

kathara lab

arp

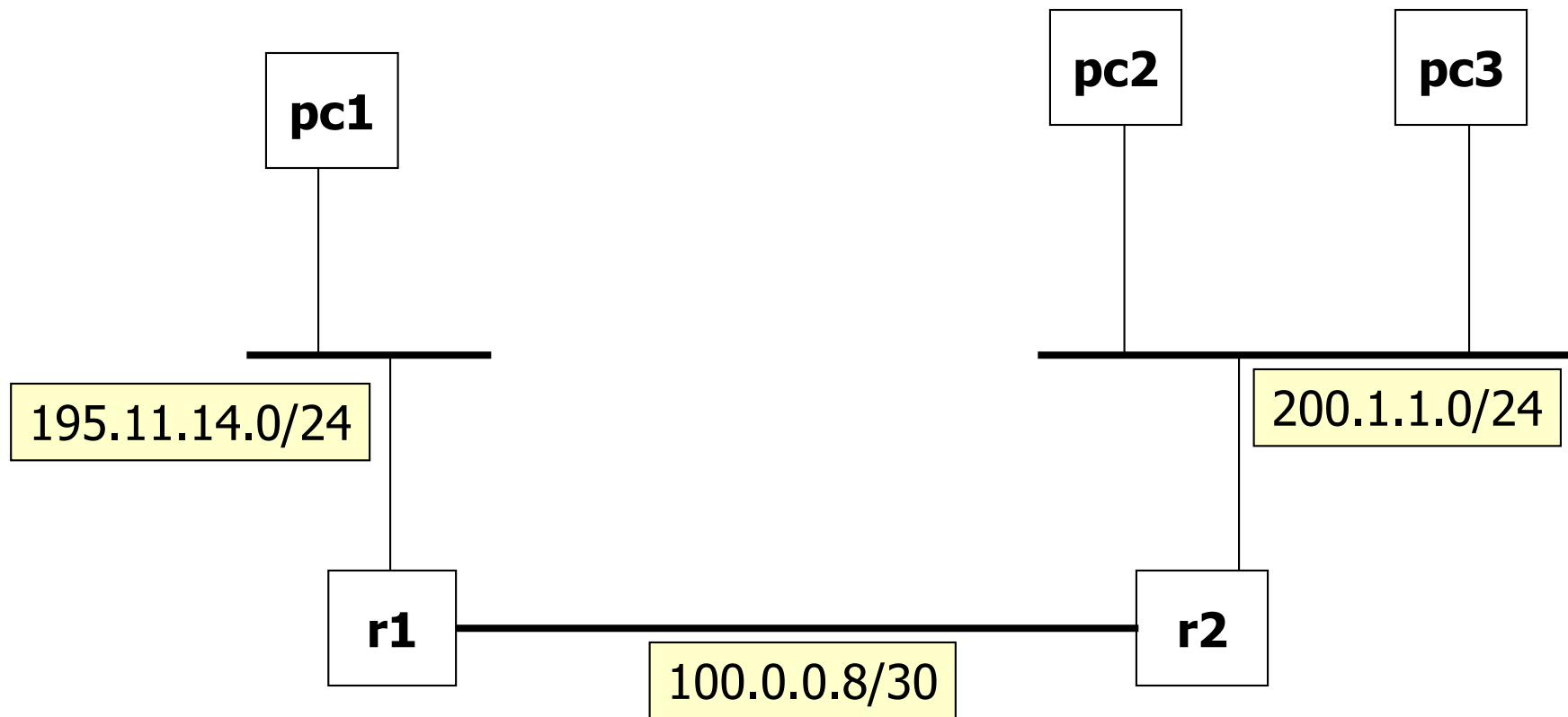
Version	1.0
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Web	http://www.kathara.org/
Description	using the address resolution protocol for local and non local traffic – kathara version of the netkit lab on arp version 2.2

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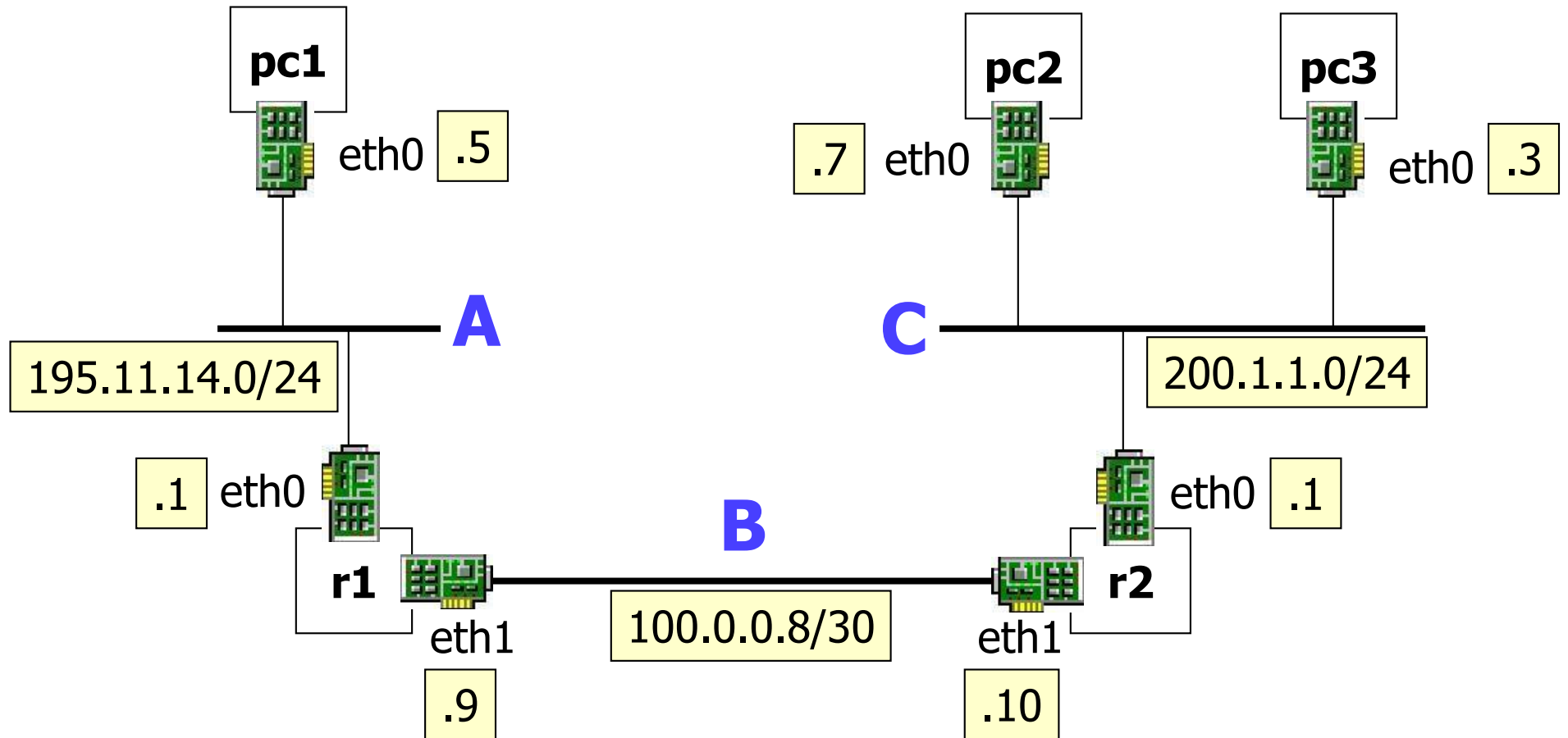
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step 1 – network topology

high level view



step 1 – network topology configuration details



step 2 – a quick look at the lab

lab.conf

```
r1[0]="A"  
r1[1]="B"  
  
r2[0]="C"  
r2[1]="B"  
  
pc1[0]="A"  
  
pc2[0]="C"  
  
pc3[0]="C"
```

pc1.startup

```
ifconfig eth0 195.11.14.5 up  
route add default gw 195.11.14.1
```

pc2.startup

```
ifconfig eth0 200.1.1.7 up  
route add default gw 200.1.1.1
```

pc3.startup

```
ifconfig eth0 200.1.1.3 up  
route add default gw 200.1.1.1
```

step 2 – a quick look at the lab

r1.startup

```
ifconfig eth0 195.11.14.1 up
ifconfig eth1 100.0.0.9 netmask 255.255.255.252 broadcast 100.0.0.11 up
route add -net 200.1.1.0 netmask 255.255.255.0 gw 100.0.0.10 dev eth1
```

r2.startup

```
ifconfig eth0 200.1.1.1 up
ifconfig eth1 100.0.0.10 netmask 255.255.255.252 broadcast 100.0.0.11 up
route add -net 195.11.14.0 netmask 255.255.255.0 gw 100.0.0.9 dev eth1
```

■ start the lab

▼ host machine

```
user@localhost:~$ cd kathara-lab_arp
user@localhost:~/kathara-lab_arp$ ./start
```

step 3 – inspecting the arp cache

ARP (8)

Linux Programmer's Manual

ARP (8)

NAME

`arp` - manipulate the system ARP cache

SYNOPSIS

```
arp [-vn] [-H type] [-i if] -a [hostname]
arp [-v] [-i if] -d hostname [pub]
arp [-v] [-H type] [-i if] -s hostname hw_addr [temp]
arp [-v] [-H type] [-i if] -s hostname hw_addr [netmask nm] pub
arp [-v] [-H type] [-i if] -Ds hostname ifa [netmask nm] pub
arp [-vnD] [-H type] [-i if] -f [filename]
```

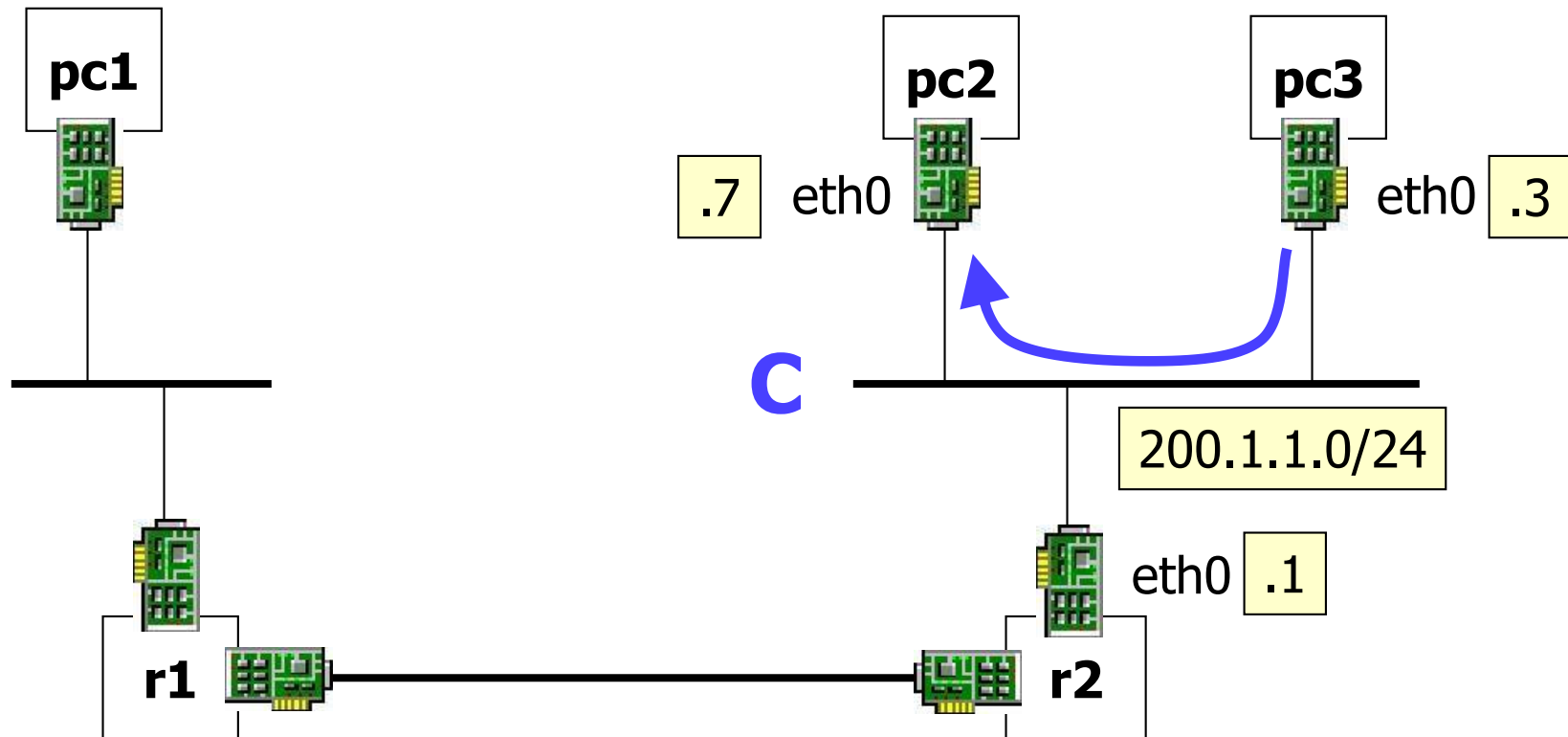
DESCRIPTION

`Arp` manipulates the kernel's ARP cache in various ways. The primary options are clearing an address mapping entry and manually setting up one. For debugging purposes, the `arp` program also allows a complete dump of the ARP cache.

.....

step 3 – inspecting the arp cache (local traffic)

- traffic within the same network does not traverse routers



step 3 – inspecting the arp cache (local traffic)

pc3

the arp cache is
initially empty

sending packets to
200.1.1.7 requires
address resolution

```
pc3:~# arp
pc3:~# ping 200.1.1.7
PING 200.1.1.7 (200.1.1.7) 56(84) bytes of data.
64 bytes from 200.1.1.7: icmp_seq=1 ttl=64 time=1.39 ms
64 bytes from 200.1.1.7: icmp_seq=2 ttl=64 time=0.542 ms

--- 200.1.1.7 ping statistics ---
2 packets transmitted, 2 received, 0% packet loss, time 1022ms
rtt min/avg/max/mdev = 0.542/0.969/1.396/0.427 ms
pc3:~# arp -n
```


Address	Hwtype	Hwaddress	Flags	Mask	Iface
200.1.1.7	ether	FE:FD:C8:01:01:07	C		eth0

```
pc3:~# █
```

address resolution
results are stored in
the arp cache

step 3 – inspecting the arp cache (local traffic)

- communications are usually bi-directional
- the receiver of the arp request learns the mac address of the other party, to avoid a new arp in opposite direction (standard behavior, see rfc 826)

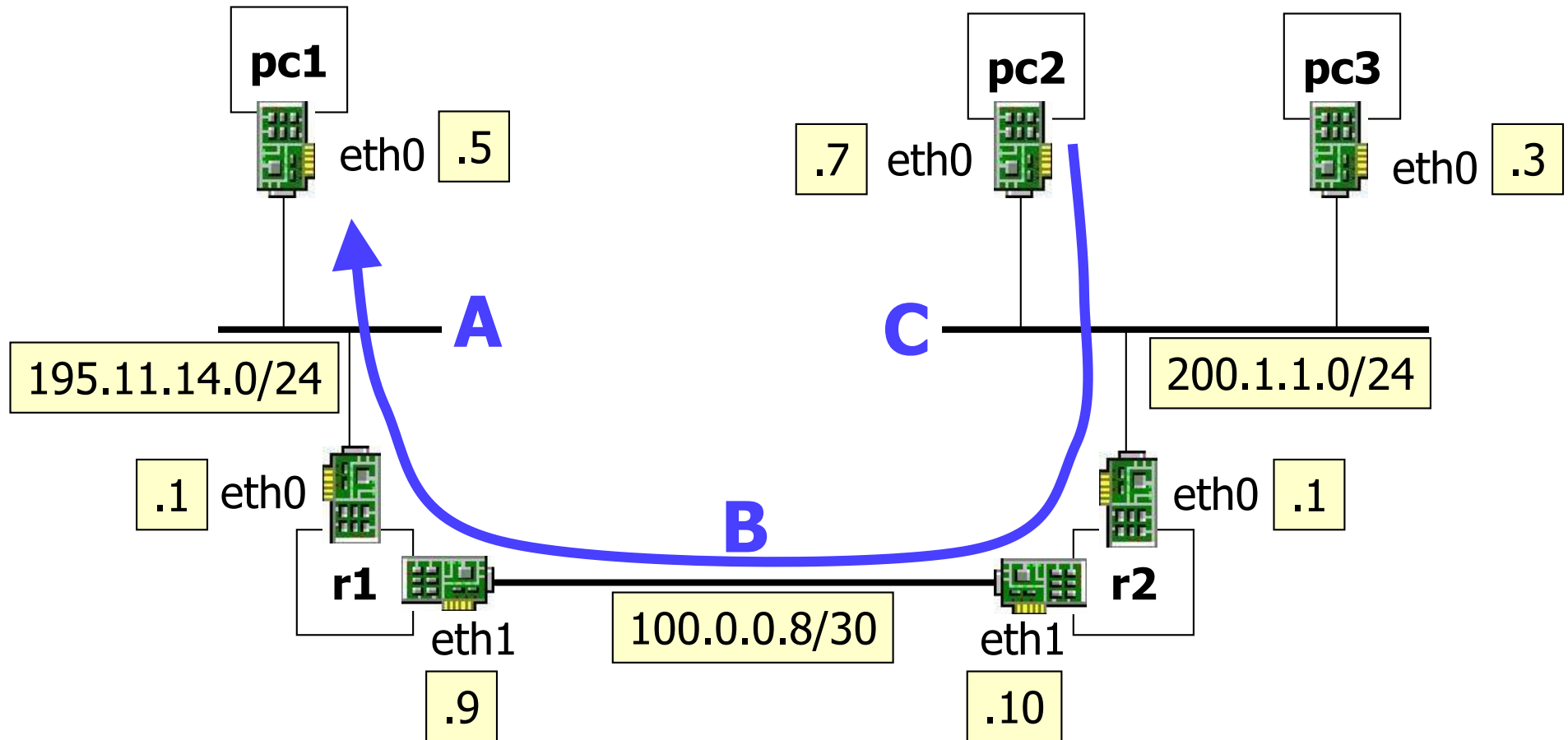


```
pc2:~# arp -n
```

Address	HWtype	HWaddress	Flags	Mask	Iface
200.1.1.3	ether	FE:FD:C8:01:01:03	C		eth0

```
pc2:~# █
```

step 4 – inspecting the arp cache (non local traffic)



step 4 – inspecting the arp cache (non local traffic)

- when ip traffic is addressed outside the local network, the sender needs the mac address of the router
- arp requests can get replies only within the local network

```
pc2:~# ping 195.11.14.5
PING 195.11.14.5 (195.11.14.5) 56(84) bytes of data.
64 bytes from 195.11.14.5: icmp_seq=1 ttl=62 time=30.4 ms
64 bytes from 195.11.14.5: icmp_seq=2 ttl=62 time=1.02 ms

--- 195.11.14.5 ping statistics ---
2 packets transmitted, 2 received, 0% packet loss, time 1013ms
rtt min/avg/max/mdev = 1.024/15.731/30.438/14.707 ms
pc2:~# arp -n
```

Address	Hwtype	Hwaddress	Flags	Mask	Iface
200.1.1.1	ether	FE:FD:C8:01:01:01	C		eth0
200.1.1.3	ether	FE:FD:C8:01:01:03	C		eth0

```
pc2:~# █
```

step 4 – inspecting the arp cache (non local traffic)

- what about routers?
- routers perform arp too (hence have arp caches)
anytime they have to send ip packets on an ethernet lan

r1

```
r1:~# arp -n
```

Address	Hwtype	Hwaddress	Flags	Mask	Iface
100.0.0.10	ether	FE:FD:64:00:00:0A	C		eth1
195.11.14.5	ether	FE:FD:C3:0B:0E:05	C		eth0

r1:~# █

Callouts: r2 (eth1), pc1

r2

```
r2:~# arp -n
```

Address	Hwtype	Hwaddress	Flags	Mask	Iface
200.1.1.7	ether	FE:FD:C8:01:01:07	C		eth0
100.0.0.9	ether	FE:FD:64:00:00:09	C		eth1

r2:~# █

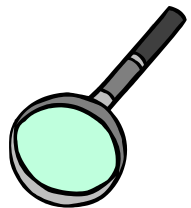
Callouts: pc2, r1 (eth1)

step5 – sniffing arp traffic

- restart the lab in order to clear arp caches

```
host machine
user@localhost:~/kathara-lab_arp$ lclean
user@localhost:~/kathara-lab_arp$ lstart
```

- get ready to sniff



=

```
virtual machine
vm:~# tcpdump -t -i eth0
```

sniff on this interface

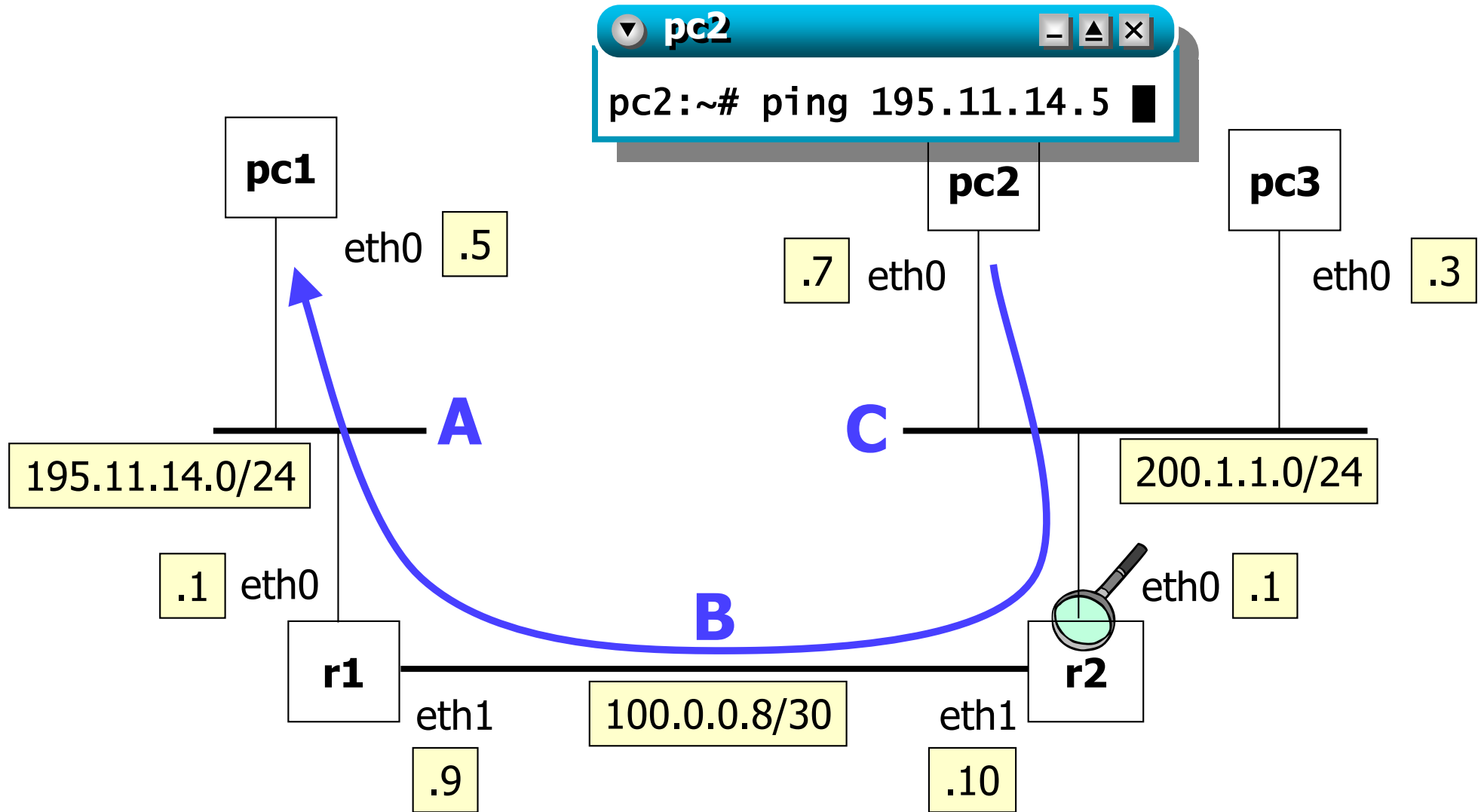
No DNS look up

No protocol name

suppress timestamps (kathara is not a simulation system, hence timestamps are not meaningful)

show link-level headers (mac addresses)

step 5 – sniffing arp traffic



step 5 – sniffing arp traffic

- on collision domain C

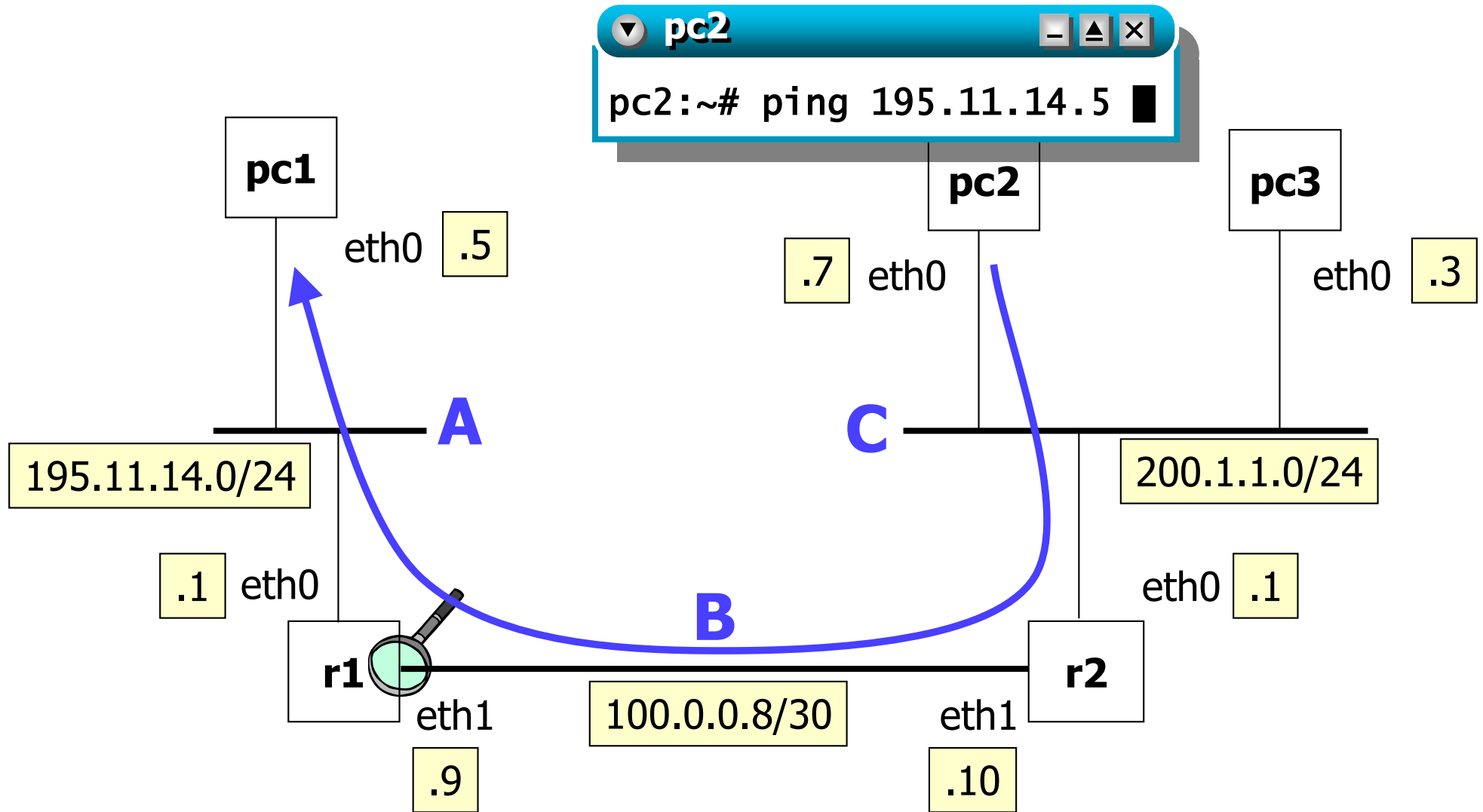
▼ r2



```
r2:~# tcpdump -t -i eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 96 bytes
fe:fd:c8:01:01:07 > Broadcast, ethertype ARP (0x0806), length 42: arp who-has
200.1.1.1 tell 200.1.1.7
fe:fd:c8:01:01:01 > fe:fd:c8:01:01:07, ethertype ARP (0x0806), length 42: arp reply
200.1.1.1 is-at fe:fd:c8:01:01:01
fe:fd:c8:01:01:07 > fe:fd:c8:01:01:01, ethertype IPv4 (0x0800), length 98: IP
200.1.1.7 > 195.11.14.5: icmp 64: echo request seq 1
fe:fd:c8:01:01:01 > fe:fd:c8:01:01:07, ethertype IPv4 (0x0800), length 98: IP
195.11.14.5 > 200.1.1.7: icmp 64: echo reply seq 1
```

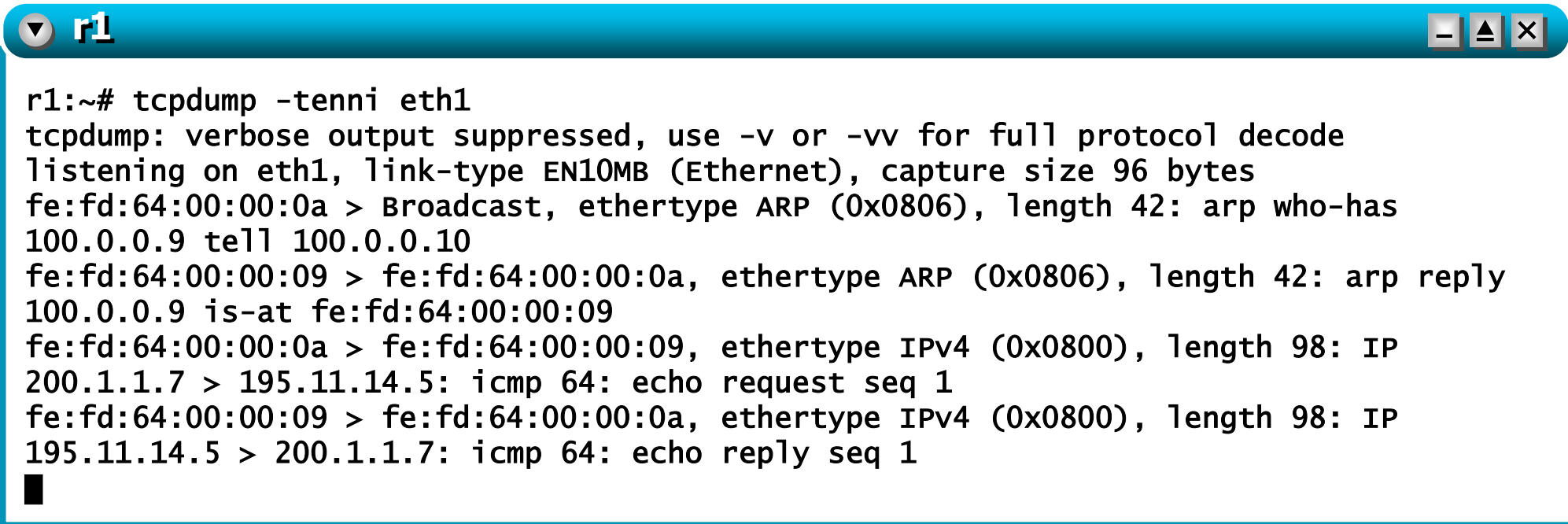
1. **pc2** asks all the stations on collision domain C: "who has 200.1.1.1?" (200.1.1.1 is **pc2**'s default gateway)
2. **r2** replies \Rightarrow both **pc2** and **r2** update their arp cache
3. **pc2** sends to **r2** the ip packet (icmp echo request) for **pc1**
4. **r2** sends to **pc2** the corresponding echo reply (generated by **pc1**)

step 5 – sniffing arp traffic



step 6 – sniffing arp traffic

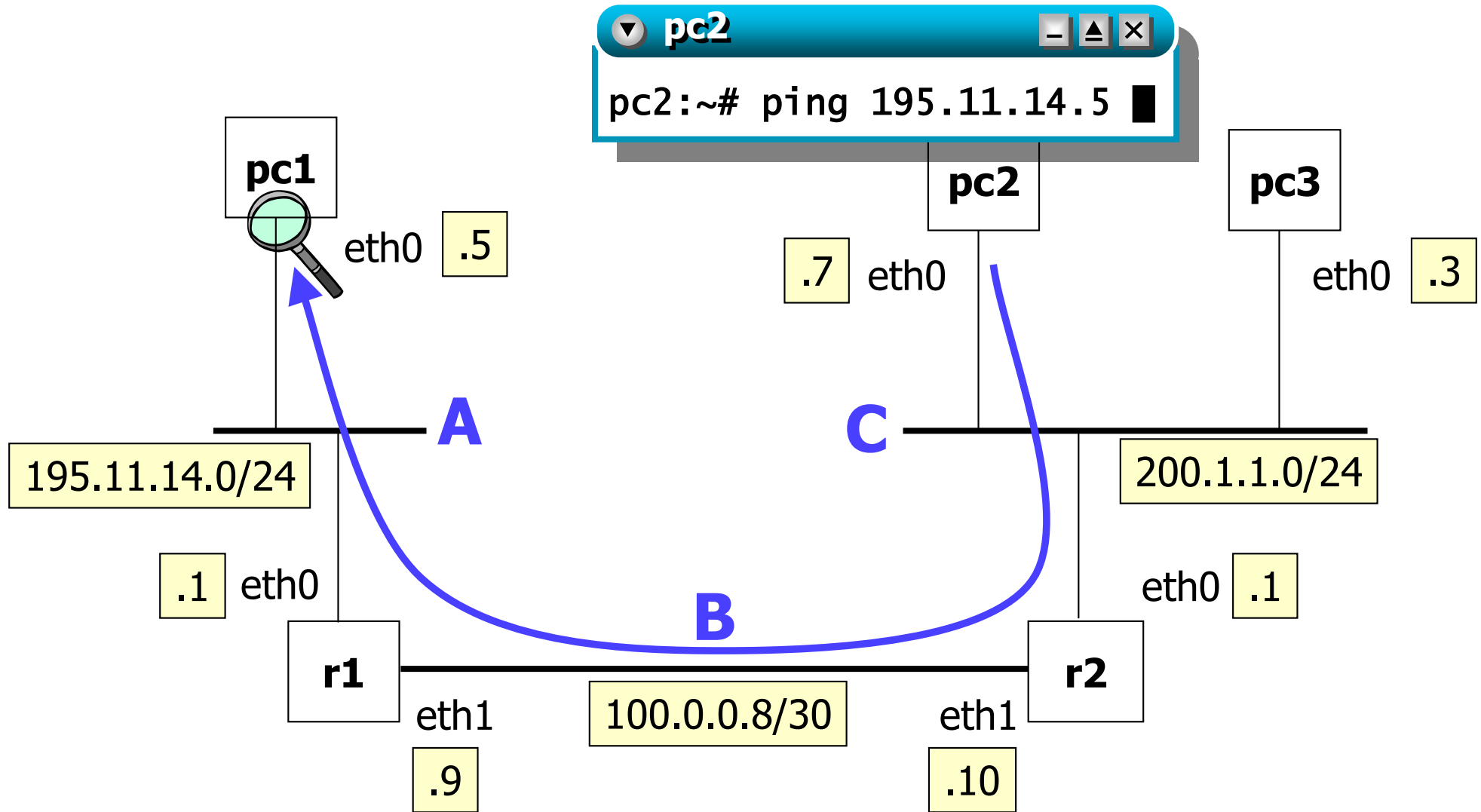
- on collision domain B



```
r1:~# tcpdump -t -n -i eth1
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth1, link-type EN10MB (Ethernet), capture size 96 bytes
fe:fd:64:00:00:0a > Broadcast, ethertype ARP (0x0806), length 42: arp who-has
100.0.0.9 tell 100.0.0.10
fe:fd:64:00:00:09 > fe:fd:64:00:00:0a, ethertype ARP (0x0806), length 42: arp reply
100.0.0.9 is-at fe:fd:64:00:00:09
fe:fd:64:00:00:0a > fe:fd:64:00:00:09, ethertype IPv4 (0x0800), length 98: IP
200.1.1.7 > 195.11.14.5: icmp 64: echo request seq 1
fe:fd:64:00:00:09 > fe:fd:64:00:00:0a, ethertype IPv4 (0x0800), length 98: IP
195.11.14.5 > 200.1.1.7: icmp 64: echo reply seq 1
■
```

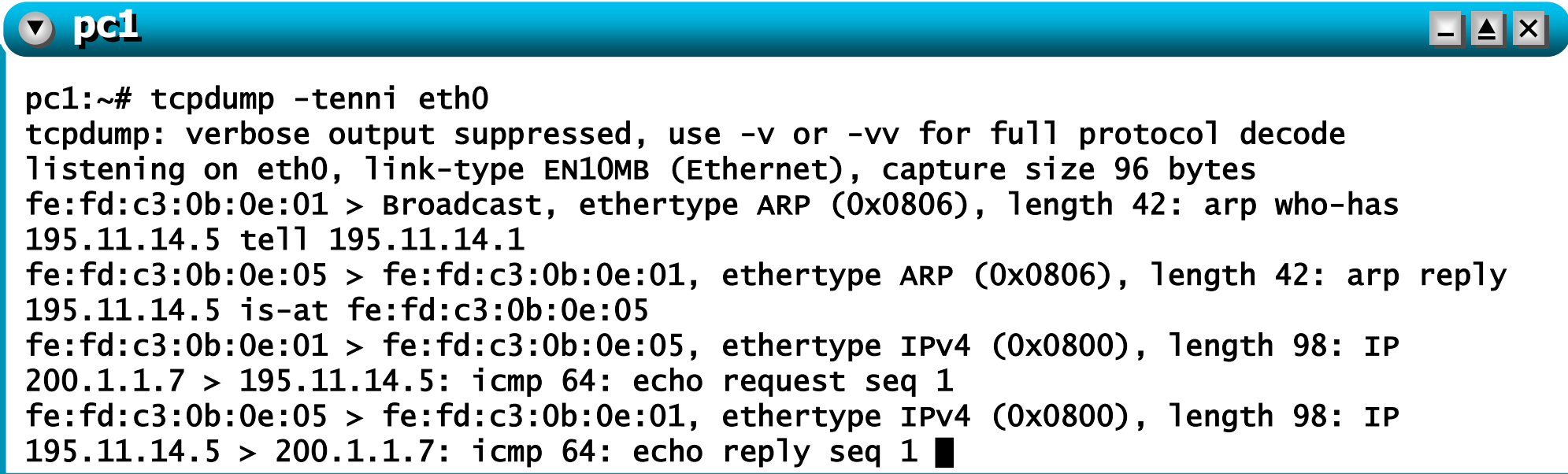
1. **r2** asks all the stations on collision domain B: “who has 100.0.0.9?” (100.0.0.9 is the next hop obtained from the routing table)
2. **r1** replies \Rightarrow both **r1** and **r2** update their arp cache
3. **r2** sends to **r1** the echo request generated by **pc2** for **pc1**
4. **r1** sends to **r2** the echo reply generated by **pc1** for **pc2**

step 5 – sniffing arp traffic



step 5 – sniffing arp traffic

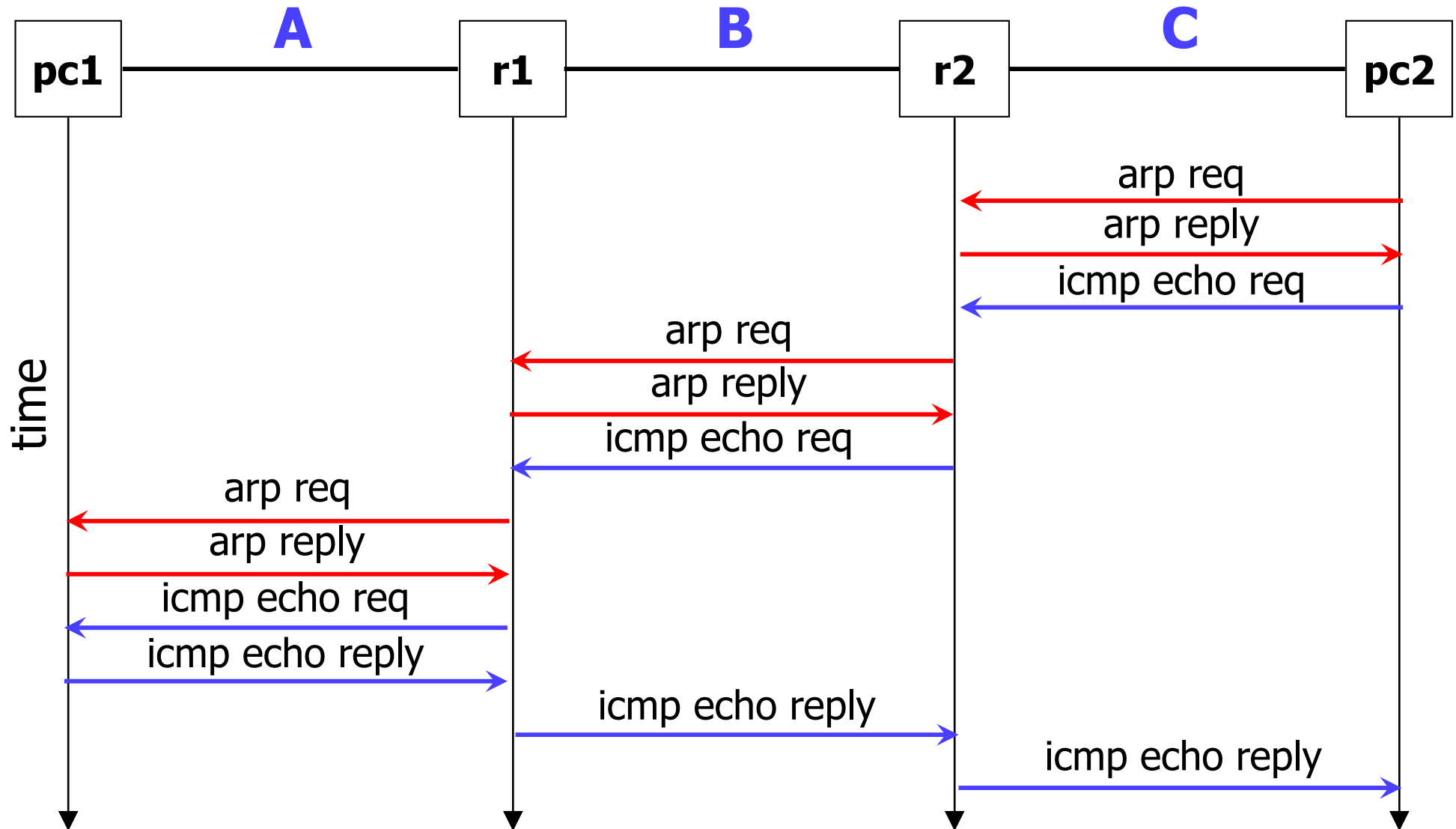
- on collision domain A



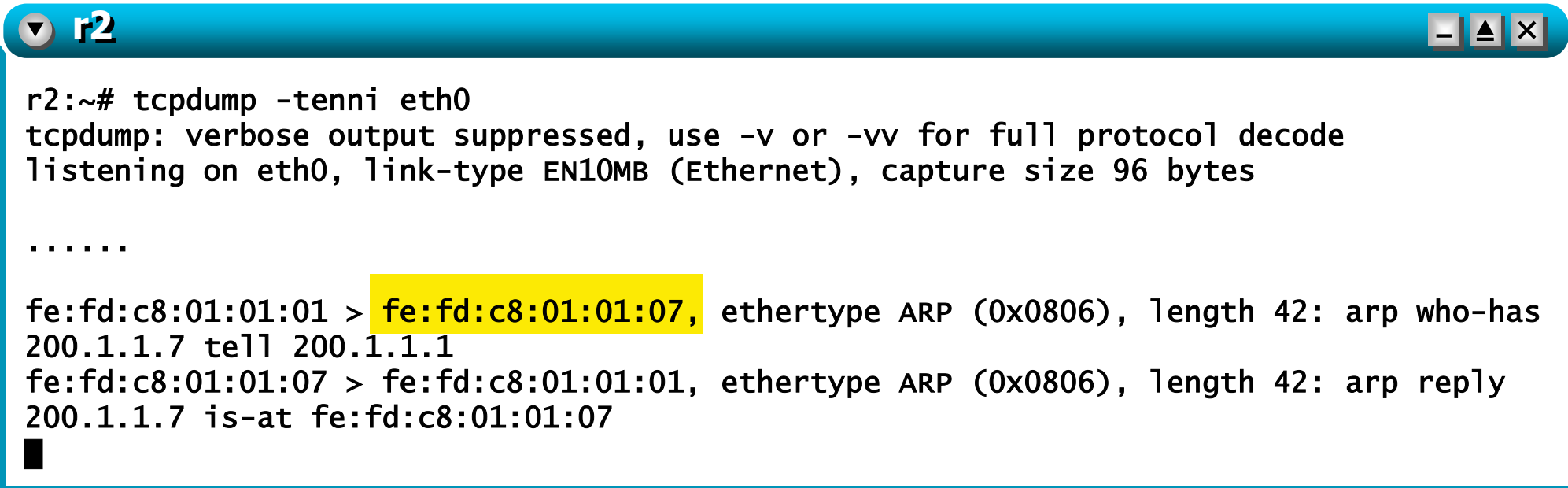
```
pc1:~# tcpdump -ttni eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 96 bytes
fe:fd:c3:0b:0e:01 > Broadcast, ethertype ARP (0x0806), length 42: arp who-has
195.11.14.5 tell 195.11.14.1
fe:fd:c3:0b:0e:05 > fe:fd:c3:0b:0e:01, ethertype ARP (0x0806), length 42: arp reply
195.11.14.5 is-at fe:fd:c3:0b:0e:05
fe:fd:c3:0b:0e:01 > fe:fd:c3:0b:0e:05, ethertype IPv4 (0x0800), length 98: IP
200.1.1.7 > 195.11.14.5: icmp 64: echo request seq 1
fe:fd:c3:0b:0e:05 > fe:fd:c3:0b:0e:01, ethertype IPv4 (0x0800), length 98: IP
195.11.14.5 > 200.1.1.7: icmp 64: echo reply seq 1 ■
```

1. **r1** asks all the stations on collision domain A: “who has 195.11.14.5?” (195.11.14.5 is the destination address of the icmp request obtained from the ip header)
2. **pc1** replies \Rightarrow both **pc1** and **r1** update their arp cache
3. **r1** sends the ip packet (echo request) to **pc1**
4. **pc1** generates the corresponding echo reply for **pc2** and sends it to **r1**

step 6 – understanding the whole picture



step 7 – arp implementation details



```
r2:~# tcpdump -ttni eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth0, link-type EN10MB (Ethernet), capture size 96 bytes

.....

fe:fd:c8:01:01:01 > fe:fd:c8:01:01:07, ethertype ARP (0x0806), length 42: arp who-has
200.1.1.7 tell 200.1.1.1
fe:fd:c8:01:01:07 > fe:fd:c8:01:01:01, ethertype ARP (0x0806), length 42: arp reply
200.1.1.7 is-at fe:fd:c8:01:01:07
█
```

- arp requests are usually in broadcast
- it may also happen that a station (router/pc) sends a unicast arp request to check if an entry of the arp cache is still valid (suggested by the standard, see rfc 826)
- unicast arp requests may be performed periodically on each entry of the arp cache, depending on the implementation

proposed exercises

- what packets can we observe over all the three collision domains as the ping from **pc2** to **pc1** continues (ignore any implementation dependent arp behavior)?

proposed exercises

- check the different error messages obtained by trying to ping an unreachable destination in the case of
 - local destination
 - non local destination
- which packets are exchanged in the local collision domain in the two cases?