K(3)]$$

$$= K(9) \oplus K(3)$$

$$= \overline{K(9)}$$

# Ex. 3 — Simple Questions

1. In ECB Mode, each block is encrypted separately with a function E and a key K, so the corruption of one encrypted block won't influence other blocks, only one block will be decrypted incorrectly.

In CBC Mode, from the second block, each block is encrypted and xor with the previous encrypted block. If one encrypted block (not the last block) is corrupted, the next block will also be influenced when applied xor with the wrong block, so two blocks will be decrypted incorrectly.

2. If IV is incremented by 1 each time, after a reset, the attacker would know the exact value of IV every time and can construct whatever plaintext he want and xor it with IV, and then input it into the block cipher. Then he can compare the ciphertext with the plaintext more efficiently. So it is not CPA secure.

3. Since p = 29 is a prime, according to Theorem 2.17, we can test the prime factors of p - 1 = 28, which are 2 and 7.

First, when q = 2,

$$2^{(p-1)/q} = 2^{28/2} = 2^{14} \equiv 28 \mod 29$$

Second, when q = 7,

$$2^{(p-1)/q} = 2^{28/7} = 2^4 \equiv 16 \mod 29$$

So

$$2^{(p-1)/d} \not\equiv 1 \bmod p$$

We can concluded that 2 is a generator of U(Z/29Z).

4. Since 1801 and 8191 are primes, it is a Legendre Symbol, and we can directly calculate  $1801^{4095} \mod 8191$  to solve it.

By applying modular exponentiation, we get the following table.

$$\begin{array}{c|cccc} i & d_i & \text{power mod } 8191 \\ \hline 11 & 1 & 1^2 \cdot 1801 \equiv 1801 \\ 10 & 1 & 1801^2 \cdot 1801 \equiv 2493 \\ 9 & 1 & 2493^2 \cdot 1801 \equiv 6873 \\ 8 & 1 & 6873^2 \cdot 1801 \equiv 7874 \\ 7 & 1 & 7874^2 \cdot 1801 \equiv 544 \\ 6 & 1 & 544^2 \cdot 1801 \equiv 557 \\ 5 & 1 & 557^2 \cdot 1801 \equiv 1193 \\ 4 & 1 & 1193^2 \cdot 1801 \equiv 4482 \\ 3 & 1 & 4482^2 \cdot 1801 \equiv 6085 \\ 2 & 1 & 6085^2 \cdot 1801 \equiv 6085 \\ 2 & 1 & 6085^2 \cdot 1801 \equiv 5027 \\ 1 & 1 & 5027^2 \cdot 1801 \equiv 4046 \\ 0 & 1 & 4046^2 \cdot 1801 \equiv 8190 \\ \hline & 1801^{4095} \equiv 8190 \bmod 8191 \\ \hline & \left(\frac{1801}{8191}\right) = -1 \\ \hline \end{array}$$

5. First, if  $\left(\frac{b^2-4ac}{p}\right)=0$ , then  $b^2-4ac=0$ , so the equation only have one solution  $x=-\frac{b}{2a}$ , and it can always mod p, thus the number of solutions satisfies  $1+\left(\frac{b^2-4ac}{p}\right)=1$ .

Second, if  $\binom{b^2-4ac}{p} \neq 0$ , then  $b^2-4ac \neq 0$ , the equation have two solutions  $x=-\frac{b\pm\sqrt{b^2-4ac}}{2a}$ , which means

$$-\frac{b \pm \sqrt{b^2 - 4ac}}{2a} \equiv x \mod p$$
$$\sqrt{b^2 - 4ac} \equiv \pm (2ax + b) \mod p$$

Then the problem becomes whether  $b^2 - 4ac$  is a square mod p.

If  $\left(\frac{b^2-4ac}{p}\right)=1$ ,  $b^2-4ac$  is a square mod p, and we can get 2 solutions mod p.

Otherwise,  $\left(\frac{b^2-4ac}{p}\right)=-1$ ,  $b^2-4ac$  is not a square mod p, and we can get no solution mod p.

In conclusion, the number of solutions mod p to the equation  $ax^2 + bx + c$  is

$$1 + \left(\frac{b^2 - 4ac}{p}\right)$$

6. According to Euler's theorem,

$$n^{p-1} \equiv 1 \mod p$$
  
 $n^{q-1} \equiv 1 \mod q$ 

Let 
$$(p-1) = k(q-1)$$
,

$$(n^{q-1})^p = n^{p-1} \equiv 1 \bmod q$$

Since gcd(n, pq) = 1, according to Chinese Reminder Theorem, we get

$$n^{p-1} \equiv 1 \mod pq$$

7. (a) Necessity: We know

$$\left(\frac{-3}{p}\right) = \left(\frac{-1}{p}\right)\left(\frac{3}{p}\right) = 1$$

First, if  $p \equiv 1 \mod 4$ ,  $\left(\frac{-1}{p}\right) = 1$ , then  $\left(\frac{3}{p}\right) = 1$ . And since  $p \not\equiv 3 \mod 4$ ,  $\left(\frac{3}{p}\right) = \left(\frac{p}{3}\right) = 1$ , so  $p \equiv 1 \mod 3$ , thus

$$p \equiv 1 \mod 12$$

Second, if  $p \equiv 3 \mod 4$ ,  $\left(\frac{-1}{p}\right) = -1$ , then  $\left(\frac{3}{p}\right) = -1$ . And since  $p \not\equiv 3 \mod 4$ ,  $\left(\frac{3}{p}\right) = -\left(\frac{p}{3}\right) = -1$ , so  $p \equiv 1 \mod 3$ , thus

$$p \equiv 7 \mod 12$$

In conclusion,  $p \equiv 1 \mod 6$ , and since p is odd prime, it can be deduced to  $p \equiv 1 \mod 3$ .

(b) Sufficiency: We know

$$p \equiv 1 \mod 3$$

So we can get

$$\left(\frac{p}{3}\right) = 1$$

First, if  $p \equiv 1 \mod 4$ ,

$$\left(\frac{-3}{p}\right) = \left(\frac{-1}{p}\right)\left(\frac{3}{p}\right) = 1 \cdot 1 = 1$$

Second, if  $p \equiv 3 \mod 4$ ,

$$\left(\frac{-3}{p}\right) = \left(\frac{-1}{p}\right) \cdot - \left(\frac{3}{p}\right) = -1 \cdot -1 = 1$$

So it is proved.

So  $\left(\frac{-3}{p}\right) = 1$  if and only if  $p \equiv 1 \mod 3$ .

8. Since  $\left(\frac{a}{p}\right) = 1$ ,

$$a^{(p-1)/2} \equiv 1 \mod p$$

And we know 2 is a prime factor of p-1, so it doesn't satisfy the condition that for all prime factors q of p-1,  $a^{(p-1)/q} \not\equiv 1 \mod p$ . Then a is not a generator of  $F_p^*$ .

## Ex. 4 — Prime vs. irreducible

- 1. In an integral domain, suppose a prime element p is reducible, so it can be implied as p=ab, where a,b are non-zero, non-invertible and not equals to the ring element 1. Let  $x=k_1a$ ,  $y=k_2b$ ,  $k_1,k_2\neq 0$ , so  $x,y\neq 0$ . If  $b\nmid k_1$  and  $a\nmid k_2$ , then  $ab\nmid k_1a$  and  $ab\nmid k_2b$ , which means  $p\nmid x$  and  $p\nmid y$ , which makes a contradiction with (\*), so we can conclude that any prime element in an integral domain is reducible number is not prime. And since any number p>1 is either a prime or
- 2. In Z, suppose an irreducible number p > 1 is not prime, but it can't be implied as p = ab, where a, b > 1. Suppose  $a \mid p$ , we can get a = 1 or a = p, which makes a contradiction with (\*\*). So we can conclude that any irreducible integer is prime in Z.
- 3. From (\*\*), we can simply get any prime number is irreducible. So if p is a prime and  $p \mid x \cdot y$ ,  $x, y \in Z$ , suppose  $p \nmid x$  and  $p \nmid y$ , we can get  $p \nmid x \cdot y$ , which makes a contradiction. So we (\*) is proved, and (\*\*) implies (\*).
- 4. From (\*), we know any prime integer is irreducible, if p is a prime and a|p, suppose  $a \neq 1$  and  $a \neq p$ , then 1 < a < p, but a|p and p is not reducible, which makes a contradiction, so (\*) implies (\*\*). And according to part 3, (\*\*) implies (\*), we can conclude that (\*) and (\*\*) are equivalent for integers.

### Ex. 5 — Primitive root mod 65537

1. Since 65537 is a prime, we can calculate  $3^{32768} \mod 65537$  and we can find that  $3^{32768} \equiv -1 \mod 65537$  (The calculation is shown in part 2), so

$$\left(\frac{3}{65537}\right) = -1$$

2. By applying modular exponentiation, we get the following table.

i	$d_i$	power mod $65537$	
15	1	$1^2 \cdot 3 \equiv 3$ $3^2 \equiv 9$ $9^2 \equiv 81$	
14	0		
13	0		
12	0	$81^2 \equiv 6561$	
11	0	$6561^{2} \equiv 54449$ $54449^{2} \equiv 61869$ $61869^{2} \equiv 19139$ $19139^{2} \equiv 15028$ $15028^{2} \equiv 282$ $282^{2} \equiv 13987$ $13987^{2} \equiv 8224$ $8224^{2} \equiv 65529$	
10	0		
9	0		
8	0		
7	0		
6	0		
5	0		
4	0		
3	0	$65529^2 \equiv 64$	
2	0	$64^2 \equiv 4096$	
1	0	$4096^2 \equiv 65281$	
0	0	$65281^2 \equiv 65536$	

So

$$3^{32768} \equiv 65536 \mod 65537$$
$$3^{32768} \not\equiv -1 \mod 65537$$

3.	According to Theorem 2.17, we can conclude that 3 is a primitive root the only prime factor of 65536 and $3^{(65537-1)/2} \not\equiv 1 \mod 65537$ .	ot mod 65537	because 2 is