



# Condition Variables, Classic Sync Problem

*CS 571: Operating Systems (Spring 2020)*  
Lecture 4

Yue Cheng

Some material taken/derived from:

- Wisconsin CS-537 materials created by Remzi Arpacı-Dusseau.

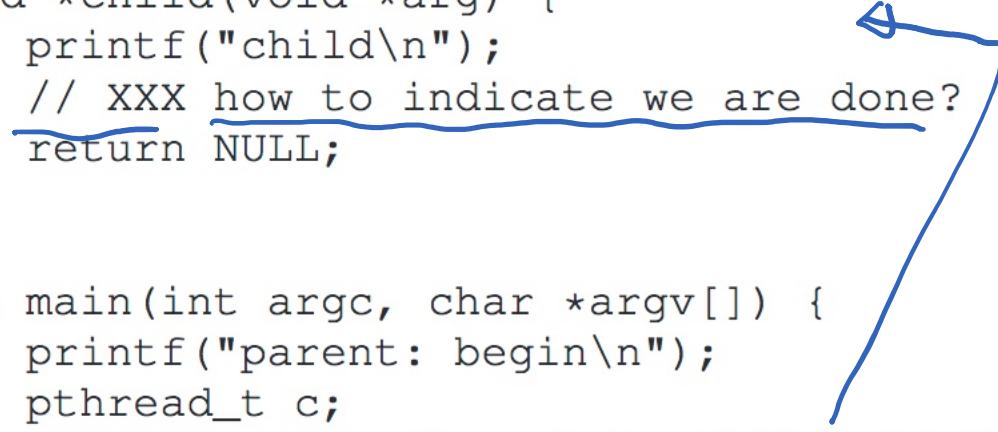
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# Condition Variables

# Condition Variables

A parent waiting for its child

```
1 void *child(void *arg) {  
2     printf("child\n");  
3     // XXX how to indicate we are done?  
4     return NULL;  
5 }  
6  
7 int main(int argc, char *argv[]) {  
8     printf("parent: begin\n");  
9     pthread_t c;  
10    Pthread_create(&c, NULL, child, NULL); // create child  
11    // XXX how to wait for child?  
12    printf("parent: end\n");  
13    return 0;  
14 }
```



# Spin-based Approach

Using a shared variable, parent spins until child set it to 1

```
1  volatile int done = 0;
2
3  void *child(void *arg) {
4      printf("child\n");
5      done = 1;
6      return NULL;
7  }
8
9  int main(int argc, char *argv[]) {
10     printf("parent: begin\n");
11     pthread_t c;
12     Pthread_create(&c, NULL, child, NULL); // create child
13     while (done == 0) 
14         ; // spin
15     printf("parent: end\n");
16     return 0;
17 }
```

# Spin-based Approach

Using a shared variable, parent spins until child set it to 1

```
1  volatile int done = 0;
2
3  void *child(void *arg) {
4      printf("child\n");
5      done = 1;
6      return NULL;
7  }
8
9  int main(int argc, char *argv[]) {
10     printf("parent: begin\n");
11     pthread_t c;
12     Pthread_create(&c, NULL, child, NULL); // create child
13     while (done == 0)
14         ; // spin
15     printf("parent: end\n");
16     return 0;
17 }
```

What's the problem of this approach?

# Condition Variables (CV)

- Definition:
  - An explicit queue that threads can put themselves when some **condition** is not as desired (by **waiting** on the condition)
  - Other thread can wake one of those waiting threads to allow them to continue (by **signaling** on the condition)
- Pthread CV

```
pthread_cond_wait(pthread_cond_t *c, pthread_mutex_t *m);  
pthread_cond_signal(pthread_cond_t *c);
```

# CV-based Approach

```
void *child(void *arg) {  
    printf("child\n");  
    thr_exit();          ??  
    return NULL;  
}
```

```
int main(int argc, char *argv[]) {  
    printf("parent: begin\n");  
    pthread_t p;  
    Pthread create(&p, NULL, child, NULL);  
    thr_join();          ??  
    printf("parent: end\n");  
    return 0;  
}
```

# Broken Implementation 1

```
void *child(void *arg) {  
    printf("child\n");  
    thr_exit();  
    return NULL;  
}  
  
int main(int argc, char *argv[]) {  
    printf("parent: begin\n");  
    pthread_t p;  
    Pthread_create(&p, NULL, child, NULL);  
    thr_join();  
    printf("parent: end\n");  
    return 0;  
}
```

```
1 → void thr_exit() {  
2     Pthread_mutex_lock(&m);  
3     Pthread_cond_signal(&c);  
4     Pthread_mutex_unlock(&m);  
5 }  
6  
7     void thr_join() {  
8         Pthread_mutex_lock(&m);  
9         Pthread_cond_wait(&c, &m);  
10        Pthread_mutex_unlock(&m);  
11    }
```

# Broken Implementation 1

```
void *child(void *arg) {  
    printf("child\n");  
    thr_exit();  
    return NULL;  
}  
  
int main(int argc, char *argv[]) {  
    printf("parent: begin\n");  
    pthread_t p;  
    Pthread_create(&p, NULL, child, NULL);  
    thr_join();  
    printf("parent: end\n");  
    return 0;  
}  
  
1 → void thr_exit() {  
2     Pthread_mutex_lock(&m);  
3     Pthread_cond_signal(&c);  
4     Pthread_mutex_unlock(&m);  
5 }  
6  
7     void thr_join() {  
8         Pthread_mutex_lock(&m);  
9         Pthread_cond_wait(&c, &m);  
10        Pthread_mutex_unlock(&m);  
11    }
```

If parent comes after child, parent sleeps forever

# Broken Implementation 1

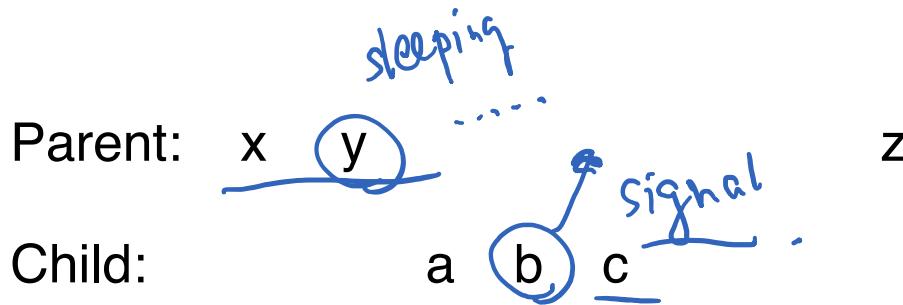
time



```
void thread_exit() {  
    Mutex_lock(&m);          // a  
    Cond_signal(&c);         // b  
    Mutex_unlock(&m);        // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);          // x  
    Cond_wait(&c, &m);       // y  
    Mutex_unlock(&m);        // z  
}
```

# Broken Implementation 1



```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    Cond_signal(&c);          // b  
    Mutex_unlock(&m);         // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);           // x  
    Cond_wait(&c, &m);        // y  
    Mutex_unlock(&m);         // z  
}
```

# Broken Implementation 1

Parent: x y z

Child: a b c

**GOOD!**

```
void thread_exit() {  
    Mutex_lock(&m);          // a  
    Cond_signal(&c);         // b  
    Mutex_unlock(&m);        // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);          // x  
    Cond_wait(&c, &m);       // y  
    Mutex_unlock(&m);        // z  
}
```

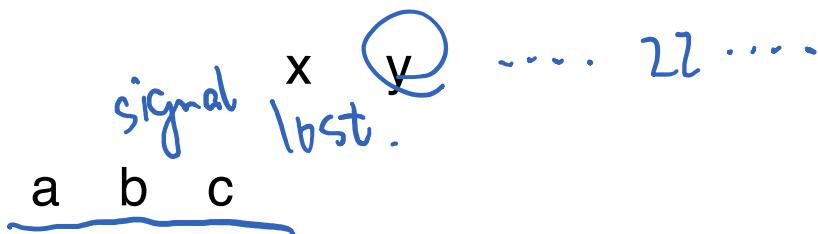
# Broken Implementation 1

```
void thread_exit() {
    Mutex_lock(&m);          // a
    Cond_signal(&c);         // b
    Mutex_unlock(&m);        // c
}
```

```
void thread_join() {
    Mutex_lock(&m);          // x
    Cond_wait(&c, &m);       // y
    Mutex_unlock(&m);        // z
}
```

# Broken Implementation 1

Parent:



Child:

a b c

```
void thread_exit() {
    Mutex_lock(&m);           // a
    Cond_signal(&c);          // b
    Mutex_unlock(&m);         // c
}
```

```
void thread_join() {
    Mutex_lock(&m);           // x
    Cond_wait(&c, &m);        // y
    Mutex_unlock(&m);         // z
}
```

# Broken Implementation 1

Parent:                    x    y    ... *sleep forever ...*

Child:    a    b    c

```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    Cond_signal(&c);          // b  
    Mutex_unlock(&m);         // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);           // x  
    Cond_wait(&c, &m);        // y  
    Mutex_unlock(&m);         // z  
}
```

# Broken Implementation 2

```
void *child(void *arg) {  
    printf("child\n");  
    thr_exit();  
    return NULL;  
}
```

```
1 void thr_exit() {  
2     done = 1;  
3     Pthread_cond_signal(&c);  
4 }  
5  
6 void thr_join() {  
7     if (done == 0)  
8         Pthread_cond_wait(&c);  
9 }
```

```
int main(int argc, char *argv[]) {  
    printf("parent: begin\n");  
    pthread_t p;  
    Pthread_create(&p, NULL, child, NULL);  
    thr_join();  
    printf("parent: end\n");  
    return 0;  
}
```

# Broken Implementation 2

```
void *child(void *arg) {  
    printf("child\n");  
    thr_exit();  
    return NULL;  
}
```

```
int main(int argc, char *argv[]) {  
    printf("parent: begin\n");  
    pthread_t p;  
    Pthread_create(&p, NULL, child, NULL);  
    thr_join();  
    printf("parent: end\n");  
    return 0;  
}
```

```
1 void thr_exit() {  
2     done = 1;  
3     Pthread_cond_signal(&c);  
4 }  
5  
6 void thr_join() {  
7     if (done == 0)  
8         Pthread_cond_wait(&c);  
9 }
```

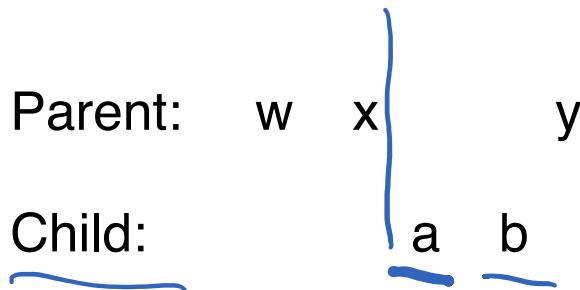
No mutual exclusion, hence child may signal before parent calls cond\_wait(). In this case, parent sleeps forever!

# Broken Implementation 2

```
void thread_exit() {  
    done = 1;          // a  
    Cond_signal(&c); // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m);      // w  
    if (done == 0)        // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);    // z  
}
```

# Broken Implementation 2



```
void thread_exit() {  
    done = 1; // a  
    Cond_signal(&c); // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m); // w  
    if (done == 0) // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m); // z  
}
```

# Broken Implementation 2

Parent: w x y ... *sleeeeeeeeep forever ...*

Child: a b

```
void thread_exit() {  
    done = 1; // a  
    Cond_signal(&c); // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m); // w  
    if (done == 0) // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m); // z  
}
```

# Broken Implementation 2

Parent: w x y ... *sleeeeeeeeep forever ...*

Child: a b

```
void thread_exit() {  
    done = 1; // a  
    Cond_signal(&c); // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m); // w  
    if (done == 0) // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m); // z  
}
```

How to fix?

# Broken Implementation 2

Parent: w x y ... *sleeeeeeeeep forever ...*

Child: a b

*Child*

```
void thread_exit() {
    done = 1;
    Cond_signal(&c);
}
Mutex_lock(&m); // a
Mutex_unlock(&m); // b
```

```
void thread_join() {
    Mutex_lock(&m);
    while (done == 0)
        Cond_wait(&c, &m);
    Mutex_unlock(&m);
}
// w
// x
// y
// z
```

# Trap 1 When Using CV



# Trap 1 When Using CV



# Trap 1 When Using CV



# Trap 1 When Using CV



Only one thread gets a signal

# Trap 2 When Using CV

Condition Variable

# Trap 2 When Using CV



# Trap 2 When Using CV

Condition Variable

# Trap 2 When Using CV



# Trap 2 When Using CV



# Trap 2 When Using CV



Signal lost if nobody waiting at that time

# Guarantee

Upon signal, there has to be **at least one** thread waiting;  
If there are threads waiting, **at least one** thread will wake



FIFO

```

1 int done    = 0;
2 pthread_mutex_t m = PTHREAD_MUTEX_INITIALIZER;
3 pthread_cond_t c  = PTHREAD_COND_INITIALIZER;
4
5 void thr_exit() {
6     Pthread_mutex_lock(&m);
7     done = 1;
8     Pthread_cond_signal(&c);
9     Pthread_mutex_unlock(&m);
10}
11
12 void *child(void *arg) {
13     printf("child\n");
14     thr_exit();
15     return NULL;
16}
17
18 void thr_join() {
19     Pthread_mutex_lock(&m);
20     while (done == 0)
21         Pthread_cond_wait(&c, &m);
22     Pthread_mutex_unlock(&m);
23}
24
25 int main(int argc, char *argv[]) {
26     printf("parent: begin\n");
27     pthread_t p;
28     Pthread_create(&p, NULL, child, NULL);
29     thr_join();
30     printf("parent: end\n");
31     return 0;
32} Y.Cheng

```

## CV-based Parent-wait-for-child Approach

```

1 int done = 0;
2 pthread_mutex_t m = PTHREAD_MUTEX_INITIALIZER;
3 pthread_cond_t c = PTHREAD_COND_INITIALIZER;
4
5 void thr_exit() {
6     Pthread_mutex_lock(&m);
7     done = 1;
8     Pthread_cond_signal(&c);
9     Pthread_mutex_unlock(&m);
10}

```

```

12 void *child(void *arg) {
13     printf("child\n");
14     thr_exit();
15     return NULL;
16 }

```

```

18 void thr_join() {
19     Pthread_mutex_lock(&m);
20     while (done == 0)
21         Pthread_cond_wait(&c, &m);
22     Pthread_mutex_unlock(&m);
23 }

```

```

25 int main(int argc, char *argv[]) {
26     printf("parent: begin\n");
27     pthread_t p;
28     Pthread_create(&p, NULL, child, NULL);
29     thr_join();
30     printf("parent: end\n");
31     return 0;
32 } Y.Cheng

```

## CV-based Parent-wait-for-child Approach

### Good Rule of Thumb

Always do 1. **wait** and 2. **signal** while holding the lock

Why: To prevent lost signal

# Classical Problems of Synchronization

- Producer-consumer problem
  - CV-based version
- Readers-writers problem
- Dining-philosophers problem

# CV-based Producer-Consumer Implementation 1

## Single CV and if statement

```
cond_t cond;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);
        if (count == 1)
            Pthread_cond_wait(&cond, &mutex);
        put(i);
        Pthread_cond_signal(&cond);
        Pthread_mutex_unlock(&mutex);
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);
        if (count == 0)
            Pthread_cond_wait(&cond, &mutex);
        int tmp = get();
        Pthread_cond_signal(&cond);
        Pthread_mutex_unlock(&mutex);
        printf("%d\n", tmp);
    }
}
```

```
1  int buffer;
2  int count = 0; // initially, empty
3
4  void put(int value) {
5      assert(count == 0);
6      count = 1;
7      buffer = value;
8  }
9
10 int get() {
11     assert(count == 1);
12     count = 0;
13     return buffer;
14 }
```

Put and Get routines  
Single buffer

# CV-based Producer-Consumer Implementation 1

## Single CV and if statement

```
cond_t cond;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

```
1  int buffer;
2  int count = 0; // initially, empty
3
4  void put(int value) {
5      assert(count == 0);
6      count = 1;
7      buffer = value;
8  }
9
10 int get() {
11     assert(count == 1);
12     count = 0;
13     return buffer;
14 }
```

Put and Get routines  
Single buffer

What's the problem of this approach?

# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) { ← C1 running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0)
            Pthread_cond_wait(&cond, &mutex); // c2
        int tmp = get(); // c3
        Pthread_cond_signal(&cond);
        Pthread_mutex_unlock(&mutex); // c4
        printf("%d\n", tmp);
    }
}
```

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1)
            Pthread_cond_wait(&cond, &mutex); // p2
        put(i); // p3
        Pthread_cond_signal(&cond); // p4
        Pthread_mutex_unlock(&mutex); // p5
    }
} // p6
```

C1

C2

P

empty

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get



# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) { ← P running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);
        Pthread_cond_signal(&cond); // p4
        Pthread_mutex_unlock(&mutex); // p5
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) { ← P running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) { ← P running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken



# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

← P running

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

P running

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1)
            Pthread_cond_wait(&cond, &mutex); // p2
        put(i); // p3
        Pthread_cond_signal(&cond); // p4
        Pthread_mutex_unlock(&mutex); // p5
    }
}

```

C2 running

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) { ← C2 running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                    // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)
            Pthread_cond_wait(&cond, &mutex); // p2
        put(i);
        Pthread_cond_signal(&cond);          // p4
        Pthread_mutex_unlock(&mutex);        // p5
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1)
            Pthread_cond_wait(&cond, &mutex); // p2
        put(i); // p3
        Pthread_cond_signal(&cond); // p4
        Pthread_mutex_unlock(&mutex); // p5
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	$T_p$ awoken

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) { ← C2 running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                    // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)
            Pthread_cond_wait(&cond, &mutex); // p2
        put(i);                            // p3
        Pthread_cond_signal(&cond);          // p4
        Pthread_mutex_unlock(&mutex);        // p5
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	$T_p$ awoken
	Ready	c6	Running		Ready	0	

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C1 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);
        Pthread_cond_signal(&cond); // p4
        Pthread_mutex_unlock(&mutex); // p5
    }
}

```

// p1  
// p2  
// p3  
// p4  
// p5  
// p6

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	$T_p$ awoken
	Ready	c6	Running		Ready	0	
c4	Running		Ready		Ready	0	Oh oh! No data

# CV-based Producer-Consumer Implementation 2



```
1 cond_t cond;
2 mutex_t mutex;
3
4 void *producer(void *arg) {
5     int i;
6     for (i = 0; i < loops; i++) {
7         Pthread_mutex_lock(&mutex); // p1
8         while (count == 1) // p2
9             Pthread_cond_wait(&cond, &mutex); // p3
10        put(i); // p4
11        Pthread_cond_signal(&cond); // p5
12        Pthread_mutex_unlock(&mutex); // p6
13    }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex); // c1
20         while (count == 0) // c2
21             Pthread_cond_wait(&cond, &mutex); // c3
22         int tmp = get(); // c4
23         Pthread_cond_signal(&cond); // c5
24         Pthread_mutex_unlock(&mutex); // c6
25         printf("%d\n", tmp);
26     }
27 }
```

# CV-based Producer-Consumer Implementation 2

```
1  cond_t  cond;
2  mutex_t mutex;
3
4  void *producer(void *arg) {
5      int i;
6      for (i = 0; i < loops; i++) {
7          Pthread_mutex_lock(&mutex);           // p1
8          while (count == 1)                  // p2
9              Pthread_cond_wait(&cond, &mutex); // p3
10         put(i);                         // p4
11         Pthread_cond_signal(&cond);       // p5
12         Pthread_mutex_unlock(&mutex);     // p6
13     }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);           // c1
20         while (count == 0)                  // c2
21             Pthread_cond_wait(&cond, &mutex); // c3
22         int tmp = get();                  // c4
23         Pthread_cond_signal(&cond);       // c5
24         Pthread_mutex_unlock(&mutex);     // c6
25         printf("%d\n", tmp);
26     }
27 }
```

Single CV and while

What's the problem of this approach?

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C1 running

```
void *producer(void *arg) {
```

```

    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get

(cond-wait())
 ↗  
 unlock()  
 // sleep ...  
 // woken ...  
 lock(..)  
 ↘

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C2 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	<u>Sleep</u>		Ready	0	Nothing to get

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

P running ←

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

P running ←

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
Ready			Sleep	p5	Running	1	<u><math>T_{c1}</math> awoken</u>

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

P sleeping ← P sleeping

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C1 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C1 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	$T_{c1}$ grabs data

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C1 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);
        Pthread_cond_signal(&cond); // p4
        Pthread_mutex_unlock(&mutex); // p5
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Nothing to get
	Ready		Sleep	p5	Running	1	
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	
c2	Running		Sleep		Sleep	1	Must sleep (full)
c4	Running		Sleep		Sleep	0	Recheck condition
c5	Running		Ready		Sleep	0	$T_{c1}$ grabs data
							Oops! Woke $T_{c2}$

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C1 sleeping

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	$T_{c1}$ grabs data
c5	Running		Ready		Sleep	0	Oops! Woke $T_{c2}$
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

C2 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	$T_{c1}$ grabs data
c5	Running		Ready		Sleep	0	Oops! Woke $T_{c2}$
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	
	Sleep	c2	Running		Sleep	0	
	Sleep	c3	Sleep		Sleep	0	Everyone asleep...

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

## C2 sleeping

```
void *producer(void *arg) {
```

```

    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	
	Ready		Sleep	p5	Running	1	Buffer now full
	Ready		Sleep	p6	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	$T_{c1}$ grabs data
c5	Running		Ready		Sleep	0	Oops! Woke $T_{c2}$
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get
	Sleep	c2	Running		Sleep	0	
	Sleep	c3	Sleep		Sleep	0	Everyone asleep...

# CV-based Producer-Consumer Implementation 3

```
1 cond_t empty, fill;
2 mutex_t mutex;
3
4 void *producer(void *arg) {
5     int i;
6     for (i = 0; i < loops; i++) {
7         Pthread_mutex_lock(&mutex);
8         while (count == 1)
9             Pthread_cond_wait(&empty, &mutex);
10        put(i);
11        Pthread_cond_signal(&fill);
12        Pthread_mutex_unlock(&mutex);
13    }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);
20         while (count == 0)
21             Pthread_cond_wait(&fill, &mutex);
22         int tmp = get();
23         Pthread_cond_signal(&empty);
24         Pthread_mutex_unlock(&mutex);
25         printf("%d\n", tmp);
26     }
27 }
```

Two CVs and while +

# CV-based Producer-Consumer Implementation 3

```
1 cond_t empty, fill;
2 mutex_t mutex;
3
4 void *producer(void *arg) {
5     int i;
6     for (i = 0; i < loops; i++) {
7         Pthread_mutex_lock(&mutex);
8         while (count == 1)
9             Pthread_cond_wait(&empty, &mutex);
10        put(i);
11        Pthread_cond_signal(&fill);
12        Pthread_mutex_unlock(&mutex);
13    }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);
20         while (count == 0)
21             Pthread_cond_wait(&fill, &mutex);
22         int tmp = get();
23         Pthread_cond_signal(&empty);
24         Pthread_mutex_unlock(&mutex);
25         printf("%d\n", tmp);
26     }
27 }
```

## Two CVs and while

Using **two CVs** to distinguish two types of threads; in order to properly signal which thread should wake up

- Producer waits on **empty**
- Consumer waits on **full**

# Readers-Writers Problem

# Readers-Writers Problem

- A data object (e.g. a file) is to be shared among several concurrent processes/threads
- A writer process/thread must have exclusive access to the data object
- **Multiple reader** processes/threads may access the shared data simultaneously without a problem

# Reader-Writer Lock

RWlock.

```
1  typedef struct _rwlock_t {
2      sem_t lock;          // binary semaphore (basic lock)
3      sem_t writelock;    // used to allow ONE writer or MANY readers
4      int   readers;      // count of readers reading in critical section
5  } rwlock_t;
6
7  void rwlock_init(rwlock_t *rw) {
8      rw->readers = 0;
9      sem_init(&rw->lock, 0, 1);
10     sem_init(&rw->writelock, 0, 1);
11 }
12
13 void rwlock_acquire_readlock(rwlock_t *rw) {
14     sem_wait(&rw->lock);
15     rw->readers++;
16     if (rw->readers == 1)
17         sem_wait(&rw->writelock); // first reader acquires writelock
18     sem_post(&rw->lock);
19 }
20
21 void rwlock_release_readlock(rwlock_t *rw) {
22     sem_wait(&rw->lock);
23     rw->readers--;
24     if (rw->readers == 0)
25         sem_post(&rw->writelock); // last reader releases writelock
26     sem_post(&rw->lock);
27 }
28
29 void rwlock_acquire_writelock(rwlock_t *rw) {
30     sem_wait(&rw->writelock);
31 }
32
33 void rwlock_release_writelock(rwlock_t *rw) {
34     sem_post(&rw->writelock);
35 }
```

# Reader-Writer Lock

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1  typedef struct _rwlock_t {
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7  void rwlock_init(rwlock_t *rw) { ←  
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Initially, # readers is 0  
binary sem lock set to 1  
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31 }  
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34     sem_post(&rw->writelock);  
35 }
```

Initially, # readers is 0  
binary sem lock set to 1  
writelock set to 1

Writer cannot  
be in CS when  
readers are!

# Readers-Writers Problem: Writer Thread

```
rwlock_acquire_writelock(rw);  
...  
write is performed  
...  
rwlock_release_writelock(rw);
```

The diagram illustrates the access pattern of a writer thread to a shared resource. It shows two lines of C code: `rwlock_acquire_writelock(rw);` at the top and `rwlock_release_writelock(rw);` at the bottom. A blue wavy line underlines the first line, and another underlines the second line. Between these underlined lines, there are three horizontal ellipses: one above the acquisition line, one below the release line, and one between them. To the right of the code, a blue arrow points downwards, labeled with the letters "CS", indicating the Critical Section being entered or exited.

# Readers-Writers Problem: Reader Thread

```
rwlock_acquire_readlock(rw)
...
read is performed
...
rwlock_release_readlock(rw)
```

The diagram illustrates the access pattern of a reader thread to a critical section (CS). It shows a sequence of operations:

- `rwlock_acquire_readlock(rw)`: A blue wavy line underlines the word `rwlock`.
- ...  
read is performed  
...*: The text is centered between two ellipses.
- `rwlock_release_readlock(rw)`: A blue wavy line underlines the word `rwlock`.

A blue arrow points from the end of the `read is performed` section down to the label `CS`, indicating that the entire read operation is contained within the critical section.

**Well, is this solution Okay?**

# Readers-Writers Problem: Reader Thread

```
rwlock_acquire_readlock(rw)
...
read is performed
...
rwlock_release_readlock(rw)
```

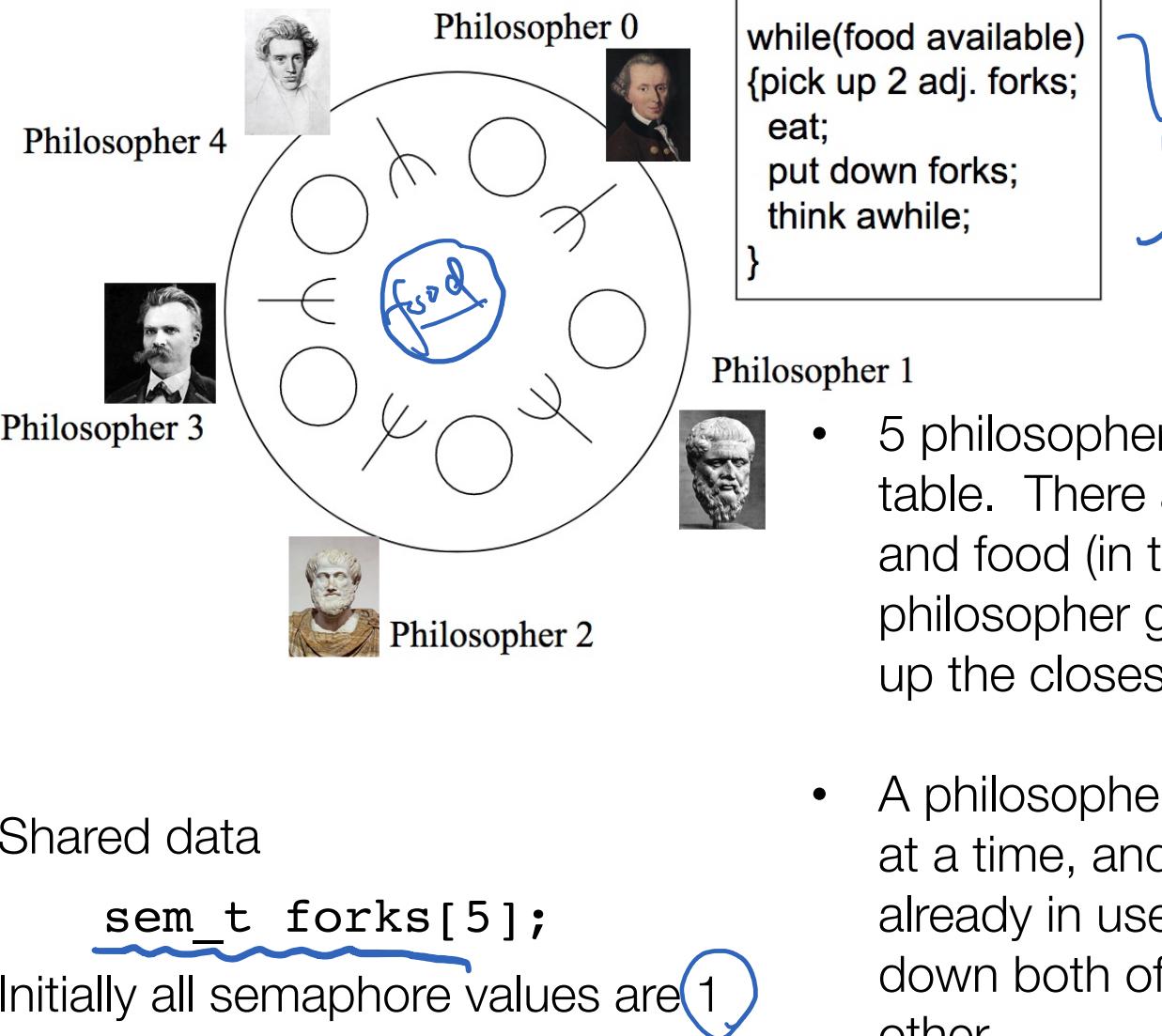
**Well, is this solution Okay?**  
A: Technically it works. But  
**starvation** may happen

# Starvation

- A process/thread that is forced to wait **indefinitely** in a synchronization program is said to be subject to **starvation**
  - In some execution scenarios, that process does not make any progress
  - Deadlocks imply starvation, but the reverse is not true

# Dining-Philosophers Problem

# Dining-Philosophers Problem



# Dining-Philosophers Problem

- The basic loop of a philosopher

```
while (1) {  
    think();  
    getforks(); → ??  
    eat(); → Critical section  
    putforks(); → ??  
}
```

# The Helper Functions

```
int left(int p) { return R; }  
int right(int p) { return (p + 1) % 5; }
```

sem\_t forks[5]

- Each fork initialized to 1

```
1 void getforks() {  
2     sem_wait(forks[_left(p)_]);  
3     sem_wait(forks[_right(p)_]);  
4 }  
5  
6 void putforks() {  
7     sem_post(forks[_left(p)_]);  
8     sem_post(forks[_right(p)_]);  
9 }
```

**Is this solution correct?**

# Simplest Example of A Deadlock

Thread 0

```
sem_wait(fork[0])  
sem_wait(fork[1])  
sem_signal(fork[0])  
sem_signal(fork[1])
```

Interleaving

Thread 1

T<sub>0</sub> [f<sub>0</sub> f<sub>1</sub>] T<sub>1</sub>

```
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_signal(fork[1])  
sem_signal(fork[0])
```

# Simplest Example of A Deadlock

Thread 0

```
sem_wait(fork[0])
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```

Interleaving

```
sem_wait(fork[0])
```

Thread 1

```
sem_wait(fork[1])
sem_wait(fork[0])
sem_signal(fork[1])
sem_signal(fork[0])
```

# Simplest Example of A Deadlock

Thread 0

```
sem_wait(fork[0])
sem_wait(fork[1])
sem_signal(fork[0])
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```

Interleaving

```
sem_wait(fork[0])
sem_wait(fork[1])
```

Thread 1

```
sem_wait(fork[1])
sem_wait(fork[0])
sem_signal(fork[1])
sem_signal(fork[0])
```

# Simplest Example of A Deadlock

Thread 0

```
sem_wait(fork[0])
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```

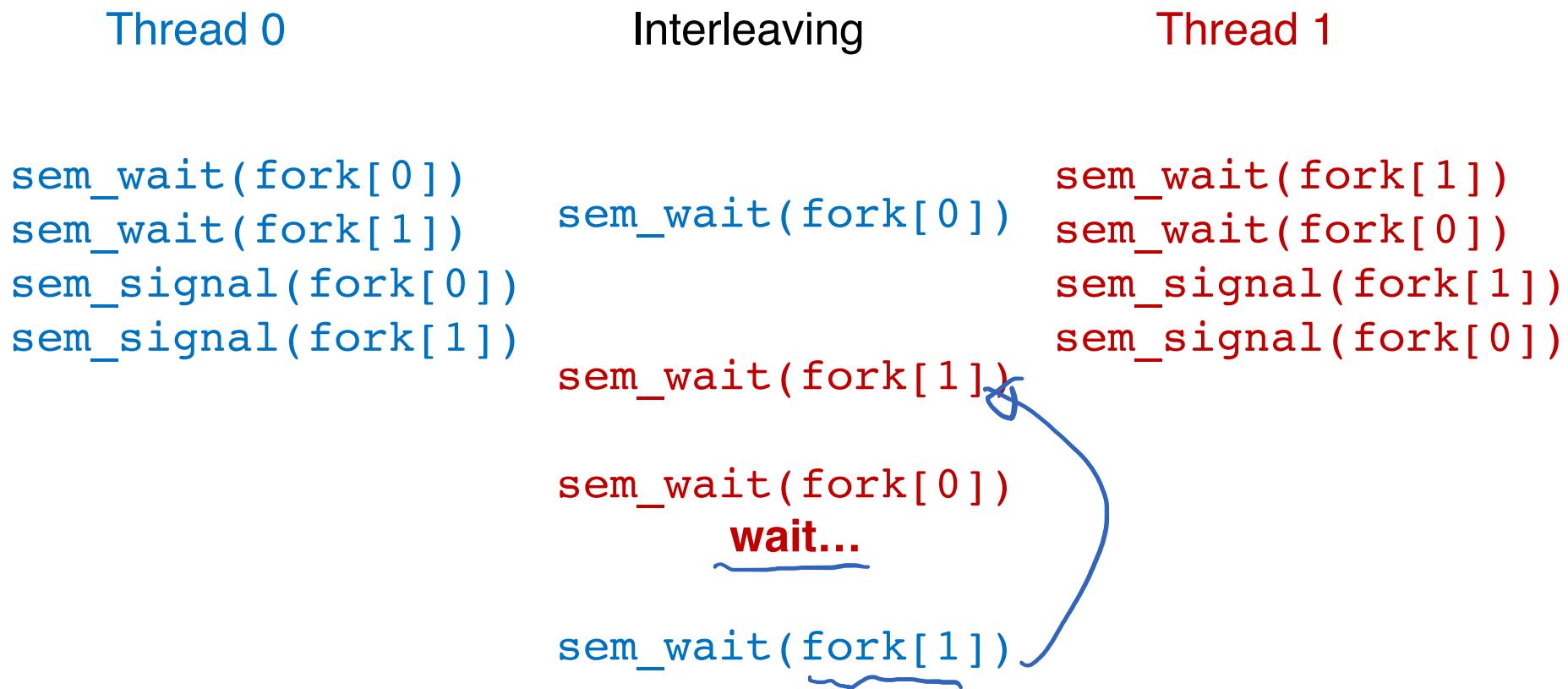
Interleaving

```
sem_wait(fork[0])
sem_wait(fork[1])
sem_wait(fork[0])
```

Thread 1

```
sem_wait(fork[1])
sem_wait(fork[0])
sem_signal(fork[1])
sem_signal(fork[0])
```

# Simplest Example of A Deadlock



# Simplest Example of A Deadlock

Thread 0

```
sem_wait(fork[0])
sem_wait(fork[1])
sem_signal(fork[0])
sem_signal(fork[1])
```

Interleaving

```
sem_wait(fork[0])
sem_wait(fork[1])
```

Thread 1

```
sem_wait(fork[1])
sem_wait(fork[0])
sem_signal(fork[1])
sem_signal(fork[0])
```

```
sem_wait(fork[0])
wait...
```

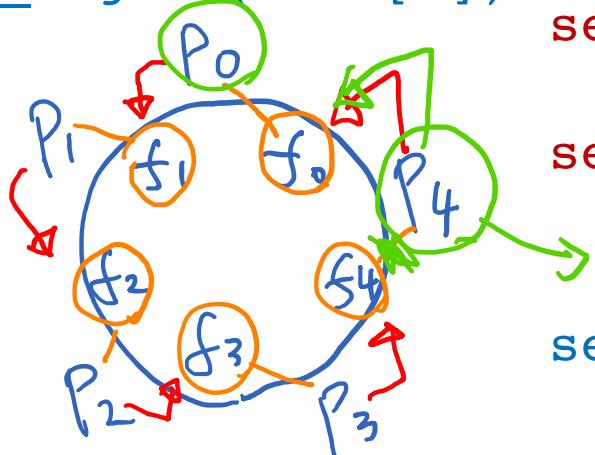
```
sem_wait(fork[1])
wait...
```

# Simplest Example of A Deadlock

Q: Would the previous 5DP implementation cause exactly the same form of a deadlock as shown below?

Thread 0

```
sem_wait(fork[0])  
sem_wait(fork[1])  
sem_signal(fork[0])  
sem_signal(fork[1])
```



Interleaving

```
sem_wait(fork[0])  
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_wait(fork[1])  
wait...  
sem_wait(fork[0])  
wait...
```

Thread 1

```
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_signal(fork[1])  
sem_signal(fork[0])
```

# Review: Conditions for Deadlocks

- Mutually exclusive access of shared resources
  - Binary semaphore `fork[ 0 ]` and `fork[ 1 ]`

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- Mutually exclusive access of shared resources
  - Binary semaphore `fork[ 0 ]` and `fork[ 1 ]`
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  - Thread 0 waits for Thread 1 to signal(`fork[ 1 ]`) and
  - Thread 1 waits for Thread 0 to signal(`fork[ 0 ]`)

# Review: Conditions for Deadlocks

- Mutually exclusive access of shared resources
  - Binary semaphore `fork[ 0 ]` and `fork[ 1 ]`
- Circular waiting
  - Thread 0 waits for Thread 1 to signal(`fork[ 1 ]`) and
  - Thread 1 waits for Thread 0 to signal(`fork[ 0 ]`)
- Hold and wait
  - Holding either `fork[ 0 ]` or `fork[ 1 ]` while waiting on the other

# Review: Conditions for Deadlocks

- Mutually exclusive access of shared resources
  - Binary semaphore `fork[ 0 ]` and `fork[ 1 ]`
- Circular waiting
  - Thread 0 waits for Thread 1 to signal(`fork[ 1 ]`) and
  - Thread 1 waits for Thread 0 to signal(`fork[ 0 ]`)
- Hold and wait
  - Holding either `fork[ 0 ]` or `fork[ 1 ]` while waiting on the other
- No preemption
  - Neither `fork[ 0 ]` and `fork[ 1 ]` can be removed from their respective holding threads

# Why 5DP is Interesting?

- How to eat with your fellows without causing deadlocks
  - Circular arguments (the **circular wait condition**)
  - Not giving up on firmly held things (**no preemption**)
  - Infinite patience with half-baked schemes (**hold some & wait for more**)
- Why starvation exists and what we can do about it?

# Why 5DP is Interesting?

- ~~How to eat with your fellows without causing deadlocks~~ — **How to mess with your fellows!**
  - Circular arguments (the **circular wait condition**)
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# Dijkstra's Solution: Break the Circular Wait Condition

- Change how forks are acquired by at least one of the philosophers
- Assume P0 – P4, 4 is the highest number

```
1 void getforks() {  
2     if (p == 4) {  
3         sem_wait(forks[right(p)]);  
4         sem_wait(forks[left(p)]);  
5     } else {  
6         sem_wait(forks[left(p)]);  
7         sem_wait(forks[right(p)]);  
8     }  
9 }
```

# Again, Starvation

- Subtle difference between deadlock and starvation
  - Once a set of processes are in a deadlock, there is **no future execution sequence** that can get them out of it!
  - In starvation, there does exist hope – some execution order may be favorable to the starving process although no guarantee it would ever occur
- Rollback and retry are prone to starvation
- Continuous arrival of higher priority process is another common starvation situation

# **Building a Semaphore w/ CV Worksheet**