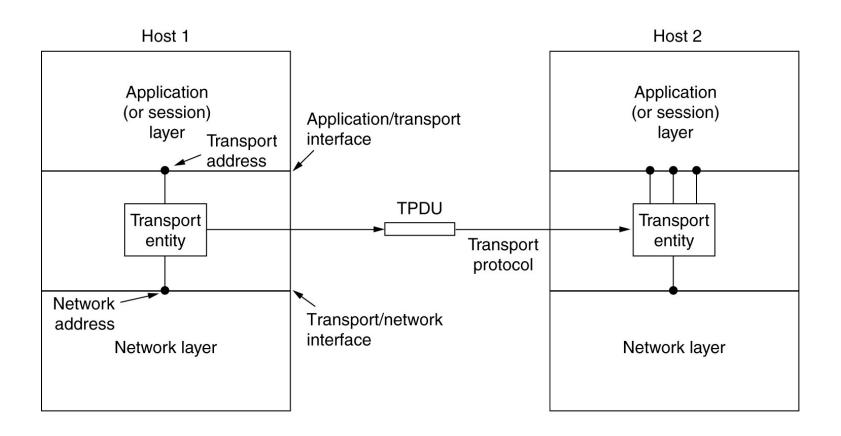
Chapter 6

The Transport Layer

The Transport Service

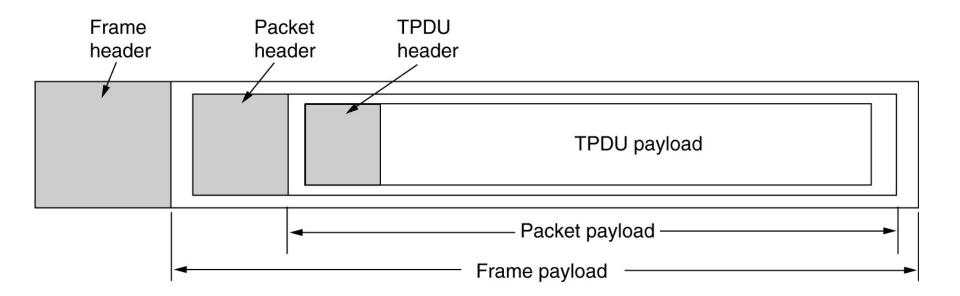
- Services Provided to the Upper Layers
- Transport Service Primitives
- Berkeley Sockets
- An Example of Socket Programming:
 - An Internet File Server
- User Datagram Protocol

Services Provided to the Upper Layers



The network, transport, and application layers.

Transport Service Primitives (2)



The nesting of TPDUs, packets, and frames.

Elements of Transport Protocols

- Addressing
- Connection Establishment
- Connection Release
- Flow Control and Buffering
- Multiplexing
- Crash Recovery

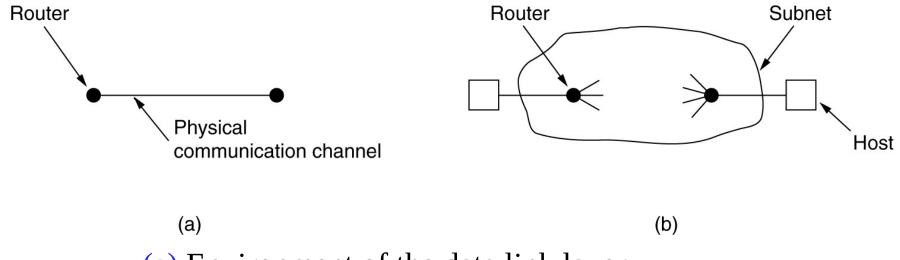
Transport Protocol

Transport Protocols resemble data link protocols:

Error control

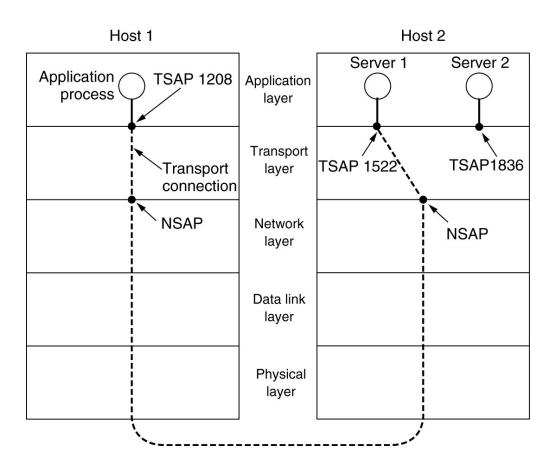
Sequencing

Flow Control



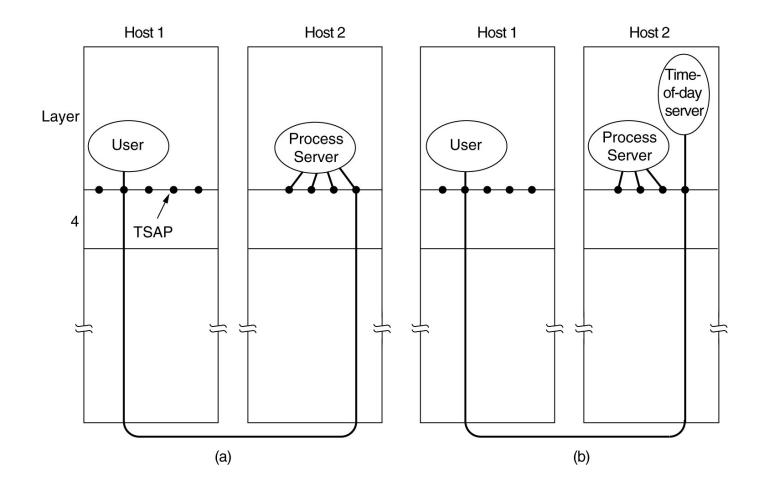
- (a) Environment of the data link layer.
- (b) Environment of the transport layer.

Addressing



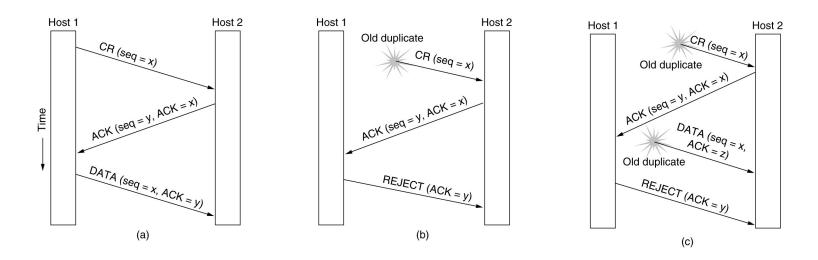
TSAPs, NSAPs and transport connections.

Connection Establishment



How a user process in host 1 establishes a connection with a time-of-day server in host 2 (xinetd).

Connection Establishment (3)

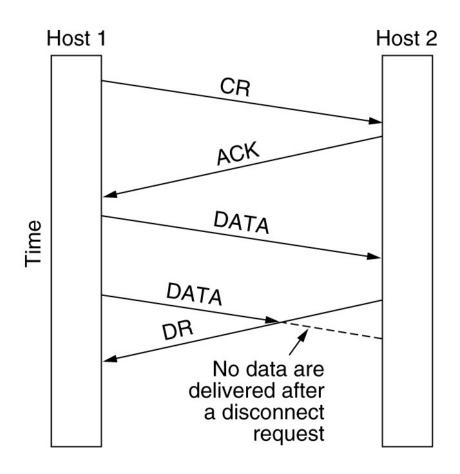


Three protocol scenarios for establishing a connection using a three-way handshake. CR denotes CONNECTION REQUEST.

- (a) Normal operation,
- (b) Old Dup CONNECTION REQUEST appearing out of nowhere.
- (c) Duplicate CONNECTION REQUEST and duplicate ACK.

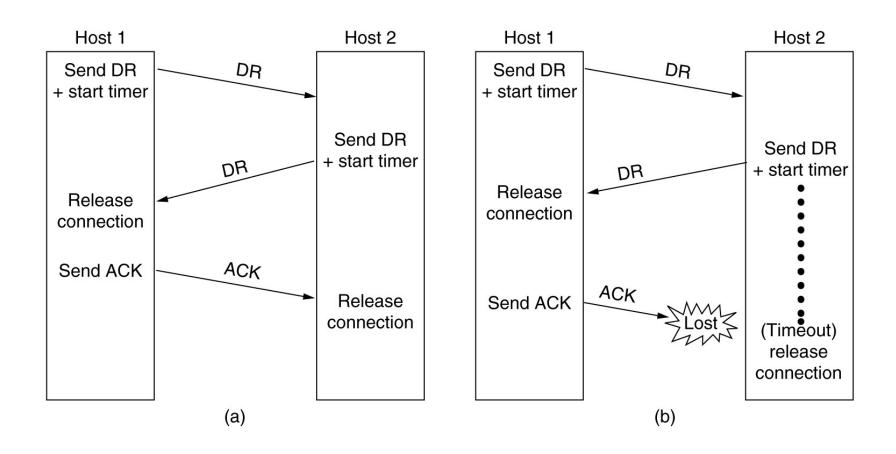
Connection Release

Asymmetric Release



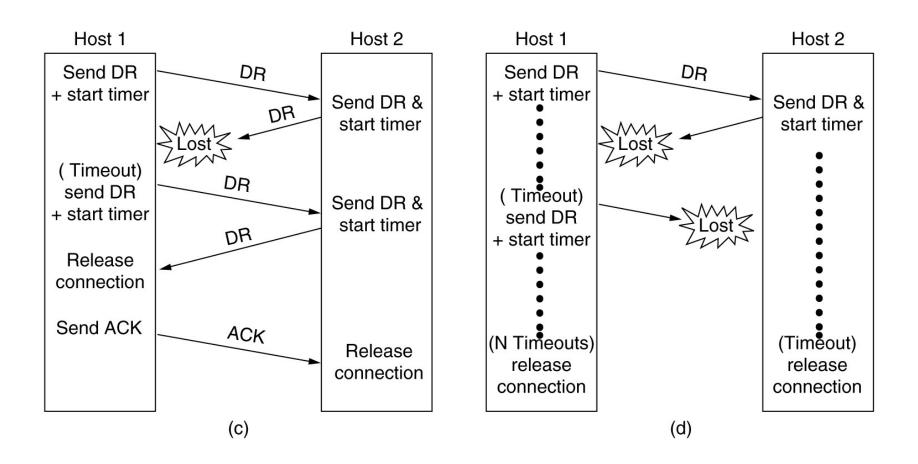
Abrupt disconnection with loss of data.

Connection Release (2)



Four protocol scenarios for releasing a connection. (a) Normal case of a three-way handshake. (b) final ACK lost.

Connection Release (3)



(c) Response lost. (d) Response lost and subsequent DRs lost.

Flow Control and Buffering

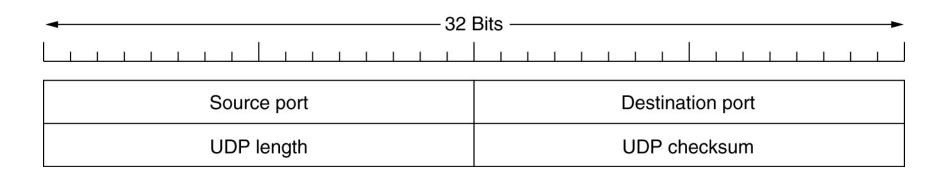
Α	Message	В	Comments
· —		_	
-	< request 8 buffers>	-	A wants 8 buffers
←	<ack = 15, buf = 4 $>$	-	B grants messages 0-3 only
-	<seq = 0, data = m0 $>$	-	A has 3 buffers left now
-	<seq = 1, data = m1>	-	A has 2 buffers left now
-	<seq = 2, data = m2 $>$	•••	Message lost but A thinks it has 1 left
←	<ack = 1, buf = 3>	•	B acknowledges 0 and 1, permits 2-4
-	<seq = 3, data = m3 $>$	\rightarrow	A has 1 buffer left
-	<seq = 4, data = m4 $>$	\rightarrow	A has 0 buffers left, and must stop
-	<seq = 2, data = m2 $>$	\rightarrow	A times out and retransmits
•	<ack = 4, buf = 0>	-	Everything acknowledged, but A still blocked
←	<ack = 4, buf = 1>	←	A may now send 5
←	<ack = 4, buf = 2 $>$	•	B found a new buffer somewhere
-	<seq = 5, data = m5 $>$	-	A has 1 buffer left
-	<seq = 6, data = m6 $>$	-	A is now blocked again
←	<ack = 6, buf = 0 $>$	←	A is still blocked
• • •	<ack = 6, buf = 4 $>$	•	Potential deadlock
		<pre></pre>	<pre></pre>

Dynamic buffer allocation. The arrows show the direction of transmission. An ellipsis (...) indicates a lost TPDU.

The Internet Transport Protocols: UDP

- Introduction to UDP
- Remote Procedure Call
- The Real-Time Transport Protocol

Introduction to UDP



The UDP header.

The Internet Transport Protocols: TCP

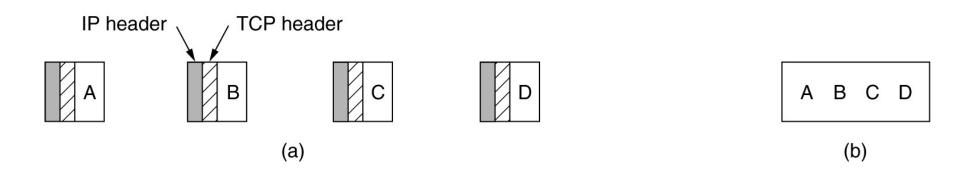
- Introduction to TCP
- The TCP Service Model
- The TCP Protocol
- The TCP Segment Header
- TCP Connection Establishment
- TCP Connection Release
- TCP Connection Management Modeling
- TCP Transmission Policy
- TCP Congestion Control
- TCP Timer Management
- Wireless TCP and UDP
- Transactional TCP

The TCP Service Model

Port	Protocol	Use
21	FTP	File transfer
23	Telnet	Remote login
25	SMTP	E-mail
69	TFTP	Trivial File Transfer Protocol
79	Finger	Lookup info about a user
80	HTTP	World Wide Web
110	POP-3	Remote e-mail access
119	NNTP	USENET news

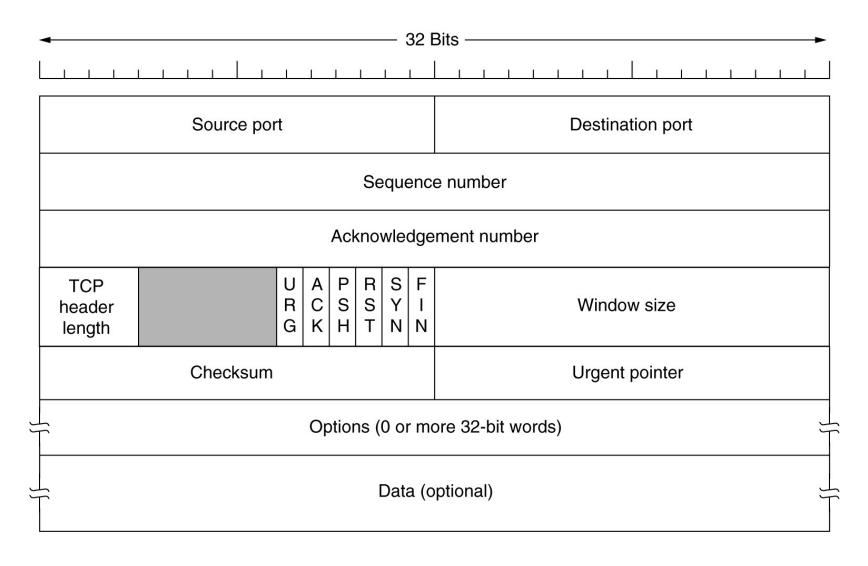
Some assigned ports.

The TCP Service Model (2)



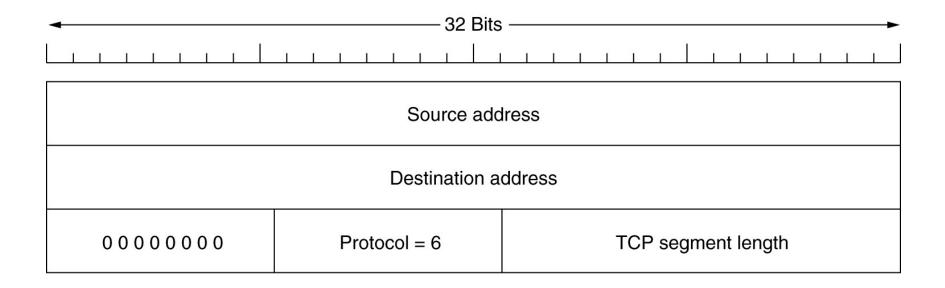
- (a) Four 512-byte segments sent as separate IP datagrams.
- (b) The 2048 bytes of data delivered to the application in a single READ CALL.

The TCP Segment Header



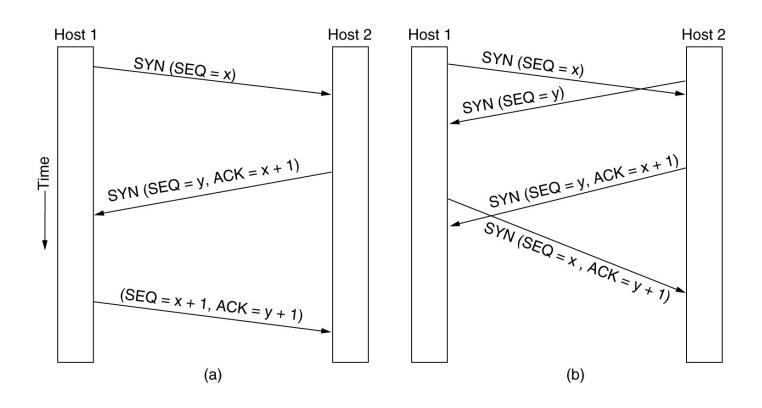
TCP Header.

The TCP Segment Header (2)



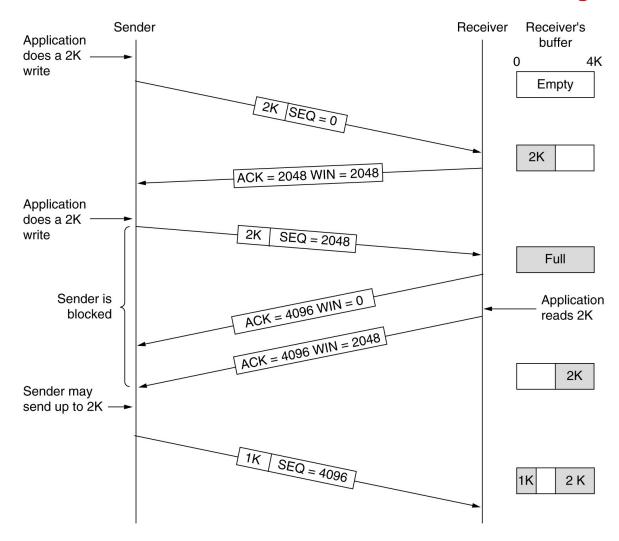
The pseudoheader included in the TCP checksum.

TCP Connection Establishment



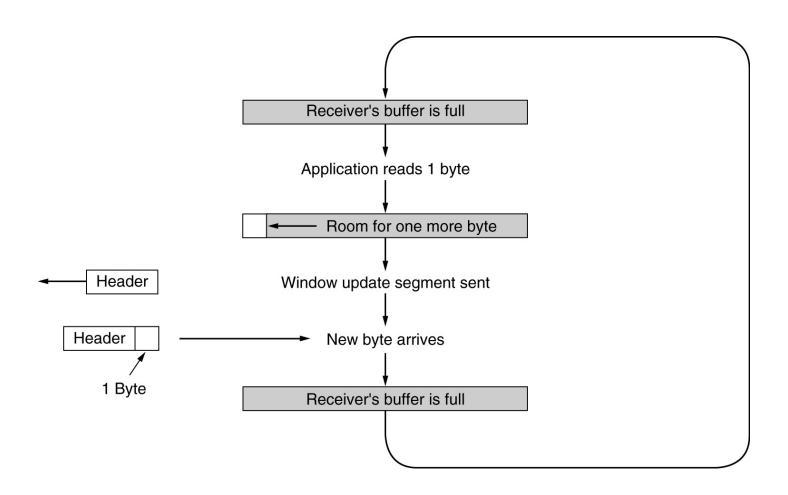
- (a) TCP connection establishment in the normal case.
- (b) Call collision.

TCP Transmission Policy



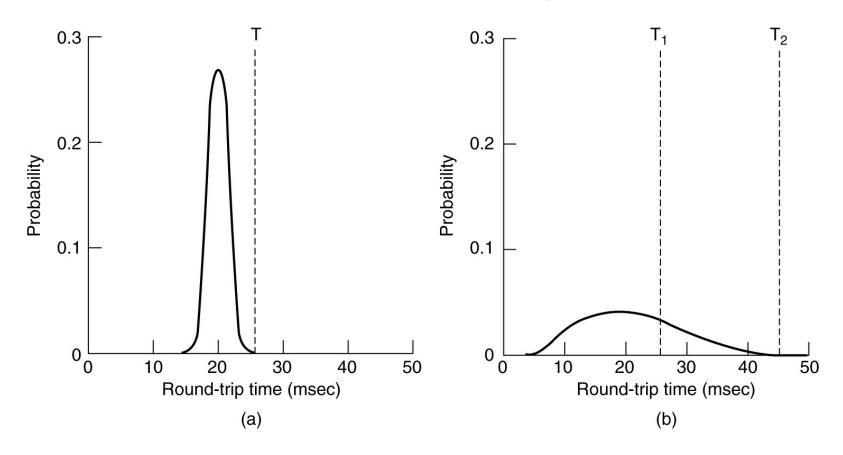
Window management in TCP.

TCP Transmission Policy (2)



Silly window syndrome.

TCP Timer Management



- (a) Probability density of ACK arrival times in the data Link layer.
- (b) Probability density of ACK arrival times for TCP.