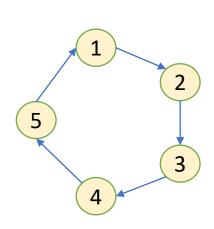
Cassandra

A distributed Open Source NoSQL Database

Cassandra CQL	vs SQL
CREATE KEYSPACE myDatabase WITH replication = {'class': 'SimpleStrategy', 'replication_factor': 1};	CREATE DATABASE myDatabase;
USE myDatabase;	"
CREATE TABLE IF NOT EXISTS myTable (id INT PRIMARY KEY); (synonyms in cql: COLUMNFAMILY=TABLE) NB table need primary key in CQL.	"
ALTER TABLE myTable ADD myField INT;	"
CREATE INDEX myIndex ON myTable (myField);	"
INSERT INTO myTable (id, myField) VALUES (1, 7);	"
SELECT * FROM myTable WHERE myField = 7;	"
SELECT COUNT(*) FROM myTable;	"
DELETE FROM myTable WHERE myField = 7;	"

CQL	SQL
- No support for things like JOIN, GROUP BY, or FOREIGN KEY . Leaving these features out is important because it makes writing and retrieving data from Cassandra much more efficient. Writes are cheap . Write everything the way you want to read it.	JOIN, GROUP BY, FOREIGN KEY
CQL does not perform a read while inserting. Without a read, there is no way to know if the data being inserted is replacing an existing record. This means that both inserts and updates are extremely fast.	
 UPDATE myTable SET myField = 2 WHERE id = 6; - However, if the row does not exist, it will still get created. Similarly as unintuitive, an INSERT statement will actually replace data if it exists. In where-clause, only primary key column can be used. Under the hood, INSERT and UPDATE are treated the same by Cassandra ("Upserts"), except for Counter columns/tables. Both INSERT and UPDATE require complete PRIMARY KEY. 	
Transaction Control Language (TCL) - Not in CQL	COMMIT – It saves the work done SAVEPOINT – It identifies a point in a transaction to which you can later roll back ROLLBACK – It restores database to original since the last COMMIT
Data Retrieval/Query Language (DRL/DQL): Simple transactions (Relation between database objects is not possible): - Where clause: only on primary key or secondary indexes! - Can use only AND operator, There are no OR and NOT operators.	Data Retrieval/Query Language (DRL/DQL): Full transactions.



П	
has	P
h(F	arti
rim	tior
⊃Ke	#

	Node	$es \rightarrow$				_
	1	2	3	4	5	
1	SE	SE	SE		λe	, e
2		DE	DE	DE	piid	yspa
3			NO	NO	NO :	Keyspace->
4	DK			DK	DK	ľ
5	UK	UK			UK	

	1	2	3	4
1	SE			
2		DE		
3			NO	
4				DK

Replicas:1

	1	2	
1	SE	SE	
2	DE	DE	
3	NO	NO	
4	DK	DK	

	1	2
1	SE	
2		DE
3	NO	
4		DK

Replicas:1

	1	
1	SE	
2	DE	ivebiicas. I
3	NO	J. L.
4	DK	

	1	2	3	4
1	SE, DE	SE, DE		
	DE	DE		
2		NO, DK	NO, DK	
		DK	DK	

userid	firstname	lastname	email
3b4c48a7-13b5-4aba-9f0a-dc75ded08a99	Barney	Rubble	rubble@hotmail.com
2506535a-4999-438d-8682-d5a739596343	Fred	Flinstone	fred@gmail.com
31e24f9d-0d27-4143-82b1-fa1a4268d028	Joe	Rockhead	joer@yahoo.com
57bae20b-3694-4975-a274-db5e856d24ab	Wilma	Flinstone	wilma@bedrock.com

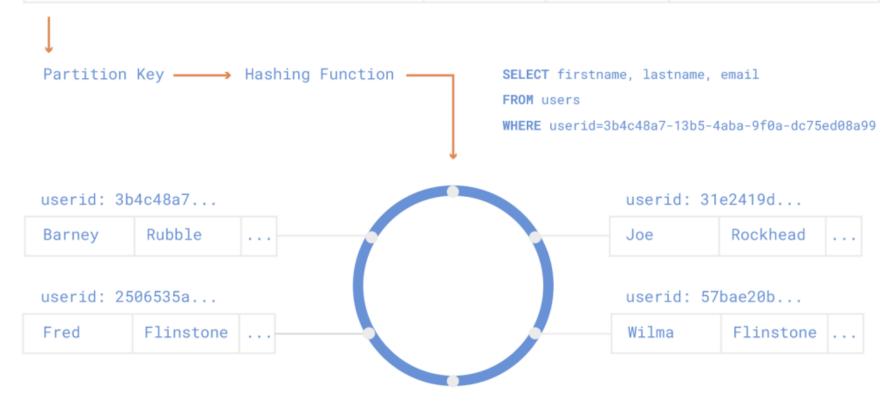


FIGURE 1 How Cassandra stores data