



Resilient Distributed Datasets: Spark

CS 675: *Distributed Systems* (Spring 2020)

Lecture 6

Yue Cheng

Some material taken/derived from:

- Matei Zaharia's NSDI'12 talk slides.
- Utah CS6450 by Ryan Stutsman.

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Announcement

- Deadline of the project report gets extended to 11:59pm next Friday, 04/03
- Doodle poll for Lab 2 demo and proposal discussion meetings (Thursday and Friday)
 - <https://doodle.com/poll/tbskp9hqaqsn7ysi>

What's good with MapReduce

- Scaled analytics to thousands of machines
- Eliminated fault-tolerance as a concern

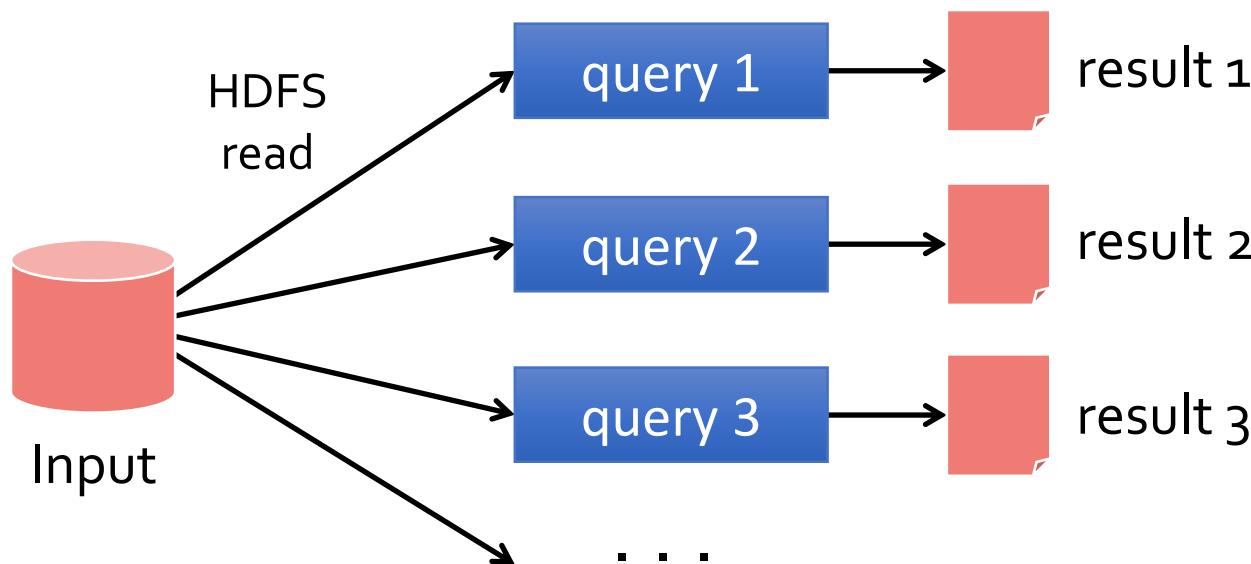
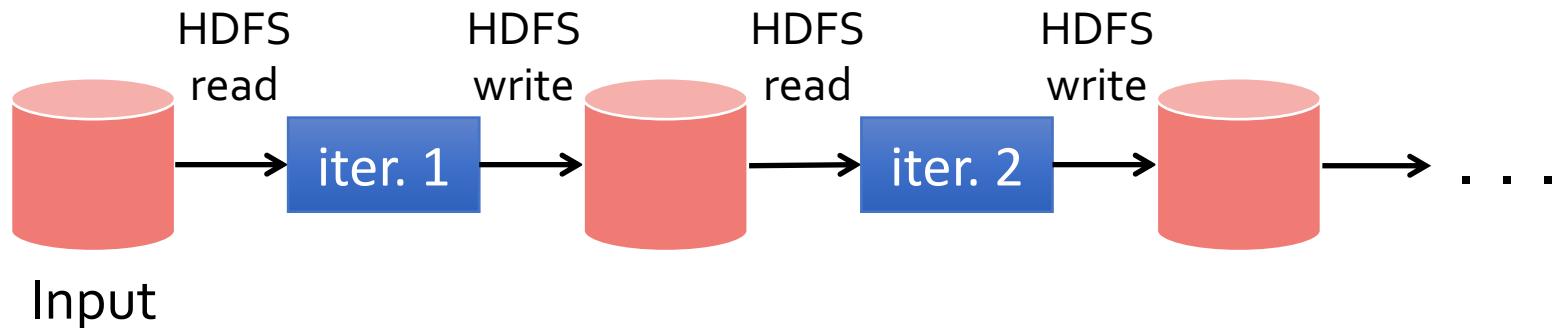
Problems with MapReduce

- Scaled analytics to thousands of machines
- Eliminated fault-tolerance as a concern
- **Not very expressive**
 - Iterative algorithms
(PageRank, Logistic Regression, Transitive Closure)
 - Interactive and ad-hoc queries
(Interactive Log Debugging)
- Lots of specialized frameworks
 - Pregel, GraphLab, PowerGraph, DryadLINQ, HaLoop...

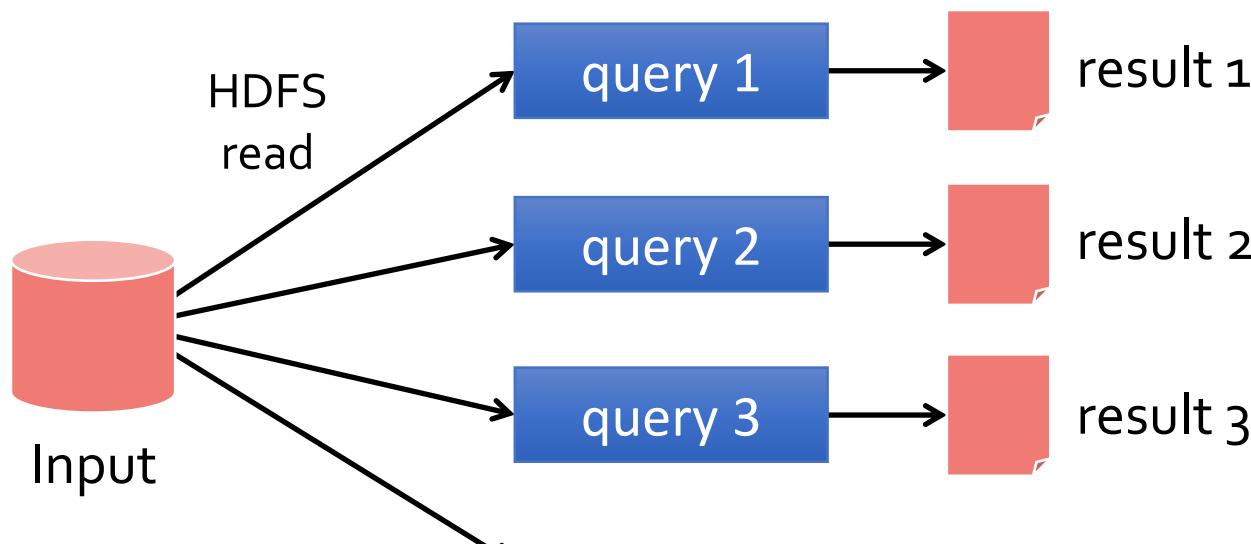
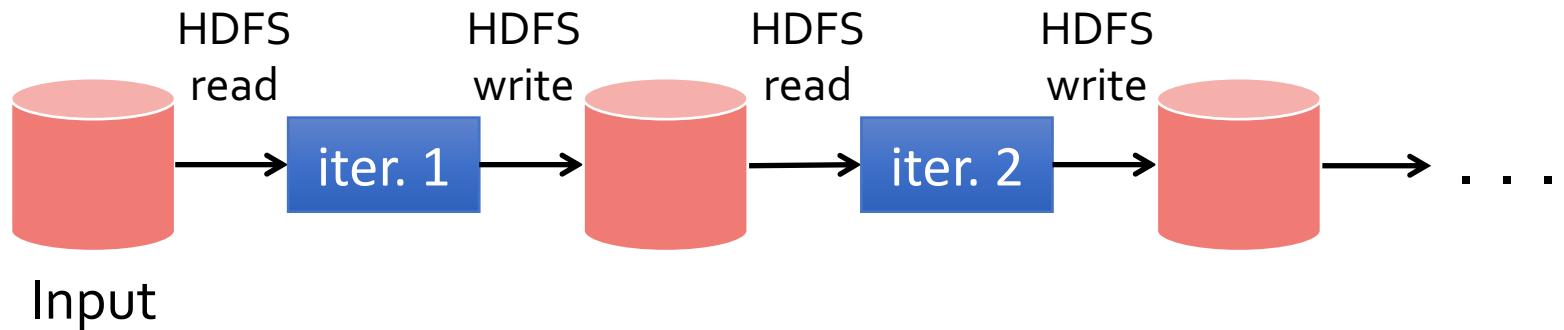
Sharing data between iterations/ops

- Only way to share data between iterations / phases is through shared storage
 - **Slow!**
- Allow operations to feed data to one another
 - Ideally, through memory instead of disk-based storage
- Need the “chain” of operations to be exposed to make this work
- Also, does this break the MR fault-tolerance scheme?
 - Retry and Map or Reduce task since idempotent

Examples

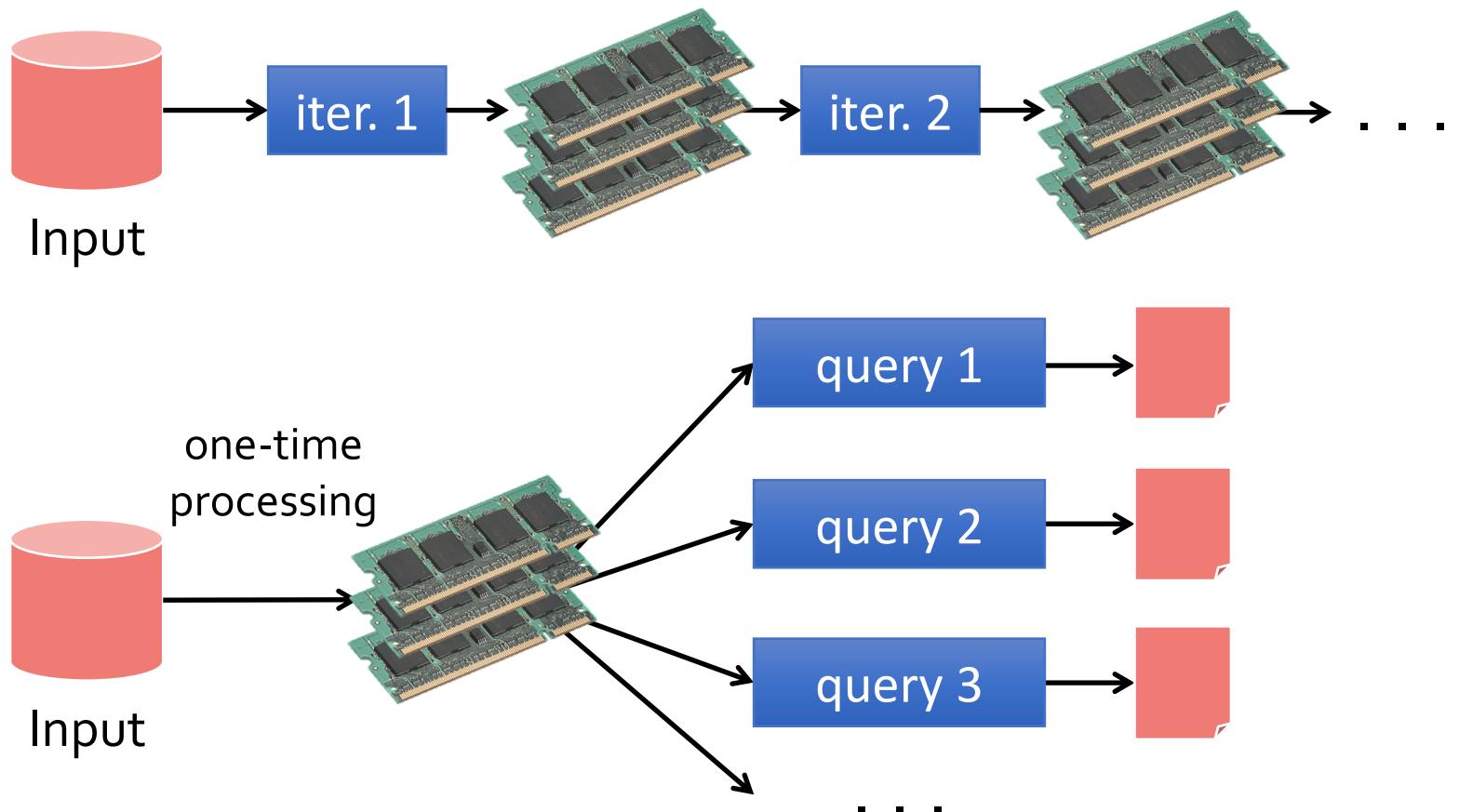


Examples

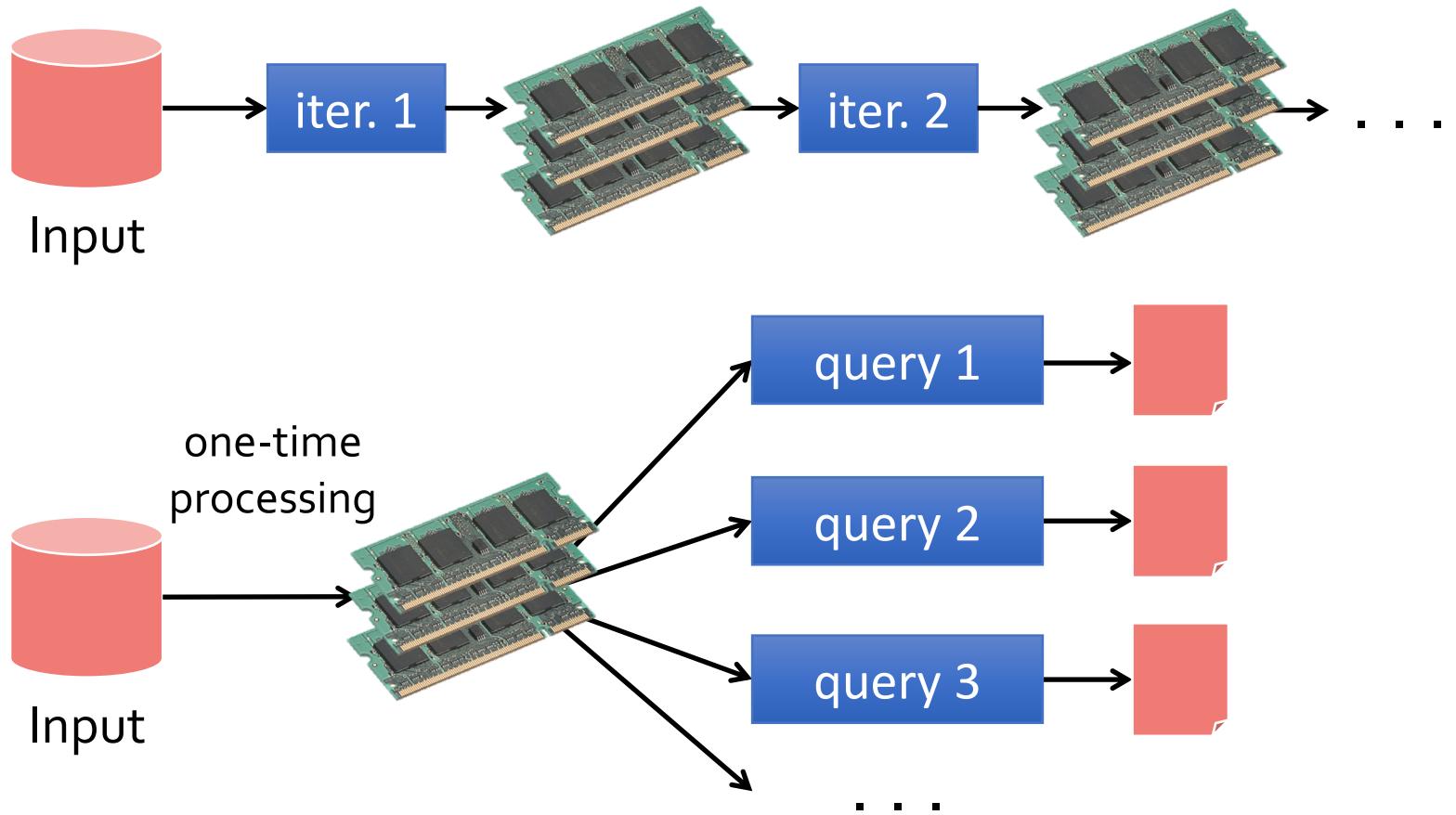


Slow due to replication and disk I/O,
but necessary for fault tolerance

Goal: In-memory data sharing



Goal: In-memory data sharing



10-100× faster than network/disk, but how to get FT?

Challenges

- How to design a distributed memory abstraction that is both **fault-tolerant** and **efficient**?

Challenges

- How to design a distributed memory abstraction that is both **fault-tolerant** and **efficient**?
- Existing storage systems allow **fine-grained** mutation to state
 - In-memory key-value stores
 - Requires replicating data or logs across nodes for fault tolerance
 - Costly for data-intensive apps
 - 10-100x slower than memory write
 - They also require costly on-the-fly replication for mutations

Challenges

- How to design a distributed memory abstraction that is both **fault-tolerant** and **efficient**?
- Existing storage systems allow **fine-grained** mutation to state

Insight: leverage similar coarse-grained approach that transforms whole dataset per operation, like MapReduce (batch processing)

- 10-100x slower than memory write
- They also require costly on-the-fly replication for mutations

Solution: Resilient Distributed Datasets (RDDs)

- Restricted form of distributed shared memory
 - **Immutable**, partitioned collections of records
 - Can only be built through *coarse-grained* deterministic *transformations* (map, filter, join, ...)
- Efficient fault recovery using *lineage*
 - Log **one operation** to apply to many elements
 - Recompute lost partitions on failure
 - No cost if nothing fails

Spark programming interface

- Scala API, exposed within interpreter as well
- RDDs
- Transformations on RDDs ($\text{RDD}_1 \rightarrow \text{RDD}_2$)
- Actions on RDDs ($\text{RDD} \rightarrow \text{output}$)
- Control over RDD partitioning (how items are split over nodes)
- Control over RDD persistence (in memory, on disk, or recompute on loss)

Transformations

Transformations
(define a new RDD)

map	flatMap
filter	union
sample	join
groupByKey	cogroup
reduceByKey	cross
sortByKey	mapValues

RDDs in terms of Scala types → Scala semantics at workers

Transformations are **lazy “thunks”**; cause no cluster action

Actions

Actions
(return a result to
driver program)

collect
reduce
count
save
lookupKey

Consumes an RDD to **produce** output
either to storage (save), or
to interpreter/Scala (count, collect, reduce)

Causes RDD lineage chain to get executed on the cluster to
produce the output
(for any missing pieces of the computation)

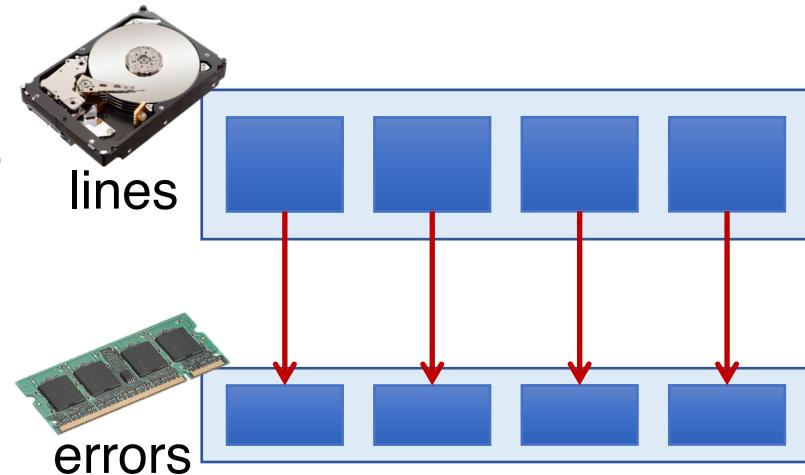
Interactive debugging

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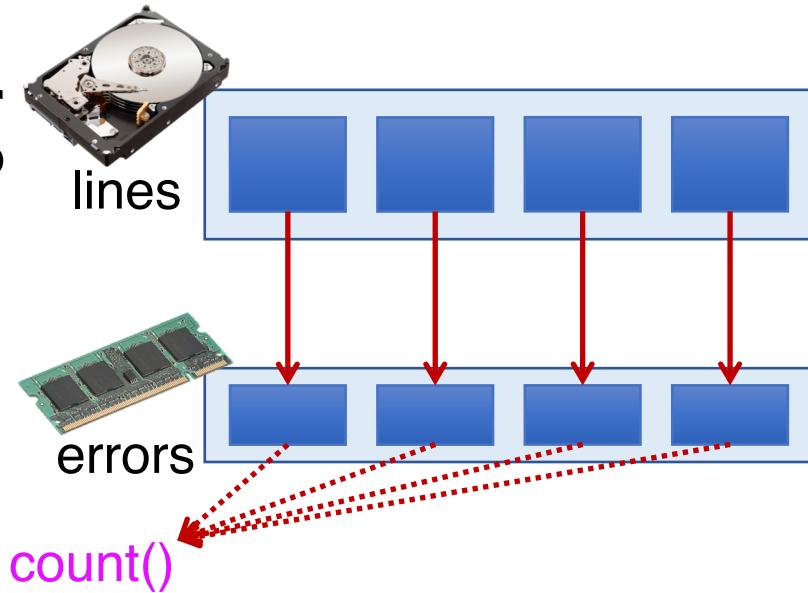
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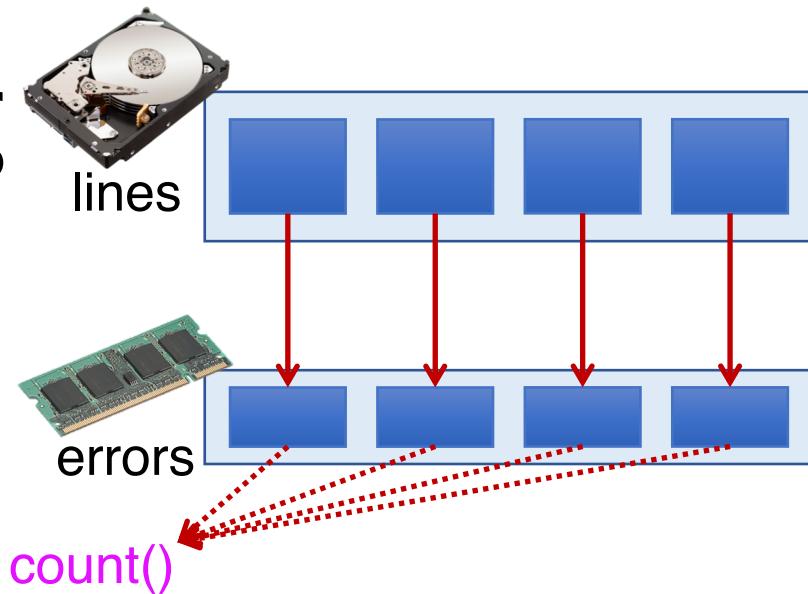


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```
errors.filter(
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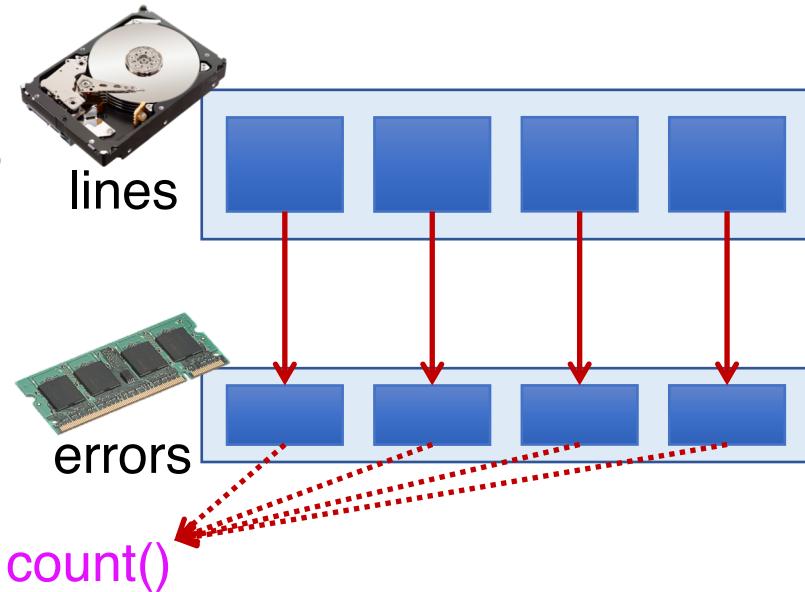


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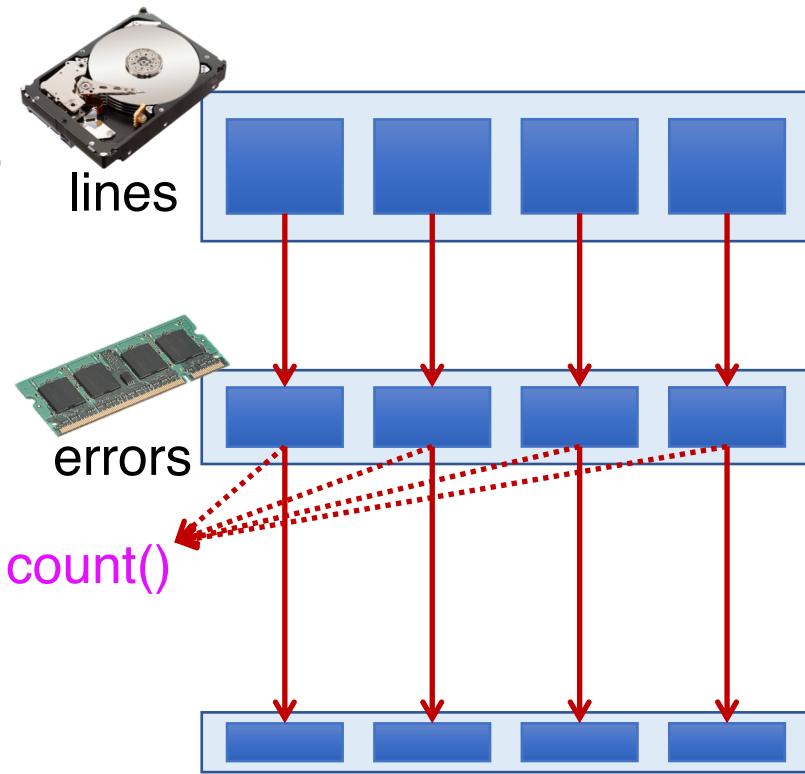


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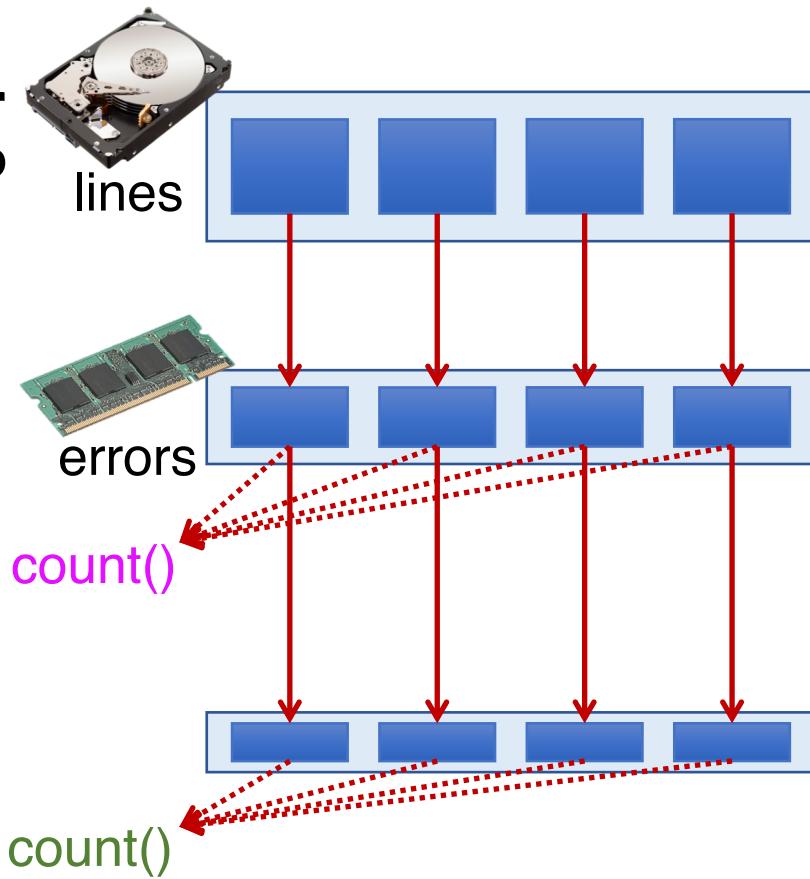


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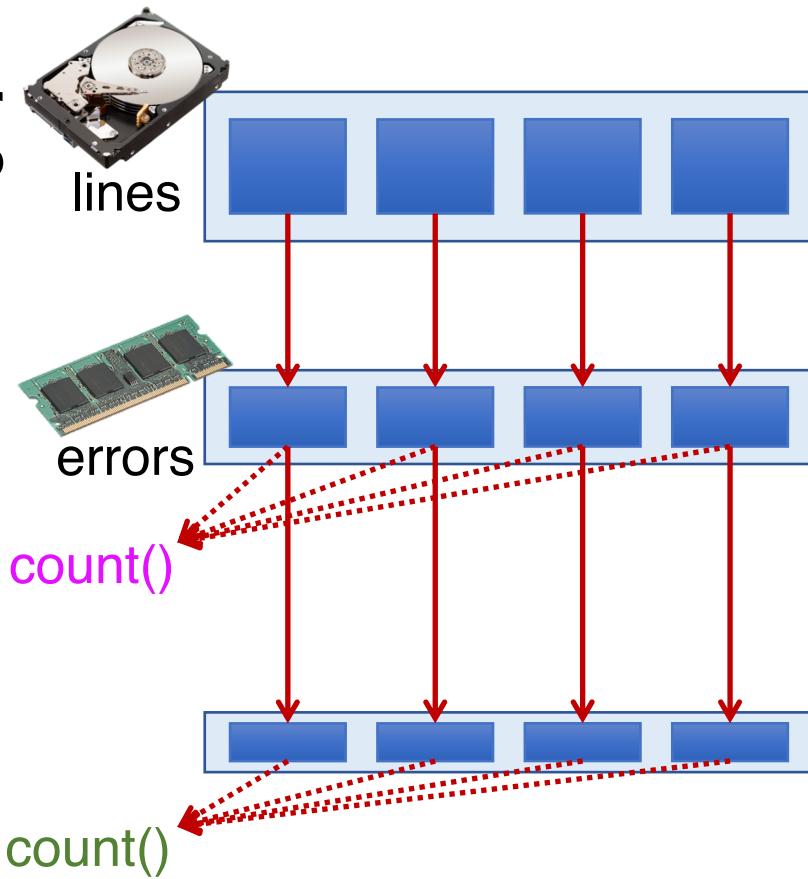


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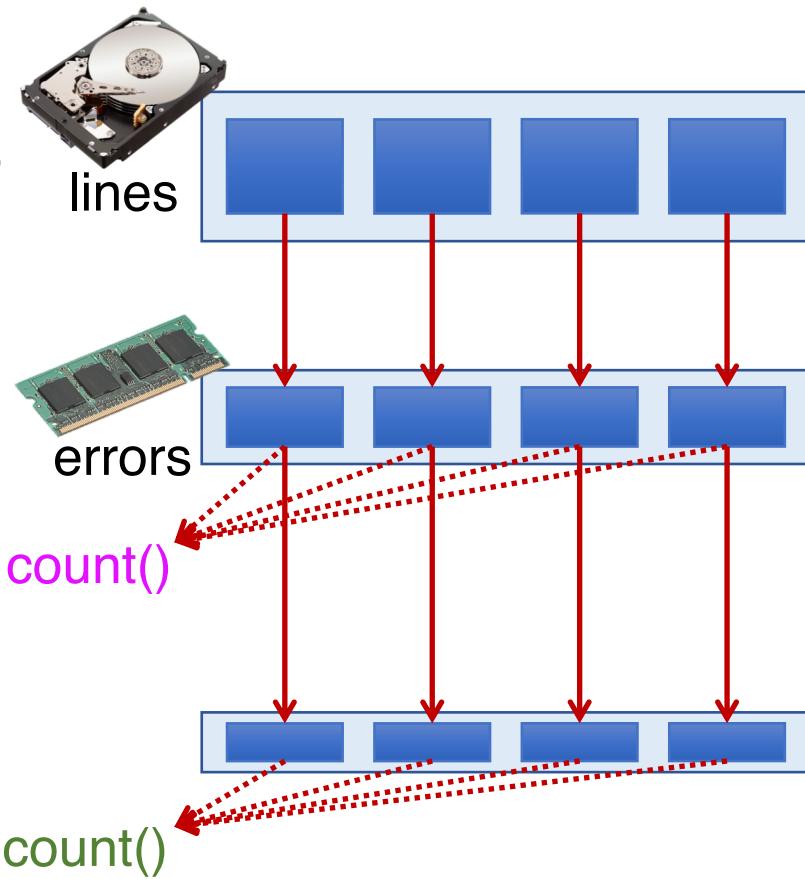


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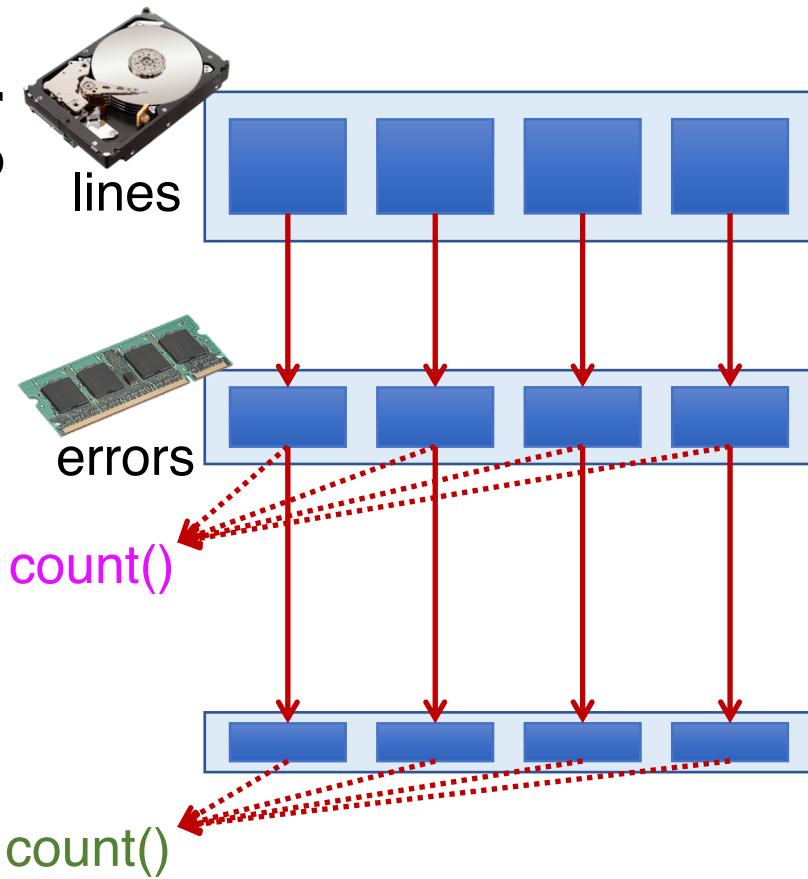


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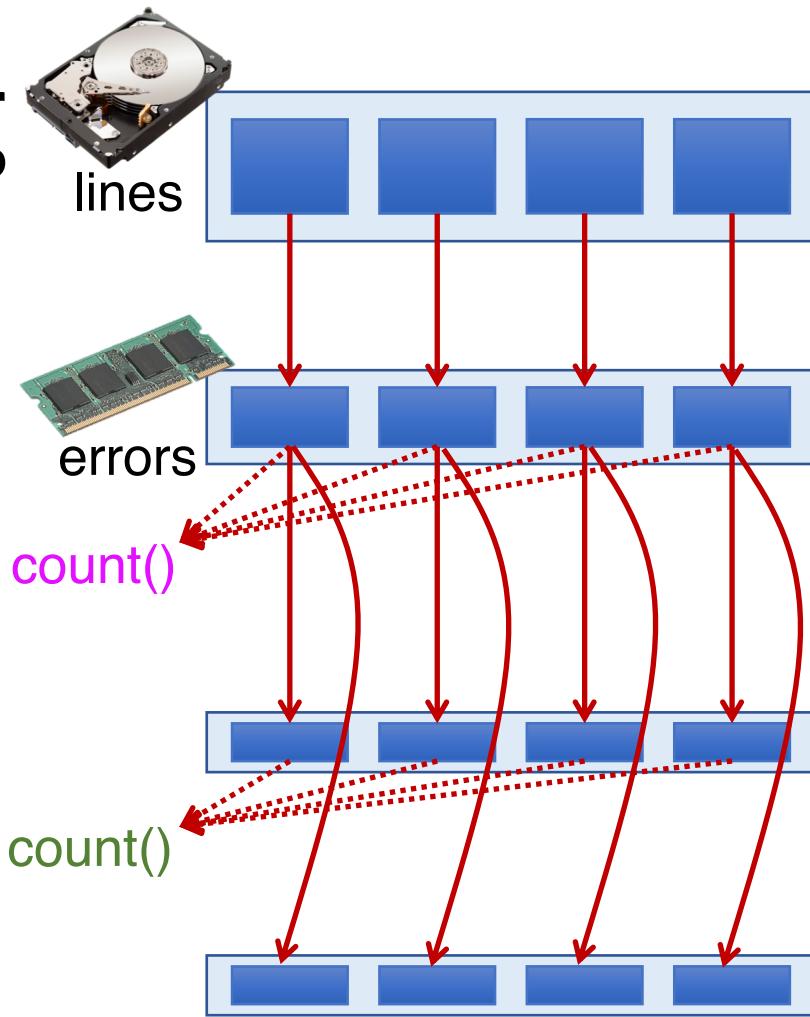


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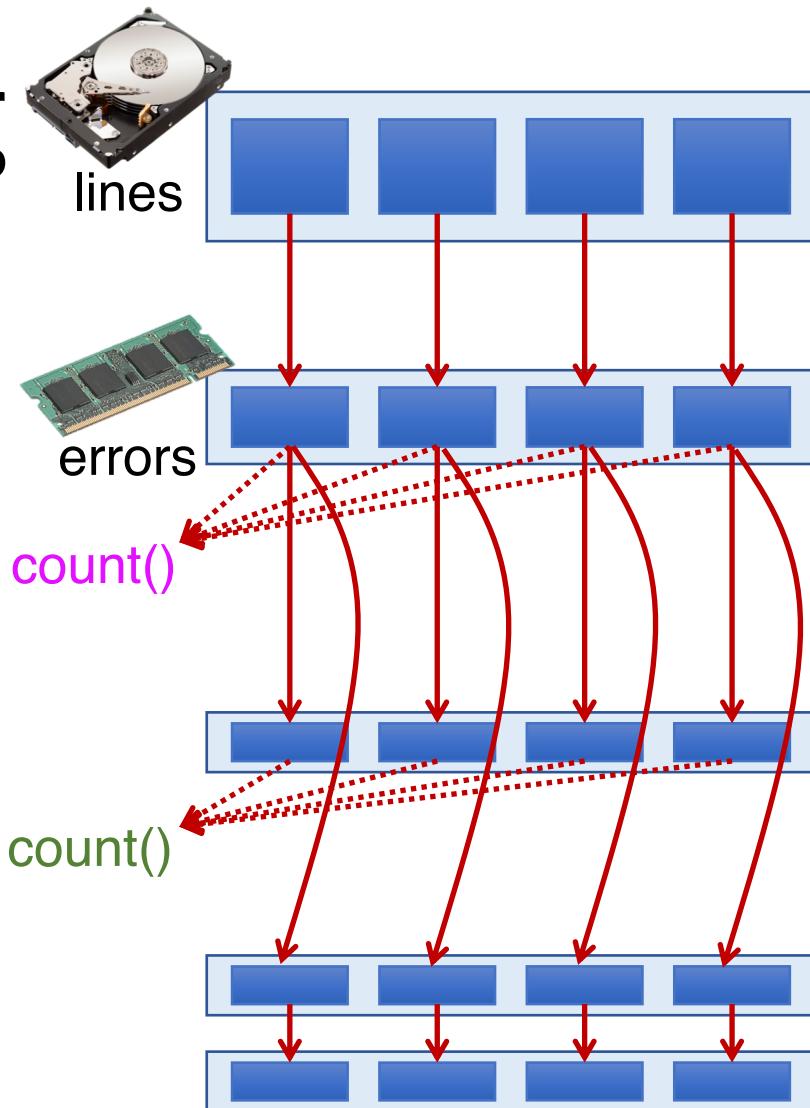


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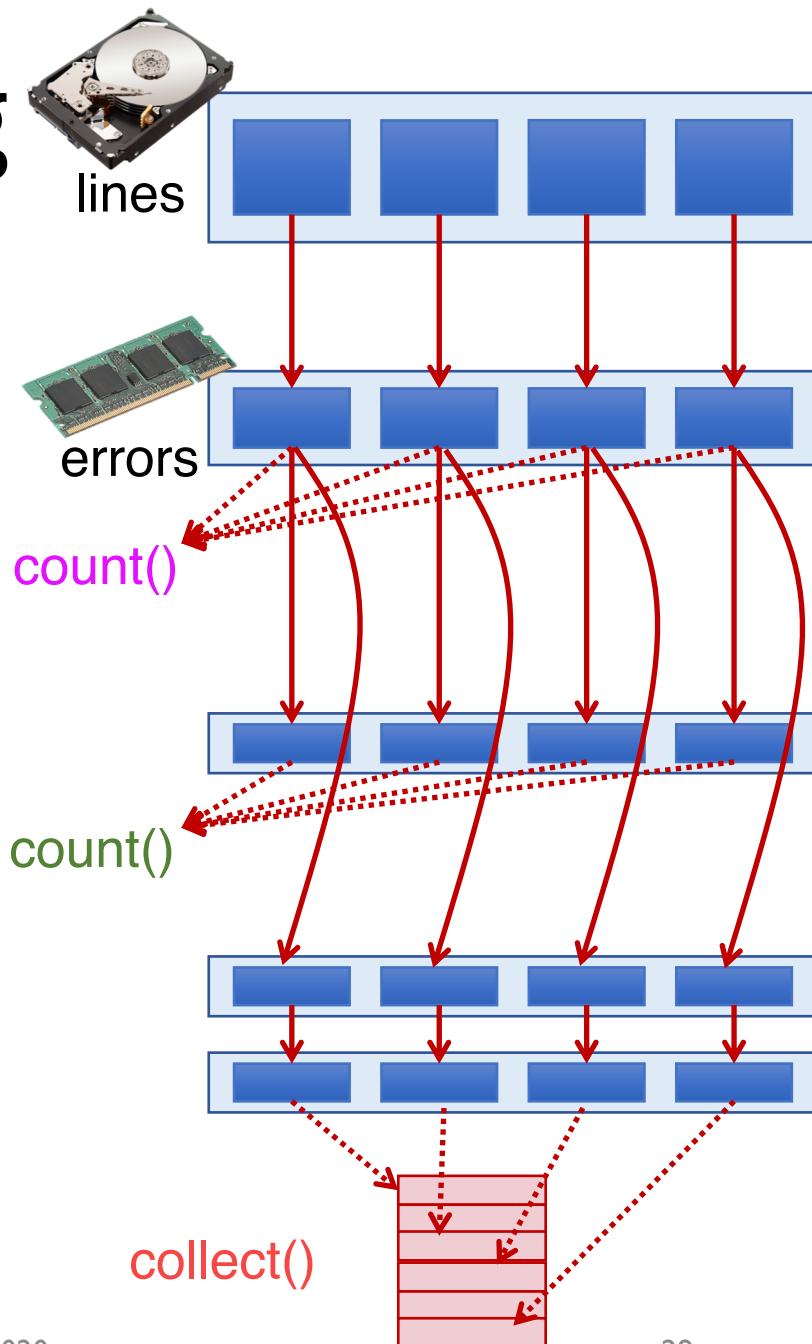


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persist()

- Not an action and not a transformation
- A scheduler hint
- Tells which RDDs the Spark schedule should materialize and whether in memory or storage
- Gives the user control over reuse/recompute/recovery tradeoffs

persist()

- Not an action and not a transformation
- A scheduler hint
- Tells which RDDs the Spark schedule should materialize and whether in memory or storage
- Gives the user control over reuse/recompute/recovery tradeoffs
- **Q:** If persist() asks for the materialization of an RDD why isn't it an action?

Lineage graph of RDDs

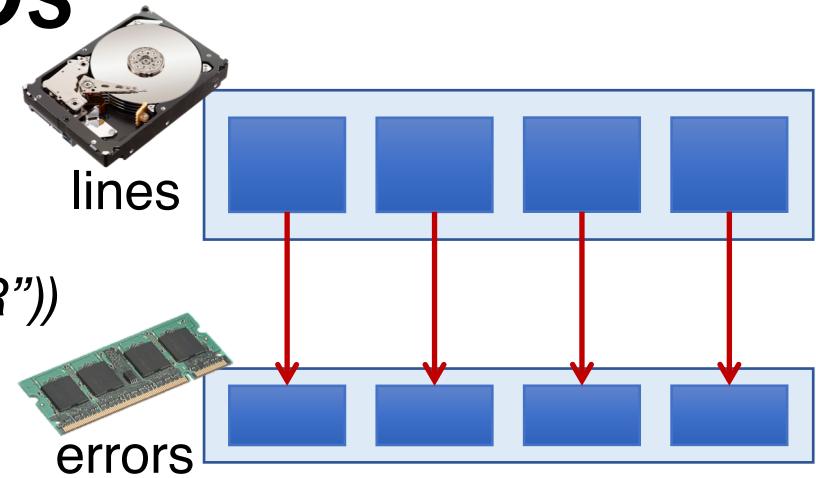
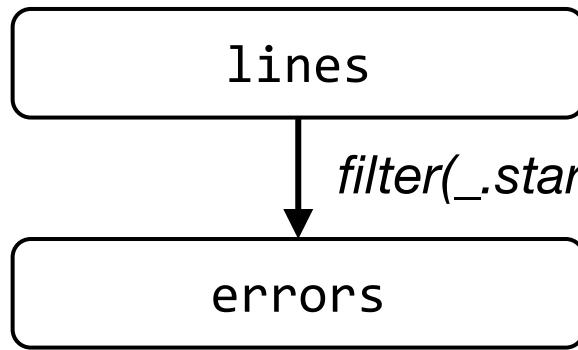
lines



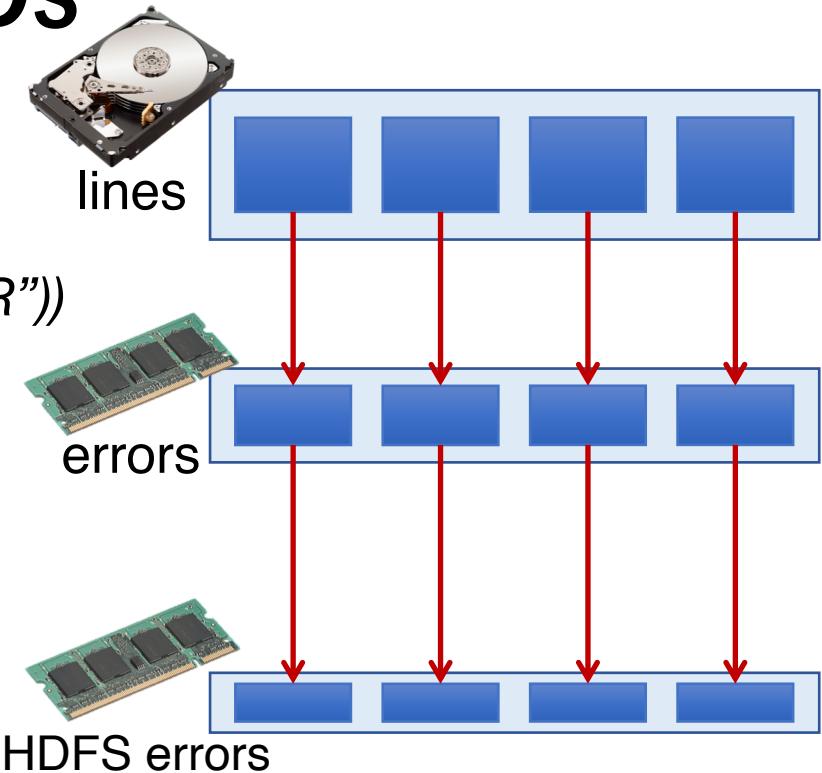
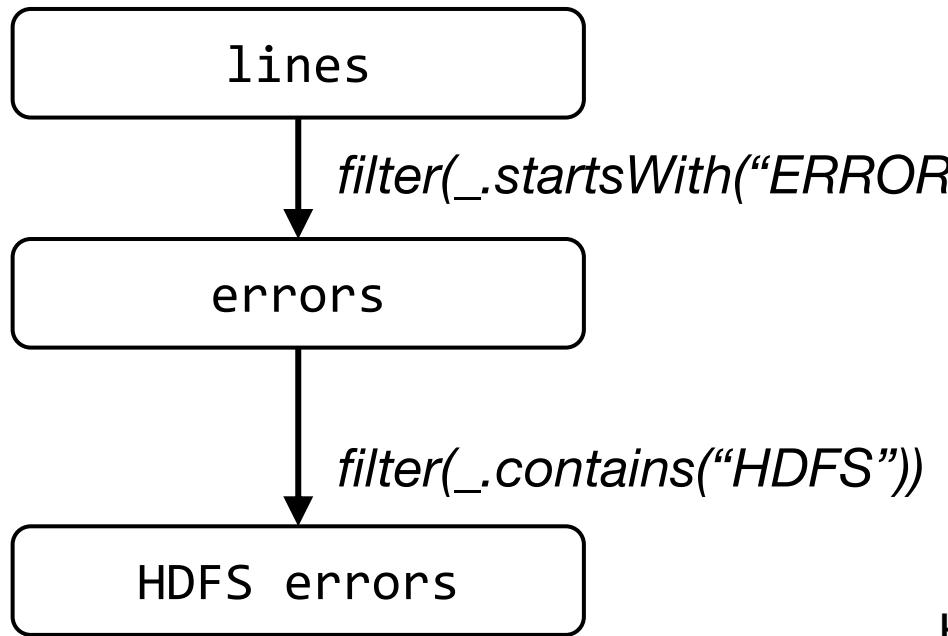
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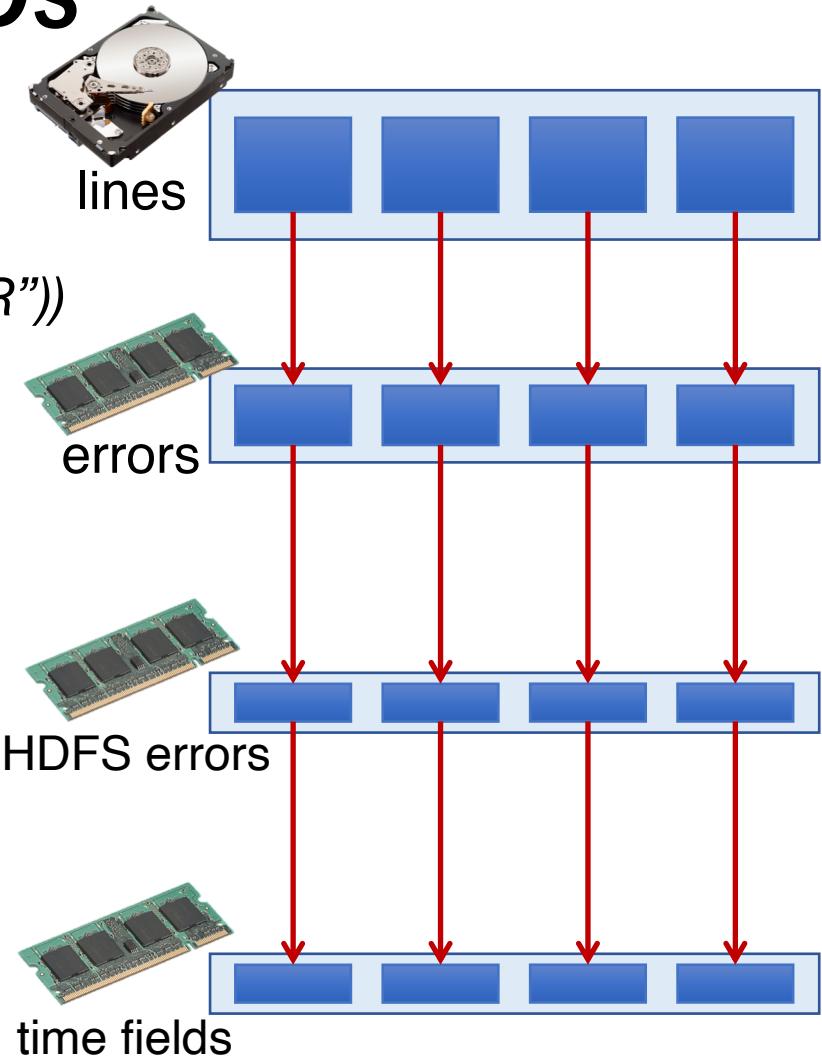
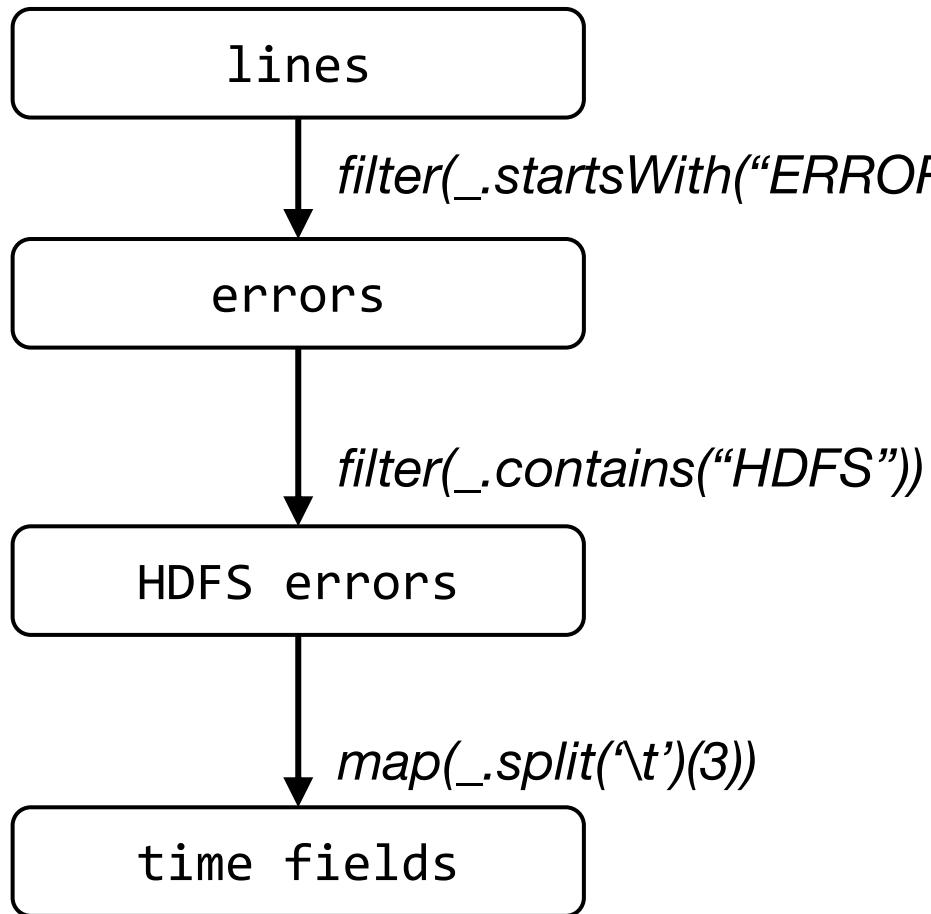
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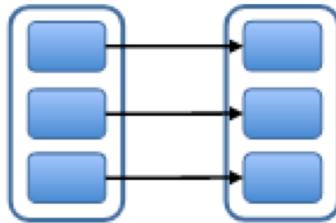


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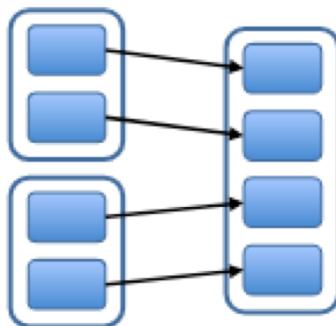


Narrow & wide dependencies

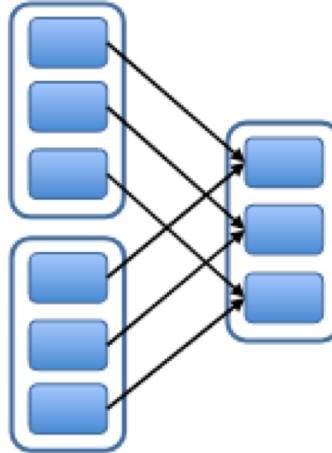
Narrow Dependencies:



map, filter

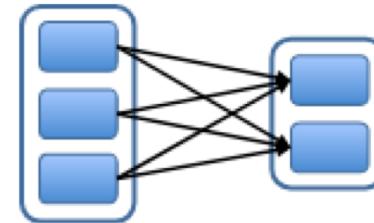


union

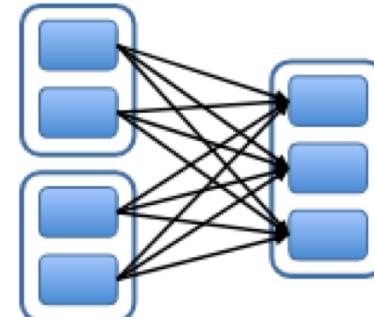


join with inputs
co-partitioned

Wide Dependencies:



groupByKey



join with inputs not
co-partitioned

Narrow: each parent partition used by at most one child partition (can partition on one machine)

Wide: multiple child partitions depend on one parent partition

Must stall for all parent data, loss of child requires whole parent RDD (not just a small # of partitions)

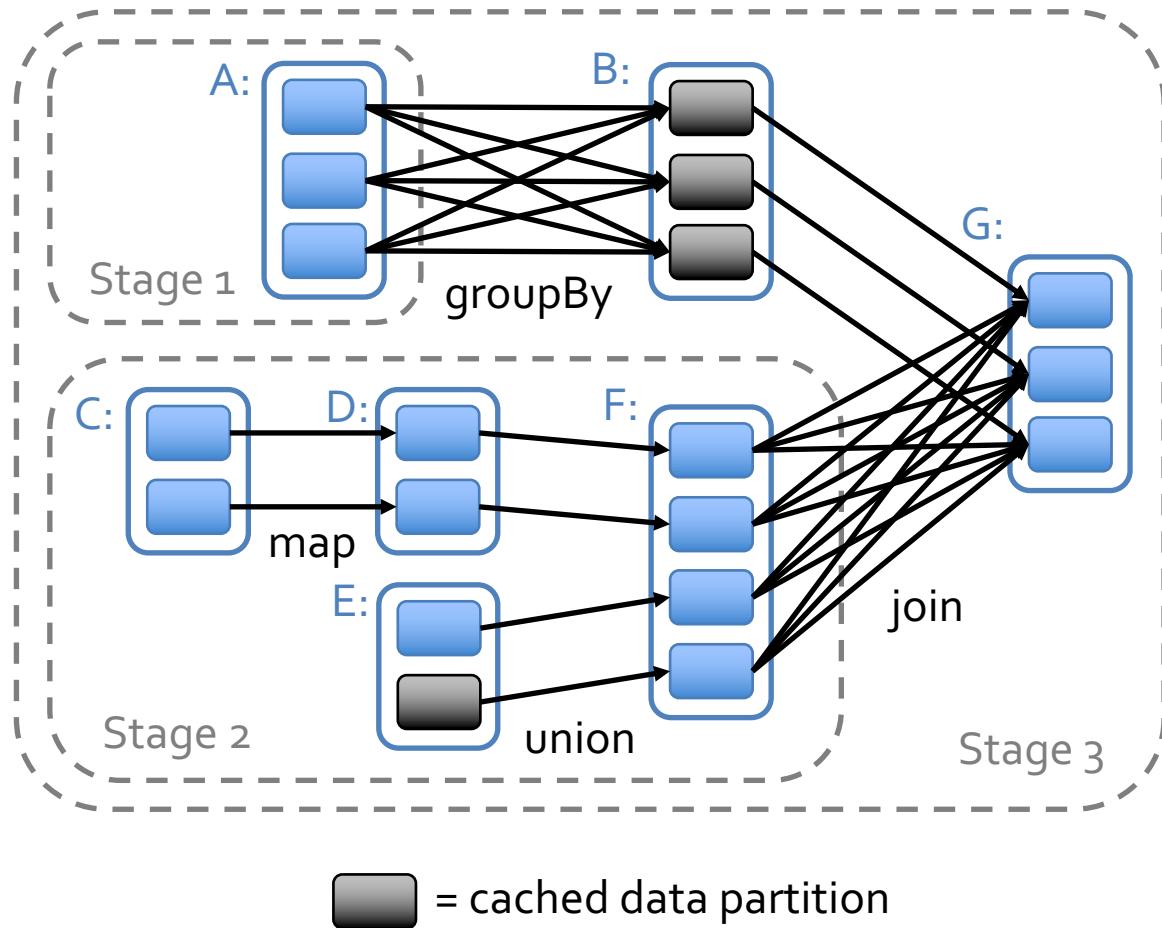
Task scheduler

Dryad-like DAGs

Pipelines functions
within a stage

Locality & data
reuse aware

Partitioning-aware
to avoid shuffles



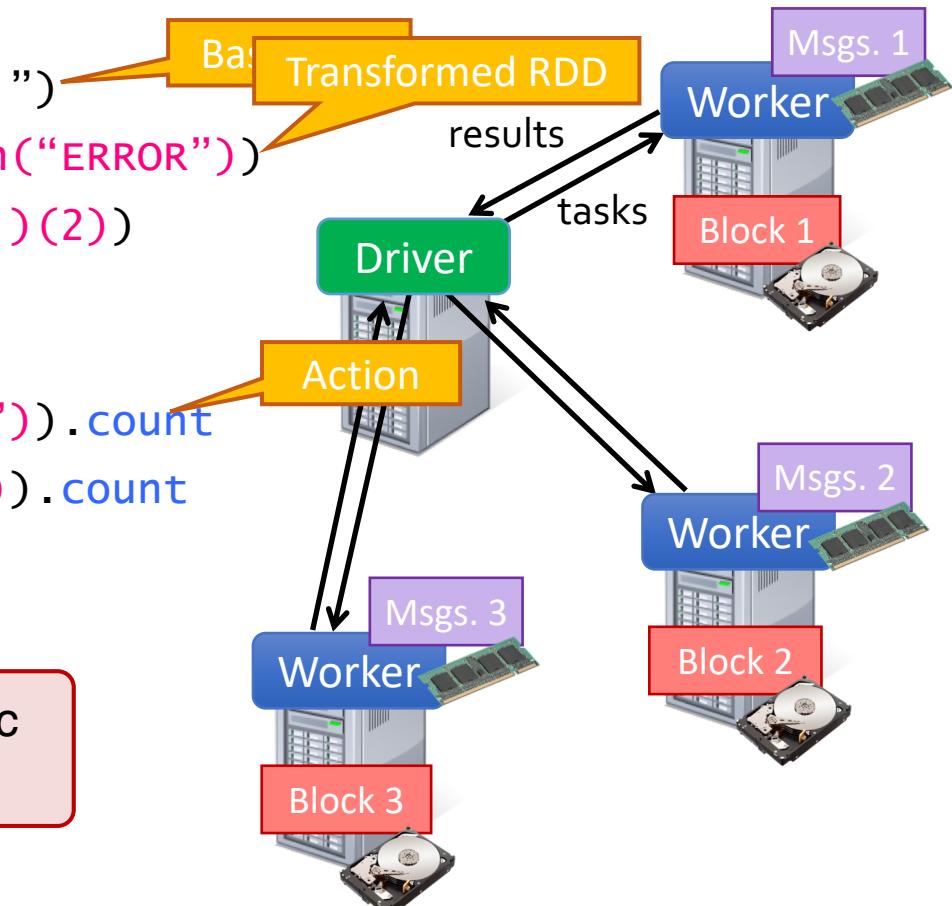
Interactive debugging (control and data flow)

Load error messages from a log into memory, then interactively search for various patterns

```
lines = spark.textFile("hdfs://...")  
errors = lines.filter(_.startsWith("ERROR"))  
messages = errors.map(_.split('\t')(2))  
messages.persist()
```

```
messages.filter(_.contains("MySQL")).count  
messages.filter(_.contains("HDFS")).count
```

Result: scaled to 1 TB data in 5-7 sec
(vs 170 sec for on-disk data)

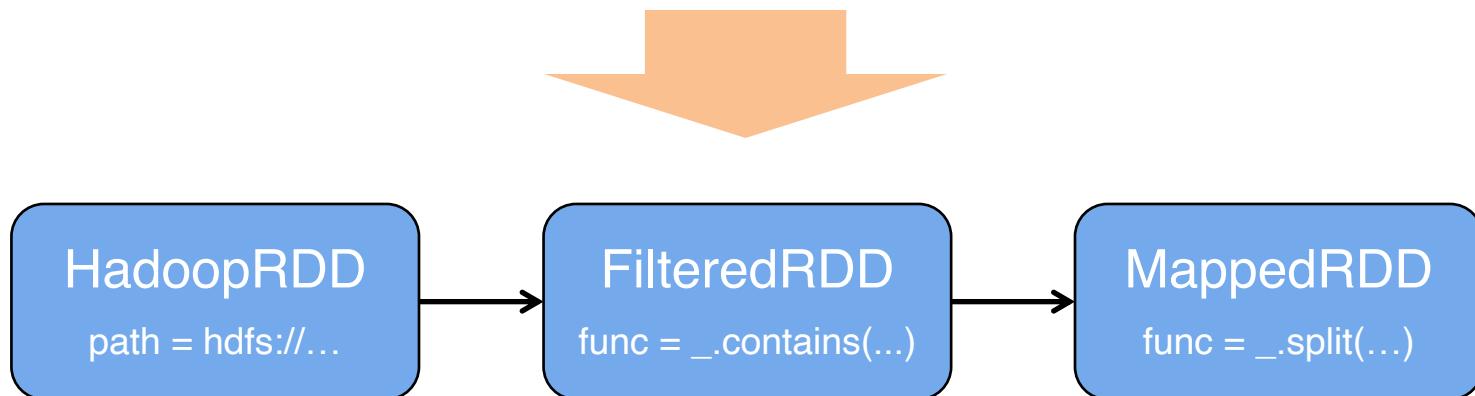


Fault recovery

- RDDs track the graph of transformations that built them (their *lineage*) to rebuild lost data

E.g.:

```
messages = textFile(...).filter(_.contains("error"))
               .map(_.split('\t'))(2)
```

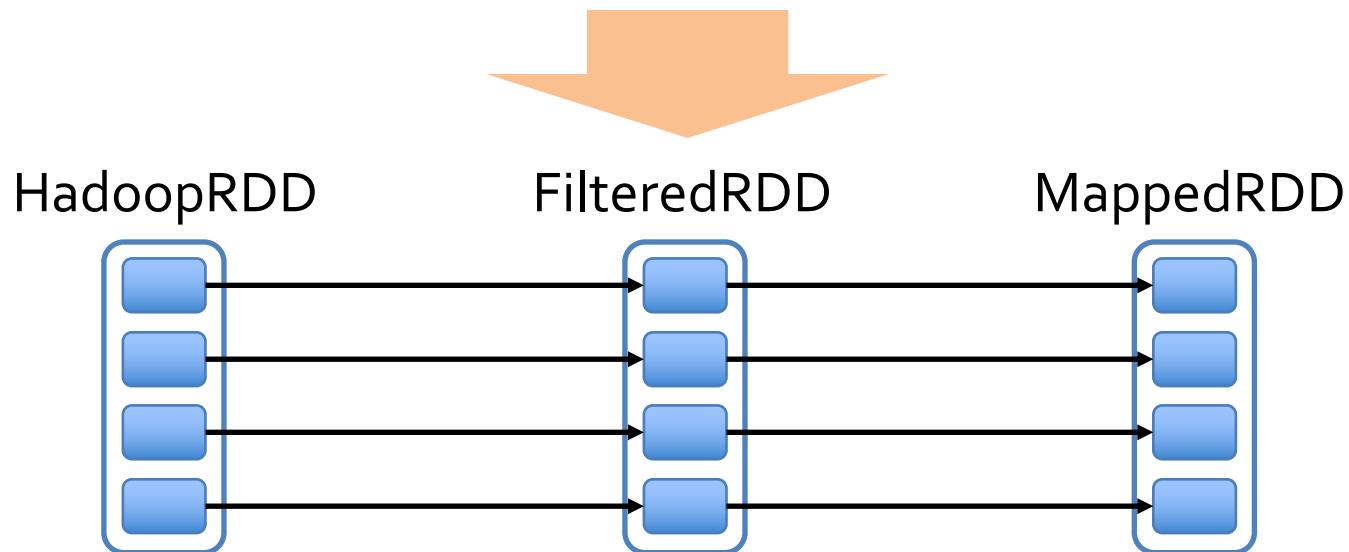


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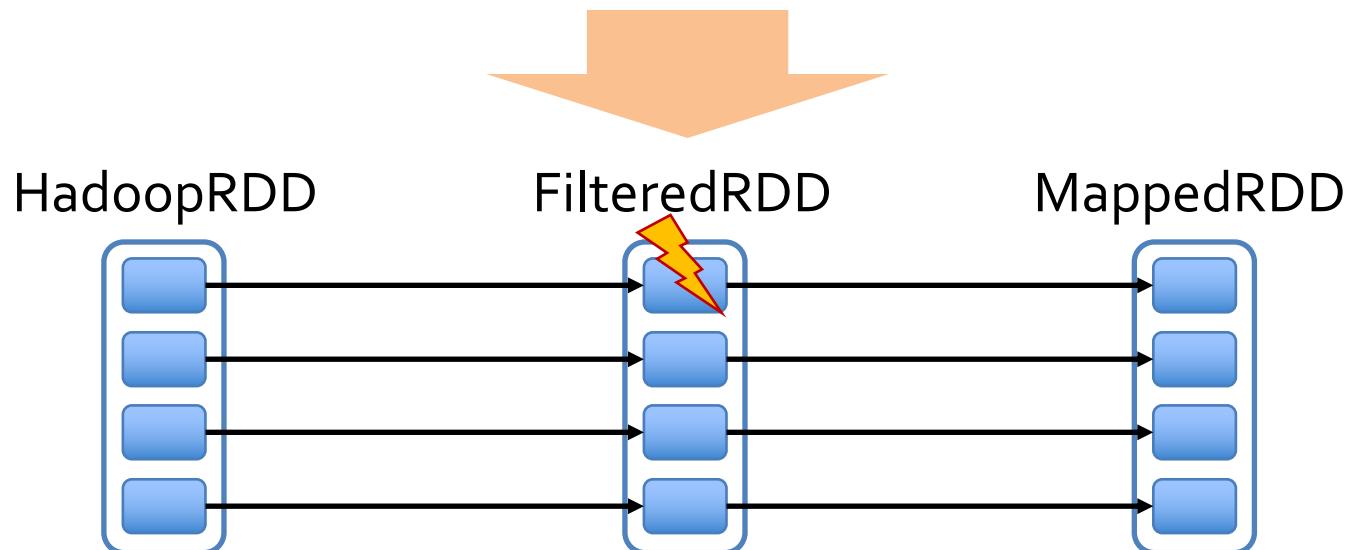


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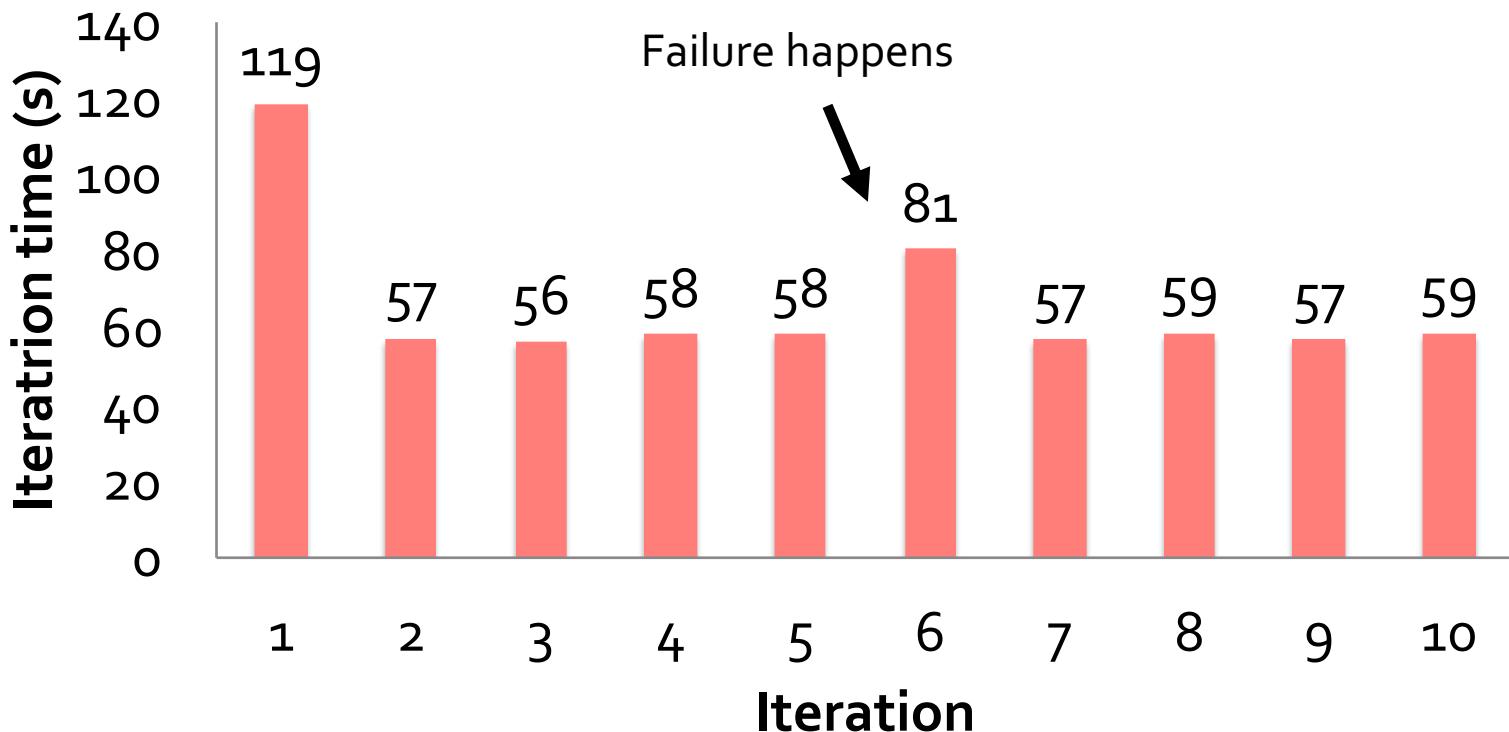
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E.g.:

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Fault recovery results



Example: PageRank

1. Start each page with a rank of 1
2. On each iteration, update each page's rank to
$$\sum_{i \in \text{neighbors}} \text{rank}_i / |\text{neighbors}_i|$$

```
links = // RDD of (url, neighbors) pairs
ranks = // RDD of (url, rank) pairs

for (i <- 1 to ITERATIONS) {
    ranks = links.join(ranks).flatMap {
        (url, (links, rank)) =>
        links.map(dest => (dest, rank/links.size))
    }.reduceByKey(_ + _)
}
```

Example: PageRank

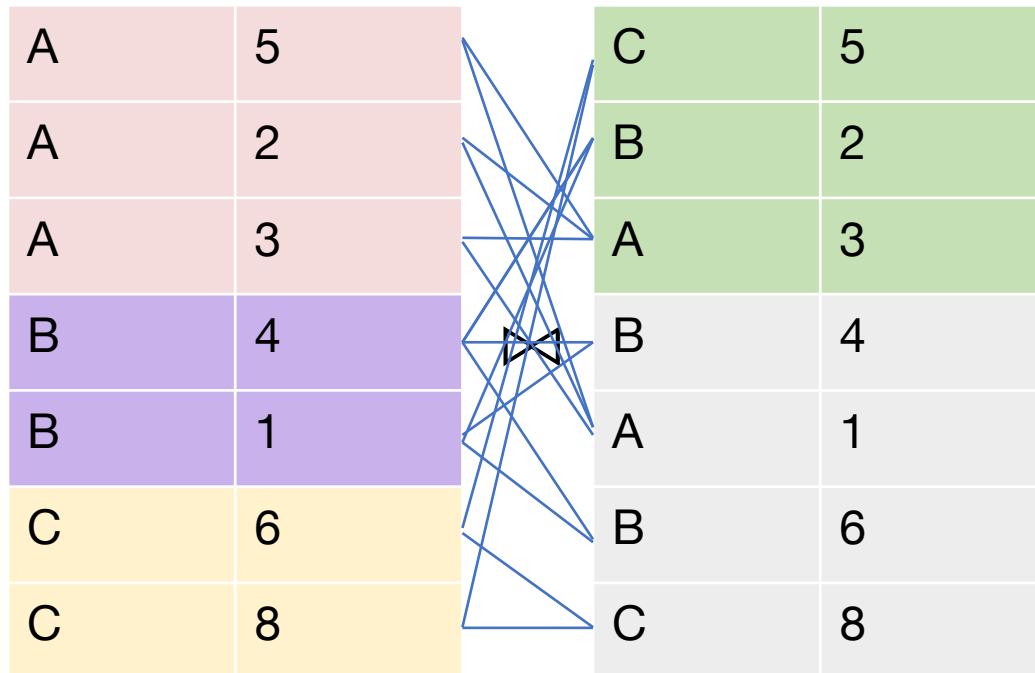
1. Start each page with a rank of 1
2. On each iteration, update each page's rank to

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```
RDD[(URL, Seq[URL])]  
links = // RDD of (url, neighbors) pairs  
ranks = // RDD of (url, rank) pairs ← RDD[(URL, Rank)]  
  
for (i <- 1 to ITERATIONS) {  
    ranks = links.join(ranks).flatMap {  
        (url, (links, rank)) =>  
            links.map(dest => (dest, rank/links.size))  
    }.reduceByKey(_ + _)  
}  
Reduce to RDD[(URL, Rank)]  
For each neighbor in links emits (URL, RankContrib)
```

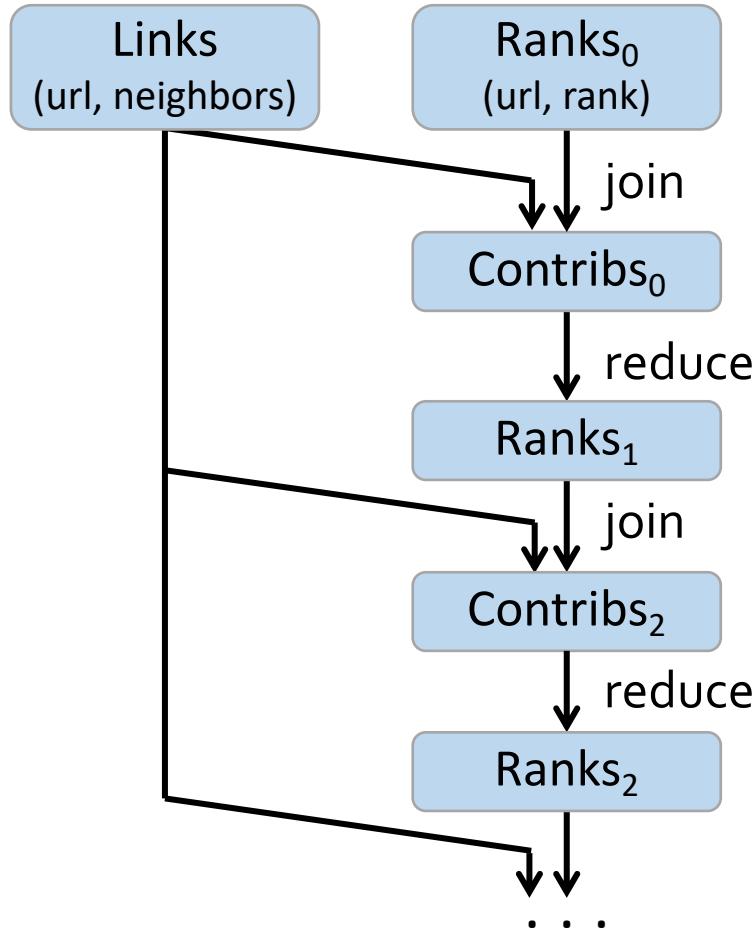
Join (\bowtie)

Alice	5	\bowtie	Alice	F	=	Alice	5	F
Bob	6		Bob	M		Bob	6	M
Claire	4		Claire	F		Claire	4	F



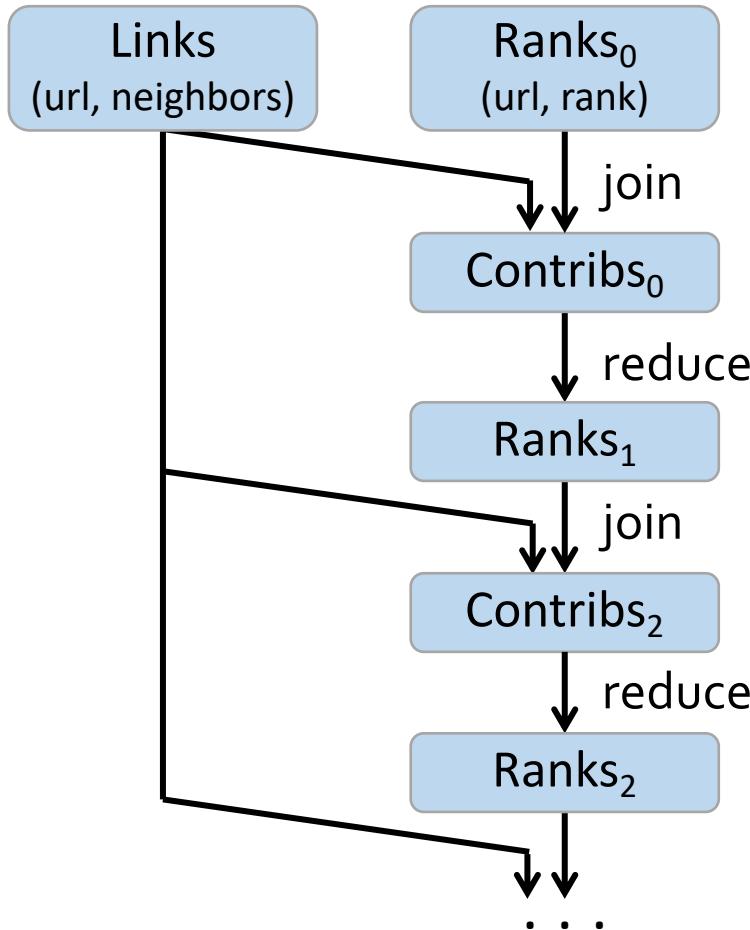
If partitioning doesn't match, then need to reshuffle to match pairs. Same problem in `reduce()` for MapReduce.

Optimizing placement



- `Links` & `ranks` repeatedly joined
- Can co-partition them (e.g. hash both on URL) to avoid shuffles
- Can also use app knowledge, e.g., hash on DNS name
- `Links = links.partitionBy(new URLPartitioner())`

Optimizing placement



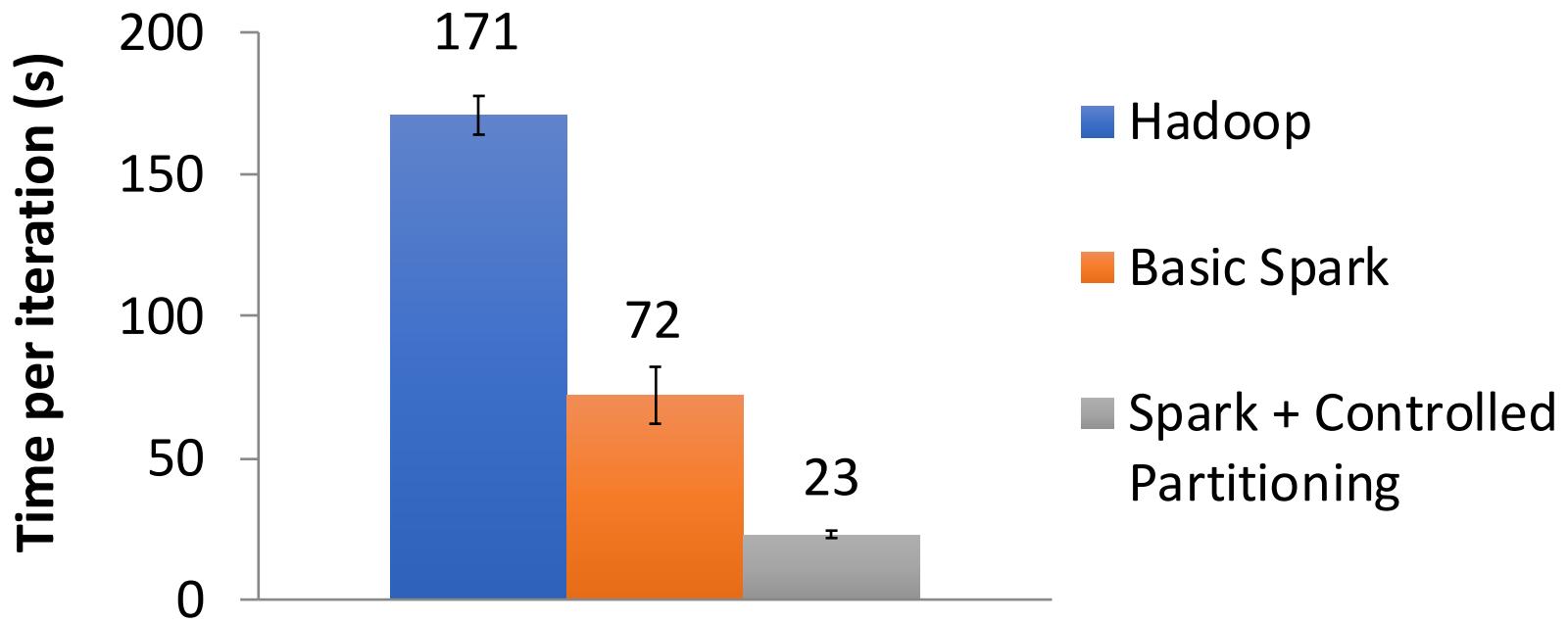
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Q: Where might we have placed `persist()`?

Co-partitioning example

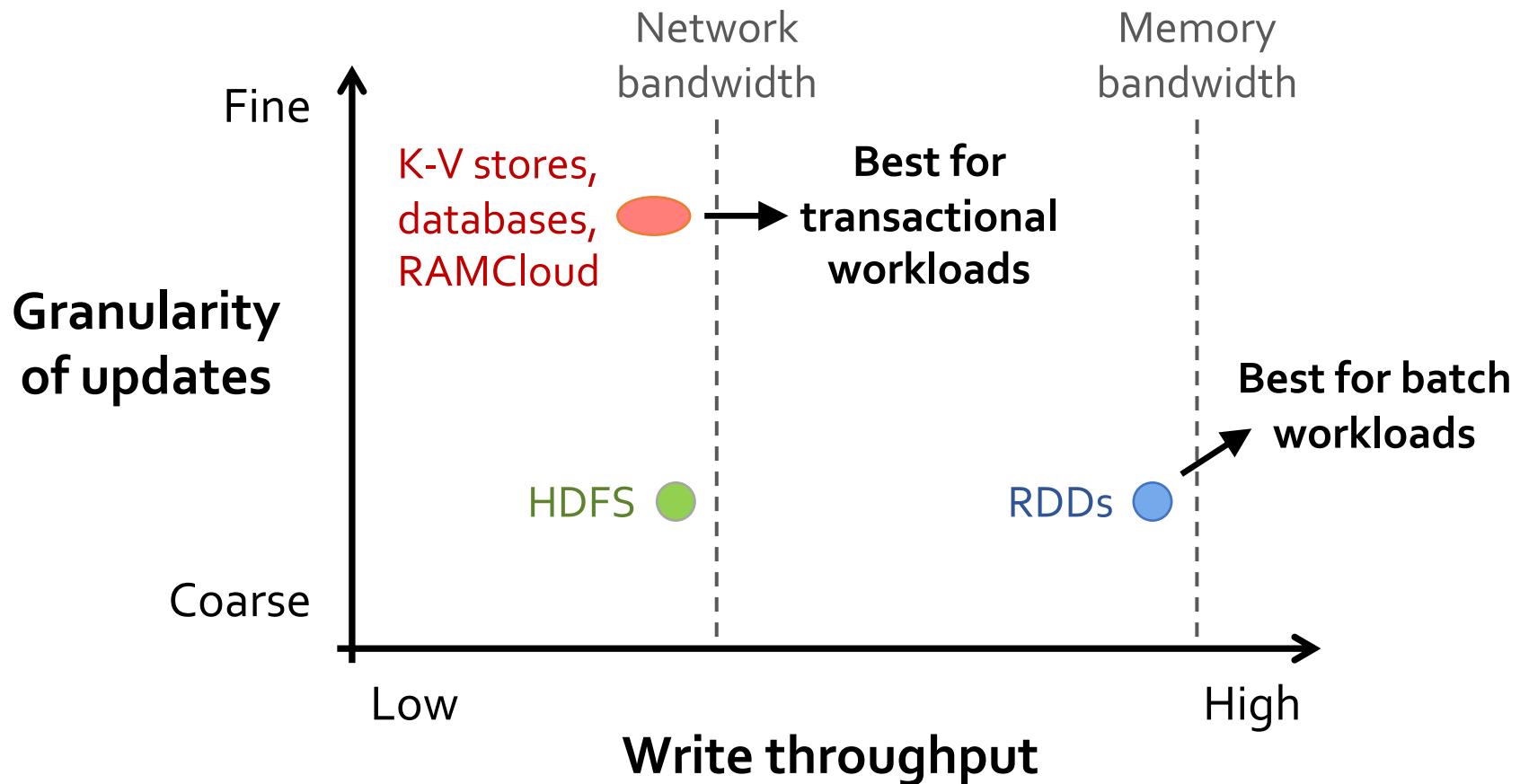
Co-partitioning can avoid shuffle on join
But, fundamentally a shuffle on `reduceByKey`
Optimization: custom partitioner on domain

PageRank performance



* Figure 10a: 30 machines on 54 GB of Wikipedia data computing PageRank

Tradeoff space



Discussion & wrap-up