## **Solution for Assignment 5**

1. Consider the following context-free grammar

```
A \rightarrow B C D \mid f B

B \rightarrow a B \mid \epsilon

C \rightarrow b C f \mid c \mid \epsilon

D \rightarrow d D \mid e \mid \epsilon

where A is the start symbol.
```

(a) (30%) Compute First sets and Follow sets for the nonterminals in the grammar.

```
Ans : First(A) = \{a,b,c,d,e,f,\epsilon\} Follow(A) = \{\$\}

First(B) = \{a,\epsilon\} Follow(B) = \{b,c,d,e,\$\}

First(C) = \{b,c,\epsilon\} Follow(C) = \{d,e,f,\$\}

First(D) = \{d,e,\epsilon\} Follow(D) = \{\$\}
```

(b) (40%) Construct the functions of the recursive-decent parser for the nonterminals.

```
Ans:
```

```
const int
     a = 1; b = 2; c = 3; d = 4; e = 5; f = 6; $ = 7;
int token = lexer();
void match(int t){
     if(token == t) token = lexer(); else error();
}
void A(){
     switch(token){
          case a:
          case b:
          case c:
          case d:
          case e:
          case $:
              B(); C(); D(); break;
          case f:
              match(f); B(); break;
          default : error();
     }
}
```

```
void B(){
    switch(token){
         case a:
             match(a); B(); break;
         case b:
         case c:
         case d:
         case e:
        case $:
             break;
        default : error();
    }
}
void C(){
    switch(token){
        case b:
             match(b); C(); match(f); break;
        case c:
             match(c); break;
         case d:
         case e:
         case f:
        case $:
             break;
        default : error();
    }
}
void D(){
    switch(token){
         case d:
             match(d); D(); break;
         case e:
             match(e); break;
        case $:
             break;
        default : error();
    }
}
```

(c) (30%) Construct the parsing table of the table-driven predictive parser for the grammar.

Ans:

	Α	В	С	D
а	$A \rightarrow B C D$	$B \rightarrow a B$		
b	$A \rightarrow B C D$	$B \rightarrow \epsilon$	$C \rightarrow b C f$	
С	$A \rightarrow B C D$	$B \rightarrow \epsilon$	$C \rightarrow c$	
d	$A \rightarrow B C D$	$B \rightarrow \epsilon$	$C \rightarrow \epsilon$	$D \rightarrow dD$
е	$A \rightarrow B C D$	$B \rightarrow \epsilon$	$C \rightarrow \epsilon$	D → e
f	$A \rightarrow f B$		$C \rightarrow \epsilon$	
\$	$A \rightarrow B C D$	$B \rightarrow \epsilon$	$C \rightarrow \epsilon$	$D \rightarrow \epsilon$