

Compilers

CS414-2015S-01

Compiler Basics & Lexical Analysis

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01-0: Syllabus

- Office Hours
- Course Text
- Prerequisites
- Test Dates & Testing Policies
- Projects
 - Teams of up to 2
- Grading Policies
- Questions?

01-1: Notes on the Class

- Don't be afraid to ask me to slow down!
- We will cover some pretty complex stuff here, which can be difficult to get the first (or even the second) time. *ASK QUESTIONS*
- While specific questions are always preferred, “I don't get it” is always an acceptable question. I am always happy to stop, re-explain a topic in a different way.
 - If you are confused, I can *guarantee* that at least one other person in the class would benefit from more explanation

01-2: Notes on the Class

- Projects are non-trivial
 - Using new tools (JavaCC)
 - Managing a large scale project
 - Lots of complex classes & advanced programming techniques.

01-3: Notes on the Class

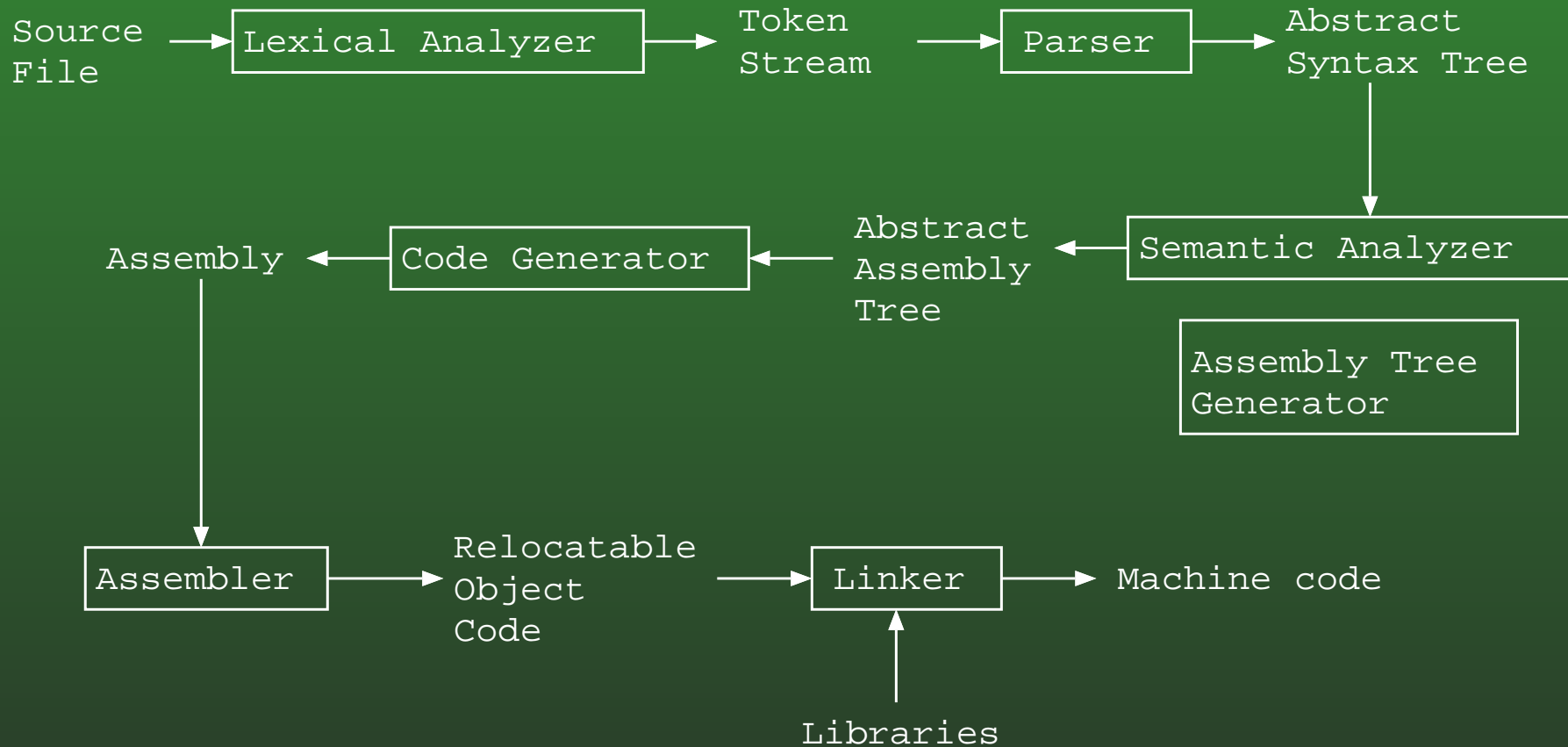
- Projects are non-trivial
 - Using new tools (JavaCC)
 - Managing a large scale project
 - Lots of complex classes & advanced programming techniques.
- *START EARLY!*
 - Projects will take longer than you think (especially starting with the semantic analyzer project)
- *ASK QUESTIONS!*

01-4: What is a compiler?



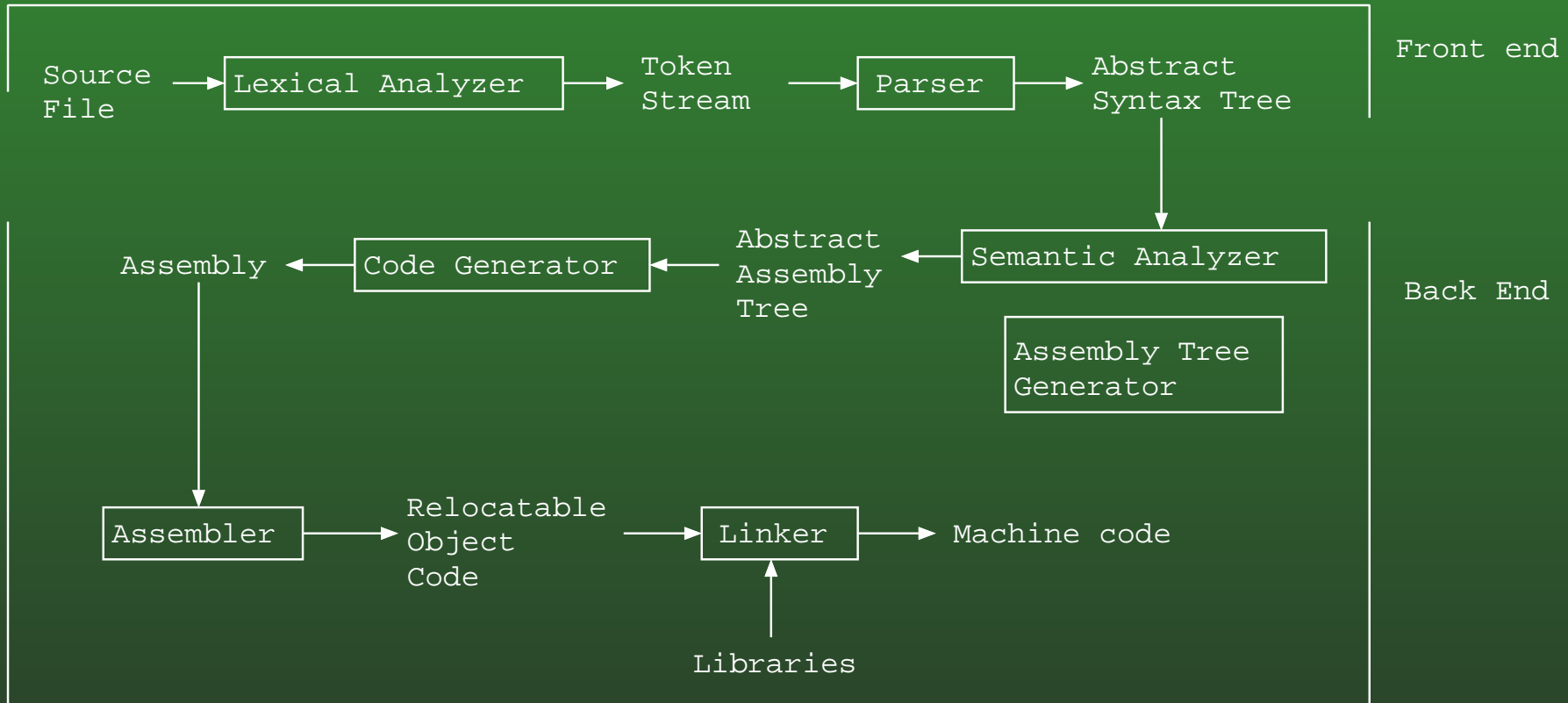
Simplified View

01-5: What is a compiler?

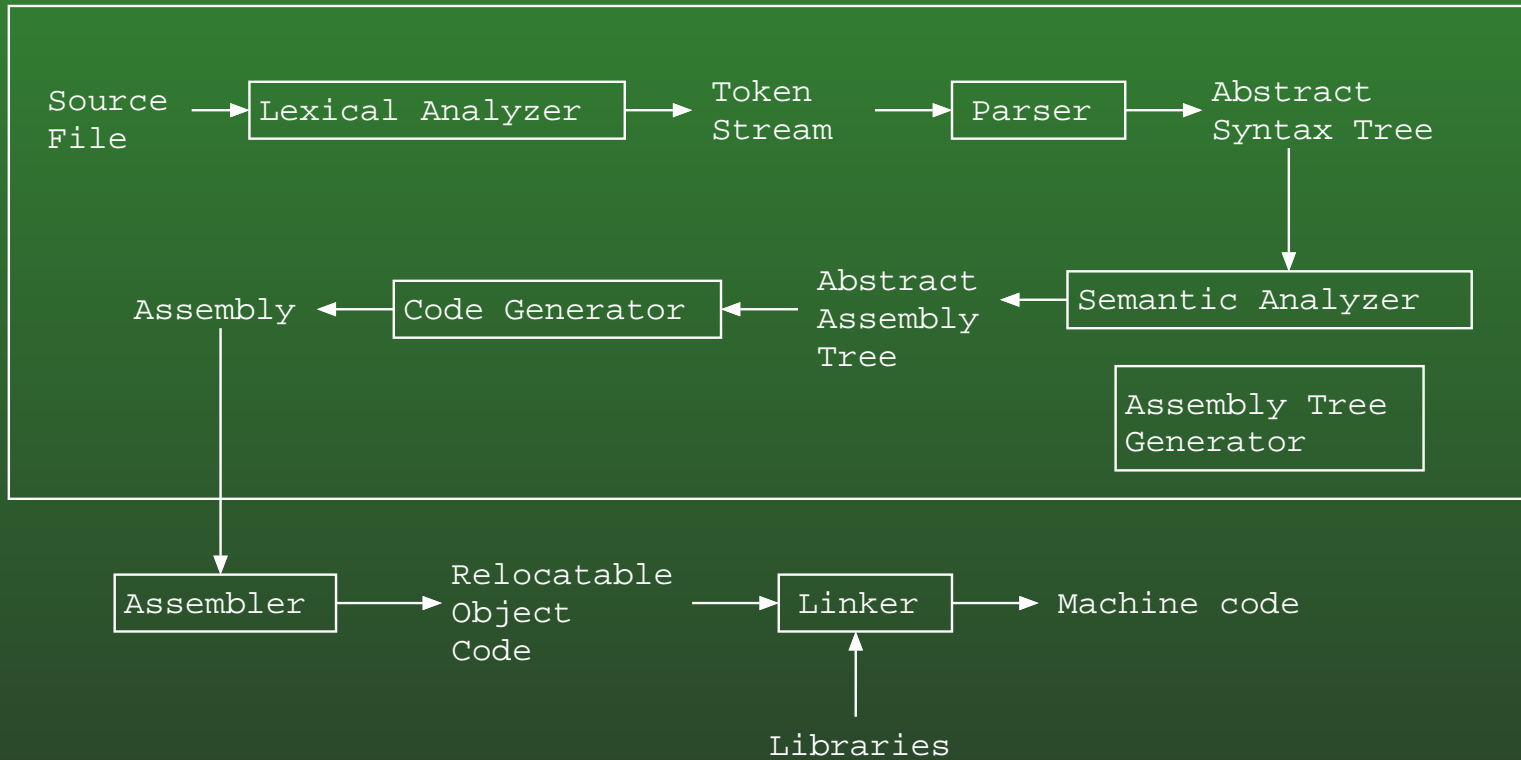


More Accurate View

01-6: What is a compiler?



01-7: What is a compiler?



Covered in
this course

01-8: Why Use Decomposition?

01-9: Why Use Decomposition?

Software Engineering!

- Smaller units are easier to write, test and debug
- Code Reuse
 - Writing a suite of compilers (C, Fortran, C++, etc) for a new architecture
 - Create a new language – want compilers available for several platforms

01-10: Lexical Analysis

- Converting input file to stream of tokens

```
void main() {  
    print(4);  
}
```

01-11: Lexical Analysis

- Converting input file to stream of tokens

void main() {	IDENTIFIER(void)
print(4);	IDENTIFIER(main)
}	LEFT-PARENTHESIS
	RIGHT-PARENTHESIS
	LEFT-BRACE
	IDENTIFIER(print)
	LEFT-PARENTHESIS
	INTEGER-LITERAL(4)
	RIGHT-PARENTHESIS
	SEMICOLON
	RIGHT-BRACE

01-12: Lexical Analysis

Brute-Force Approach

- Lots of nested if statements

```
if (c = nextchar() == 'P') {  
    if (c = nextchar() == 'R') {  
        if (c = nextchar() == 'O') {  
            if (c = nextchar() == 'G') {  
                /* Code to handle the rest of either  
                 PROGRAM or any identifier that starts  
                 with PROG  
                */  
            } else if (c == 'C') {  
                /* Code to handle the rest of either  
                 PROCEDURE or any identifier that starts  
                 with PROC  
                */  
            }  
        }  
    }  
}
```

...

01-13: Lexical Analysis

Brute-Force Approach

- Break the input file into words, separated by spaces or tabs
 - This can be tricky – not all tokens are separated by whitespace
 - Use string comparison to determine tokens

01-14: Deterministic Finite Automata

- Set of states
- Initial State
- Final State(s)
- Transitions

DFA for else, end, identifiers

Combine DFA

01-15: DFAs and Lexical Analyzers

- Given a DFA, it is easy to create C code to implement it
- DFAs are easier to understand than C code
 - Visual – almost like structure charts
- ... However, creating a DFA for a complete lexical analyzer is still complex

01-16: Automatic Creation of DFAs

We'd like a tool:

- Describe the tokens in the language
- Automatically create DFA for tokens
- Then, automatically create C code that implements the DFA

We need a method for describing tokens

01-17: Formal Languages

- **Alphabet Σ :** Set of all possible symbols (characters) in the input file
 - Think of Σ as the set of symbols on the keyboard
- **String w :** Sequence of symbols from an alphabet
- **String length $|w|$** Number of characters in a string: $|\text{car}| = 3$, $|\text{abba}| = 4$
 - **Empty String ϵ :** String of length 0: $|\epsilon| = 0$
- **Formal Language:** Set of strings over an alphabet

Formal Language \neq Programming language – Formal Language is only a set of strings.

01-18: Formal Languages

Example formal languages:

- Integers $\{0, 23, 44, \dots\}$
- Floating Point Numbers $\{3.4, 5.97, \dots\}$
- Identifiers $\{\text{foo}, \text{bar}, \dots\}$

01-19: Language Concatenation

- **Language Concatenation** Given two formal languages L_1 and L_2 , the concatenation of L_1 and L_2 , $L_1 L_2 = \{xy | x \in L_1, y \in L_2\}$

For example:

$\{\text{fire, truck, car}\} \{\text{car, dog}\} =$
 $\{\text{firecar, firedog, truckcar, truckdog, carcar, cardog}\}$

01-20: Kleene Closure

Given a formal language L :

$$L^0 = \{\epsilon\}$$

$$L^1 = L$$

$$L^2 = LL$$

$$L^3 = LLL$$

$$L^4 = LLLL$$

$$L^* = L^0 \cup L^1 \cup L^2 \cup \dots \cup L^n \cup \dots$$

01-21: Regular Expressions

Regular expressions are use to describe formal languages over an alphabet Σ :

Regular Expression	Language
--------------------	----------

ϵ	$L[\epsilon] = \{\epsilon\}$
------------	------------------------------

$a \in \Sigma$	$L[a] = \{a\}$
----------------	----------------

(MR)	$L[MR] = L[M]L[R]$
--------	--------------------

$(M R)$	$L[(M R)] = L[M] \cup L[R]$
---------	-----------------------------

(M^*)	$L[(M^*)] = L[M]^*$
---------	---------------------

01-22: r.e. Precedence

From highest to Lowest:

Kleene Closure *

Concatenation

Alternation |

$$ab^*c|e = (a(b^*)c) \mid e$$

01-23: Regular Expression Examples

all strings over $\{a,b\}$

binary integers (with leading zeroes)

all strings over $\{a,b\}$ that

begin and end with a

all strings over $\{a,b\}$ that

contain aa

all strings over $\{a,b\}$ that

do not contain aa

01-24: Regular Expression Examples

all strings over $\{a,b\}$

$(a|b)^*$

binary integers (with leading zeroes)

$(0|1)(0|1)^*$

all strings over $\{a,b\}$ that

$a(a|b)^*a$

begin and end with a

all strings over $\{a,b\}$ that

$(a|b)^*aa(a|b)^*$

contain aa

all strings over $\{a,b\}$ that

$b^*(abb^*)^*(a|\epsilon)$

do not contain aa

01-25: Reg. Exp. Shorthand

$[a,b,c,d] = (a|b|c|d)$

$[d-g] = [d,e,f,g] = (b|e|f|g)$

$[d-f,M-O] = [d,e,f,M,N,O]$
 $= (d|e|f|M|N|O)$

$(\alpha)? = \text{Optionally } \alpha \text{ (i.e., } (\alpha | \epsilon))$

$(\alpha)^+ = \alpha(\alpha)^*$

01-26: Regular Expressions & Unix

- Many unix tools use regular expressions
- Example: `grep '<reg exp>' filename`
 - Prints all lines that contain a match to the regular expression
 - Special characters:
 - `^` beginning of line
 - `$` end of line
 - (grep examples on other screen)

01-27: JavaCC Regular Expressions

- All characters & strings must be in quotation marks
 - "else"
 - "+"
 - ("a" | "b")
- All regular expressions involving * must be parenthesized
 - ("a")*, not "a"*

01-28: JavaCC Shorthand

`["a","b","c","d"]` = `("a"|"b"|"c"|"d")`
`["d"-"g"]` = `["d","e","f","g"]` = `("b"|"e"|"f"|"g")`
`["d"-"f","M"-"O"]` = `["d","e","f","M","N","O"]`
= `("d"|"e"|"f"|"M"|"N"|"O")`
`(α)?` = Optionally α (i.e., `(α | ϵ)`)
`(α)+` = $\alpha(\alpha)^*$
`(~["a","b"])` = Any character *except* "a" or "b".
Can only be used with `[]` notation
`~(a(a|b)*b)` is not legal

01-29: r.e. Shorthand Examples

Regular Expression	Language
	{if}
	Set of legal identifiers
	Set of integer literals
	(leading zeroes allowed)
	Set of real literals

01-30: r.e. Shorthand Examples

Regular Expression	Language
"if"	{if}
["a"-"z"](["0"-"9","a"-"z"])*	Set of legal identifiers
["0"-"9"]	Set of integer literals (leading zeroes allowed)
(["0"-"9"]+"."(["0"-"9"]*)) ((["0"-"9"])*"."["0"-"9"]+)	Set of real literals

01-31: Lexical Analyzer Generator

JavaCC is a Lexical Analyzer Generator and a Parser Generator

- Input: Set of regular expressions (each of which describes a type of token in the language)
- Output: A lexical analyzer, which reads an input file and separates it into tokens

01-32: Structure of a JavaCC file

```
options{
    /* Code to set various options flags */
}

PARSER_BEGIN(foo)

public class foo {
    /* This segment is often empty */
}

PARSER_END(foo)

TOKEN_MGR_DECLS :
{
    /* Declarations used by lexical analyzer */
}

/* Token Rules & Actions */
```

01-33: Token Rules in JavaCC

- Tokens are described by rules with the following syntax:

```
TOKEN :  
{  
    <TOKEN_NAME: RegularExpression>  
}
```

- TOKEN_NAME is the name of the token being described
- RegularExpression is a regular expression that describes the token

01-34: Token Rules in JavaCC

- Token rule examples:

```
TOKEN :  
{  
    <ELSE: "else">  
}
```

```
TOKEN :  
{  
    <INTEGER_LITERAL: (["0"-"9"])+>  
}
```

01-35: Token Rules in JavaCC

- Several different tokens can be described in the same TOKEN block, with token descriptions separated by |.

```
TOKEN :  
{  
    <ELSE: "else">  
|    <INTEGER_LITERAL: ("0"-"9")+>  
|    <SEMICOLON: ";">  
}
```

01-36: getNextToken

- When we run javacc on the input file `foo.jj`, it creates the class `fooTokenManager`
- The class `fooTokenManager` contains the static method `getNextToken()`
- Every call to `getNextToken()` returns the next token in the input stream.

01-37: getNextToken

- When getNextToken is called, a regular expression is found that matches the next characters in the input stream.
- What if more than one regular expression matches?

```
TOKEN :  
{  
    <ELSE: "else">  
|    <IDENTIFIER: ("a"-"z")+>  
}
```

01-38: getNextToken

- When more than one regular expression matches the input stream:
 - Use the longest match
 - “elsed” should match to IDENTIFIER, not to ELSE followed by the identifier “d”
 - If two matches have the same length, use the rule that appears first in the .jj file
 - “else” should match to ELSE, not IDENTIFIER

01-39: JavaCC Example

```
PARSER_BEGIN(simple)
public class simple {

}
PARSER_END(simple)
```

```
TOKEN :
{
    <ELSE: "else">
|    <SEMICOLON: ";">
|    <FOR: "for">
|    <INTEGER_LITERAL: ("0"-"9")+>
|    <IDENTIFIER: ["a"-"z"](["a"-"z", "0"-"9"])*>
}
```

else;ford for

01-40: SKIP Rules

- Tell JavaCC what to ignore (typically whitespace) using SKIP rules
- SKIP rule is just like a TOKEN rule, except that no TOKEN is returned.

```
SKIP:
{
    < regularexpression1 >
|
    < regularexpression2 >
|
    ...
|
    < regularexpressionn >
}
```

01-41: Example SKIP Rules

```
PARSER_BEGIN(simple2)
public class simple2 {
}
PARSER_END(simple2)
```

```
SKIP :
{
    < " " >
|
    < "\n" >
|
    < "\t" >
}
```

```
TOKEN :
{
    <ELSE: "else">
|
    <SEMICOLON: ";">
|
    <FOR: "for">
|
    <INTEGER_LITERAL: ("0"-"9")+>
|
    <IDENTIFIER: ["A"-"Z"](["A"-"Z", "0"-"9"])*>
}
```

01-42: JavaCC States

- Comments can be dealt with using SKIP rules
- How could we skip over 1-line C++ Style comments?

```
// This is a comment
```

01-43: JavaCC States

- Comments can be dealt with using SKIP rules
- How we could skip over 1-line C++ Style comments:

```
// This is a comment
```

- Using a SKIP rule

```
SKIP :  
{  
    < "//" (~["\n"])* "\n" >  
}
```

01-44: JavaCC States

- Writing a regular expression to match multi-line comments (using `/*` and `*/`) is much more difficult
- Writing a regular expression to match nested comments is impossible (take Automata Theory for a proof :))
- What can we do?
 - Use JavaCC States

01-45: JavaCC States

- We can label each TOKEN and SKIP rule with a “state”
- Unlabeled TOKEN and SKIP rules are assumed to be in the default state (named DEFAULT, unsurprisingly enough)
- Can switch to a new state after matching a TOKEN or SKIP rule using the : NEWSTATE notation

01-46: JavaCC States

```
SKIP :
{
    < " " >
|
    < "\\n" >
|
    < "\\t" >
}

SKIP :
{
    < "/*" > : IN_COMMENT
}

<IN_COMMENT>

SKIP :
{
    < "*/" > : DEFAULT
|
    < ~[] >
}

TOKEN :
{
    <ELSE: "else">
|
    ... (etc)
}
```


01-47: **Actions in TOKEN & SKIP**

- We can add Java code to any SKIP or TOKEN rule
- That code will be executed when the SKIP or TOKEN rule is matched.
- Any methods / variables defined in the TOKEN_MGR_DECLS section can be used by these actions

01-48: Actions in TOKEN & SKIP

```
PARSER_BEGIN(remComments)
public class remComments { }
PARSER_END(remComments)
```

```
TOKEN_MGR_DECLS :
{
    public static int numcomments = 0;
}
```

```
SKIP :
{
    < "/*" > : IN_COMMENT
}
```

```
SKIP :
{
    < "//" (~["\n"])* "\n" > { numcomments++; }
}
```

01-49: Actions in TOKEN & SKIP

```
<IN_COMMENT>
SKIP :
{
    < "*/" > { numcomments++; SwitchTo(DEFAULT);}
}
```

```
<IN_COMMENT>
SKIP :
{
    < ~[] >
}
```

```
TOKEN :
{
    <ANY: ~[]>
}
```

01-50: Tokens

- Each call to getNextToken returns a “Token” object
- Token class is automatically created by javaCC.
- Variables of type Token contain the following public variables:
 - `public int kind;` The type of token. When javacc is run on the file `foo.jj`, a file `fooConstants.java` is created, which contains the symbolic names for each constant

```
public interface simplejavaConstants {  
    int EOF = 0;  
    int CLASSSS = 8;  
    int D0 = 9;  
    int ELSE = 10;  
    ...  
}
```

01-51: Tokens

- Each call to getNextToken returns a “Token” object
- Token class is automatically created by javaCC.
- Variables of type Token contain the following public variables:
 - `public int beginLine, beginColumn, endLine, endColumn;` The location of the token in the input file

01-52: Tokens

- Each call to getNextToken returns a “Token” object
- Token class is automatically created by javaCC.
- Variables of type Token contain the following public variables:
 - `public String image;` The text that was matched to create the token.

01-53: Generated TokenManager

```
class TokenTest {
    public static void main(String args[]) {
        Token t;
        Java.io.InputStream infile;
        pascalTokenManager tm;
        boolean loop = true;

        if (args.length < 1) {
            System.out.print("Enter filename as command line argument");
            return;
        }
        try {
            infile = new Java.io.FileInputStream(args[0]);
        } catch (Java.io.FileNotFoundException e) {
            System.out.println("File " + args[0] + " not found.");
            return;
        }
        tm = new sjavaTokenManager(new SimpleCharStream(infile));
```

01-54: Generated TokenManager

```
t = tm.getNextToken();
while(t.kind != sjavaConstants.EOF) {
    System.out.println("Token : "+ t + " : ");
    System.out.println(pascalConstants.tokenImage[t.kind]);
}
}
}
```


01-55: **Lexer Project**

- Write a .jj file for simpleJava tokens
- Need to handle all whitespace (tabs, spaces, end-of-line)
- Need to handle nested comments (to an arbitrary nesting level)

01-56: Project Details

- JavaCC is available at <https://javacc.dev.java.net/>
- To compile your project

```
% javacc simplejava.jj
% javac *.java
```
- To test your project

```
% java TokenTest <test filename>
```
- To submit your program: Create a branch:

<https://www.cs.usfca.edu/svn/<username>/cs414/lexer/>