# Hihocoder 题目泛做解题报告

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# 目录

1	第一	部分	1	
	1.1	1001	1	
	1.2	1014	1	
	1.3	1015	3	
	1.4	1032	4	
2	Tree			
	2.1	Divide And Conquer Tree	6	
	2.2	Link Tree	9	
	2.3	Segment Tree	12	
	2.4	Splay Tree	17	
	2.5	Treap	20	
3	Geometry 22			
	3.1	Basic Struct and Algorithm	22	
	3.2	Polygon Area	23	
	3.3	Half Plane Intersection	23	
4	Math 24			
	4.1	China Remainder Theory	24	
	4.2	Decompose	24	
	4.3	Euler Phi	26	
	4.4	Extend GCD	26	
	4.5	Integer Inverse	27	
	4.6	Line Mod	27	
	4.7	Log Mod	27	
	4.8	Lucas	28	
	4.9	Miller Rabin	29	
	4.10	Mul Mod	29	
	4.11	Pollard Rho	30	
	4.12	Pow Mod	30	
	4.13	Power Mod	30	
	4.14	Primitive Root	31	
	4.15	Square Mod	32	
5	Others 32			
	5.1	Exact Cover	32	
	5.2	Matrix Fast Power	35	
	5.3	Polynomial	36	

# 1 第一部分

#### 1.1 1001

Default , Hello World :)

```
#include<cstdio>
using namespace std;

int main(){
    int a, b;
    //this is a very very typical problem
    //I have ever sloved, so I am proud of myself
    while(~scanf("%d%d", &a, &b)){
        printf("%d\n", a+b);
    }
    return 0;
}
```

#### 1.2 1014

- 经典字典树 (trie)
- Trie 是一种基于字符串前缀保存的数据结构,将每个字符串以一棵树的形式保存,而相同的前缀 会在同一颗树的路径上.
- 关于 Trie 常见的操作有: Build, Update, Query. Update | Query 的时候由于仅仅按照字符串的 长度进行逐位操作,所以复杂度为 O(d), build 的复杂度为 O(n\*d).
- Trie 的扩展有基于 Trie 的贪心 Xor(经典问题如:查找 N 个数最大的两个异或值,查找区间最大的异或和等).

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
char s[20];

typedef struct Trie{
   int val, cnt;
   Trie *next[26];
   Trie(){
      for(int i = 0; i < 26; ++i){
        next[i] = NULL;
      }
}</pre>
```

```
}Trie;
Trie* buildTrieTree(Trie *root, char *dict){
   Trie *p = root, *q;
   int i = 0;
   while(dict[i] != '0'){
      int k = dict[i] - 'a';
      if(p->next[k] == NULL){
          q = new Trie();
         for(int j = 0; j < 26; ++j){
     //
            q->next[j] = NULL;
     // }
          q->val = k;
          q->cnt = 0;
          p->next[k] = q;
      p = p->next[k];
      p->cnt += 1;
      ++i;
   }
   return root;
}
int query(Trie *root, char *dict){
   int p = 0;
   while(dict[p] != '0'){
      int k = dict[p] - 'a';
      if(root->next[k] != NULL){
          root = root->next[k];
      else{
         return 0;
      ++p;
   }
   return root->cnt;
}
int main(){
   int n, m;
   scanf("%d", &n);
   Trie *root = new Trie();
 // for(int i = 0; i < 26; ++i){
 // root->next[i] = NULL;
 // }
```

```
for(int i = 0; i < n; ++i){
    scanf("%s", s);
    root = buildTrieTree(root, s);
}

scanf("%d", &m);
for(int i = 0; i < m; ++i){
    scanf("%s", s);
    int ans = query(root, s);
    printf("%d\n", ans);
}
return 0;
}</pre>
```

#### 1.3 1015

- 经典 KMP
- KMP 太经典了, 关于匹配串与模式串的匹配问题, 利用的是最长前缀与后缀相等来减少不必要的 匹配, 将匹配的复杂度从 O(n\*m) 降到 O(n+m).
- KMP 的 next 数组可以做很多扩展,最常见的是求循环节。注意理解 next 数组的本质,本质上 next 是多次迭代的,而每次迭代始终保持前缀与后缀相等,这个性质同样重要。

```
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int N = 1e4 + 1;
char a[N], b[N*100];
int p[N];
void GetNextArray(int n, char *b){
   int i, j;
   p[0] = -1;
   j = -1;
   for(i = 1; i < n; ++i){</pre>
       while(j >= 0 && b[j+1] != b[i]){
          j = p[j];
       if(b[j+1] == b[i]){
          ++j;
       }
       p[i] = j;
   }
```

```
}
void KMP(int n1, int n2, char *a, char *b){
   int j = -1, ans = 0;
   for(int i = 0; i < n1; ++i){</pre>
       while(j >= 0 && b[j+1] != a[i]){
          j = p[j];
       if(b[j+1] == a[i]){
          ++j;
       if(j == n2-1){
          ++ans;
          j = p[j];
       }
   printf("%d\n", ans);
}
int main(){
   int n;
   scanf("%d", &n);
   for(int i = 0; i < n; ++i){
       scanf("%s%s", a, b);
       int n1 = strlen(a), n2 = strlen(b);
       GetNextArray(n1, a);
       KMP(n2, n1, b, a);
   }
   return 0;
}
```

#### 1.4 1032

- 经典 Manacher-最长回文串
- 字符串最长回文串最简单的就是  $O(n^2)$  的暴力匹配 (枚举中心点), 当然也可以求正向串与逆向串的最长公共子序列, 也是  $O(n^2)$ . 也可以通过枚举 i 点的前缀与后缀,基于后缀数组与高度数组,求后缀与前缀的最长公共前缀 (LCP) 来解决,但是复杂度与效率仍然不及 Manacher.
- Manacher 的算法细节不表,说一下大概思路:在遍历字符串的过程中使用记录两个变量值 *mx*, *idx* 分别表示迭代到当前位置,最远的回文串能达到的位置以及其对应回文中心的下标。考虑对称性, 迭代到 i 的时候,考虑 i 关于 idx 的对称点 j,这样通过判断可以将 j 的回文长度加入到 i 的回文串中,这部分计算之前已经产生,所以节省了大量重复计算。最终比较一下所有位置的回文串长度即可。

• 算法的复杂度为O(n)。有很多细节,比如预处理,位置比较等等,可以参考:\_https://www.felix021.com/blog/read.

```
#include <cstdio>
#include <iostream>
#include <cstring>
#include <algorithm>
using namespace std;
#define minab(a, b) ((a) < (b) ? (a) : (b))
const int N = 1000010;
int p[N<<1];</pre>
char s[N<<1];</pre>
string st;
int mx, idx;//the Rightest postion can be reached now, and the idx
void manacher(char *str) {
   mx = 0, idx = 0;
   int ans = 0;
   memset(p, 0, sizeof(p));
   int size = (int)strlen(str);
   for(int i = 1; str[i]; ++i) {
       p[i] = mx > i ? minab(p[idx*2 - i], mx - i) : 1;
       \label{eq:while(p[i] + i < size) && (i - p[i] >= 0) && str[p[i]+i] == str[i-p[i]]) } \\ 
          ++p[i];
       }
       if(i + p[i] > mx) {
          mx = i + p[i];
          idx = i;
   // printf("p[i] = %d\n", p[i]);
       if(str[i] == '#') {
          if(((p[i]) / 2) * 2 > ans) {
              ans = (p[i]) / 2 * 2;
          }
       }
       else {
           if((p[i] - 1)/2 * 2 + 1 > ans) {
              ans = (p[i] - 1) / 2 * 2 + 1;
          }
       }
   printf("%d\n", ans);
}
```

```
int main() {
   int n;
   scanf("%d", &n);
   for(int i = 0; i < n; ++i) {
      cin >> st;
      for(int j = 0; j < st.size(); ++j) {
        s[j<<1|1] = '#';
        s[(j<<1) + 2] = st[j];
      }
      s[st.size() << 1 | 1] = '#';
      s[(st.size() << 1) + 2] = '0';
      // printf("%s\n", s+1);
      manacher(s+1);
   }
   return 0;
}</pre>
```

#### 1.5 1033

- 数位 DP
- 数位 DP 是一类按位进行 DP 统计问题。按位统计之后,发现很多重复子问题,于是就可以愉快的 DP 了.
- 数位 DP 的写法比较固定,一般都是用记忆化搜索的写法(参考我的代码. 注意需要考虑前导 0, 以及每次枚举的数字的范围(能不能到 9)
- 对于

```
// Author:tcOops
// Level -> CF/TC -> Yellow
// > -> Ag
// -> F/L/A/G
// -> Latency \[ \text{ 2017/5/15} \]

#include <cstdio>
#include <cstring>
#include <algorithm>
using namespace std;
const int N = 21;
const int MOD = 1e9 + 7;
const int OFFSET = 210;

int dig[N];
long long base[N];
```

```
struct google {
   long long n, sum;
   google() {
      n = -1, sum = 0;
   }
};
google dp[N][420];
google solve(int len, long long sum, bool isPrefix0, bool limit) {
   //printf("%d %d %lld\n", len, digital, sum);
   google ans, res;
   ans.n = 0, ans.sum = 0;
   if(len == 0) {
       if(sum == 0) {
          ans.n = 1;
      return ans;
   }
   if(!limit && isPrefix0 && dp[len][sum+OFFSET].n != -1) {
      return dp[len][sum+OFFSET];
   }
   int up = limit ? dig[len] : 9;
   for(int i = 0; i <= up; ++i) {</pre>
       if(!isPrefix0) {
          if(i == 0) {
              res = solve(len-1, sum, false, limit && (i == up));
          }
          else {
             res = solve(len-1, i-sum, true, limit && (i == up));
       }
       else {
          long long new_sum = i - sum;
          res = solve(len-1, new_sum, true, limit && (i == up));
       }
       ans.n += res.n;
       ans.sum = ((ans.sum + res.sum) % MOD + (res.n*i) % MOD * base[len-1] % MOD + MOD) % MOD;
   }
```

```
//printf("%lld %lld\n", ans.n, ans.sum);
   if(!limit && isPrefix0) {
       dp[len][sum+OFFSET] = ans;
   }
   return ans;
}
long long go(long long n, long long k) {
   if(n <= 0) {
       return 0;
   }
   int len = 0;
   while(n) {
       dig[++len] = n\%10;
      n /= 10;
   }
  // dig[len++] = 0;
   return solve(len, k, false, true).sum;
}
int main () {
   long long l, r, k;
   base[0] = 1;
   for(int i = 1; i < N; ++i) {</pre>
       base[i] = (base[i-1] * 10) \% MOD;
   }
   scanf("%lld %lld %lld", &1, &r, &k);
   printf("lld\n", (go(r, k) - go(l-1, k) + MOD) % MOD);
   return 0;
}
```

# 2 Tree

# 2.1 Divide And Conquer Tree

```
//hdu 4812 D Tree
#include <iostream>
#include <cstdio>
#include <cstring>
#include <vector>
```

```
#pragma comment(linker,"/STACK:102400000,102400000")
using namespace std;
const int maxn = 1e5 + 10;
const int md = 1e6 +3;
int N,K;
vector<int > edge[maxn];
void add_edge(int from,int to) {
 edge[from].push_back(to);
void init() {
 for(int i = 1;i <= N;i ++) edge[i].clear();</pre>
int vi[maxn];
int vis[maxn];
int root;
int mi;
int son[maxn];
int hash[md + 10];
int vers[md + 10];
int verc;
pair<int , int > ans;
int fastpow(int x,int y) {
 int ret = 1 ,mul = x;
 while(y) {
   if(y & 1 ) ret = 1LL * mul * ret % md;
   mul = 1LL * mul* mul % md;
   y >>= 1;
 }
 return ret;
int comm[md + 10];
void inv1() {
 for(int i = 0;i < md;i ++) {</pre>
   comm[i] = fastpow(i,md - 2);
 }
int inv(int t) {
 return comm[t];
void getroot(int t,int sz) {
 vis[t] = true;
 son[t] = 1;
 int mx = 0;
 for(int i = 0;i < edge[t].size();i ++) {</pre>
   int nxt = edge[t][i];
   if(!vis[nxt]) {
```

```
getroot(nxt,sz);
     son[t] += son[nxt];
     mx = max(mx, son[nxt]);
   }
 }
 mx = max(mx,sz - son[t]);
 if(mx <= mi) {</pre>
   root = t;
   mi = mx;
 vis[t] = false;
void dfs(int t,int mul,int ri) {
 vis[t] = true;
 //query
 mul =1LL * mul * vi[t] % md;
 if(1LL * mul * ri % md == K) {
   pair<int ,int > tmp = pair<int ,int > (min(root,t),max(root,t));
   if(tmp < ans) ans = tmp;</pre>
 int q = 1LL* inv(1LL * mul * ri % md) * K % md;
 if(vers[q] == verc && hash[q]!= 0 ) {
   pair<int ,int > tmp = pair<int ,int > (min(t,hash[q]),max(t,hash[q]));
   if(tmp < ans) ans = tmp;</pre>
 }
 son[t] = 1;
 for(int i = 0;i < edge[t].size();i ++) {</pre>
   int nxt = edge[t][i];
   if(!vis[nxt]) {
     dfs(nxt,mul,ri);
     son[t] += son[nxt];
   }
 }
 //set
 if(vers[mul] != verc ) {
   vers[mul] = verc;
   hash[mul] = t;
 hash[mul] = min(hash[mul],t);
 vis[t] = false;
}
void work(int t,int sz) {
 mi = sz;
 getroot(t,sz);
 // dfs
 int rt = root;
```

```
vis[rt] =true;
 verc ++;
 for(int i = 0;i < edge[root].size();i ++) {</pre>
   int nxt = edge[rt][i];
   if(!vis[nxt]) {
     dfs(nxt,1,vi[rt] % md);
   }
 }
 for(int i = 0;i < edge[rt].size();i ++) {</pre>
   int nxt = edge[rt][i];
   if(!vis[nxt]) {
     work(nxt,son[rt]);
   }
 }
}
int main() {
 inv1();
 verc = 0;
 while(scanf("%d%d",&N,&K) != EOF) {
   init();
   for(int i = 1;i <= N;i ++) {</pre>
     scanf("%d",&vi[i]);
   }
   for(int i = 0;i < N - 1;i ++) {</pre>
     int u,v;
     scanf("%d%d",&u,&v);
     add_edge(u,v);
     add_edge(v,u);
   }
   memset(vis,0,sizeof(vis));
   ans = pair<int ,int > (N+1,N+1);
   work(1,N);
   if(ans.first == N+1 && ans.second == N + 1) {
     puts("No solution");
   } else {
     printf("%d %d\n",ans.first,ans.second);
   }
 }
}
```

#### 2.2 Link Tree

```
//HDU 3966
//operation1 path c1 to c2 plus k
//operation2 path c1 to c2 minus k
```

```
#include <iostream>
#include <cstdio>
#include <algorithm>
#include <vector>
#include <cstring>
#pragma comment(linker, "/STACK:1024000000,1024000000")
using namespace std;
#define lc (o<<1)</pre>
#define rc (o<<1|1)
int N,M,P;
const int maxn = 100010;
vector<int > edge[maxn];
int ai[maxn];
void add_edge(int from,int to) {
 edge[from].push_back(to);
void init() {
 for(int i = 1;i <= N;i ++) edge[i].clear();</pre>
int son[maxn]; // size of children
int fa[maxn];
int wn[maxn]; //index in segment
int wcnt;
int vis[maxn];
int dep[maxn]; // depth
int top[maxn]; // link fa
//Tree link
void dfs1(int t,int d) {
 vis[t] = true;
 dep[t] = d;
 son[t] = 1;
 for(int i = 0;i < edge[t].size();i ++) {</pre>
   int nxt = edge[t][i];
   if(!vis[nxt]) {
     fa[nxt] = t;
     dfs1(nxt,d + 1);
     son[t] += nxt;
   }
 vis[t] = false;
void dfs2(int t) {
 vis[t] = true;
 wn[t] = wcnt ++;
 bool first = true;
 int index = -1;
```

```
for(int i = 0;i < edge[t].size();i ++) {</pre>
   int nxt = edge[t][i];
   if(!vis[nxt]) {
     if(first) {
       first =false;
       index = nxt;
     if(son[nxt] > son[index]) {
       index= nxt;
     }
   }
 }
 if(!first ) {
   top[index] = top[t];
   dfs2(index);
   for(int i = 0;i < edge[t].size();i ++) {</pre>
     int nxt = edge[t][i];
     if(!vis[nxt] && nxt != index) {
       top[nxt] = nxt;
      dfs2(nxt);
     }
   }
 }
 vis[t] = false;
//segment tree
int addv[maxn << 2];</pre>
void add(int o,int l,int r,int y1,int y2,int v) {
 if(y1 <= 1 && r <= y2) {</pre>
   addv[o] += v;
 } else {
   int m = (1 + r) >> 1;
   if(y1 <= m) add(lc,1,m,y1,y2,v);</pre>
   if(m < y2) add(rc,m+1,r,y1,y2,v);</pre>
 }
void query(int o,int 1,int r,int x,int & ans) {
 if(1 == r && r == x) {
   ans += addv[o];
 } else {
   int m = (1 + r) >> 1;
   ans += addv[o];
   if(x <= m ) {
     query(lc,1,m,x,ans);
   } else {
     query(rc,m+1,r,x,ans);
```

```
}
 }
}
void init_seg() {
 memset(addv,0,sizeof(addv));
}
char buff[5];
int main() {
 while(~scanf("%d%d%d",&N,&M,&P) ) {
   init();
   for(int i = 1;i <= N;i ++) {</pre>
     scanf("%d",&ai[i]);
   for(int i = 0;i < M;i ++) {</pre>
     int u,v;
     scanf("%d%d",&u,&v);
     add_edge(u,v);
     add_edge(v,u);
   dfs1(1,1);
   wcnt = 0;
   top[1] = 1;
   dfs2(1);
   init_seg();
   while(P --) {
     scanf("%s",buff);
     if(buff[0] == 'I' || buff[0] == 'D') {
       int c1,c2,k;
       scanf("%d%d%d",&c1,&c2,&k);
       if(buff[0] == 'D') k = - k;
       /// query path
       while(top[c1] != top[c2]) {
         int f1 = top[c1];
         int f2 = top[c2];
         if(dep[f1] < dep[f2]) {</pre>
          swap(f1,f2);
          swap(c1,c2);
         add(1,0,N - 1,wn[f1],wn[c1],k);
         c1 = fa[f1];
       if(dep[c1] < dep[c2]) {</pre>
         swap(c1,c2);
       add(1,0,N - 1,wn[c2],wn[c1],k);
     } else if(buff[0] == 'Q') {
```

```
int d;
    scanf("%d",&d);
    int ans = 0;
    query(1,0,N-1,wn[d],ans);
    ans += ai[d];
    printf("%d\n",ans);
    }
}
```

## 2.3 Segment Tree

```
//HDU 4578
//segment plus mul power sum
#include <cstdio>
#include <algorithm>
using namespace std;
#define lc (o<<1)</pre>
#define rc (o<<1|1)
const int maxn = 100010;
const int md = 10007;
int sumv1[maxn<<2], sumv2[maxn<<2], sumv3[maxn<<2];</pre>
int addv[maxn<<2], setv[maxn<<2], timv[maxn<<2];</pre>
void pushdown(int o) {
 if (setv[o] >= 0) {
   setv[lc] = setv[rc] = setv[o];
   addv[lc] = addv[rc] = 0;
   timv[lc] = timv[rc] = 1;
   setv[o] = -1;
 if (timv[o] != 1) {
   addv[lc] *= timv[o];
   addv[lc] %= md;
   addv[rc] *= timv[o];
   addv[rc] %= md;
   timv[lc] *= timv[o];
   timv[lc] %= md;
   timv[rc] *= timv[o];
   timv[rc] %= md;
   timv[o] = 1;
 if (addv[o] > 0) {
   addv[lc] += addv[o];
   addv[lc] %= md;
```

```
addv[rc] += addv[o];
   addv[rc] %= md;
   addv[o] = 0;
 }
}
void maintain(int o,int l,int r) {
 if (1 == r) {
   if (setv[o] != -1) {
     sumv1[o] = setv[o];
     setv[o] = -1;
   if (timv[o] != 1) {
     sumv1[o] *= timv[o];
     timv[o] = 1;
     sumv1[o] %= md;
   if (addv[o] > 0) {
     sumv1[o] += addv[o];
     sumv1[o] %= md;
     addv[o] = 0;
   sumv2[o] = sumv1[o] * sumv1[o] % md;
   sumv3[o] = sumv1[o] * sumv2[o] % md;
 } else {
   sumv1[o] = (sumv1[lc] + sumv1[rc]) % md;
   sumv2[o] = (sumv2[lc] + sumv2[rc]) % md;
   sumv3[o] = (sumv3[lc] + sumv3[rc]) % md;
   if (setv[o] != -1) {
     sumv1[o] = setv[o] * (r - 1 + 1) % md;
     sumv2[o] = setv[o] * setv[o] % md * (r - 1 + 1) % md;
     sumv3[o] = setv[o] * setv[o] % md * setv[o] % md * (r - 1 + 1) % md;
   }
   if (timv[o] != 1) {
     sumv1[o] *= timv[o];
     sumv1[o] %= md;
     sumv2[o] *= timv[o] * timv[o] % md;
     sumv2[o] %= md;
     sumv3[o] *= timv[o] * timv[o] % md * timv[o] % md;
     sumv3[o] %= md;
   if (addv[o] > 0) {
     int tmp1 = sumv1[o];
     sumv1[o] += addv[o] * (r - 1 + 1) % md;
     sumv1[o] %= md;
     int tmp2 = sumv2[o];
     int tmp3 = sumv3[o];
```

```
sumv2[o] = (tmp2 + 2*tmp1%md * addv[o]%md + addv[o] * addv[o] %md* (r - 1 +1)%md) % md;
               sumv3[o] = tmp3 + 3 * tmp2/md * addv[o] % md + 3 * tmp1 % md * addv[o]/md * addv[o] % md + 3 * tmp2/md * addv[o] % md * addv[o] %
                            addv[o] * addv[o] % md * addv[o] % md * (r - 1 + 1) %md;
               sumv3[o] %= md;
         }
     }
}
void setq(int o,int l,int r,int y1,int y2,int v) {
     if (y1 <= 1 && r <= y2) {</pre>
          setv[o] = v;
          addv[o] = 0;
          timv[o] = 1;
     } else {
         pushdown(o);
          int m = (1 + r) >> 1;
         if (y1 <= m) setq(lc,1,m,y1,y2,v);</pre>
          else maintain(lc,1,m);
          if (m < y2) setq(rc,m+1,r,y1,y2,v);</pre>
          else maintain(rc,m+1,r);
     maintain(o,1,r);
}
void addq(int o,int l,int r,int y1,int y2,int v) {
     if (y1 <= 1 && r <= y2) {</pre>
          addv[o] += v;
          addv[o] %= md;
     } else {
         pushdown(o);
          int m = (1 + r) >> 1;
         if (y1 <= m ) addq(lc,1,m,y1,y2,v);</pre>
          else maintain(lc,1,m);
          if (m < y2) addq(rc,m+1,r,y1,y2,v);</pre>
          else maintain(rc,m+1,r);
     }
    maintain(o,1,r);
void timq(int o,int l,int r,int y1,int y2,int v) {
     if (y1 <= 1 && r <= y2) {</pre>
          timv[o] *= v;
         timv[o] %= md;
          addv[o] *= v;
          addv[o] %= md;
     } else {
          pushdown(o);
          int m = (1 + r) >> 1;
         if (y1 <= m) timq(lc,1,m,y1,y2,v);</pre>
```

```
else maintain(lc,1,m);
   if (m < y2) timq(rc,m+1,r,y1,y2,v);</pre>
   else maintain(rc,m+1,r);
 maintain(o,1,r);
}
int ans1, ans2, ans3;
void query(int o,int l,int r,int y1,int y2,int add,int ti) {
 if (setv[o] > 0) {
   add = ti * addv[o] % md + add;
   ti = ti * timv[o] % md;
   int len = min(r,y2) - max(y1,1) + 1;
   int tmp1 = setv[o] * len % md * ti % md;
   int tmp2 = setv[o] * setv[o] % md * len % md * ti%md * ti %md;
   int tmp3 = setv[o] * setv[o] % md * setv[o] % md * len % md *ti %md* ti % md* ti % md;
   int _sum1 = tmp1 + add * len % md;
   _sum1 %= md;
   int _sum2 = (tmp2 + 2* tmp1 * add % md + add * add % md * len % md) % md;
   int _sum3 = (tmp3 + 3 * tmp2 * add % md + 3 * tmp1 * add % md * add % md + len * add % md *
        add % md *add % md) % md;
   ans1 = (ans1 + _sum1) \% md;
   ans2 = (ans2 + \_sum2) % md;
   ans3 = (ans3 + sum3) \% md;
   return ;
 if (y1 <= 1 && r <= y2) {</pre>
   int tmp1 = sumv1[o] * ti % md;
   int tmp2 = sumv2[o] * ti % md * ti % md;
   int tmp3 = sumv3[o] * ti % md * ti % md * ti % md;
   int _sum = tmp1 + add * (r - 1 + 1) % md;
   int _sum2 = tmp2 + 2* tmp1 * add % md + add * add % md * (r - 1 + 1) % md;
   int _sum3 = tmp3 + 3 * tmp2 % md * add % md + 3 * tmp1 % md * add % md * add % md + add * add
       % md * add % md * (r-1+1) % md;
   _sum %= md;
   _sum2 %= md;
   _sum3 %= md;
   ans1 = (ans1 + _sum) % md;
   ans2 = (ans2 + _sum2) \% md;
   ans3 = (ans3 + _sum3) \% md;
 } else {
   int m = (1 +r ) >> 1;
   if (y1 <= m) query(lc,1,m,y1,y2,(ti * addv[o] % md + add) % md,ti * timv[o] % md);</pre>
   if (m < y2) query(rc,m+1,r,y1,y2,(ti * addv[o] % md + add) % md,ti * timv[o] % md);</pre>
 }
}
```

```
void init(int o,int l,int r) {
 setv[o] = -1;
 timv[o] = 1;
 addv[o] = 0;
 sumv1[o] = sumv2[o] = sumv3[o] = 0;
 if (1 == r) {
 } else {
   int m = (1 + r) >> 1;
   init(lc,1,m);
   init(rc,m+1,r);
 }
}
int main() {
 int N,M;
 while (scanf("%d%d",&N,&M)==2 && N && M) {
   init(1,1,N);
   while (M --) {
     int cmd,x,y,c;
     scanf("%d%d%d%d",&cmd,&x,&y,&c);
     if(cmd == 1) {
       c %= md;
       addq(1,1,N,x,y,c);
     } else if(cmd == 2) {
       c %= md;
      timq(1,1,N,x,y,c);
     } else if(cmd == 3) {
       c %= md;
       setq(1,1,N,x,y,c);
     } else if(cmd == 4) {
       ans1 = ans2 = ans3 = 0;
       query(1,1,N,x,y,0,1);
       if(c == 1) {
        printf("%d\n",ans1);
       } else if(c == 2){
        printf("%d\n",ans2);
       } else if(c == 3) {
        printf("%d\n",ans3);
      }
     }
   }
 }
}
```

# 2.4 Splay Tree

```
#include <cstdio>
#include <iostream>
using namespace std;
struct Node {
 Node* ch[2];
 int v, s, flip;
 void maintain() {
   s = 1 + ch[0] -> s + ch[1] -> s;
 void pushdown() {
   if (flip) {
     flip = 0;
     swap(ch[0], ch[1]);
     ch[0]->flip ^= 1;
     ch[1]->flip ^= 1;
   }
 int cmp(int k) const {
   int d = k - ch[0] ->s;
   if (d == 1) return -1;
   return d <= 0 ? 0 : 1;
 }
};
Node* null = new Node();
void rotate(Node* &o, int d) {
 Node* k = o \rightarrow ch[d^1];
 o->ch[d^1] = k->ch[d];
 k \rightarrow ch[d] = o;
 o->maintain();
 k->maintain();
 o = k;
void splay(Node* &o, int k) {
 o->pushdown();
 int d = o->cmp(k);
 if (d == 1) k -= o -> ch[0] -> s + 1;
 if (d != -1) {
   Node* p = o->ch[d];
   p->pushdown();
   int d2 = p - cmp(k);
   int k2 = (d2 == 0) ? k : k - p -> ch[0] -> s - 1;
   if (d2 != -1) {
     splay(p->ch[d2], k2);
     if (d == d2) {
       rotate(o, d^1);
```

```
} else {
       rotate(o->ch[d], d);
   }
   rotate(o, d^1);
 }
}
Node* merge(Node* left, Node* right) { // make sure left != null
 splay(left, left->s);
 left->ch[1] = right;
 left->maintain();
 return left;
void split(Node* o, int k, Node* &left, Node* &right) { // make sure 1 <= k <= o->s
 splay(o, k);
 left = o;
 right = o->ch[1];
 o->ch[1] = null;
 left->maintain();
}
const int maxn = 300000 + 10;
struct SS {
 int n;
 Node seq[maxn];
 Node* root;
 Node* build(int sz) {
   if (!sz) return null;
   Node* L = build(sz/2);
   Node* o = &seq[++n];
   o->v = n-1;
   o->flip = 0;
   o \rightarrow ch[0] = L;
   o \rightarrow ch[1] = build(sz - sz/2 - 1);
   o->maintain();
   return o;
 void init(int sz) {
   n = 0;
   null->s = null->flip = 0;
   root = build(sz);
 void print(Node *o) {
   if (o != null) {
     o->pushdown();
     print(o->ch[0]);
     if (o->v) {
```

```
if (o->v != 1) putchar(' ');
      printf("%d", o->v);
     print(o->ch[1]);
   }
 }
} ss;
int n, m, a, b, c;
char op[10];
int main() {
 while (scanf("%d%d",&n,&m) == 2 && n != -1 && m != -1) {
   ss.init(n+1);
   Node *t1, *t2, *t3;
   while(m--){
     scanf("%s",op);
     if(op[0]=='C'){ // split [a,b], put it after c}
       scanf("%d%d%d",&a,&b,&c);
       split(ss.root, b+1, t1, t2);
       split(t1, a, t1, t3);
      ss.root = merge(t1, t2);
       split(ss.root, c+1, t1, t2);
       ss.root = merge(merge(t1, t3), t2);
     } else { // flip [a,b]
       scanf("%d%d",&a,&b);
       split(ss.root, b+1, t1, t3);
       split(t1, a, t1, t2);
      t2->flip ^= 1;
       ss.root = merge(merge(t1, t2), t3);
     }
   }
   ss.print(ss.root);
   puts("\n");
}
```

## 2.5 Treap

```
struct Node {
  Node *ch[2]; // 0-left 1-right
  int r, v, s; // rank, val, #node
  Node(int v): v(v) {
    ch[0] = ch[1] = NULL;
    r = rand();
    s = 1;
}
```

```
int cmp(int x) const {
   if (x == v) return -1;
   return x < v ? 0 : 1;</pre>
 void maintain() { // maintain #node
   s = 1;
   if (ch[0] != NULL) s += ch[0]->s;
   if (ch[1] != NULL) s += ch[1]->s;
 }
};
void rotate(Node* &o, int d) {
 Node* k = o \rightarrow ch[d^1];
 o->ch[d^1] = k->ch[d];
 k \rightarrow ch[d] = o;
 o->maintain();
 k->maintain();
 o = k;
void insert(Node* &o, int x) {
 if (o == NULL) {
   o = new Node(x);
 } else {
   int d = o \rightarrow cmp(x);
   if (d != -1) { // same ele won't be inserted
     insert(o->ch[d], x);
     if (o->ch[d]->r > o->r) rotate(o, d^1);
   }
 }
 o->maintain();
void remove(Node* &o, int x) {
 if (o == NULL) return ; // ele to be removed not exist
 int d = o \rightarrow cmp(x);
 if (d == -1) {
   Node* ret = o;
   if (o->ch[0] != NULL && o->ch[1] != NULL) {
     int d2 = (o->ch[0]->r > o->ch[1]->r ? 1 : 0);
     rotate(o, d2);
     remove(o->ch[d2], x);
   } else {
     if (o->ch[0] == NULL) o = o->ch[1];
     else o = o \rightarrow ch[0];
     delete ret;
   }
 } else {
   remove(o->ch[d], x);
```

```
if (o) o->maintain();
int find(Node* o, int x) {
 while (o != NULL) {
   int d = o \rightarrow cmp(x);
   if (d == -1) return 1;
   else o = o->ch[d];
 return 0;
}
int kth_big(Node* o, int k) {
 if (o == NULL || k <= 0 || k > o->s) return 0;
 int s = o > ch[1] == NULL ? 0 : o > ch[1] -> s;
 if (k == s+1) return o->v;
 else if (k <= s) return kth_big(o->ch[1], k);
 else return kth_big(o->ch[0], k-s-1);
int kth_small(Node* o, int k) {
 if (o == NULL || k <= 0 || k > o->s) return 0;
 int s = o \rightarrow ch[0] == NULL ? 0 : o \rightarrow ch[0] \rightarrow s;
 if (k == s) return o->v;
 else if (k < s) return kth_small(o->ch[0], k);
 else return kth_small(o->ch[1], k-s-1);
void merge(Node* &src, Node* &dest) {
 if (src == NULL) return ;
 merge(src->ch[0], dest);
 merge(src->ch[1], dest);
 insert(dest, src->v);
 delete src;
 src = NULL;
void clear(Node* &o) {
 if (o == NULL) return ;
 clear(o->ch[0]);
 clear(o->ch[1]);
 delete o;
 o = NULL;
}
```

# 3 Geometry

## 3.1 Basic Struct and Algorithm

```
struct Point {
 double x, y;
 Point(double x=0, double y=0):x(x),y(y){}
};
typedef Point Vector;
Vector operator + (const Vector &A, const Vector &B) { return Vector(A.x+B.x, A.y+B.y); }
Vector operator - (const Point &A, const Point &B) { return Vector(A.x-B.x, A.y-B.y); }
Vector operator * (const Vector &A, double p) { return Vector(A.x*p, A.y*p); }
double Dot(const Vector &A, const Vector &B) { return A.x*B.x + A.y*B.y; }
double Cross(const Vector &A, const Vector &B) { return A.x*B.y - A.y*B.x; }
double Length(const Vector &A) { return sqrt(Dot(A, A)); }
Vector Normal(const Vector &A) { double L = Length(A); return Vector(-A.y/L, A.x/L); }
struct Line {
 Point P;
 Vector v;
 double ang;
 Line() {}
 Line(Point P, Vector v):P(P),v(v){ ang = atan2(v.y, v.x); }
 bool operator < (const Line &L) const {</pre>
   return ang < L.ang;</pre>
 }
};
// if p is on the left side of L
bool OnLeft(const Line &L, const Point &p) {
 return Cross(L.v, p-L.P) > 0;
}
// intersection of line $a$ and $b$
Point GetLineIntersection(const Line &a, const Line &b) {
 Vector u = a.P-b.P;
 double t = Cross(b.v, u) / Cross(a.v, b.v);
 return a.P+a.v*t;
}
```

#### 3.2 Polygon Area

```
double PolygonArea(vector<Point> p) {
  int n = p.size();
  double area = 0;
  for(int i = 1; i < n-1; i++)</pre>
```

```
area += Cross(p[i]-p[0], p[i+1]-p[0]);
return area/2;
}
```

#### 3.3 Half Plane Intersection

```
const double eps = 1e-6;
// intersection of areas (leftside of lines)
vector<Point> HalfplaneIntersection(vector<Line> L) {
 int n = L.size();
 sort(L.begin(), L.end());
 int first, last;
 vector<Point> p(n);
 vector<Line> q(n);
 vector<Point> ans;
 q[first=last=0] = L[0];
 for(int i = 1; i < n; i++) {</pre>
   while(first < last && !OnLeft(L[i], p[last-1])) last--;</pre>
   while(first < last && !OnLeft(L[i], p[first])) first++;</pre>
   q[++last] = L[i];
   if(fabs(Cross(q[last].v, q[last-1].v)) < eps) {</pre>
     if(OnLeft(q[last], L[i].P)) q[last] = L[i];
   if(first < last) p[last-1] = GetLineIntersection(q[last-1], q[last]);</pre>
 while(first < last && !OnLeft(q[first], p[last-1])) last--;</pre>
 if(last - first <= 1) return ans;</pre>
 p[last] = GetLineIntersection(q[last], q[first]);
 for(int i = first; i <= last; i++) ans.push_back(p[i]);</pre>
 return ans;
}
```

# 4 Math

# 4.1 China Remainder Theory

```
// china remainder theory, no matter whether gcd(m[i],m[j])=1
LL CRT(const vector<LL>&m, const vector<LL> &b){
  bool flag = false;
  LL x, y, i, d, result, a1, m1, a2, m2, Size = m.size();
  m1 = m[0], a1 = b[0];
  for(int i = 1; i < Size; i++){</pre>
```

```
m2 = m[i], a2 = b[i];
d = exgcd(m1, m2, x, y);
if ((a2 - a1) % d != 0) flag = true;
result = (mul_mod(x, (a2 - a1) / d, m2) % m2 + m2) % m2;
LL tmp = m1;
m1 = m1 / d * m2;
a1 = (a1 + mul_mod(tmp, result, m1)) % m1;
a1 = (a1 % m1 + m1) % m1;
}
if (flag) return -1;
else return a1;
}
```

#### 4.2 Decompose

```
// eg: poj 3471
const int maxn = 10000000;
const int maxp = 700000; // about maxn/log(maxn)
struct Factor{ // factor as p^num
 int p, num;
};
struct DeComposer {
 DeComposer() { gen_primes(); }
 bool vis[maxn+5];
 int pn, prime[maxp];
 void sieve() {
   int m = (int)sqrt(maxn+0.5);
   memset(vis,0,sizeof(vis));
   for(int i=2;i<=m;++i)if(!vis[i])</pre>
     for(int j=i*i;j<=maxn;j+=i)vis[j]=1;</pre>
 void gen_primes() {
   sieve();
   pn = 0;
   for (int i = 2; i <= maxn; ++ i) {</pre>
     if (!vis[i]) prime[pn++] = i;
   }
 }
 int fcn;
 Factor fc[64]; // x = p1^a1 * p2^a2 * ...
 int fn, factor[maxp]; // all y satisify y|x
 void decompose2(int x,int d){
   if(d==fcn){
     factor[fn++] = x;
   } else {
```

```
for(int i = 0; i <= fc[d].num; ++ i) {</pre>
       decompose2(x, d+1);
       x *= fc[d].p;
   }
 }
 void decompose1(int x) {
   fcn = 0;
   for(int i = 0; i < pn && prime[i] * prime[i] <= x; ++ i) if (x % prime[i] == 0) {</pre>
     fc[fcn].p = prime[i];
     fc[fcn].num = 0;
     while(x % prime[i] == 0) {
      fc[fcn].num ++;
      x /= prime[i];
     }
     fcn ++;
   }
   if (x > 1) {
     fc[fcn].p = x;
    fc[fcn].num = 1;
     fcn ++;
   }
 }
 void decompose(int x){
   decompose1(x);
   fn = 0;
   decompose2(1,0);
} dc_solver;
```

## 4.3 Euler Phi

```
// #x that x<=n && gcd(x,n)==1
int euler_phi(int n) {
  int m = (int)sqrt(n+0.5);
  int ans = n;
  for (int i = 2; i <= m; ++ i) if (n % i == 0) {
    ans = ans / i * (i-1);
    while (n%i == 0) n /= i;
  }
  if (n > 1) ans = ans / n * (n-1);
  return ans;
}
int phi[maxn];
void phi_table(int n) {
```

```
for (int i = 2; i <= n; ++ i) phi[i] = 0;
phi[1] = 1;
for (int i = 2; i <= n; ++ i) {
   if (!phi[i]) {
      for (int j = i; j <= n; j += i) {
        if (!phi[j]) phi[j] = j;
        phi[j] = phi[j] / i * (i-1);
      }
   }
   phi[i] += phi[i-1];
}</pre>
```

#### 4.4 Extend GCD

```
// a * x + b * y = d, |x| + |y| get the minimum
LL exgcd(LL a, LL &d, LL &d, LL &x, LL &y){
  if (a) { x = 0; y = 1; return a; }
  else { exgcd(b, a%b, d, y, x); y -= x*(a/b); }
}
```

## 4.5 Integer Inverse

```
LL inv1(LL a, LL n) { // a^-1 under n
    LL d, x, y;
    gcd(a,n,d,x,y);
    return d == 1 ? (x+n)%n : -1;
}
LL inv2(LL a, LL p) { // in case that p is a prime
    return pow_mod(a, p-2, p);
}
```

#### 4.6 Line Mod

```
// ax = b (mod n)
// let d = gcd(a,n), use exgcd to solve ax + ny = d
// if b|d, then there are #ans=d, otherwise, no solution
vector<LL> line_mod(LL a, LL b, LL n) {
   LL x, y;
   exgcd(a,n,x,y);
   vector<LL>ans;
   ans.clear();
```

```
if(b%d==0){
    x%=n; x+=n; x%=n;
    ans.push_back(x*(b/d)%(n/d));
    for(LL i=1;i<d;++i){
        ans.push_back((ans[0]+i*n/d)%n);
    }
}
return ans;
}</pre>
```

# 4.7 Log Mod

```
// eg: hdu 2815
// d*a^(x-c) = b \pmod{n}, make sure that (a,n) = 1 and (d,n) = 1
map<LL,LL>f;
LL log_mod(LL a, LL b, LL n, LL c, LL d) {
 LL m, v, e=1, i, x, y, dd;
  m = ceil( sqrt(n + 0.5) );
  f.clear();
  f[1] = m;
  for(i = 1; i < m; ++ i) {</pre>
    e = e*a%n;
   if (!f[e]) f[e] = i;
  }
  e = (e*a)%n;
  for (i = 0; i < m; ++ i) {</pre>
    exgcd(d,n,dd,x,y);
   x = (x*b%n + n) % n;
   if (f[x]) {
     LL num = f[x];
      return c + i*m + (num==m ? 0 : num);
    d = (d*e) \% n;
 return -1;
}
// a^x = b (mod n), no restriction
LL log_mod(LL a, LL b, LL n) {
  b%=n;
  LL c = 0, d = 1, t;
  \label{eq:while} \begin{aligned} & \text{while}((\texttt{t=\_gcd}(\texttt{a},\texttt{n}))\,!=1) \{ \end{aligned}
   if(b%t) return -1;
   c++;
   n/=t;
    b/=t;
```

```
d=d*a/t%n;
if(d==b)return c;
}
return log_mod(a,b,n,c,d);
}
```

#### 4.8 Lucas

```
// C(n,m) % p, make sure p is prime, p <= 10<sup>5</sup>
// n = n[k] * p^k + n[k-1] * p^(k-1) + ... + n[0]
// m = m[k] * p^k + m[k-1] * p^(k-1) + ... + m[0]
// then, C(n,m) = C(n[k],m[k])*C(n[k-1],m[k-1])*..*C(n[0],m[0]) (mod p)
// C(n,m) = C(n/p, m/p) * C(n/p, m/p) (mod p)
// eg: hdu3037
LL Lucas(LL n, LL m, LL p) {
 LL ret = 1;
 while(n && m) {
   LL np = n%p, mp = m%p;
   if(np < mp) return 0;</pre>
   ret = ret * factorial(np) % p * reverse(factorial(mp), p) % p * reverse(factorial(np-mp), p) %
       p;
   n /= p;
   m \neq p;
 return ret;
}
```

#### 4.9 Miller Rabin

```
// prime test
bool Witness(LL n, LL a) {
    LL m = n-1, j = 0;
    while(!(m&1)) m >>= 1, j ++;
    LL ans = pow_mod(a, m, n);
    while (j --) {
        LL tmp = mul_mod(ans, ans, n);
        if (tmp == 1 && ans != 1 && ans != n-1) return 1;
        ans = tmp;
    }
    return ans != 1;
}
bool Miller_Rabin(LL n) {
    if (n < 2) return 0;</pre>
```

```
if (n == 2) return 1;
if (!(n&1)) return 0;
for (int i = 0; i < max_test; ++ i) {
    ll a = rand() % (n-2) + 2;
    if (Witness(n,a)) return 0;
}
return 1;
}</pre>
```

#### 4.10 Mul Mod

```
// x*y % n
LL mul_mod(LL x, LL y, LL n) {
    LL T = floor(sqrt(n) + 0.5);
    LL t = T * T - n;
    LL a = x / T, b = x % T;
    LL c = y / T, d = y % T;
    LL e = a * c / T, f = a * c % T;
    LL v = ((a*d + b*c) % n + e*t) % n;
    LL g = v / T, h = v % T;
    LL ret = (((f+g)*t % n + b*d) % n + h*T) % n;
    return (ret % n + n) % n;
}
```

#### 4.11 Pollard Rho

#### 4.12 Pow Mod

```
// a^x % n
LL pow_mod(LL a, LL x, LL n) {
    LL ret = 1, mul = a;
    while (x) {
        if (x&1) ret = mul_mod(ret, mul, n);
        mul = mul_mod(mul, mul, n);
        x >>= 1;
    }
    return ret;
}
```

#### 4.13 Power Mod

```
// x^n = a \pmod{p}, make sure that p is prime
// let g be a primitive root of p, x = g^y, a = g^m
// use log_mod to get m, g^(yn) = g^m \pmod{p}
// thus yn = m \pmod{p-1}, use exgcd to solve and get back
vector<int> power_mod(int a, int n, int p) {
 int g = primitive_root(p);
 LL m = log_mod(g, a, p);
 vector<int>ret;
 if(a==0){
   ret.push_back(0);
   return ;
 if(m==-1)return ret;
 LL A=n, B=p-1, C=m, x, y;
 LL d = exgcd(A,B,x,y);
 if(C%d!=0)return ret;
 x=x*(C/d)%B;
 LL delta=B/d;
 for(int i=0;i<d;++i){</pre>
   x=((x+delta)\%B+B)\%B;
   ret.push_back((int)pow_mod(g,x,p));
 sort(ret.begin(),ret.end());
 ret.erase(unique(ret.begin(),ret.end()), ret.end());
 return ret;
}
```

#### 4.14 Primitive Root

```
// eg: SGU 511
struct PR {
 // make sure that p is prime
 // if p = 2, solve the prob. without PR
 int divs[N+5];
 int primitive_root(const int p) {
   if (p == 2) return 1;
   int cnt = 0, m = p-1;
   for (int i = 2; i*i <= m; ++ i) if (m%i == 0) {</pre>
     divs[cnt++] = i;
     if (i*i < m) divs[cnt++] = m/i;</pre>
   int r = 2, j = 0;
   while (1) {
    for (j = 0; j < cnt; ++ j) {</pre>
       if (fastpow(r, divs[j], p) == 1) break;
     if (j >= cnt) return r;
     r ++;
   return -1;
} pr_solver;
```

#### 4.15 Square Mod

```
// x*x = a \pmod{n}, make sure that n is prime
// be careful there is a single sol. when n = 2
// otherwise, x and n-x are both okay
// eg: ural 1132
LL modsqr(LL a, LL n) {
 LL b, k, i, x;
 if (n == 2) return a % n;
 if (pow_mod(a, (n-1)/2, n) == 1) {
   if (n%4 == 3) {
     x = pow_mod(a, (n+1)/4, n);
   }else{
     for(b=1; pow_mod(b, (n-1)/2, n) == 1; b ++);
     i = (n-1)/2;
    k = 0;
     do {
      i/=2:
      k/=2;
      if((pow_mod(a,i,n) * pow_mod(b,k,n)+1) %n == 0) {
```

```
k += (n-1)/2;
}
while(i%2 == 0);
x = (pow_mod(a,(i+1)/2,n) * pow_mod(b,k/2,n)) %n;
}
if(x*2 > n) x = n-x;
return x;
}
return -1;
}
```

## 5 Others

#### 5.1 Exact Cover

```
// la 2659
#include <cstdio>
#include <vector>
using namespace std;
const int MROW = 16*16*16 + 5;
const int MCOL = 16*16*4 + 5;
const int NODE = 16*16*16*4 + 5;
struct DLX {
 int n, sz;
 int S[MCOL];
 int row[NODE], col[NODE];
 int ansd, ans[MROW];
 int L[NODE], R[NODE], U[NODE], D[NODE];
 void init(int n) {
   this \rightarrow n = n;
   for (int i = 0; i <= n; ++ i) {</pre>
     U[i] = D[i] = i;
     L[i] = i-1; R[i] = i+1;
     S[i] = 0;
   R[n] = 0; L[0] = n;
   sz = n+1;
 void addRow(int r, const vector<int> &columns) {
   int first = sz;
   for (int i = 0; i < columns.size(); ++ i) {</pre>
     int c = columns[i];
     L[sz] = sz-1; R[sz] = sz+1;
     D[sz] = c; U[sz] = U[c];
```

```
D[U[c]] = sz; U[c] = sz;
     row[sz] = r; col[sz] = c;
     S[c] ++; sz ++;
   R[sz-1] = first; L[first] = sz-1;
 #define FOR(i,A,s) for(int i=A[s];i!=s;i=A[i])
 void remove(int c) {
   L[R[c]] = L[c]; R[L[c]] = R[c];
   FOR(i,D,c)
     FOR(j,R,i) \{ U[D[j]] = U[j]; D[U[j]] = D[j]; -- S[col[j]]; \}
 void restore(int c) {
   FOR(i,U,c)
     FOR(j,L,i) { ++S[col[j]]; U[D[j]]=j; D[U[j]]=j; }
   L[R[c]] = c; R[L[c]] = c;
 bool dfs(int d) {
   if (R[0] == 0) {
     ansd = d;
     return 1;
   int c = R[0];
   FOR(i,R,0) if(S[i]<S[c]) c=i;</pre>
   remove(c);
   FOR(i,D,c) {
     ans[d] = row[i];
     FOR(j,R,i) remove(col[j]);
     if(dfs(d+1)) return 1;
     FOR(j,L,i) restore(col[j]);
   restore(c);
   return 0;
 bool solve(vector<int>&v) {
   v.clear();
   if (!dfs(0)) return 0;
   for (int i = 0; i < ansd; ++ i) v.push_back(ans[i]);</pre>
   return 1;
 }
} dlx;
char data[18][18];
bool input() {
 for (int i = 0; i < 16; ++ i) {</pre>
   if (scanf("%s",data[i]) == EOF) return 0;
 }
```

```
return 1;
}
enum { SLOT=0, ROW, COL, BLOK };
int encode(int i, int j, int k) {
 return i*256 + j*16 + k + 1;
int block(int i, int j) {
 return 4*(i/4) + (j/4);
void decode(int x, int &a, int &b, int &c) {
 x --;
 c = x \% 16; x /= 16;
 b = x \% 16; x /= 16;
 a = x;
}
vector<int>columns;
void solve() {
 dlx.init(16*16*4);
 for (int i = 0; i < 16; ++ i) {</pre>
   for (int j = 0; j < 16; ++ j) {
     for (int k = 0; k < 16; ++ k) {
       if (data[i][j] == '-' || data[i][j] == k+'A') {
         columns.clear();
         columns.push_back(encode(SLOT, i, j));
         columns.push_back(encode(ROW, i, k));
         columns.push_back(encode(COL, j, k));
         columns.push_back(encode(BLOK, block(i,j), k));
         dlx.addRow(encode(i,j,k), columns);
       }
     }
   }
 }
 columns.clear();
 dlx.solve(columns);
 for (int i = 0; i < columns.size(); ++ i) {</pre>
   int r, c, v;
   decode(columns[i], r, c, v);
   data[r][c] = char('A' + v);
 for (int i = 0; i < 16; ++ i) {</pre>
   printf("%s\n", data[i]);
 }
int main() {
 int kcase = 0;
 while (input()) {
```

```
if (kcase) puts("");
  kcase ++;
  solve();
}
```

#### 5.2 Matrix Fast Power

```
struct Matrix {
 int n, a[N][N];
 Matrix operator * (const Matrix &b) const {
   Matrix ret; ret.clear();
   ret.n = n;
   for (int i = 0; i < n; ++ i) {</pre>
     for (int k = 0; k < n; ++ k) if (a[i][k]) {</pre>
       for (int j = 0; j < n; ++ j) {
        ret.a[i][j] += a[i][k] * b.a[k][j];
        ret.a[i][j] %= mod;
      }
     }
   }
   return ret;
 void clear() {
   memset(a,0,sizeof(a));
 }
};
Matrix matrix_one(int n) {
 Matrix ret; ret.clear();
 ret.n = n;
 for (int i = 0; i < n; ++ i) {</pre>
   ret.a[i][i] = 1;
 }
 return ret;
Matrix matrix_pow(Matrix x, int n) {
 Matrix ret = matrix_one(x.n), mul = x;
 while (n) {
   if (n&1) ret = ret * mul;
   mul = mul * mul;
   n >>= 1;
 }
 return ret;
```

## 5.3 Polynomial

```
// eg: UVALive 4305
const int MAXN = 500;
const double EPS = 1e-10;
inline int sgn(const double &a) { return a > EPS ? 1 : (a < -EPS ? -1 : 0); }</pre>
struct Polynomial {
 double data[MAXN];
 int n;
 Polynomial() {}
 Polynomial(int _n) : n(_n) {
   memset(data, 0, sizeof(data));
 Polynomial(double *_data, int _n) {
   memset(data, 0, sizeof(data));
   n = _n;
   for (int i = n; i >= 0; i--) data[i] = _data[i];
 Polynomial operator + (const Polynomial &a) {
   Polynomial c(max(n, a.n));
   for (int i = c.n; i >= 0; i--) c.data[i] = data[i] + a.data[i];
   while (sgn(c.data[c.n]) == 0 \&\& c.n) c.n--;
   return c;
 Polynomial operator - (const Polynomial &a) {
   Polynomial c(max(n, a.n));
   for (int i = c.n; i >= 0; i--) c.data[i] = data[i] - a.data[i];
   while (sgn(c.data[c.n]) == 0 && c.n) c.n--;
   return c;
 Polynomial operator * (const Polynomial &a) {
   Polynomial c(n + a.n);
   for (int i = n; i >= 0; i--) for (int j = a.n; j >= 0; j--) c.data[i + j] += data[i] *
       a.data[j];
   return c;
 Polynomial operator / (const Polynomial &a) {
   if (n < a.n) return *this;</pre>
   else {
     Polynomial c(n - a.n);
     for (int i = c.n; i >= 0; i--) c.data[i] = data[i + a.n];
     for (int i = c.n; i >= 0; i--) {
       c.data[i] /= a.data[a.n];
      for (int j = i - 1; a.n - i + j \ge 0 && j \ge 0; j - 0) c.data[j] -= c.data[i] * a.data[a.n -
           i + j];
     }
```

```
return c;
   }
 Polynomial operator % (const Polynomial &a) {
   Polynomial c = *this - *this / a * a;
   while (sgn(c.data[c.n]) == 0 && c.n) c.n--;
   return c;
 }
 bool iszero() {
   return n == 0 && sgn(data[0]) == 0;
 bool isconst() {
   return n > 0;
 Polynomial derivative() {
   Polynomial a(n - 1);
   for (int i = n - 1; i >= 0; i--) a.data[i] = data[i + 1] * (double)(i + 1);
   return a;
 Polynomial integral() {
   Polynomial a(n + 1);
   for (int i = n + 1; i >= 1; i--) a.data[i] = data[i - 1] / (double)i;
   return a;
 }
 void show() {
   for (int i = n; i >= 0; i--) {
     printf("%.6f", data[i], i);
     if (i != 0) printf(" x");
     if (i != 1 && i != 0) printf(" ^ %d", i);
     if (i != 0) printf(" + ");
     else printf("\n");
   }
 }
};
Polynomial gcd(Polynomial a , Polynomial b) {
 if (b.iszero()) return a;
 else return gcd(b, a % b);
}
```