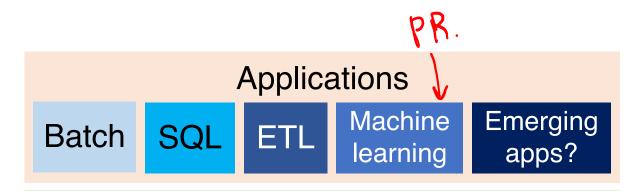
Resilient Distributed Datasets, Spark

DS 5110: Big Data Systems (Spring 2023) Lecture 4

Yue Cheng

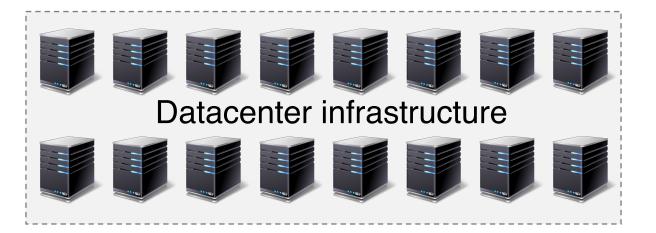




Scalable computing engines



Scalable storage systems



What's good with MapReduce

- Scaled analytics to thousands of machines
- Eliminated fault tolerance as a concern

Problems with MapReduce

- Scaled analytics to thousands of machines
- Eliminated fault tolerance as a concern
- Not very expressive
 - Iterative algorithms (PageRank, Logistic Regression, ...)
- Interactive and ad-hoc queries (Interactive Log Debugging)
- Lots of specialized frameworks
 - Pregel, GraphLab, PowerGraph, DryadLINQ, HaLoop...

Sharing data between stages/iterations

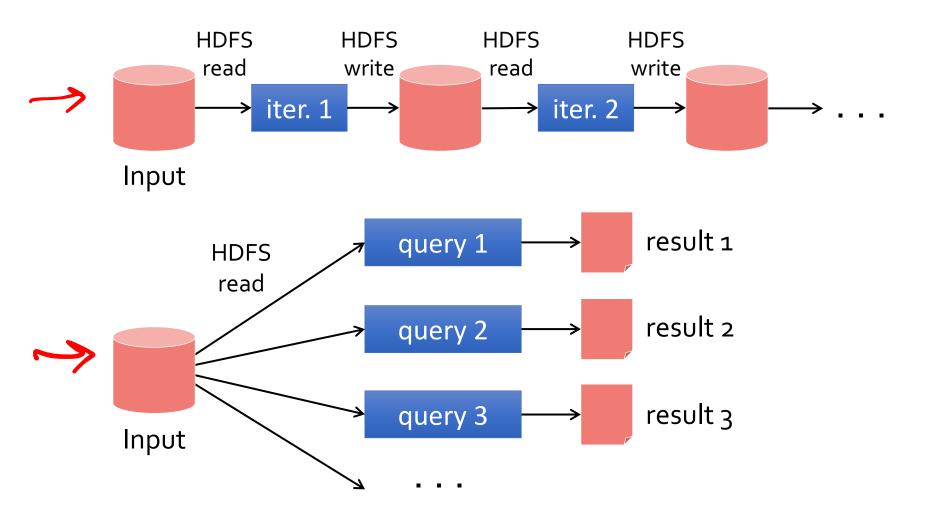
- Only way to share data between iterations / phases is through shared storage

 M R
 - Slow!
- Allow operations to feed data to one another
 - Ideally, through memory instead of disk-based storage

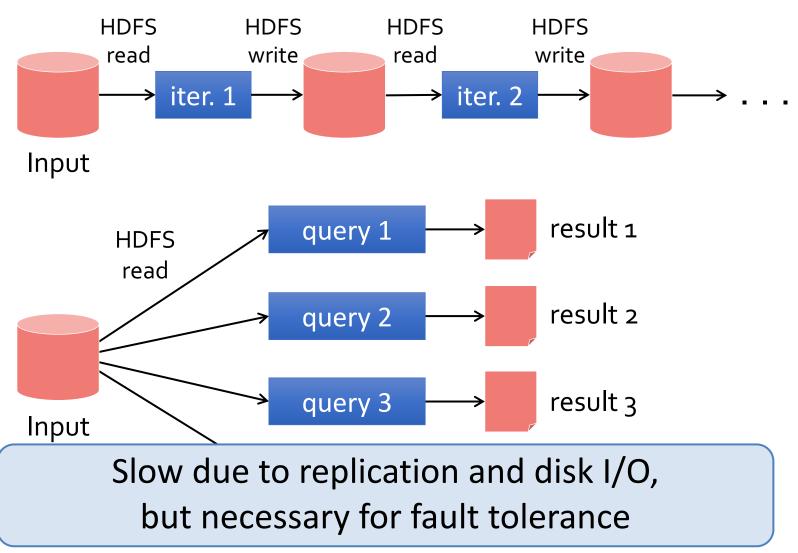
Sharing data between stages/iterations

- Only way to share data between iterations / phases is through shared storage
 - Slow!
- Allow operations to feed data to one another
 - Ideally, through memory instead of disk-based storage
- Need the "chain" of operations to be exposed to make this work
- Problem to solve: Would this break the MR faulttolerance scheme?
 - Retry and Map or Reduce task since idempotent

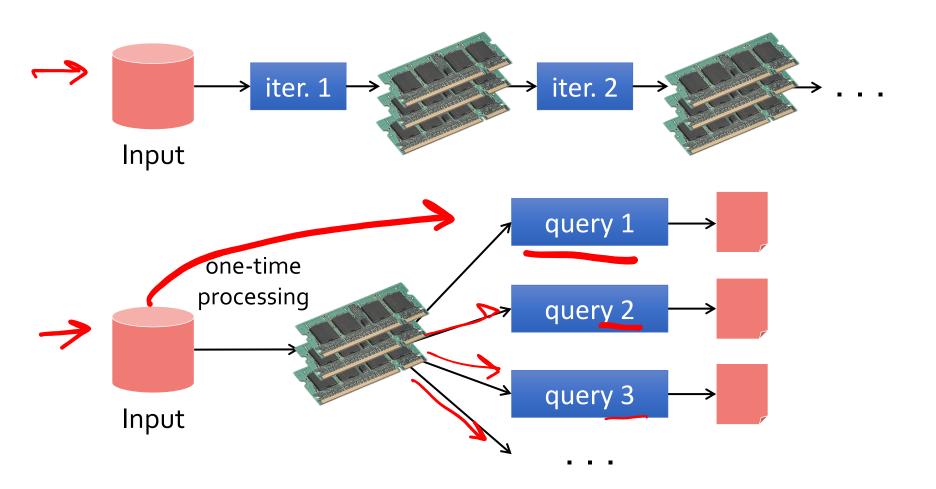
Examples



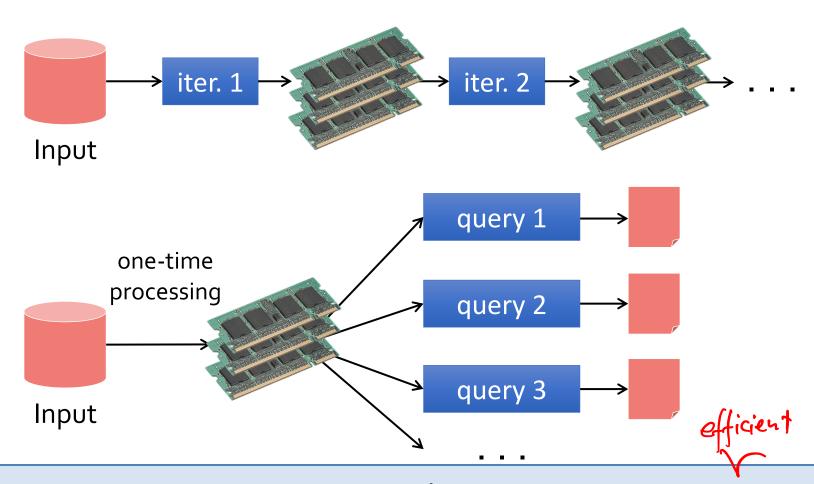
Examples



Goal: In-memory data sharing



Goal: In-memory data sharing



10-100× faster than network/disk, but how to get FT?

Challenges

 How to design a distributed memory abstraction that is both fault-tolerant and efficient?

Challenges

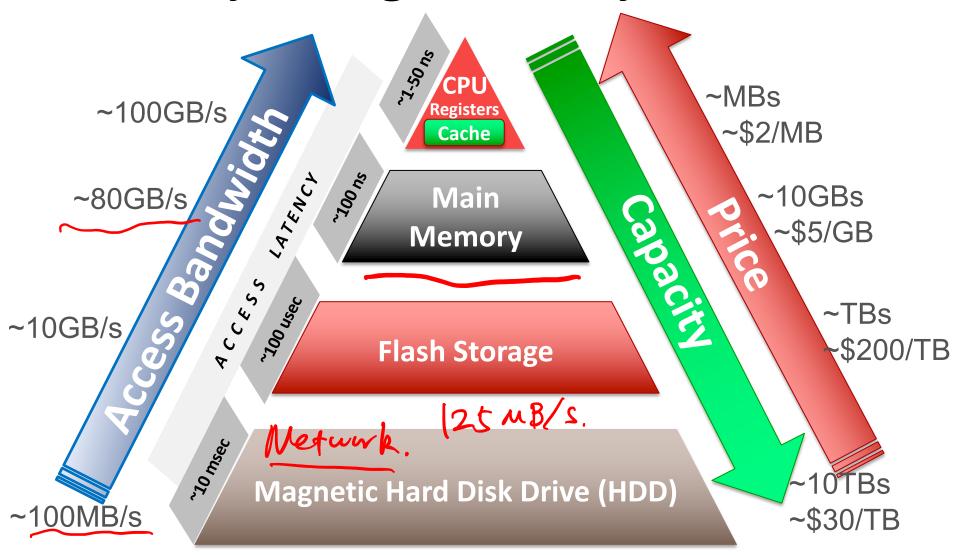
 How to design a distributed memory abstraction that is both fault-tolerant and efficient?

- Existing storage systems allow fine-grained mutation to state
 - nutation to state

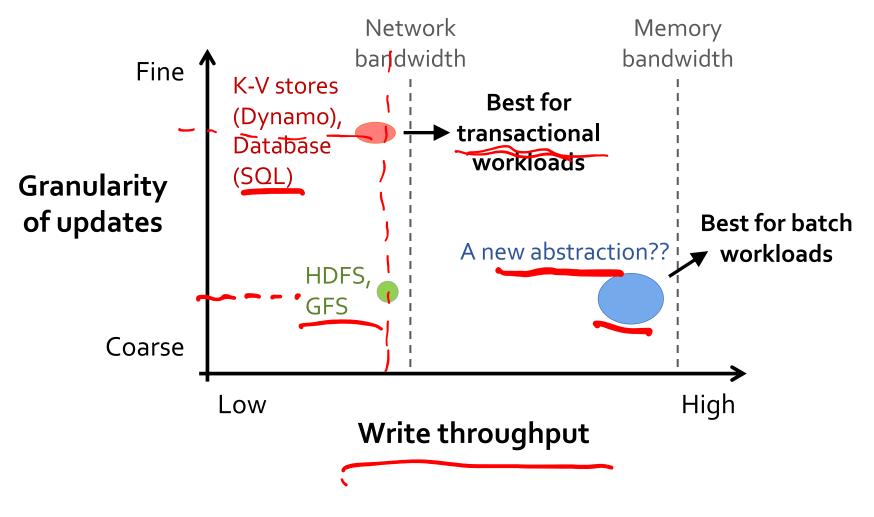
 In-memory key-value stores

 range query
 - Requires replicating data or logs across nodes for fault tolerance
 - Costly for data-intensive apps
 - 10-100x slower than memory write
 - They also require costly on-the-fly replication for mutations

Memory-storage hierarchy



Tradeoff space



Challenges

 How to design a distributed memory abstraction that is both fault-tolerant and efficient?

 Existing storage systems allow fine-grained mutation to state

Insight: leverage similar coarse-grained approach that transforms whole dataset per operation, like MapReduce (batch processing)

- 10-100x slower than memory write
- They also require costly on-the-fly replication for mutations

Solution: Resilient Distributed Datasets (RDDs)

- Restricted form of distributed shared memory
 - Immutable, partitioned collections of records
 - Can only be built through coarse-grained, deterministic transformations (map, filter, join, ...)

- Efficient fault recovery using lineage
 - Log one operation to apply to many elements
 - Recompute lost partitions on failure
 - No cost if nothing fails

Spark programming interface

Scala API, now have multi-language bindings such as Python, Java, etc.

Managing RDDs

- Transformations on RDDs (RDD₁ → RDD₂)
- Actions on RDDs (RDD → output)
- Control over RDD partitioning (how items are split over nodes)
- Control over RDD persistence (in memory, on disk, or recompute on loss)

Transformations

Transformations (define a new RDD)

map filter sample groupByKey reduceByKey sortByKey flatMap union join cogroup cross mapValues

RDDs in terms of Scala types -> Scala semantics at workers

Transformations are lazy operations; cause no cluster action

Actions

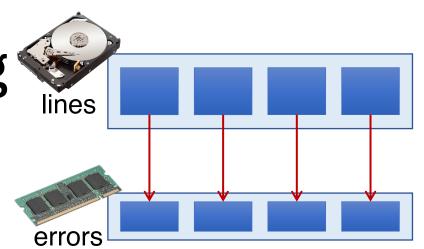
Actions (return a result to driver program)

collect reduce count save lookupKey

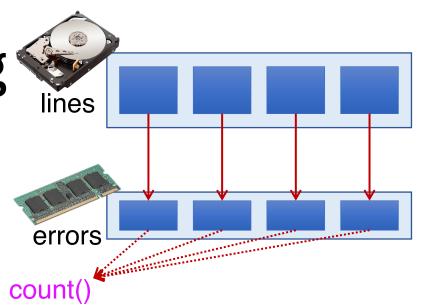
Consumes an RDD to **produce** output either to storage (save), or to interpreter/Scala (count, collect, reduce)

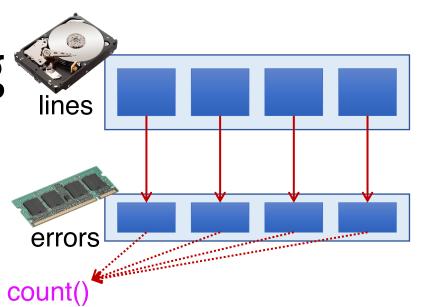
Causes RDD lineage chain to **get executed** on the cluster to produce the output (for any missing pieces of the computation)

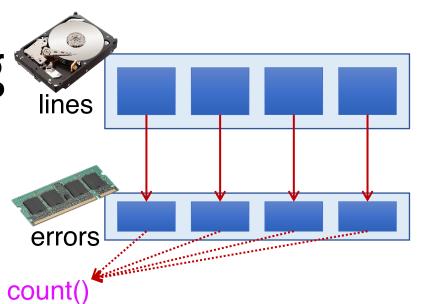
errors.count()

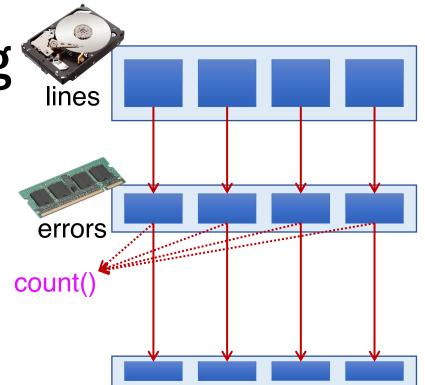


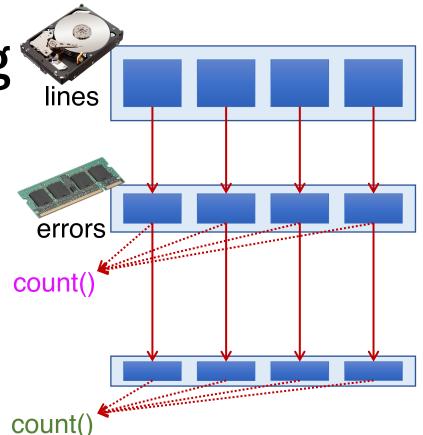
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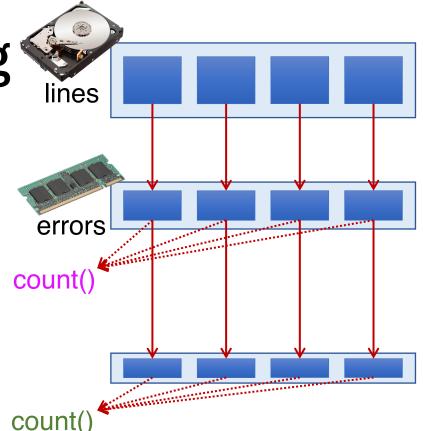


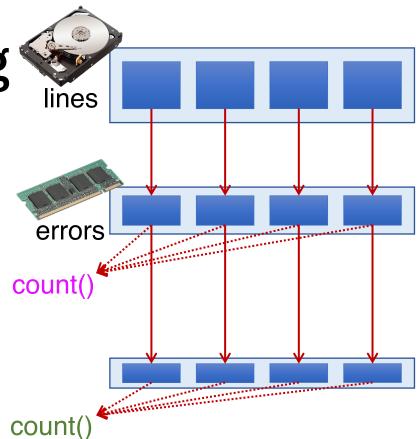




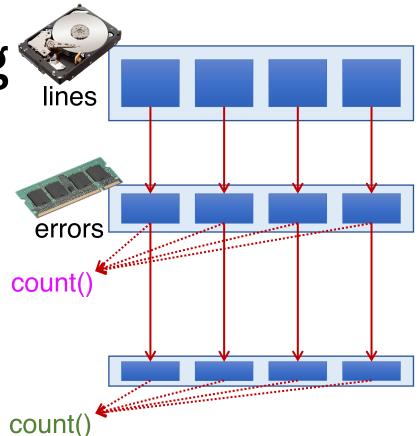




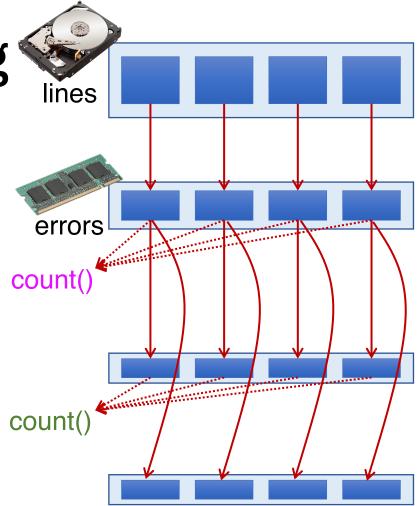




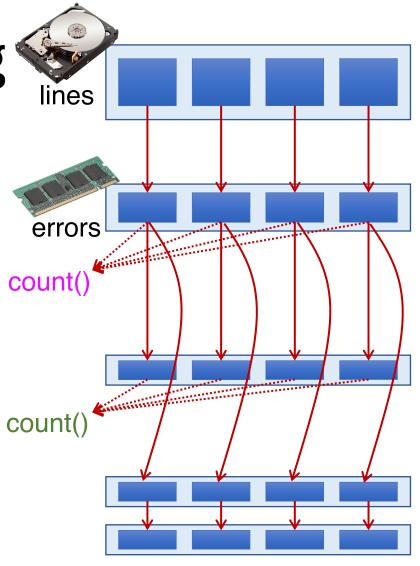
```
lines = textFile("hdfs://foo.log")
errors = lines.filter(
           .startsWith("ERROR")
errors.persist()
errors.count()
errors.filter(
     _.contains("MySQL")).count()
errors.filter(
     _.contains("HDFS"))
     _.map(_.split("\t")(3))
     .collect()
```



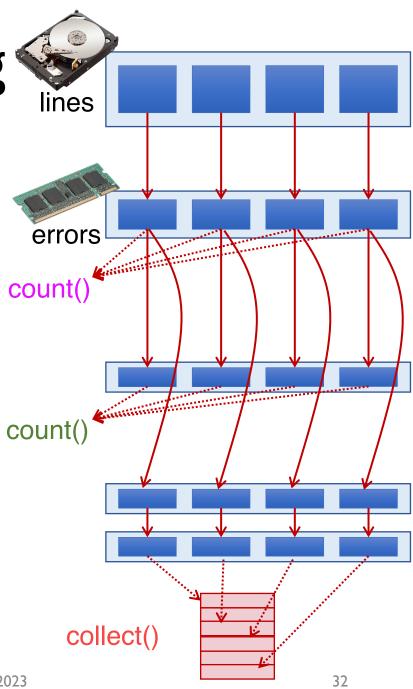
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```



persist()

- Not an action nor a transformation
- A scheduler hint

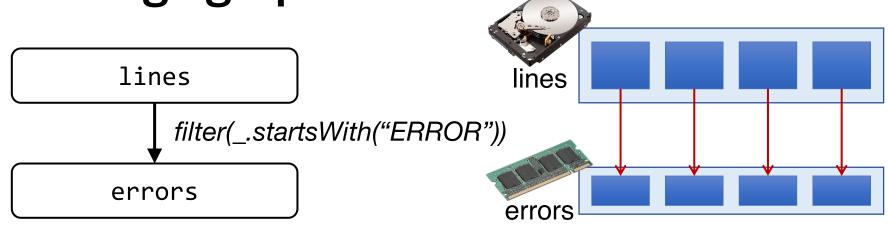
- Tells which RDDs the Spark schedule should materialize and whether in memory or storage
- Gives the user control over reuse/recompute/recovery tradeoffs

Lineage graph of RDDs

lines



Lineage graph of RDDs

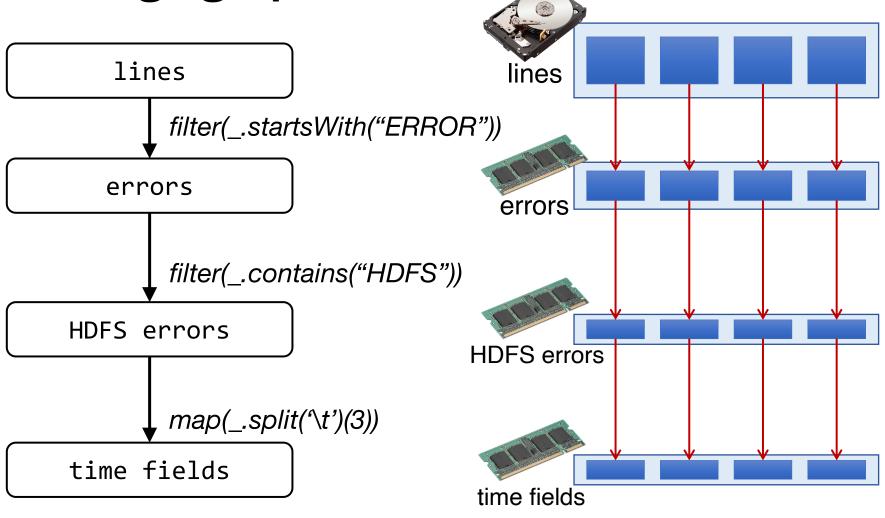


Lineage graph of RDDs lines lines filter(_.startsWith("ERROR")) errors errors filter(_.contains("HDFS"))

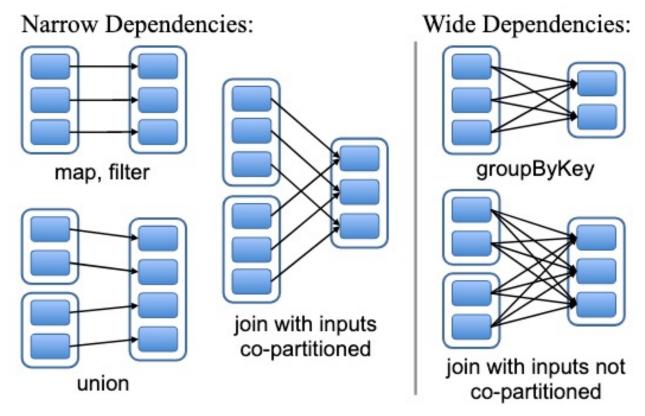
HDFS errors

HDFS errors

Lineage graph of RDDs



Narrow & wide dependencies



Narrow: each parent partition used by at most one child partition (can partition on one machine)

Wide: multiple child partitions depend on one parent partition

Must stall for all parent data, loss of child requires whole parent RDD (not just a small # of partitions)

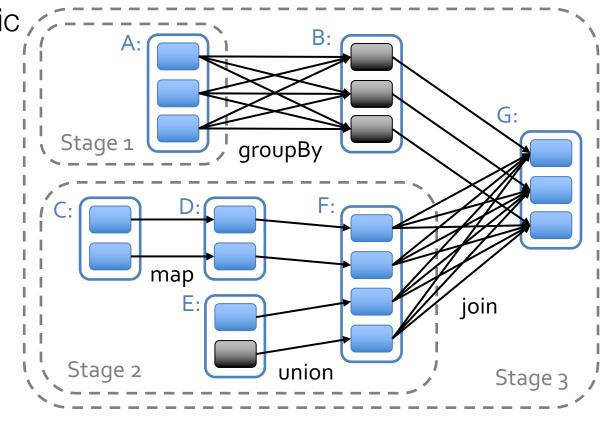
Task scheduler

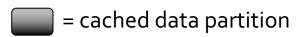
DAGs (directed acyclic graphs)

Pipelines functions within a stage

Locality & data reuse aware

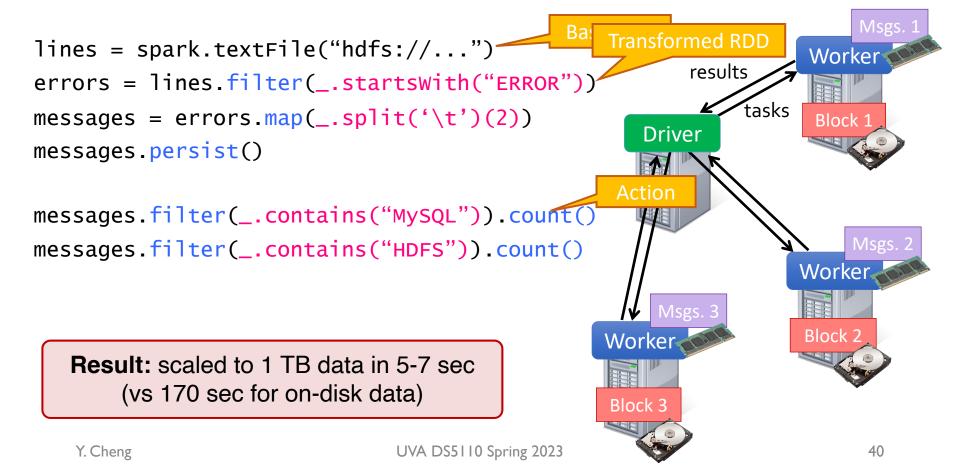
Partitioning-aware to avoid shuffles





Interactive debugging (control and data flow)

Load error messages from a log into memory, then interactively search for various patterns

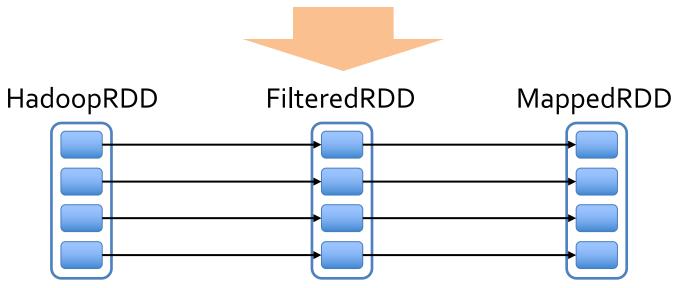


Fault recovery

 RDDs track the graph of transformations that built them (their *lineage*) to rebuild lost data

Fault recovery

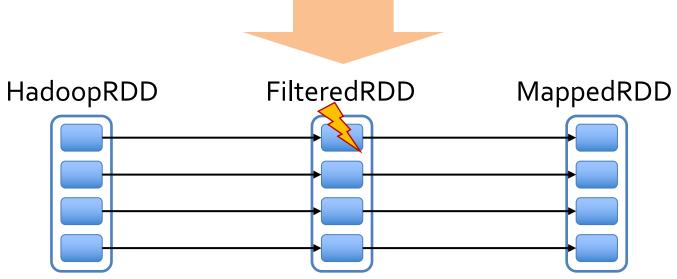
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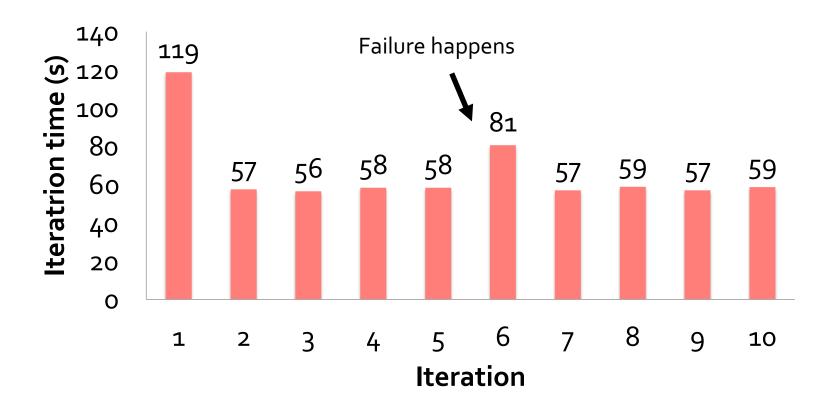
Fault recovery

 RDDs track the graph of transformations that built them (their lineage) to rebuild lost data



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Fault recovery results



Example: PageRank

- 1. Start each page with a rank of 1
- 2. On each iteration, update each page's rank to $\Sigma_{i \in neighbors}$ rank_i / |neighbors_i|

```
links = // RDD of (url, neighbors) pairs
ranks = // RDD of (url, rank) pairs

for (i <- 1 to ITERATIONS) {
   ranks = links.join(ranks).flatMap {
      (url, (links, rank)) =>
         links.map(dest => (dest, rank/links.size))
   }.reduceByKey(_ + _)
}
```

Example: PageRank

- 1. Start each page with a rank of 1
- 2. On each iteration, update each page's rank to

```
\Sigma_{i \in neighbors} \; rank_i \, / \; |neighbors_i|
```

Join (⋈)

Alice	5
Bob	6
Claire	4

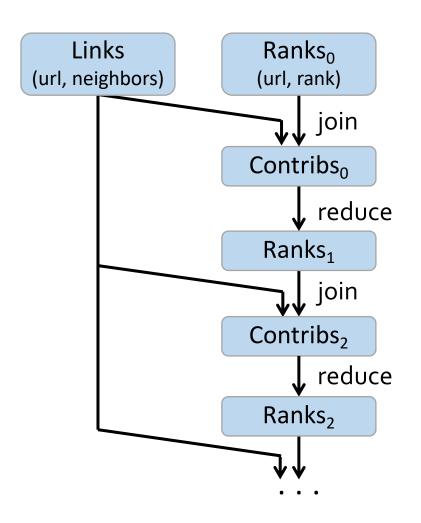
Alice	F
Bob	M
Claire	F

Alice	5	F
Bob	6	М
Claire	4	F

Α	5	С	5
Α	2	В	2
Α	3	Α	3
В	4	В	4
В	1	Α	1
С	6	В	6
С	8	С	8

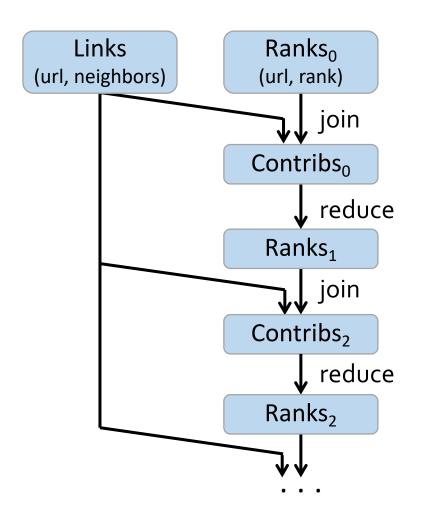
If partitioning doesn't match, then need to reshuffle to match pairs. Same problem in reduce() for MapReduce.

Optimizing placement



- links & ranks repeatedly joined
- Can co-partition them (e.g. hash both on URL) to avoid shuffles
- Can also use app knowledge,
 e.g., hash on DNS name

Optimizing placement



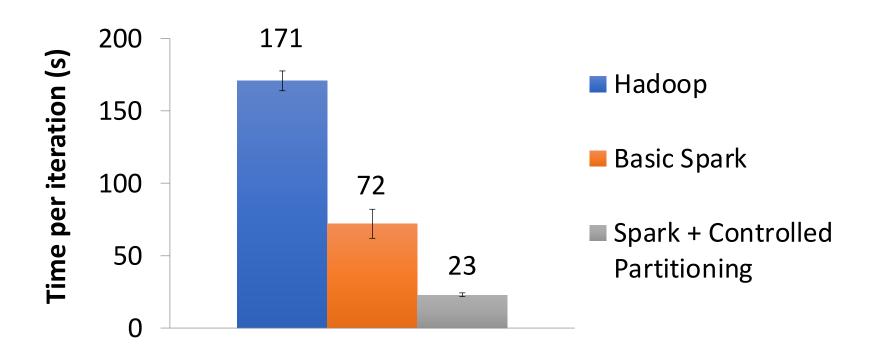
- links & ranks repeatedly joined
- Can co-partition them (e.g. hash both on URL) to avoid shuffles
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Q: Where might we have placed **persist()**?

Co-partitioning example

Co-partitioning can avoid shuffle on join
But, fundamentally a shuffle on reduceByKey
Optimization: custom partitioner on domain

PageRank performance



^{*} Figure 10a: 30 machines on 54 GB of Wikipedia data computing PageRank