Scheduling: FIFO and SJF

CS 571: Operating Systems (Spring 2020)
Lecture 4

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Some material taken/derived from:
- Wisconsin CS-537 materials created by Remzi Arpaci-Dusseau.

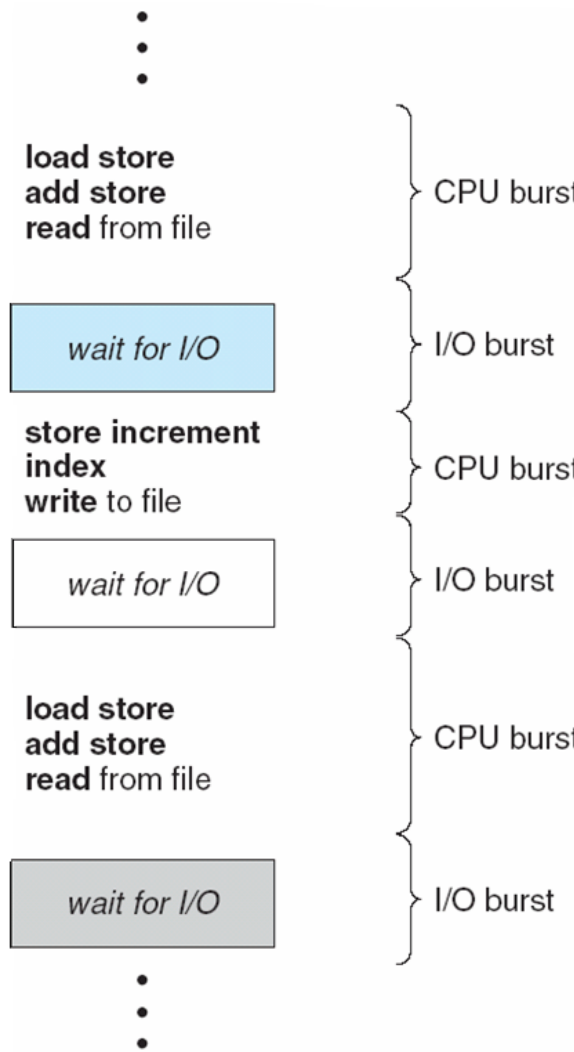
Outline

- Basic concept
- Scheduling criteria
- Scheduling algorithms
  - First In, First Out (FIFO)
  - Shortest Job First (SFJ)
  - Shortest Time-to-Completion First (STCF)
  - Round Robin (RR)
  - Priority
  - Multi-Level Feedback Queue (MLFQ)

Basic Concepts

- During its lifetime, a process goes through a sequence of CPU and I/O bursts
- The CPU scheduler (a.k.a. short-term scheduler) will select one of the processes in the ready queue for execution
- The CPU scheduler algorithm may have tremendous effects on the system performance
  - Interactive systems: Responsiveness
  - Real-time systems: Not missing the deadlines

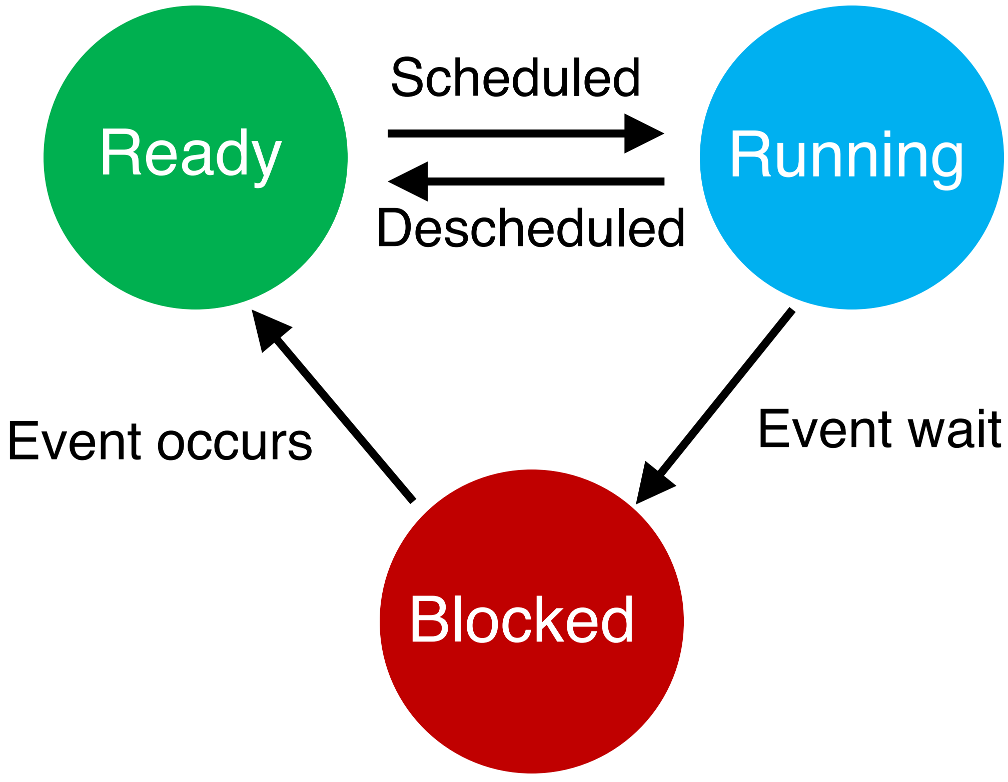
Alternating Sequence of CPU and I/O Bursts



When to Schedule?

- Under the simple process state transition model, CPU scheduler can be potentially invoked at five different points:
  1. When a process switches from the new state to the ready state
  2. When a process switches from the running state to the waiting (or blocked) state
  3. When a process switches from the running state to the ready state
  4. When a process switches from the waiting state to the ready state
  5. When a process terminates

Process State Transitions



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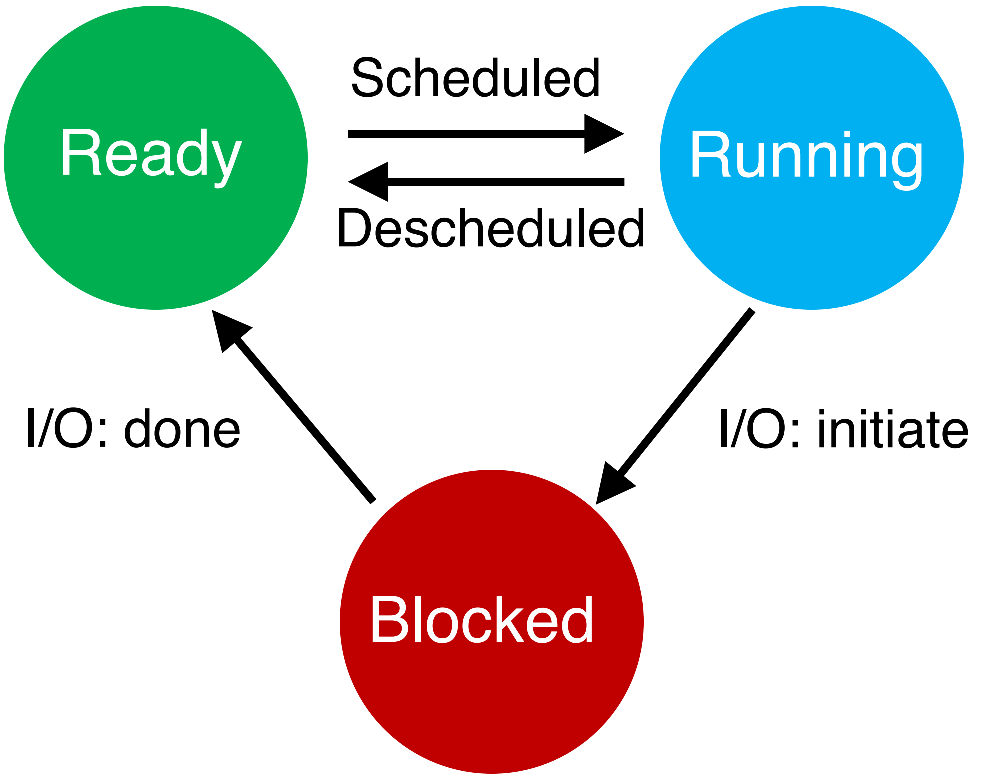


Dispatcher

- Dispatcher module gives control of the CPU to the process selected by the short-term scheduler; this involves:
  - switching context
  - switching to user mode
  - jumping to the proper (previously saved) location in the user program to restart that program

- Scheduler → Policy: When and how to schedule
- Dispatcher → Mechanism: Actuator following the commands of the scheduler

Scheduling Metrics

- To compare the performance of scheduling algorithms
    - CPU utilization – percentage of time CPU is busy executing jobs
    - Throughput – # of processes that complete their execution per time unit
    - Turnaround time – amount of time to execute a particular process
    - Waiting time – amount of time a process has been waiting in the ready queue
    - Response time – amount of time it takes from when a request was submitted until the first response is produced, not the complete output

Optimization Goals

- To **maximize**:
  - Maximize the CPU utilization
  - Maximize the throughput

- To **minimize**:
  - Minimize the (average) turnaround time
  - Minimize the (average) waiting time
  - Minimize the (average) response time

Waiting Time

- Waiting time definition

$$T_{waiting} = T_{start} - T_{arrival}$$

- Average waiting time = Sum($T_{waiting}$)/ #processes

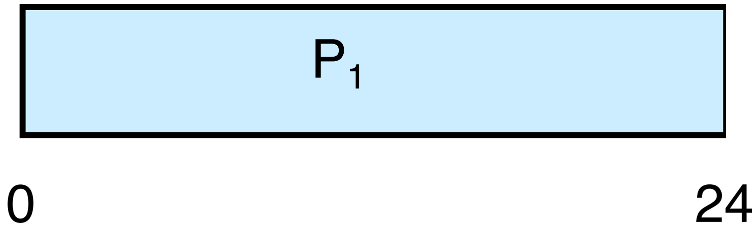
- For now, we assume
  - Average waiting time is the performance measure
  - Only one CPU burst (e.g., in milliseconds or ms) per process
  - Only CPU, No I/O
  - All processes arrive at the same time
  - Once started, each process runs to completion

First In, First Out (FIFO)

First-In-First-Out (FIFO)

Process	Burst Time
P_1	24

- Suppose that the processes arrive in order: P_1 , P_2 , P_3
  The Gantt Chart for the schedule:



First-In-First-Out (FIFO)

Process	Burst Time
P_1	24
P_2	3

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Process	Burst Time
P_1	24
P_2	3
P_3	3

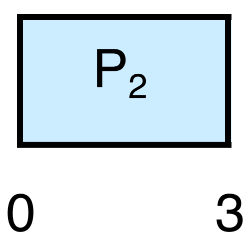
- Suppose that the processes arrive in order: P_1 , P_2 , P_3
  The Gantt Chart for the schedule:

- Waiting time for $P_1 = 0$; $P_2 = 24$; $P_3 = 27$
- Average waiting time: 17

FIFO (cont.)

- Suppose that the processes arrive in order P_2 , P_3 , P_1

- The Gantt chart for the schedule:



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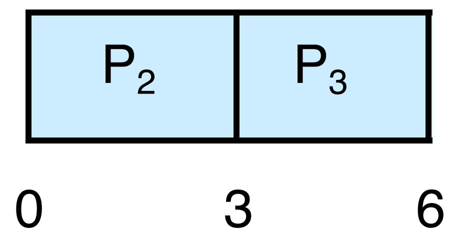

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FIFO (cont.)

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- The Gantt chart for the schedule:

- Waiting time for $P_1 = 6;$ $P_2 = 0;$ $P_3 = 3$
- Average waiting time: (6 + 0 + 3)/3 = 3

FIFO (cont.)

- Suppose that the processes arrive in order P_2 , P_3 , P_1
- The Gantt chart for the schedule:

- Waiting time for $P_1 = 6$; $P_2 = 0$; $P_3 = 3$
- Average waiting time: $(6 + 0 + 3)/3 = 3$
- Problems:
  - **Convoy effect** (short processes behind long processes)
  - Non-preemptive: Not suitable for time-sharing systems

Shortest Job First (SJF)

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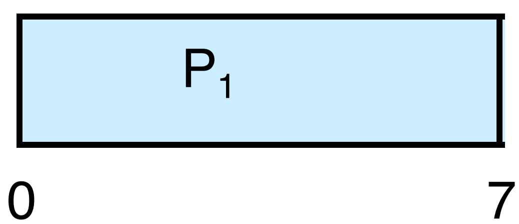
- Associate with each process the length of its next CPU burst
- The CPU is assigned to the process with the smallest (next) CPU burst (run_time)

- Two schemes (modes):
  - Non-preemptive
  - Preemptive: Also known as the Shortest Time-to-Completion First (STCF)

Example for Non-Preemptive SJF

Process	Arrival Time	Burst Time
P_1	0.0	7
P_2	2.0	4
P_3	4.0	1
P_4	5.0	4

- SJF (non-preemptive)



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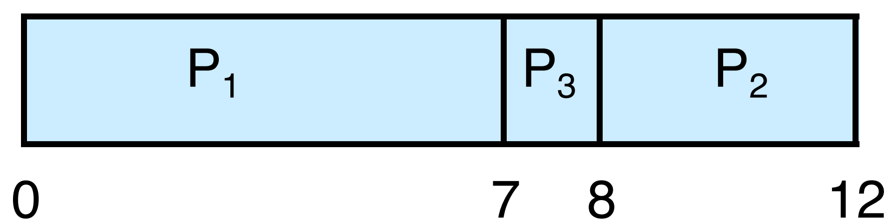

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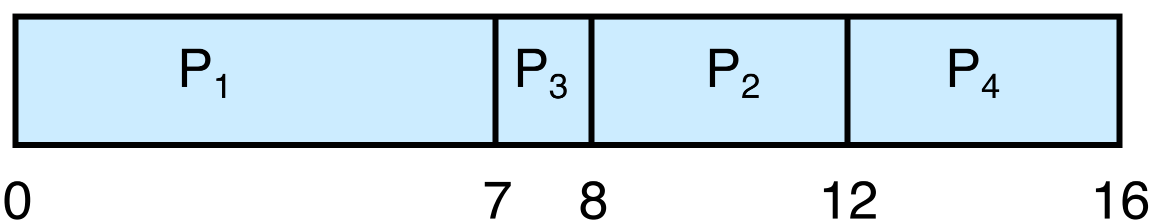

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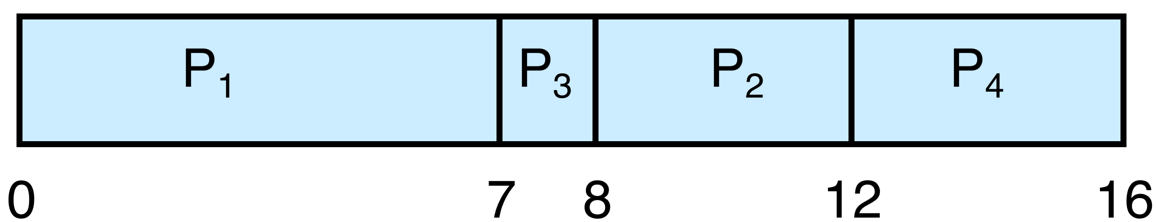


- Average waiting time = (0 + 6 + 3 + 7)/4 = 4

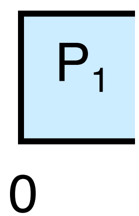
Example for Preemptive SJF (STCF)

Process	Arrival Time	Burst Time	Left Time
P_1	0.0	7	

Example for Preemptive SJF (STCF)

Process	Arrival Time	Burst Time	Left Time
P_1	0.0	7	5
P_2	2.0	4	

- SJF (preemptive)



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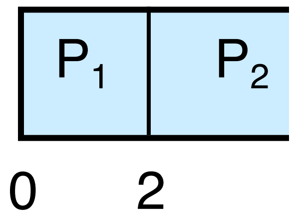

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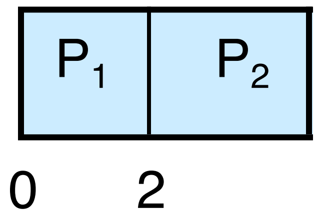

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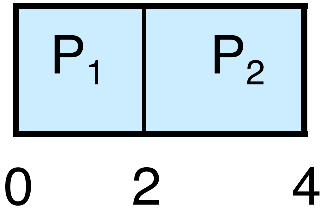

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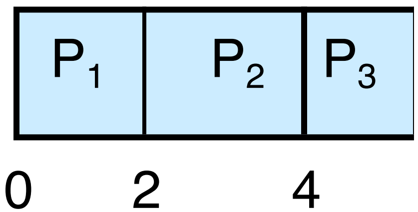

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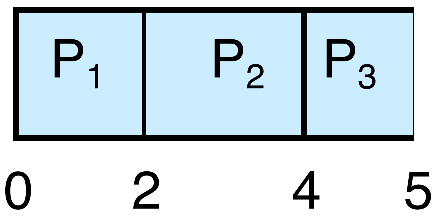

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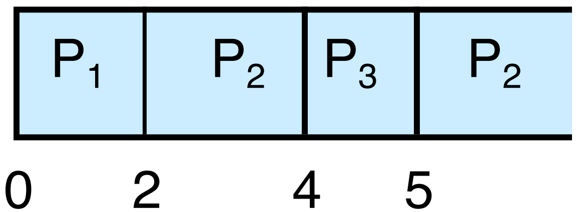

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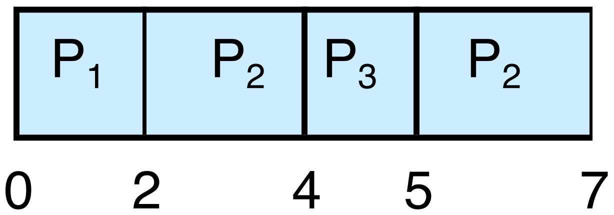

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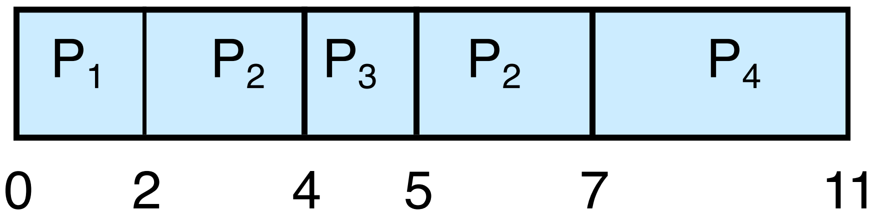

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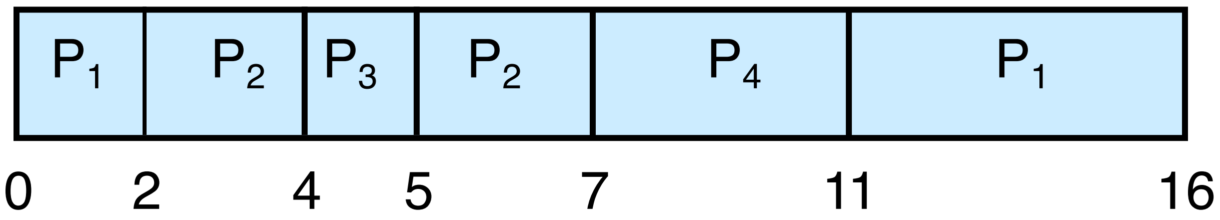

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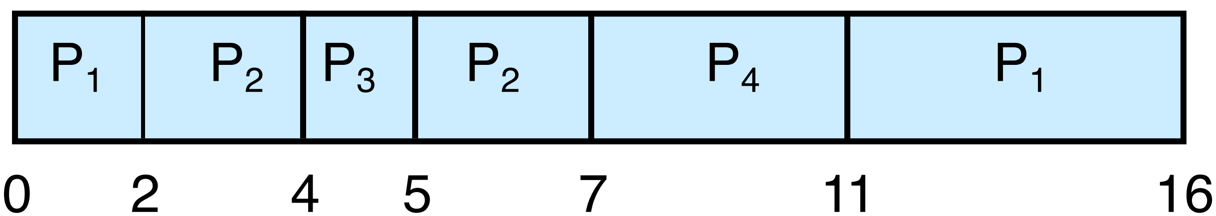

- Average waiting time = (9 + 1 + 0 + 2)/4 = 3