

# Reductions, NP, and NP-Completeness

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CS170 Reviews Session

# Time Complexity Classes

Decision problems: problems that only requires a yes/no for an answer.

- ▶ Does  $G$  has a path from  $s$  to  $t$  with length less than  $k$ ? (P)
- ▶ Is the boolean formula  $\phi$  satisfiable? (NP-Complete)

P: The set of **decision problems** that can be decided in polynomial time.

NP: The set of **decision problems** with candidate solutions (formally known as certificate/witness/proof) that can be verified in polynomial time.

We only care about decision problems because they're easier to deal with and they're as hard as search and optimization problems.

## Quick Digression: Problems (probably) not in NP

UNSAT: boolean formula  $\phi$  is unsatisfiable.

Can't come up with a polynomial size certificate for this problem.

# NP-Hard and NP-Complete

NP-Hard: decision problems that are *at least as hard* as every problem in NP.

NP-Completeness: intersection of NP and NP-Hard.

How to prove the notion of “at least as hard”? Reductions!

# Reductions

$$A \leq_P B$$

Problem A can be reduced to Problem B in polynomial time if:

- ▶ there exists a  $f$  such that an input  $w$  of problem A is an accepting instance of A if and only if  $f(w)$  is an accepting instance of B.
- ▶  $f$  can be computed in polynomial time.

If  $A \leq_P B$ , then A is an “easier” problem than B.

- ▶ Intuitively: if I want to lift a feather, I can to use my method to lift a stone to lift this feather. Therefore, the task of lifting feathers is easier than lifting stones.
- ▶ If I solved B, then I solved A, so A is the easier problem.

# Prove a problem is NP-Complete

Prove A is in NP.

- ▶ Design the poly time verifier.

Prove A is NP-Hard.

- ▶ Give a poly time reduction from a NP-Complete problem.

## Quick Digression

World's first NP-Complete problem: boolean satisfiability problem (Cook-Levin Theorem).

Richard Karp gave 21 NP-Complete problems using tree of reductions from SAT a year after Stephen Cook's paper.

# Prove this is NP-Complete

Does a boolean formula has at least two satisfying assignments?



# Reduction

Given that the directed graph Rudrata (Hamiltonian) Path problem is NP-Complete, prove that the undirected graph Rudrata Path problem is also NP-Complete.

# Reduction

Reduce from 3-SAT to Clique directly to prove it's NP-Complete.