

A wide-angle photograph of a mountain range at dusk or dawn. The peaks are covered in snow, and the sky is filled with soft, pinkish-orange clouds. A full moon is visible in the upper right corner.

# The Case for Learned Index Structures

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# Outline

- Background on Traditional Index Structures in DBMS
  - Why do we even need index structures at all?
  - B-Tree & CDF Model
- Learned Index Structures (LIS)
  - Naïve Approach (A single NNR)
  - Reclusive Model Index (RMI)
- Conclusion

# Fundamental Building Blocks of Database Systems

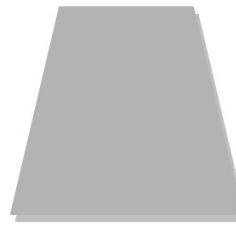


# Index Structures

B-Tree



HashMap



Blooming Filter



- Purpose of having these structures in DBMS?
- Hit: Tradeoff between speed and storage
- Scarifies storage reduce the # of blocks to read

# Accessing Data Without Index Structures

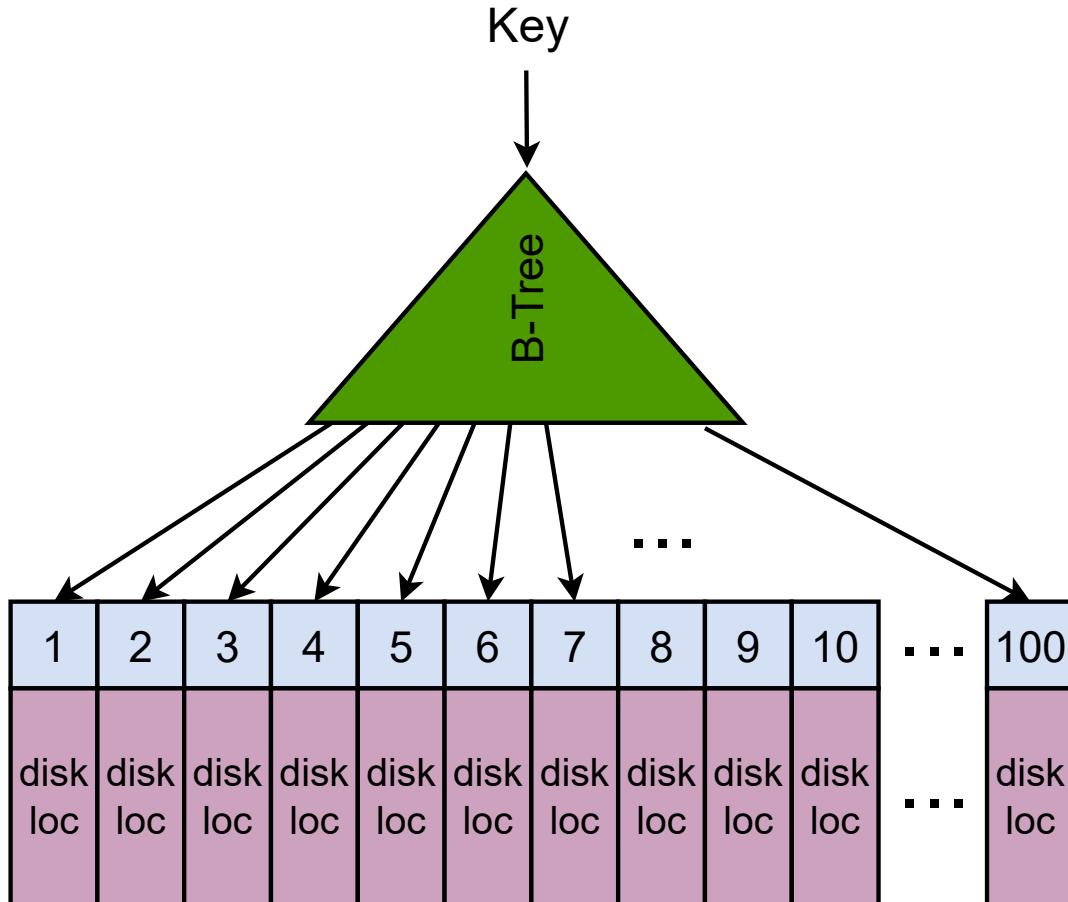
UVA DS School Info Table

SID	s_name	s_major	...	...	GPA
1	student_1	CS	...	...	4.0
2	student_2	DS	...	...	4.0

:

SID	s_name	s_major	...	...	GPA
98	student_98	DA	...	...	4.0
99	student_99	CS	...	...	4.0

# Indexing SID using B-Tree



Prepare a B-Tree index structure for SID from 1 to 100

1	2	3	4	5	6	7	8	9	10	...	100
---	---	---	---	---	---	---	---	---	----	-----	-----

Time:  $O(\log n)$   
Space:  $O(n)$

# B-Tree Operations

- Operations: *INSERT()*, *LOOKUP()*, *DELETE()*, *UPDATE()*
- *LOOKUP()* walk through example
  - Visualization : <https://people.ksp.sk/~kuko/gnarley-trees/Btree.html#>

- No assumption about data distribution
- Knowing data distribution may increase performance significantly from both speed and storage.

Indexing all integers from 1 to 100

1	2	3	4	5	6	7	8	9	10	...	100
---	---	---	---	---	---	---	---	---	----	-----	-----

array[lookup - 1]

Indexing all even integers from 2 to 100

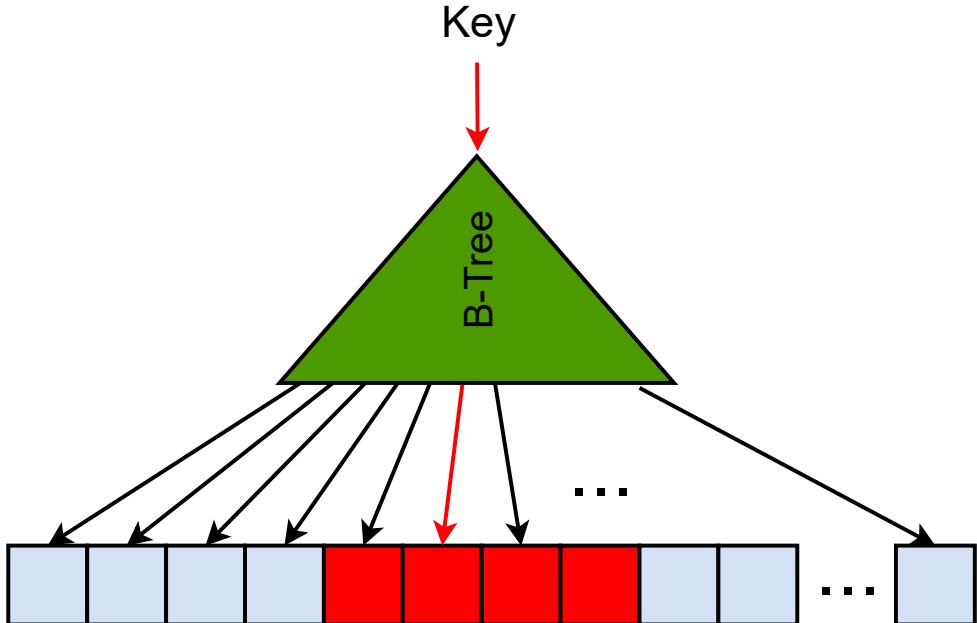
2	4	6	8	10	12	14	16	18	20	...	100
---	---	---	---	----	----	----	----	----	----	-----	-----

array[ (lookup - 2) / 2 ]

# A B-Tree is A Model

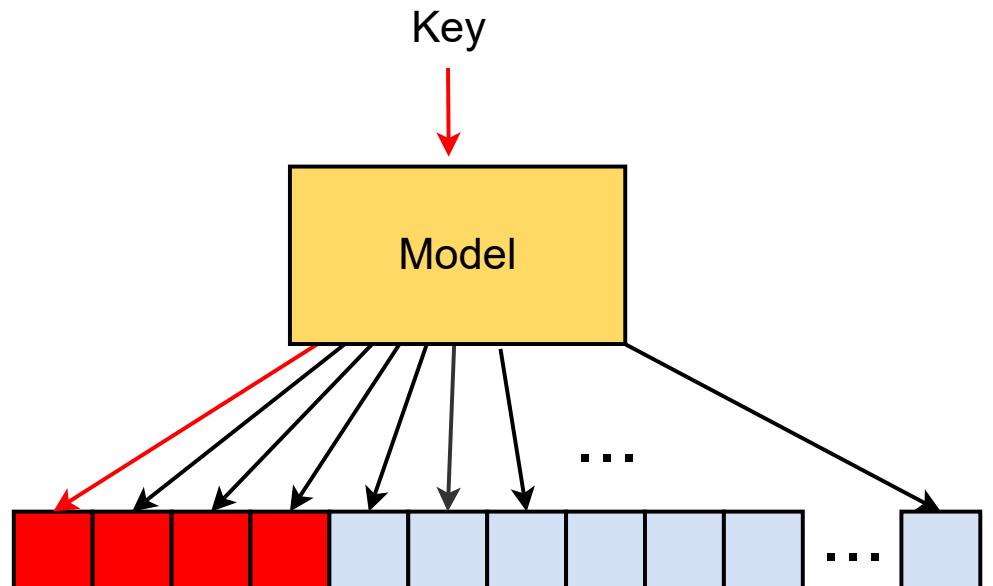
B-Tree

1. Locate the pos of input key
2. Binary search within a page size



Alternative view

1. Locate the pos of input key
2. Binary search within error boundaries



Assuming data are stored in dense array in sorted order

# Modeling B-Tree Functionality using CDF

- B-Tree indexes the data in a sorted order
  - $Pos = B\text{-Tree}(key)$
- CDF gives the probability of  $X$  that will have a value less than or equal to  $x$ 
  - $Pos = CDF(Key) * N$
- Data distribution visualization: <https://statdist.com/>

If we can learn the CDF model of a given dataset, we can replace the B-Tree index structure with learned model

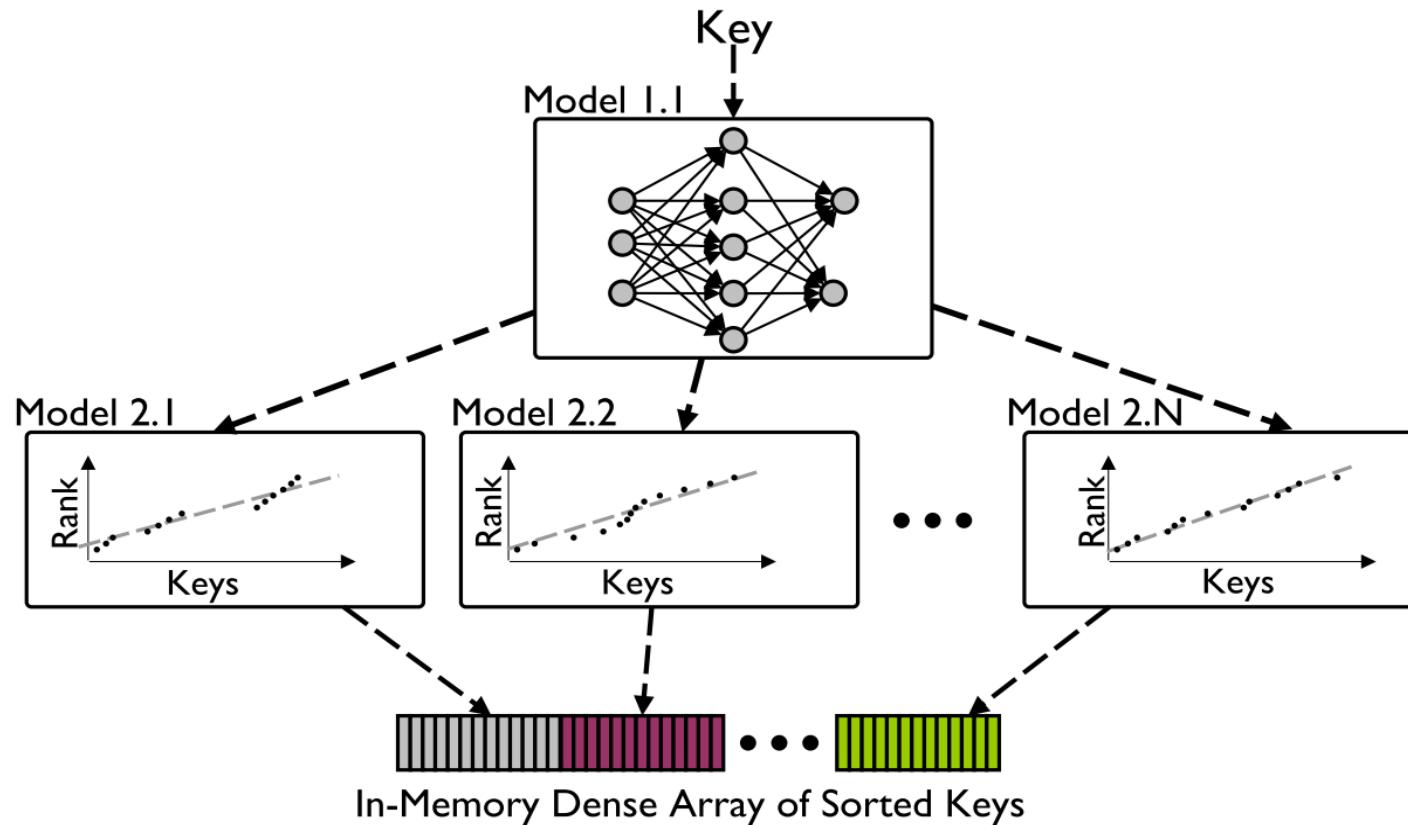
# Which ML Model?

- $Pos = CDF(Key) * N$ 
  - Approximate the position given a key inside a sorted array
  - Learn the relationship between *Key* and *Pos* -> *Regression*
  - Position range: 0 to  $N-1$
  - Key range: smallest to largest value
- Which regression model should we pick and why?
  - Linear Regression, Neural Network Regression, and etc.
- Discussion (2 mins)

# Naïve LIS

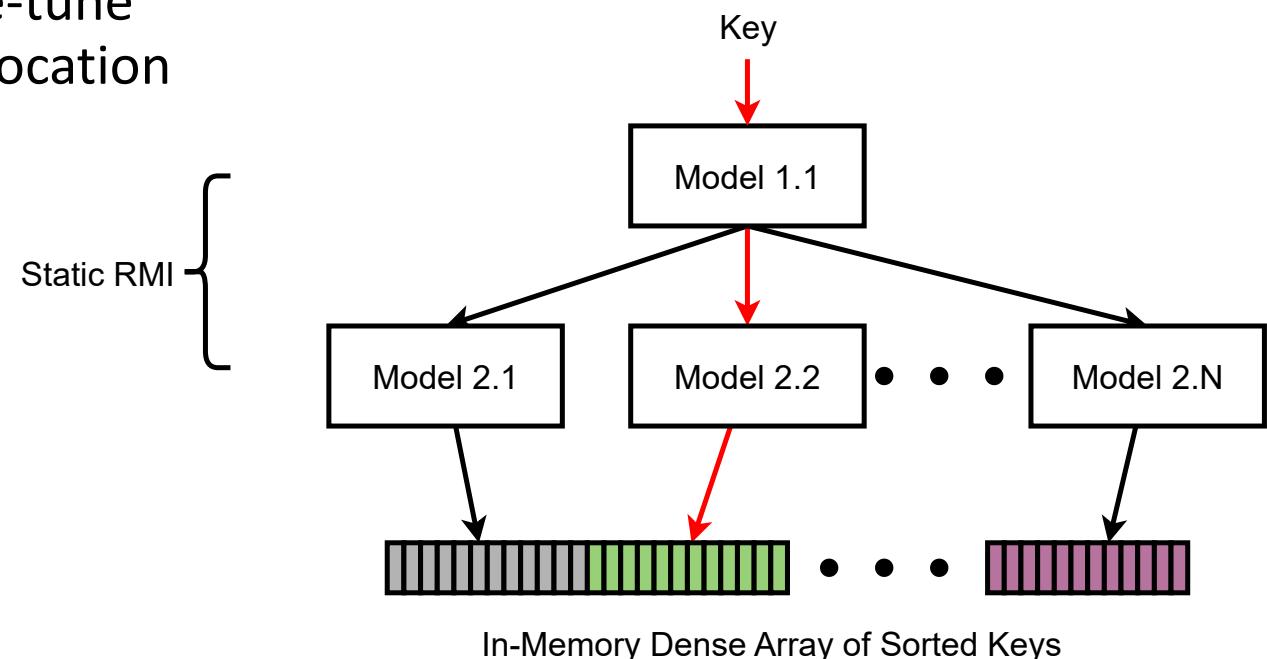
- Use a single neural network regression model
- Good at approximate the general shape of a CDF
- Large error at the last mile of predicting the actual position
- Solution: RMI

# Two-stage-RMI

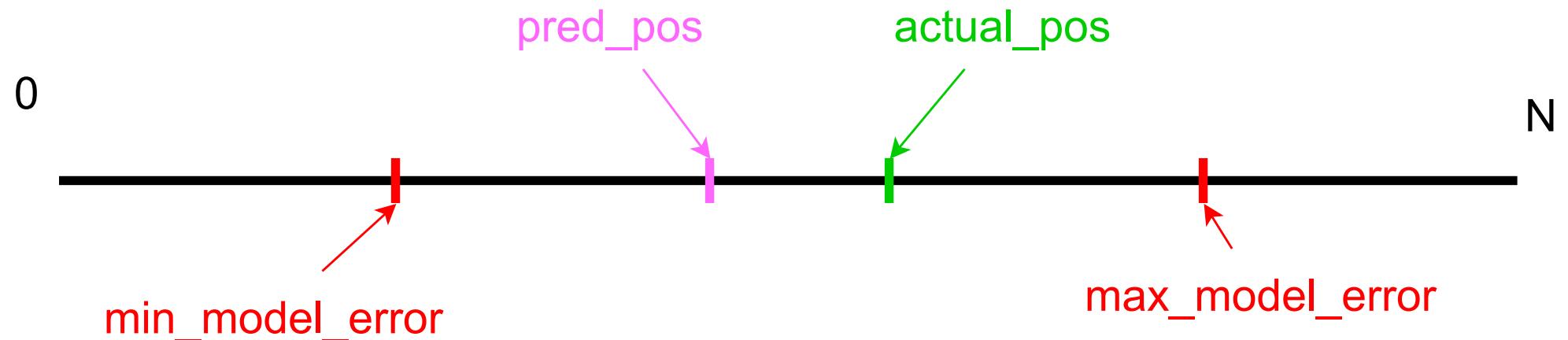


# Lookup() on A Two-stage-RMI

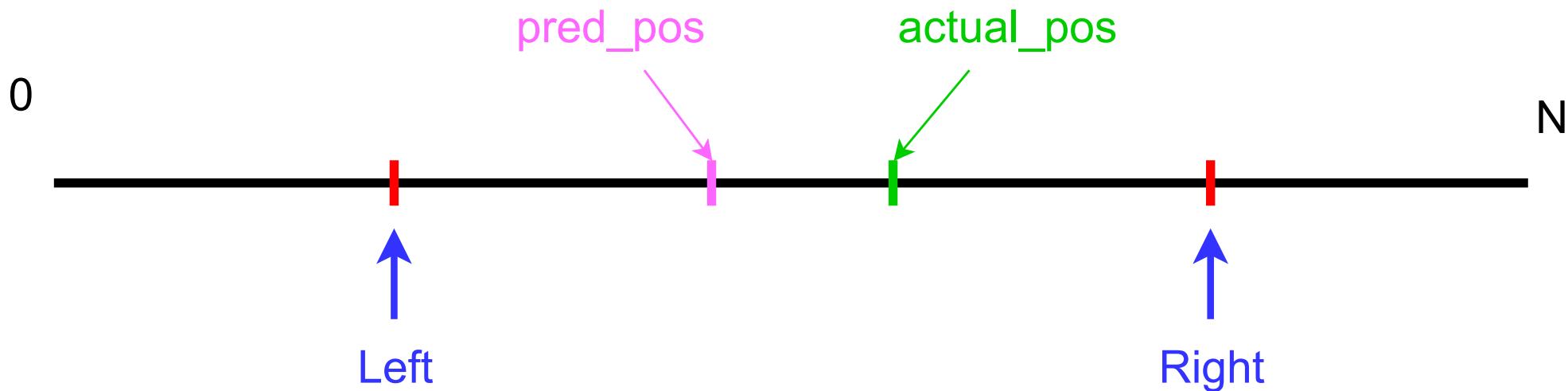
- Two stage RMI
- A higher stage model directs a lookup operation to a lower stage model to fine-tune the precision of the predicted memory location
- `learned_index_lookup(key)`
  - `ret_1 = first_stage_lookup(key)`
  - `ret_2 = second_stage_lookup(ret_1)`
  - `predicted_pos = array[ret_2]`



# Local Search



# Binary Search



# Conclusion

- DBMS index structures
  - B-Tree & CDF model
- LIS
  - Abstracting functionality of B-Tree using regression models
  - Naïve approach: using single regression model
  - RMI: a hierarchical architecture
  - Performing a local search if the predicted pos is off actual pos
  - Using binary search or other search algorithms to find the actual key