

01-0: **Syllabus**

- Office Hours
- Course Text
- Prerequisites
- Test Dates & Testing Policies
- Projects
  - Teams of up to 2
- Grading Policies
- Questions?

01-1: **Notes on the Class**

- Don't be afraid to ask me to slow down!
- We will cover some pretty complex stuff here, which can be difficult to get the first (or even the second) time.  
*ASK QUESTIONS*
- While specific questions are always preferred, "I don't get it" is always an acceptable question. I am always happy to stop, re-explain a topic in a different way.
  - If you are confused, I can *guarantee* that at least one other person in the class would benefit from more explanation

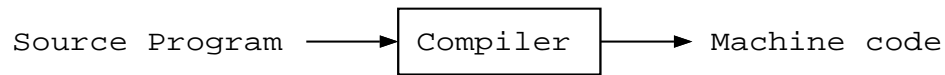
01-2: **Notes on the Class**

- Projects are non-trivial
  - Using new tools (JavaCC)
  - Managing a large scale project
  - Lots of complex classes & advanced programming techniques.

01-3: **Notes on the Class**

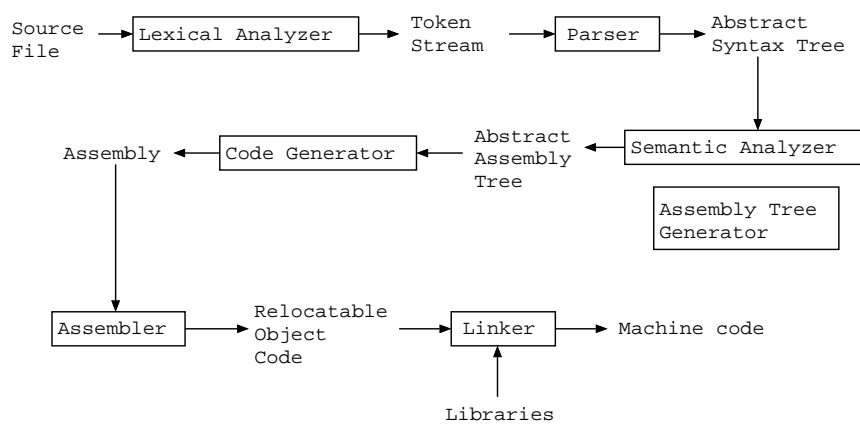
- Projects are non-trivial
  - Using new tools (JavaCC)
  - Managing a large scale project
  - Lots of complex classes & advanced programming techniques.
- *START EARLY!*
  - Projects will take longer than you think (especially starting with the semantic analyzer project)
- *ASK QUESTIONS!*

## 01-4: What is a compiler?



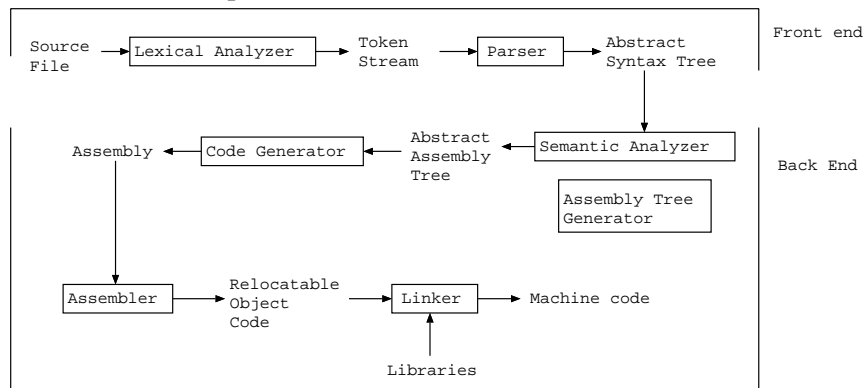
## Simplified View

## 01-5: What is a compiler?

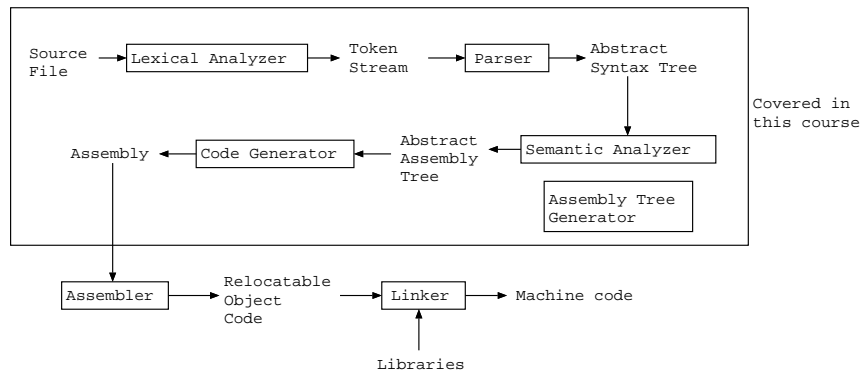


## More Accurate View

## 01-6: What is a compiler?



## 01-7: What is a compiler?



#### 01-8: Why Use Decomposition?

#### 01-9: Why Use Decomposition?

Software Engineering!

- Smaller units are easier to write, test and debug
- Code Reuse
  - Writing a suite of compilers (C, Fortran, C++, etc) for a new architecture
  - Create a new language – want compilers available for several platforms

#### 01-10: Lexical Analysis

- Converting input file to stream of tokens

```
void main() {
    print(4);
}
```

#### 01-11: Lexical Analysis

- Converting input file to stream of tokens

```
void main() {      IDENTIFIER(void)
    print(4);      IDENTIFIER(main)
}                  LEFT-PARENTHESIS
                   RIGHT-PARENTHESIS
                   LEFT-BRACE
                   IDENTIFIER(print)
                   LEFT-PARENTHESIS
                   INTEGER-LITERAL(4)
                   RIGHT-PARENTHESIS
                   SEMICOLON
                   RIGHT-BRACE
```

#### 01-12: Lexical Analysis

Brute-Force Approach

- Lots of nested if statements

```

if (c = nextchar() == 'P') {
    if (c = nextchar() == 'R') {
        if (c = nextchar() == '0') {
            if (c = nextchar() == 'G') {
                /* Code to handle the rest of either
                 PROGRAM or any identifier that starts
                 with PROG
                */
            } else if (c == 'C') {
                /* Code to handle the rest of either
                 PROCEDURE or any identifier that starts
                 with PROC
                */
            }
        }
    }
}
...

```

### 01-13: Lexical Analysis

#### Brute-Force Approach

- Break the input file into words, separated by spaces or tabs
  - This can be tricky – not all tokens are separated by whitespace
  - Use string comparison to determine tokens

### 01-14: Deterministic Finite Automata

- Set of states
- Initial State
- Final State(s)
- Transitions

DFA for else, end, identifiers

#### Combine DFA 01-15: DFAs and Lexical Analyzers

- Given a DFA, it is easy to create C code to implement it
- DFAs are easier to understand than C code
  - Visual – almost like structure charts
- ... However, creating a DFA for a complete lexical analyzer is still complex

### 01-16: Automatic Creation of DFAs

We'd like a tool:

- Describe the tokens in the language
- Automatically create DFA for tokens
- Then, automatically create C code that implements the DFA

We need a method for describing tokens

### 01-17: Formal Languages

- **Alphabet**  $\Sigma$ : Set of all possible symbols (characters) in the input file
  - Think of  $\Sigma$  as the set of symbols on the keyboard
- **String**  $w$ : Sequence of symbols from an alphabet

- **String length**  $|w|$  Number of characters in a string:  $|\text{car}| = 3$ ,  $|\text{abba}| = 4$

- **Empty String**  $\epsilon$ : String of length 0:  $|\epsilon| = 0$

- **Formal Language**: Set of strings over an alphabet

Formal Language  $\neq$  Programming language – Formal Language is only a set of strings.

#### 01-18: Formal Languages

Example formal languages:

- Integers  $\{0, 23, 44, \dots\}$
- Floating Point Numbers  $\{3.4, 5.97, \dots\}$
- Identifiers  $\{\text{foo}, \text{bar}, \dots\}$

#### 01-19: Language Concatenation

- **Language Concatenation** Given two formal languages  $L_1$  and  $L_2$ , the concatenation of  $L_1$  and  $L_2$ ,  $L_1L_2 = \{xy | x \in L_1, y \in L_2\}$

For example:

$\{\text{fire}, \text{truck}, \text{car}\} \{\text{car}, \text{dog}\} =$   
 $\{\text{firecar}, \text{firedog}, \text{truckcar}, \text{truckdog}, \text{carcar}, \text{cardog}\}$

#### 01-20: Kleene Closure

Given a formal language  $L$ :

$$\begin{aligned} L^0 &= \{\epsilon\} \\ L^1 &= L \\ L^2 &= LL \\ L^3 &= LLL \\ L^4 &= LLLL \end{aligned}$$

$$L^* = L^0 \cup L^1 \cup L^2 \cup \dots \cup L^n \cup \dots$$

#### 01-21: Regular Expressions

Regular expressions are used to describe formal languages over an alphabet  $\Sigma$ :

Regular Expression	Language
$\epsilon$	$L[\epsilon] = \{\epsilon\}$
$a \in \Sigma$	$L[a] = \{a\}$
$(MR)$	$L[MR] = L[M]L[R]$
$(M R)$	$L[(M R)] = L[M] \cup L[R]$
$(M^*)$	$L[(M^*)] = L[M]^*$

#### 01-22: r.e. Precedence

From highest to Lowest:

Kleene Closure  $*$   
 Concatenation  
 Alternation  $|$

$$\text{ab}^*\text{c}|\text{e} = (\text{a}(\text{b}^*)\text{c}) \mid \text{e}$$

#### 01-23: Regular Expression Examples

all strings over  $\{a,b\}$   
 binary integers (with leading zeroes)  
 all strings over  $\{a,b\}$  that  
     begin and end with a  
 all strings over  $\{a,b\}$  that  
     contain aa  
 all strings over  $\{a,b\}$  that  
     do not contain aa

#### 01-24: Regular Expression Examples

all strings over $\{a,b\}$	$(a b)^*$
binary integers (with leading zeroes)	$(0 1)(0 1)^*$
all strings over $\{a,b\}$ that begin and end with a	$a(a b)^*a$
all strings over $\{a,b\}$ that contain aa	$(a b)^*aa(a b)^*$
all strings over $\{a,b\}$ that do not contain aa	$b^*(abb^*)^*(a \epsilon)$

#### 01-25: Reg. Exp. Shorthand

$[a,b,c,d]$	$= (a b c d)$
$[d-g]$	$= [d,e,f,g] = (b e f g)$
$[d-f,M-O]$	$= [d,e,f,M,N,O]$
	$= (d e f M N O)$
$(\alpha)?$	$= \text{Optionally } \alpha \text{ (i.e., } (\alpha   \epsilon))$
$(\alpha)^+$	$= \alpha(\alpha)^*$

#### 01-26: Regular Expressions & Unix

- Many unix tools use regular expressions
- Example: `grep '<reg exp>' filename`
  - Prints all lines that contain a match to the regular expression
  - Special characters:
    - `^` beginning of line
    - `$` end of line
  - (grep examples on other screen)

#### 01-27: JavaCC Regular Expressions

- All characters & strings must be in quotation marks
  - `"else"`
  - `"+"`
  - `("a" | "b")`
- All regular expressions involving `*` must be parenthesized
  - `("a")*`, not `"a"*`

#### 01-28: JavaCC Shorthand

<code>["a","b","c","d"]</code>	<code>= ("a b c d")</code>	
<code>["d","g"]</code>	<code>= ["d","e","f","g"] = ("b e f g")</code>	
<code>["d","f","M","O"]</code>	<code>= ["d","e","f","M","N","O"]</code>	
	<code>= ("d e f M N O")</code>	
<code>(α)?</code>	<code>= Optionally α (i.e., (α   ε))</code>	01-29: <b>r.e. Shorthand Examples</b>
<code>(α)+</code>	<code>= α(α)*</code>	
<code>(~["a","b"])</code>	<code>= Any character <i>except</i> "a" or "b".</code>	
	<code>Can only be used with [] notation</code>	
	<code>~(a(a—b)*b) is not legal</code>	

Regular Expression	Language
	<code>{if}</code>
	Set of legal identifiers
	Set of integer literals
	(leading zeroes allowed)
	Set of real literals

01-30: **r.e. Shorthand Examples**

Regular Expression	Language
<code>"if"</code>	<code>{if}</code>
<code>["a"-"z"](["0"-"9","a"-"z"])*</code>	Set of legal identifiers
<code>["0"-"9"]</code>	Set of integer literals
	(leading zeroes allowed)
<code>(["0"-"9"]+"."(["0"-"9"]*)) </code>	Set of real literals
<code>((["0"-"9"])*"."(["0"-"9"])+)</code>	

01-31: **Lexical Analyzer Generator**

**JavaCC** is a Lexical Analyzer Generator and a Parser Generator

- Input: Set of regular expressions (each of which describes a type of token in the language)
- Output: A lexical analyzer, which reads an input file and separates it into tokens

01-32: **Structure of a JavaCC file**

```
options{
    /* Code to set various options flags */
}

PARSER_BEGIN(foo)

public class foo {
    /* This segment is often empty */
}

PARSER_END(foo)

TOKEN_MGR_DECLS :
{
    /* Declarations used by lexical analyzer */
}

/* Token Rules & Actions */
```

01-33: **Token Rules in JavaCC**

- Tokens are described by rules with the following syntax:

```
TOKEN :
{
    <TOKEN_NAME: RegularExpression>
}
```

- `TOKEN_NAME` is the name of the token being described

- `RegularExpression` is a regular expression that describes the token

#### 01-34: Token Rules in JavaCC

- Token rule examples:

```
TOKEN :
{
    <ELSE: "else">
}
```

```
TOKEN :
{
    <INTEGER_LITERAL: ( ["0"-"9"] ) +>
}
```

#### 01-35: Token Rules in JavaCC

- Several different tokens can be described in the same `TOKEN` block, with token descriptions separated by `|`.

```
TOKEN :
{
    <ELSE: "else">
|
    <INTEGER_LITERAL: ( ["0"-"9"] ) +>
|
    <SEMICOLON: ";">
}
```

#### 01-36: `getNextToken`

- When we run `javacc` on the input file `foo.jj`, it creates the class `fooTokenManager`
- The class `fooTokenManager` contains the static method `getNextToken()`
- Every call to `getNextToken()` returns the next token in the input stream.

#### 01-37: `getNextToken`

- When `getNextToken` is called, a regular expression is found that matches the next characters in the input stream.
- What if more than one regular expression matches?

```
TOKEN :
{
    <ELSE: "else">
|
    <IDENTIFIER: ( ["a"-"z"] ) +>
}
```

#### 01-38: `getNextToken`

- When more than one regular expression matches the input stream:
  - Use the longest match
    - “`elsed`” should match to `IDENTIFIER`, not to `ELSE` followed by the identifier “`d`”



- If two matches have the same length, use the rule that appears first in the .jj file
- “else” should match to ELSE, not IDENTIFIER

### 01-39: JavaCC Example

```
PARSER_BEGIN(simple)
public class simple {

}
PARSER_END(simple)

TOKEN :
{
    <ELSE: "else">
    | <SEMICOLON: ";">
    | <FOR: "for">
    | <INTEGER_LITERAL: ({"0"-"9"})*>
    | <IDENTIFIER: {"a"-"z"} ({"a"-"z", "0"-"9"})*>
}
```

else;ford for 01-40: **SKIP Rules**

- Tell JavaCC what to ignore (typically whitespace) using SKIP rules
- SKIP rule is just like a TOKEN rule, except that no TOKEN is returned.

SKIP:

```
{
    < regularexpression1 >
    | < regularexpression2 >
    | ...
    | < regularexpressionn >
}
```

### 01-41: Example SKIP Rules

```
PARSER_BEGIN(simple2)
public class simple2 {
}
PARSER_END(simple2)

SKIP :
{
    < " " >
    | < "\n" >
    | < "\t" >
}

TOKEN :
{
    <ELSE: "else">
    | <SEMICOLON: ";">
    | <FOR: "for">
    | <INTEGER_LITERAL: ({"0"-"9"})*>
    | <IDENTIFIER: {"A"-"Z"} ({"A"-"Z", "0"-"9"})*>
}
```

### 01-42: JavaCC States

- Comments can be dealt with using SKIP rules
- How could we skip over 1-line C++ Style comments?

```
// This is a comment
```

### 01-43: JavaCC States

- Comments can be dealt with using SKIP rules
- How we could skip over 1-line C++ Style comments:

```
// This is a comment
```

- Using a SKIP rule

```
SKIP :
{
    < " //" (~ ["\n"]) * "\n" >
}
```

#### 01-44: JavaCC States

- Writing a regular expression to match multi-line comments (using /\* and \*/) is much more difficult
- Writing a regular expression to match nested comments is impossible (take Automata Theory for a proof :) )
- What can we do?
  - Use JavaCC States

#### 01-45: JavaCC States

- We can label each TOKEN and SKIP rule with a “state”
- Unlabeled TOKEN and SKIP rules are assumed to be in the default state (named DEFAULT, unsurprisingly enough)
- Can switch to a new state after matching a TOKEN or SKIP rule using the : NEWSTATE notation

#### 01-46: JavaCC States

```
SKIP :
{
    < " " >
    | < "\n" >
    | < "\t" >
}
SKIP :
{
    < "/*" > : IN_COMMENT
}
<IN_COMMENT>
SKIP :
{
    < "*/" > : DEFAULT
    | < "[^]" >
}
TOKEN :
{
    <ELSE: "else">
    | ... (etc)
}
```

#### 01-47: Actions in TOKEN & SKIP

- We can add Java code to any SKIP or TOKEN rule
- That code will be executed when the SKIP or TOKEN rule is matched.
- Any methods / variables defined in the TOKEN\_MGR\_DECLS section can be used by these actions

#### 01-48: Actions in TOKEN & SKIP

```

PARSER_BEGIN(remComments)
public class remComments { }
PARSER_END(remComments)

TOKEN_MGR_DECLS :
{
    public static int numcomments = 0;
}

SKIP :
{
    < "/*" > : IN_COMMENT
}

SKIP :
{
    < "//" (~["\n"])* "\n" > { numcomments++; }
}

```

#### 01-49: Actions in TOKEN & SKIP

```

<IN_COMMENT>
SKIP :
{
    < "/*" > { numcomments++; SwitchTo(DEFAULT); }
}

<IN_COMMENT>
SKIP :
{
    < "[" >
}

TOKEN :
{
    <ANY: ~[]>
}

```

#### 01-50: Tokens

- Each call to getNextToken returns a “Token” object
- Token class is automatically created by javaCC.
- Variables of type Token contain the following public variables:
  - public int kind; The type of token. When javacc is run on the file foo.jj, a file fooConstants.java is created, which contains the symbolic names for each constant

```

public interface simplejavaConstants {
    int EOF = 0;
    int CLASS$ = 8;
    int DO = 9;
    int ELSE = 10;
    ...
}

```

#### 01-51: Tokens

- Each call to getNextToken returns a “Token” object
- Token class is automatically created by javaCC.
- Variables of type Token contain the following public variables:
  - public int beginLine, beginColumn, endLine, endColumn; The location of the token in the input file

#### 01-52: Tokens

- Each call to getNextToken returns a “Token” object
- Token class is automatically created by javaCC.
- Variables of type Token contain the following public variables:

- `public String image;` The text that was matched to create the token.

### 01-53: Generated TokenManager

```
class TokenTest {
    public static void main(String args[]) {
        Token t;
        Java.io.InputStream infile;
        pascalTokenManager tm;
        boolean loop = true;

        if (args.length < 1) {
            System.out.print("Enter filename as command line argument");
            return;
        }
        try {
            infile = new Java.io.FileInputStream(args[0]);
        } catch (Java.io.FileNotFoundException e) {
            System.out.println("File " + args[0] + " not found.");
            return;
        }
        tm = new sJavaTokenManager(new SimpleCharStream(infile));
```

### 01-54: Generated TokenManager

```
        t = tm.getNextToken();
        while (t.kind != sJavaConstants.EOF) {
            System.out.println("Token : " + t + " : ");
            System.out.println(pascalConstants.tokenImage[t.kind]);
        }
    }
}
```

### 01-55: Lexer Project

- Write a .jj file for simpleJava tokens
- Need to handle all whitespace (tabs, spaces, end-of-line)
- Need to handle nested comments (to an arbitrary nesting level)

### 01-56: Project Details

- JavaCC is available at <https://javacc.dev.java.net/>
- To compile your project

```
% javacc simplejava.jj
% javac *.java
```

- To test your project

```
% java TokenTest <test filename>
```

- To submit your program: Create a branch:

<https://www.cs.usfca.edu/svn/<username>/cs414/lexer/>