

## Problem

True or false?

$$999! < 500^{999} \quad (1)$$

## Solution

We will use Geometric Mean  $\leq$  Arithmetic Mean, i.e. for non-negative  $x$ , and  $y$ ,

$$\sqrt{x \times y} \leq \frac{x + y}{2} \quad (2)$$

with equality iff  $x = y$ .

Proof:

$$(a - b)^2 \geq 0, \quad \text{with equality iff } a = b \quad (3)$$

$$\therefore a^2 + b^2 \geq 2ab \quad (4)$$

$$\therefore \frac{x + y}{2} \geq \sqrt{xy}, \quad \text{where } x = a^2, y = b^2 \quad (5)$$

Now split each term in  $n!$  into a  $\sqrt{\cdot}$  pair, rearrange and regroup, before applying the GM $\leq$ AM inequality on each:

$$n! = \sqrt{n \times 1} \quad \sqrt{(n-1) \times 2} \quad \dots \quad \sqrt{2 \times (n-1)} \quad \sqrt{1 \times n} \quad (6)$$

$$< \frac{n+1}{2} \quad \frac{n+1}{2} \quad \dots \quad \frac{n+1}{2} \quad \frac{n+1}{2} \quad (7)$$

$$= \left( \frac{n+1}{2} \right)^n \quad (8)$$

(the inequality  $\leq$  has become strict  $<$  because at least one of the term pairs are different).

Set  $n = 999$  to answer the problem with the affirmative:

$$999! < \left( \frac{999+1}{2} \right)^{999} = 500^{999} \quad (9)$$

## Discussion

How tight is this bound? Not very! It is made from a product of  $n$  terms, each larger than the term they replace. Further, the terms “on the ends” consist of pairs of numbers which are most different, and from the proof for the inequality one can see that these have the weakest bound (or conversely, the bound is tightest when the two numbers are most similar, becoming exact when the two numbers are equal).

Is there are a different way to arrange the the  $\sqrt{\cdot}$  pairs which creates a tighter bound?

The above approach makes pairs as follows:

$$\begin{array}{ccccccc} 1 & 2 & 3 & \dots & n-2 & n-1 & n \\ \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow & \downarrow \\ n & n-1 & n-2 & \dots & 3 & 2 & 1 \end{array}$$

We can bring the pairs numerically closer to one another by offsetting by 1,

$$\begin{array}{cccccc|c}
 1 & 2 & 3 & \dots & n-2 & n-1 & n \\
 \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow & \downarrow \\
 n-1 & n-2 & n-3 & \dots & 2 & 1 & n
 \end{array}$$

Or for example, by offsetting by 3,

$$\begin{array}{cccccc|ccc}
 1 & 2 & 3 & \dots & n-3 & & n-2 & n-1 & n \\
 \downarrow & \downarrow & \downarrow & & \downarrow & & \downarrow & \downarrow & \downarrow \\
 n-3 & n-4 & n-5 & \dots & 1 & & n & n-1 & n-2
 \end{array}$$

This creates two groups:

- $n-3$  pairs which individually sum to  $n-2$ , and
- 3 pairs which individually sum to  $2n-2$ .

So the GM-AM inequality approach produces

$$n! = \sqrt{(n-3) \times 1} \sqrt{(n-4) \times 2} \dots \sqrt{1 \times (n-3)} \quad (10)$$

$$\times \sqrt{n \times (n-2)} \sqrt{(n-1) \times (n-1)} \sqrt{(n-2) \times n} \quad (11)$$

$$< \left(\frac{n-2}{2}\right)^{n-3} \times \left(\frac{2n-2}{2}\right)^3 \quad (12)$$

If we offset by an arbitrary  $k$ ,

$$\begin{array}{cccccc|cccc}
 1 & 2 & 3 & \dots & n-k & & n-k+1 & \dots & n-1 & n \\
 \downarrow & \downarrow & \downarrow & & \downarrow & & \downarrow & & \downarrow & \downarrow \\
 n-k & n-k-1 & n-k-2 & \dots & 1 & & n & \dots & n-k+2 & n-k+1
 \end{array}$$

then we again create two groups:

- $n-k$  pairs which individually sum to  $n-k+1$
- $k$  pairs which individually sum to  $2n-k+1$

Hence the GM-AM inequality approach produces

$$n! < \left(\frac{n-k+1}{2}\right)^{n-k} \times \left(\frac{2n-k+1}{2}\right)^k \quad (13)$$

Splitting down the middle for an even  $n$ , i.e.  $n = 2k$ , then

$$n! < \left(\frac{n-\frac{n}{2}+1}{2}\right)^{\frac{n}{2}} \times \left(\frac{2n-\frac{n}{2}+1}{2}\right)^{\frac{n}{2}} \quad (14)$$

$$= \left(\frac{2n-n+2}{4}\right)^{\frac{n}{2}} \times \left(\frac{4n-n+2}{4}\right)^{\frac{n}{2}} \quad (15)$$

$$= \left(\frac{(n+2)(3n+2)}{16}\right)^{\frac{n}{2}} \quad (16)$$

For large  $n$ , this behaves like

$$\left(\frac{3n^2}{16}\right)^{\frac{n}{2}} = \frac{n^n}{\left(\frac{4}{\sqrt{3}}\right)^n} \quad (17)$$

which compares favorably to the large  $n$  behavior of the original bound,

$$\left(\frac{n+1}{2}\right)^n \approx \frac{n^n}{2^n}, \quad \text{for large } n \quad (18)$$

So our “large  $n$ ” upper bound for  $n!$  has reduced (improved) by a factor of

$$\left(\frac{n^n}{\left(\frac{4}{\sqrt{3}}\right)^n}\right) / \left(\frac{n^n}{2^n}\right) = \left(\frac{2\sqrt{3}}{4}\right)^n = \left(\frac{\sqrt{3}}{2}\right)^n \approx (0.87)^n \quad (19)$$

Could this be improved further? What is the smallest  $\alpha$  such that

$$n! \lesssim \alpha^n n^n \quad \text{for large } n \quad (20)$$

Restating without the  $\lesssim$ , what is the smallest  $\alpha$  such that

$$n! < \alpha^n (n^m + a_{m-1}n^{m-1} + \dots + a_1n + a_0)^{\frac{n}{m}} \quad (21)$$

$$= \alpha^n (P_m(n))^{\frac{n}{m}} \quad (22)$$

for some set of coefficients  $a_0 \dots a_{m-1}$  making the polynomial  $P_m$  of degree  $m$ . Note the coefficient of  $n^m$  in  $P_m$  is by definition equal to 1.

The solution to the original problem found an  $\alpha_1 = 0.5$  with  $P_1(n) = n + 1$ .

Our mid-point pairing approach reduced this to  $\alpha_2 = \sqrt{3}/4 \approx 0.43$  with  $P_2(n) = n^2 + \frac{8}{3}n + \frac{4}{3}$ .

The largest  $m$  is when  $m = n$ , so  $P_n(n) = n^n + \dots$  and is raised to the power of 1, i.e.,

$$n! < \alpha_n^n (n^n + a_{n-1}n^{n-1} + \dots + a_0) \quad (23)$$

*To be continued...*