

# Fundamental about Ramsey Numbers, Ramsey Theory and Their Combinatorial Significance

## 1 Introduction

### 1.1 Historical Context

The concept of Ramsey numbers emerged from the foundational work of Frank P. Ramsey in 1930, who was investigating conditions under which order inevitably arises in sufficiently large structures. His original result, published as part of a paper on logic, stated that in any partition of the elements of a sufficiently large set, a single partition class must contain a structured subset of a given size. Although Ramsey's focus was in logic, the combinatorial implications of his work laid the groundwork for what we now call **Ramsey Theory**.

Ramsey numbers, as a specific branch of Ramsey Theory, capture the essence of this phenomenon in the context of graph theory. They quantify the point at which chaos (in the form of arbitrary coloring) can no longer obscure the emergence of structure (in the form of monochromatic subgraphs). The study of these numbers has become central to combinatorics and has deep connections to fields as diverse as number theory, computer science, and even theoretical physics.

The formalization of Ramsey numbers as a distinct concept occurred much later, with key contributions from mathematicians such as Paul Erdos and George Szekeres. Erdos, in particular, introduced probabilistic methods that revolutionized the understanding of Ramsey numbers and their growth rates. Despite nearly a century of progress, the field remains rich with unsolved problems and opportunities for exploration, especially as computational methods have expanded the scope of what can be calculated or conjectured.

### 1.2 Problem Statement

At its core, the study of Ramsey numbers involves a simple yet profound question: *How large must a graph be to ensure that it contains a certain level of structure, regardless of how its edges are colored?* Specifically, given two integers  $s$  and  $t$ , the Ramsey number  $R(s, t)$  is defined as the smallest integer  $n$  such that any edge-coloring of a complete graph  $K_n$  with two colors (commonly red and blue) contains either a monochromatic  $K_s$  in red or a monochromatic  $K_t$  in blue.

Formally, we can write:

$$R(s, t) = \min\{n \mid \forall \text{ red-blue edge-colorings of } K_n, \exists \text{ a red } K_s \text{ or a blue } K_t\}.$$

This definition extends naturally to more colors. For example, the multicolor Ramsey number  $R_k(s_1, s_2, \dots, s_k)$  generalizes the problem to  $k$ -colorings of the edges of  $K_n$ .

Even in this seemingly straightforward framework, calculating Ramsey numbers is an extraordinarily challenging task. While simple cases such as  $R(2, 2) = 2$  or  $R(3, 3) = 6$  can be determined with elementary arguments, the value of  $R(5, 5)$ —the smallest graph size where a monochromatic  $K_5$  in either color is guaranteed—is not known. Current bounds on  $R(5, 5)$  suggest that it lies between 43 and 46, but the exact value remains elusive despite significant computational and theoretical efforts.

### 1.3 Importance of Ramsey Numbers

The study of Ramsey numbers is not merely a theoretical exercise. It provides insights into fundamental questions about structure and randomness in mathematics and beyond. For example:

- **In combinatorics**, Ramsey numbers offer a lens into the interplay between global constraints and local patterns, underpinning broader results in extremal graph theory.
- **In computer science**, they inform problems in distributed systems, where ensuring consistent behavior in the face of adversarial conditions often involves Ramsey-like guarantees.
- **In logic and set theory**, Ramsey-type arguments appear in the study of infinite sets and partition relations, connecting finite Ramsey numbers to infinite generalizations.
- **In physics**, concepts from Ramsey theory have been applied to study phase transitions and critical thresholds in complex systems.

Moreover, the computational challenge of determining exact Ramsey numbers has driven the development of sophisticated algorithms and heuristics. As such, Ramsey numbers sit at the intersection of pure mathematics and practical computation, serving as a benchmark problem for understanding complexity in combinatorial settings.

### 1.4 Scope and Goals

This thesis focuses on finite Ramsey numbers, primarily in the context of graph theory. While Ramsey theory encompasses a vast array of topics, from infinite structures to applications in other disciplines, this work hones in on:

1. The properties and bounds of **classical Ramsey numbers**  $R(s, t)$  for two colors.
2. The extension to **multicolor Ramsey numbers**  $R_k(s_1, s_2, \dots, s_k)$ .
3. Key results, methods, and conjectures in the determination of Ramsey numbers.
4. Computational approaches and their role in advancing the understanding of Ramsey numbers.

The goals of this thesis are to:

- Provide a comprehensive overview of known results and open problems in the study of Ramsey numbers.
- Explore the theoretical tools, such as probabilistic methods and combinatorial bounds, that underpin current knowledge.
- Discuss the computational challenges and breakthroughs associated with Ramsey numbers, including heuristic and algorithmic approaches.

By focusing solely on Ramsey numbers, this thesis aims to present a detailed and focused treatment of their mathematical significance, theoretical challenges, and ongoing developments.

## 1.5 Structure of the Thesis

The thesis is organized as follows:

- **Introduction** (this section): Lays the groundwork by defining Ramsey numbers, explaining their significance, and outlining the goals of the thesis.
- **Preliminaries**: Covers basic definitions, key concepts, and simple examples that serve as the foundation for understanding Ramsey numbers.
- **Ramsey Numbers**: Fundamental knowledge about Ramsey Numbers: Ramsey Theorem, its property, and basic bounds.
- **Value of Ramsey Numbers**: The exact value or the range of Ramsey Numbers, and the proof for the known exact value numbers.
- **Bounds on Ramsey Numbers**: Discusses upper and lower bounds, including results from probabilistic methods and combinatorial arguments.
- **Multicolor Ramsey Numbers**: Extends the discussion to  $k$ -color Ramsey numbers, highlighting additional challenges and results.
- **Applications and Generalizations**: Explores the broader significance of Ramsey numbers and connections to other areas of mathematics.
- **Open Problems**: Identifies unresolved questions and potential avenues for future research.

This structure ensures a logical progression from foundational concepts to advanced topics, providing a thorough exploration of Ramsey numbers and their role in mathematics.

## 2 Preliminaries

### 2.1 Definitions

Before delving into the main results, we establish some basic definitions and concepts essential to the study of Ramsey numbers.

1. **Graph:** A graph  $G$  is defined as an ordered pair  $G = (V, E)$ , where  $V$  is a set of vertices, and  $E$  is a set of edges, where each edge  $e \in E$  connects two distinct vertices from  $V$ .
2. **Subgraph:** A subgraph of a graph  $G = (V, E)$  is a graph  $G' = (V', E')$  where  $V' \subseteq V$  and  $E' \subseteq E$ .
3. **Complete Graph ( $K_n$ ):** A graph with  $n$  vertices in which every pair of vertices is connected by an edge. The graph  $K_n$  contains  $\binom{n}{2}$  edges.
4. **Edge Coloring:** An assignment of colors to the edges of a graph. For Ramsey numbers, we typically consider 2-colorings (e.g., red and blue) or  $k$ -colorings.
5. **Monochromatic Subgraph:** A subgraph where all edges are of the same color.
6. **Clique Number:** The size of the largest complete subgraph in a given graph.
7. **Chromatic Number:** The minimum number of colors required to color the vertices of a graph such that no two adjacent vertices share the same color.

The Ramsey number  $R(s, t)$  is the smallest integer  $n$  such that any red-blue edge coloring of the complete graph  $K_n$  contains either a red  $K_s$  or a blue  $K_t$ . For multicolor problems,  $R_k(s_1, s_2, \dots, s_k)$  generalizes this definition to  $k$ -color edge colorings.

## 3 Ramsey Numbers

Of all divisions of Ramsey theory, one of the most researched and well-known is that of Ramsey numbers. In this section, we will establish a formal defintion for Ramsey's Theorem, and give important properties of Ramsey Numbers.

### 3.1 Ramsey's Theorem

The cornerstone of Ramsey theory is the following result, which ensures the existence of Ramsey numbers for all pairs of positive integers  $s$  and  $t$ .

**Theorem 1** (Ramsey's Theorem). *For all integers  $s, t \geq 1$ , the Ramsey number  $R(s, t)$  exists. In particular, for any  $n \geq R(s, t)$ , any 2-coloring of the edges of the complete graph  $K_n$  contains either a red  $K_s$  or a blue  $K_t$ .*

*Proof.* The proof is by induction, and it is similar to the above part Bound of Ramsey Numbers. If  $s = 1$  or  $t = 1$ , then  $R(s, t) = 1$ , as any  $K_n$  with  $n \geq 1$  trivially contains a monochromatic  $K_1$ . Assume  $R(s-1, t)$  and  $R(s, t-1)$  exist. Define  $n = R(s-1, t) + R(s, t-1)$ . Consider a 2-coloring of  $K_n$ , select an arbitrary vertex  $v$ . Partition the remaining vertices into two sets:  $V_1$ , the vertices connected to  $v$  by red edges, and  $V_2$ , the vertices connected to  $v$  by blue edges. By the pigeonhole principle, either  $|V_1| \geq R(s-1, t)$  or  $|V_2| \geq R(s, t-1)$ . If  $|V_1| \geq R(s-1, t)$ , the subgraph induced by  $V_1$  contains a red  $K_{s-1}$ , and adding  $v$  forms a red  $K_s$ . Similarly, if  $|V_2| \geq R(s, t-1)$ , the subgraph induced by  $V_2$  contains a blue  $K_t$ . Thus,  $n$  satisfies the conditions for  $R(s, t)$ , completing the proof.  $\square$

### 3.2 Important Properties of Ramsey Numbers

The following theorem provide a foundation for understanding the properties and bounds of Ramsey numbers.

**Theorem 2.** *For all integers  $s, t \geq 1$ , the following inequalities hold:*

$$R(s, t) \geq R(s-1, t) \quad \text{and} \quad R(s, t) \geq R(s, t-1). \quad (1)$$

*Proof.*

1.  $R(s, t) \geq R(s-1, t)$ : By the definition of  $R(s, t)$ ,  $R(s, t)$  is the smallest integer  $n$  such that any red-blue coloring of the edges of  $K_n$  contains either a monochromatic  $K_s$  in red, or a monochromatic  $K_t$  in blue. Similarly,  $R(s-1, t)$  is the smallest integer  $n'$  such that any red-blue coloring of the edges of  $K_{n'}$  contains either a monochromatic  $K_{s-1}$  in red, or a monochromatic  $K_t$  in blue. If  $n \geq R(s, t)$ , then any red-blue coloring of  $K_n$  guarantees the existence of a monochromatic  $K_s$  in red or  $K_t$  in blue. Since  $K_s$  contains  $K_{s-1}$  as a subgraph, the existence of  $K_s$  also implies the existence of  $K_{s-1}$ . Thus,  $n \geq R(s, t)$  also satisfies the conditions for  $R(s-1, t)$ . Therefore,  $R(s, t) \geq R(s-1, t)$ .
2.  $R(s, t) \geq R(s, t-1)$ : By the definition of  $R(s, t)$ ,  $R(s, t)$  is the smallest integer  $n$  such that any red-blue coloring of the edges of  $K_n$  contains either a monochromatic  $K_s$  in red, or a monochromatic  $K_t$  in blue. Similarly,  $R(s, t-1)$  is the smallest integer  $n'$  such that any red-blue coloring of the edges of  $K_{n'}$  contains either a monochromatic  $K_s$  in red, or a monochromatic  $K_{t-1}$  in blue. If  $n \geq R(s, t)$ , then any red-blue coloring of  $K_n$  guarantees the existence of a monochromatic  $K_s$  in red or  $K_t$  in blue. Since  $K_t$  contains  $K_{t-1}$  as a subgraph, the existence of  $K_t$  also implies the existence of  $K_{t-1}$ . Thus,  $n \geq R(s, t)$  also satisfies the conditions for  $R(s, t-1)$ . Therefore,  $R(s, t) \geq R(s, t-1)$ .

Combining these two results, we conclude that Ramsey numbers are monotonic:

$$R(s, t) \geq R(s-1, t) \quad \text{and} \quad R(s, t) \geq R(s, t-1).$$

$\square$

**Theorem 3.** *For all integers  $s, t \geq 1$ ,*

$$R(s, t) = R(t, s). \quad (2)$$

*Proof.* By definition,  $R(s, t)$  is the smallest  $n$  such that every red-blue coloring of the edges of  $K_n$  contains either a red  $K_s$ , or a blue  $K_t$ . Consider any red-blue coloring of  $K_n$ . If we swap the roles of red and blue in the coloring, then a red  $K_s$  becomes a blue  $K_s$ , and a blue  $K_t$  becomes a red  $K_t$ . This transformation does not change the structure of the graph; it simply interchanges the colors. Thus, any statement about finding a red  $K_s$  or a blue  $K_t$  applies equally to finding a blue  $K_s$  or a red  $K_t$ . If  $n$  satisfies the condition for  $R(s, t)$ , then  $n$  also satisfies the condition for  $R(t, s)$ . Therefore,  $R(s, t) \leq R(t, s)$ . The argument can be repeated in reverse. Specifically, if  $n$  satisfies  $R(t, s)$ , then it also satisfies  $R(s, t)$ . Hence,  $R(t, s) \leq R(s, t)$ . Since  $R(s, t) \leq R(t, s)$  and  $R(t, s) \leq R(s, t)$ , we conclude that:

$$R(s, t) = R(t, s).$$

□

**Theorem 4.** *Ramsey numbers satisfy the following recurrence relation:*

$$R(s, t) \leq R(s - 1, t) + R(s, t - 1), \quad (3)$$

with the base cases  $R(1, t) = R(s, 1) = 1$  for all  $s, t \geq 1$ .

*Proof.* By definition,  $R(s, t)$  is the smallest  $n$  such that every red-blue coloring of the edges of  $K_n$  contains either a red  $K_s$ , or a blue  $K_t$ . Let  $n = R(s - 1, t) + R(s, t - 1)$ , and consider any red-blue edge coloring of  $K_n$ . Choose an arbitrary vertex  $v$  in  $K_n$ . The vertex  $v$  connects to  $n - 1$  other vertices, which we partition into two subsets based on the color of the edge connecting  $v$ :  $V_1$  is the set of vertices connected to  $v$  by red edges, and  $V_2$  is the set of vertices connected to  $v$  by blue edges. By the pigeonhole principle, one of the sets  $V_1$  or  $V_2$  must have at least:

$$|V_1| \geq R(s - 1, t) \quad \text{or} \quad |V_2| \geq R(s, t - 1).$$

Now we will consider cases

1.  $|V_1| \geq R(s - 1, t)$ : By the definition of  $R(s - 1, t)$ , the subgraph induced by  $V_1$  contains either a red  $K_{s-1}$ , or a blue  $K_t$ . If  $V_1$  contains a red  $K_{s-1}$ , adding  $v$  to this subgraph forms a red  $K_s$ . If  $V_1$  contains a blue  $K_t$ , we are done.
2.  $|V_2| \geq R(s, t - 1)$ : By the definition of  $R(s, t - 1)$ , the subgraph induced by  $V_2$  contains either a red  $K_s$ , or a blue  $K_{t-1}$ . If  $V_2$  contains a blue  $K_{t-1}$ , adding  $v$  to this subgraph forms a blue  $K_t$ . If  $V_2$  contains a red  $K_s$ , we are done.

In any cases, the graph  $K_n$  contains either a red  $K_s$ , or a blue  $K_t$ . Hence,  $n = R(s - 1, t) + R(s, t - 1)$  satisfies the conditions for  $R(s, t)$ , which implies:

$$R(s, t) \leq R(s - 1, t) + R(s, t - 1).$$

□

### 3.3 Erdos-Szekeres Bound

A key result that provides an explicit upper bound on Ramsey numbers is the following.

**Theorem 5.** *For all integers  $s, t \geq 2$ ,*

$$R(s, t) \leq \binom{s+t-2}{s-1}. \quad (4)$$

*Proof.* The inequality holds if  $s = 2$  or  $t = 2$  (in fact, we have equality since  $R(s, 2) = R(2, s) = s$ ). Assume now that  $s > 2$ ,  $t > 2$  and (4) holds for every pair  $(s', t')$  with  $2 \leq s' + t' < s + t$ . Then by (1) we have

$$\begin{aligned} R(s, t) &\leq R(s-1, t) + R(s, t-1) \\ &\leq \binom{s+t-3}{s-2} + \binom{s+t-3}{s-1} = \binom{s+t-2}{s-1}. \end{aligned}$$

□

This bound provides a combinatorial estimate for Ramsey numbers but is often not tight, especially for large values of  $s$  and  $t$ .

## 4 Value of Ramsey Numbers

Since most of the talk about Ramsey numbers has been in relation to the general case  $R(s, t)$ , in this section, we will give some actual values for Ramsey numbers. It's important to note that there is a distinction between a known Ramsey number and a Ramsey number for which only upper or lower bounds are known, and the table below will give the values and the bound with  $s, t \leq 10$ , and the exact values will be noted in green while the value with upper and lower bound will be noted in yellow.

		s									
	R(s, t)	1	2	3	4	5	6	7	8	9	10
t	1	1	1	1	1	1	1	1	1	1	1
	2	1	2	3	4	5	6	7	8	9	10
	3	1	3	6	9	14	18	23	28	36	40 - 41
	4	1	4	9	18	25	36 - 40	49 - 58	59 - 79	73 - 105	92 - 135
	5	1	5	14	25	43 - 46	59 - 85	80 - 133	101 - 193	133 - 282	149 - 381
	6	1	6	18	36 - 40	59 - 85	102 - 160	115 - 270	134 - 423	183 - 651	204 - 944
	7	1	7	23	49 - 58	80 - 133	115 - 270	205 - 492	219 - 832	252 - 1368	292 - 2119
	8	1	8	28	59 - 79	101 - 193	134 - 423	219 - 832	282 - 1518	329 - 2662	343 - 4402
	9	1	9	36	73 - 105	133 - 282	183 - 651	252 - 1368	329 - 2662	565 - 4956	581 - 8675
	10	1	10	40 - 41	92 - 135	149 - 381	204 - 944	292 - 2119	343 - 4402	581 - 8675	798 - 16064

## 4.1 Proofs of Edge Case Ramsey Numbers

Now that we have a basic understanding of the Ramsey numbers that have been established, it is time to present actual evidence for these values. Obviously, we cannot spend time going over every proof of every Ramsey number that has been discovered. Rather, we will focus on two general proofs that establish all Ramsey numbers  $R(s, t)$  where  $s \leq 2$ , or  $t \leq 2$ , and then move on to proofs of the two other main diagonal Ramsey numbers that have been identified in the next part.

### 4.1.1 $R(1, u) = 1$

We start with the simplest Ramsey numbers, which can be inferred from above table that  $R(u, 1) = R(1, u) = 1$ .

**Theorem 6.**  $R(u, 1) = R(1, u) = 1$

*Proof.* By definition,  $R(u, 1)$  is the smallest  $n$  such that any red-blue edge coloring of  $K_n$  contains either a red  $K_u$ , or a blue  $K_1$ . However, since a  $K_1$  is a single vertex, then a blue  $K_1$  always exists in any graph because a single vertex does not depend on any edges, thus exactly one vertex is enough, hence  $R(u, 1) = R(1, u) = 1$ .  $\square$

#### 4.1.2 $R(2, u) = u$

Taking one step up in complexity, we analyze all Ramsey numbers with  $s = 2, t = u$  or  $s = u, t = 2$  and get  $R(u, 2) = R(2, u) = u$ .

**Theorem 7.**  $R(u, 2) = R(2, u) = u$

*Proof.* Consider  $K_{u-1}$ , a complete graph with  $u - 1$  vertices. Assign all edges the same color (e.g., all red). This coloring does not contain any blue  $K_u$ , as the graph has only  $u - 1$  vertices. Additionally, since  $K_2$  is just a single edge, a red  $K_2$  always exists trivially, irrespective of the graph size. Thus,  $R(2, u) > u - 1$ , so  $R(2, u) \geq u$ .

Now consider  $K_u$ , a complete graph with  $u$  vertices. In any red-blue edge coloring of  $K_u$ , there will be either: At least one red  $K_2$  (a single red edge) must exist, as  $K_2$  is formed by any pair of connected vertices, or if no red  $K_2$  exists, then all edges must be blue. In that case, the entire  $K_u$  becomes a blue  $K_u$ , which contains a blue clique of size  $u$ . Therefore,  $K_u$  satisfies the conditions of  $R(2, u)$ , implying that  $R(2, u) \leq u$ .

Combine together, we achieve  $R(u, 2) = R(2, u) = u$ .  $\square$

## 4.2 Proofs of Known Ramsey Number

Now that we have proved some values for simpler Ramsey numbers, we will focus in on known Ramsey numbers of the other form  $R(u, u)$ . When  $s = t = u$ , the Ramsey number  $R(u, u)$  is called a *diagonal Ramsey number*. These numbers are of particular interest, and their exact values are known only for small  $u$ . We also will discuss about some off-diagonal Ramsey numbers, which can be extremely useful for bounding the others Ramsey numbers.

#### 4.2.1 $R(3, 3) = 6$

The first Ramsey number of this form to consider is  $R(3, 3)$ , the number involved in the Party Problem.

**Theorem 8.**  $R(3, 3) = 6$

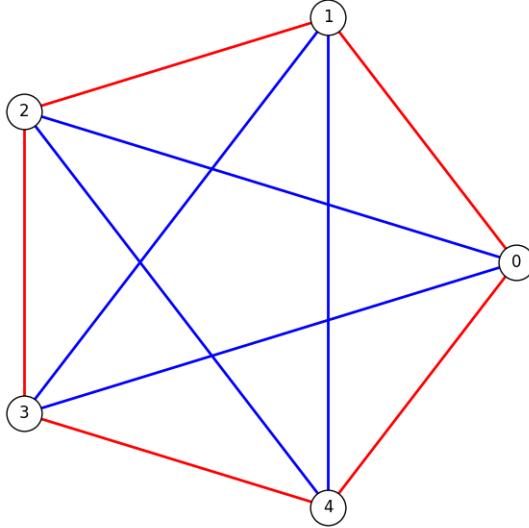
*Proof.* By definition,  $R(3, 3)$  is the smallest  $n$  such that every red-blue coloring of the edges of  $K_n$  contains either a red  $K_3$ , or a blue  $K_3$ . We will prove that  $R(3, 3) = 6$ .

- $R(3, 3) > 5$ : To show that  $R(3, 3) > 5$ , we provide a red-blue coloring of  $K_5$  that avoids monochromatic  $K_3$ . Label the vertices of  $K_5$  as  $v_1, v_2, v_3, v_4, v_5$ . Color the edges as follows:

- Color the edges of a  $K_5$  cyclically alternating between red and blue. For instance, let:

$(0, 1), (1, 2), (2, 3), (3, 4), (4, 0)$  be red, and

$(0, 2), (0, 3), (1, 3), (1, 4), (2, 4)$  be blue.



- Verify that there is no monochromatic  $K_3$ : Each triangle formed by any three vertices will always have edges of both colors.
- $R(3, 3) \leq 6$ : To show that  $R(3, 3) \leq 6$ , we argue that any red-blue edge coloring of  $K_6$  must contain a monochromatic  $K_3$ . Consider any vertex  $v$  in  $K_6$ . This vertex connects to the other 5 vertices via edges, each of which is either red or blue. By the **Pigeonhole Principle**, at least 3 of these edges must be of the same color (either red or blue), as there are only two colors and 5 edges. Without loss of generality, let  $v$  connect to three vertices  $v_1, v_2, v_3$  via red edges. If any of the edges between  $v_1, v_2$ , and  $v_3$  is also red, we form a red  $K_3$  (triangle). If not, all edges between  $v_1, v_2$ , and  $v_3$  are blue, forming a blue  $K_3$ .

Since  $5 < R(3, 3) \leq 6$ , thus  $R(3, 3) = 6$ . □

By Erdos-Szekeres bound, we have  $R(s, t) \leq R(s - 1, t) + R(s, t - 1)$ . However, we can make it tighter a little bit for a specific case.

**Theorem 9.** *If  $R(s - 1, t)$  and  $R(s, t - 1)$  are both even, then we have  $R(s, t) \leq R(s - 1, t) + R(s, t - 1) - 1$*

*Proof.* Suppose  $R(s - 1, t) = 2p$  and  $R(s, t - 1) = 2q$ . Take a graph of  $2p + 2q - 1$  vertices and a vertex  $A$ . There are  $2p + 2q - 2$  edges ending at  $A$ . Then, consider the following cases:

1.  $2p$  or more edges end at  $A$ ,
2.  $2q$  or more edges end at  $A$ ,
3.  $2p - 1$  red edges end at  $A$  and  $2q - 1$  blue edges end at  $A$ .

For the first case, consider the set  $T_1$  of the vertices at the further ends of the  $2p$  or more segments. Since the number of vertices in  $T_1$  is greater than or equal to  $R(s - 1, t)$ , there is either a red  $K_{s-1}$  or a blue  $K_t$ . However, if there is a red  $K_{s-1}$ , then the set  $T_1 \cup \{A\}$  is a red  $K_s$ . Thus, the theorem holds in this case.

The same argument shows that the theorem holds for the second case as well.

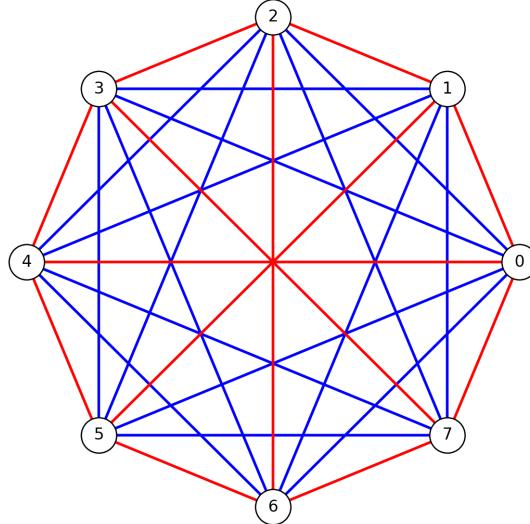
The third case cannot hold for every vertex  $A$  of the graph. Indeed, if it did, there would be  $(2p + 2q - 1)(2p - 1)$  red endpoints, which is an odd number. However, every edge has two endpoints, so this number should be even. This means that there exists at least one vertex for which either case 1 or case 2 holds. Since the theorem was shown for these two cases, it holds for the third case too.  $\square$

We currently have a tighter bound, now we can use it to find a new Ramsey numbers.

**Theorem 10.**  $R(3, 4) = 9$

*Proof.* We will do the same way to prove  $R(3, 4) = 9$  as  $R(3, 3) = 6$ . First, we want to use Theorem 8 for a bound, since both  $R(2, 4) = 4$  and  $R(3, 3) = 6$  are both even, then

$$R(3, 4) \leq R(2, 4) + R(3, 3) - 1 = 4 + 6 - 1 = 9.$$

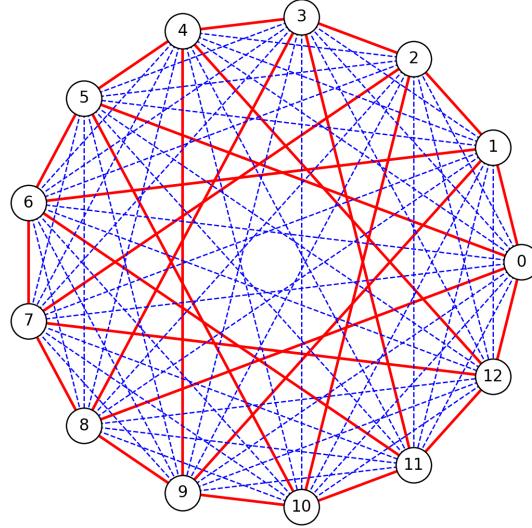


We now show that  $R(3, 4) > 8$ , by constructing a contradicted graph with 8 vertices  $\{0, 1, \dots, 7\}$ . Build a red regular graph with edges  $(0, 1), (1, 2), \dots, (6, 7), (7, 0)$  and red edges  $(0, 4), (1, 5), (2, 6), (3, 7)$ . The complement of this graph will be colored blue. The combination of two graphs will form a counter example for us: There is no red  $K_3$  and no blue  $K_4$ , thus  $R(3, 4) > 8$ . Since  $8 < R(3, 4) \leq 9$  then  $R(3, 4) = 9$ .  $\square$

**Theorem 11.**  $R(3, 5) = 14$

*Proof.* We will do the same way to prove  $R(3, 5) = 14$  as  $R(3, 4) = 9$ . First, we have a bound (note that we cannot use Theorem 8 since  $R(2, 5) = 5$  and  $R(3, 4) = 9$  are not both even)

$$R(3, 5) \leq R(2, 5) + R(3, 4) = 5 + 9 = 14.$$



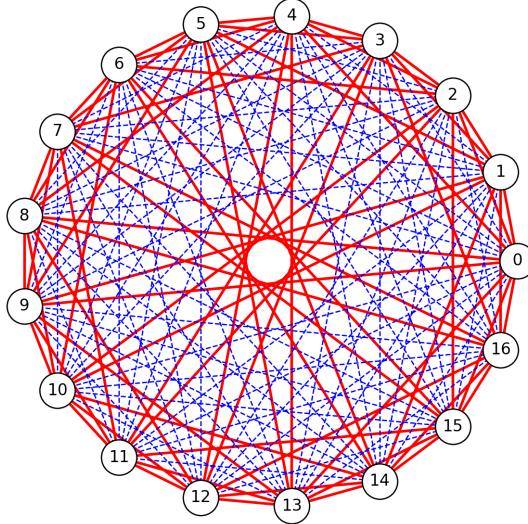
We now show that  $R(3, 5) > 13$ , by constructing a contradicted graph with 13 vertices  $\{0, 1, \dots, 12\}$ . Build a graph with red edges, that connect vertex  $u$  and  $v$  if and only if the difference of the numbers of  $u$  and  $v$  in modulo 13 is 1, 5, 8 or 12. The complement of this graph will be colored blue. The combination of two graphs will form a counter example for us: There is no red  $K_3$  and no blue  $K_5$ , thus  $R(3, 5) > 13$ . Since  $13 < R(3, 5) \leq 14$  then  $R(3, 5) = 14$ .  $\square$

We will now go to the last diagonal Ramsey number that we currently know the exact value of it, which is  $R(4, 4)$

**Theorem 12.**  $R(4, 4) = 18$

*Proof.* We will do the same way to prove  $R(4, 4) = 18$  as  $R(3, 4) = 9$ . First, we have a bound

$$R(4, 4) \leq R(3, 4) + R(4, 3) = 2R(3, 4) = 18$$



We now show that  $R(4, 4) > 17$ , by constructing a contradicted graph with 17 vertices  $\{0, 1, \dots, 17\}$ . Build a graph with red edges, that connect vertex  $u$  and  $v$  if and only if the difference of the numbers of  $u$  and  $v$  in modulo 17 is 1, 2, 4, 8, 9, 13, 15, or 16. The complement of this graph will be colored blue. The combination of two graphs will form a counter example for us: There is no red  $K_4$  and no blue  $K_4$ , thus  $R(4, 4) > 17$ . Since  $17 < R(4, 4) \leq 18$  then  $R(4, 4) = 18$ .  $\square$

### 4.3 Computation Based Ramsey Numbers

Beside some Ramsey numbers we have proved above, there are a few more numbers that which are known, those are  $R(3, 6) = 18, R(3, 7) = 23, R(3, 8) = 28, R(3, 9) = 36$  and  $R(4, 5) = 25$ . However, in order to find these numbers, beside the finder ingenuity, they must use a great amount of computational power, since the method we used above, which we have to bound and check, is no longer useful, due to the large number of cases and graphs.

## 5 Bounds on Ramsey Numbers

Determining the exact value of Ramsey numbers is a notoriously difficult problem. Instead, much of Ramsey theory focuses on establishing upper and lower bounds for these numbers. This section explores the key methods and results related to these bounds.

### 5.1 Aforementioned Bounds

This section will recall every bounds we have already mentioned, before we use it to have another simple bounds.

- For all  $s, t$ ,  $R(s, t) \leq R(s - 1, t) + R(s, t - 1)$

- For all  $s, t$  such that  $R(s-1, t)$  and  $R(s, t-1)$  are both even,  $R(s, t) \leq R(s-1, t) + R(s, t-1) - 1$
- Erdos-Szekeres bound: For all  $s, t$ ,  $R(s, t) \leq \binom{s+t-2}{s-1}$

## 5.2 Diagonal Ramsey Numbers

We want to distinguish *diagonal Ramsey numbers*  $R(s) = R(s, s)$  and *off-diagonal Ramsey numbers*  $R(s, t)$ ,  $s \neq t$ . It is not surprising that the diagonal Ramsey numbers are of greatest interest, and they are also the hardest to estimate. Note that a graph is *trivial* if it is either complete or empty, the diagonal Ramsey number  $R(s)$  is the minimal integer  $n$  such that every graph of order  $n$  has a trivial subgraph of order  $s$ .

## 5.3 Erdos Bound and some bounding methods

The initial bounds for diagonal Ramsey number  $R(u, u)$  were first established in a 1947 paper by Erdos, in which it was proved that, for  $u \geq 3$ , then

$$2^{u/2} < R(u, u) < 4^{u-1}.$$

We will prove each part of the statement. The simpler one, the upper bound  $R(u, u) < 4^{u-1}$  will be proved first.

**Theorem 13.**  $R(u, u) < 4^{u-1}$

*Proof.* We have  $R(u, u) \leq \binom{2u-2}{u-1}$ , that is

$$\begin{aligned} R(u, u) &\leq \frac{(2u-2)!}{(u-1)!(u-1)!} = \frac{(2u-2)(2u)\dots(2)\cdot(2u-3)(2u-5)\dots(1)}{(u-1)!(u-1)(u-2)\dots(1)} \\ &= 2^{u-1} \frac{2u-3}{u-1} \frac{2u-5}{u-2} \dots \frac{3}{2} \frac{1}{1} < 2^{u-1} \cdot 2^{u-1} = 4^{u-1} \end{aligned}$$

□

However, the other part is much more difficult, so we might use another method to solve it, which is probabilistic methods.

### 5.3.1 Probabilistic Methods

Introduced by Erdos, probabilistic methods revolutionized Ramsey theory by providing non-constructive lower bounds. The essence of this approach lies in demonstrating that a random coloring of the edges of  $K_n$  avoids monochromatic subgraphs with high probability.

**Theorem 14.** If  $\binom{n}{k} 2^{1-\binom{k}{2}} < 1$ , then  $R(k, k) > n$

*Proof.* We want to prove that  $R(k, k) > n$ , then we need to show there exists a coloring of the edges of  $K_n$ , such that there is no monochromatic  $K_k$ . If we assign the color of the edges in  $K_n$  randomly, and each edge will be color independently, then for all edge  $e$ ,

$$\mathcal{P}(\text{edge } e \text{ is red}) = \mathcal{P}(\text{edge } e \text{ is blue}) = \frac{1}{2}$$

There is  $\binom{n}{k}$  copies of  $K_k$  in  $K_n$ , and let  $A_i$  be the event when  $K_k$  is monochromatic (either red or blue), then

$$\mathcal{P}(A_i) = 2 \cdot \left(\frac{1}{2}\right)^{\binom{k}{2}} = 2^{1-\binom{k}{2}}$$

Then we have

$$\mathcal{P}(\exists \text{ a monochromatic } K_k) = \mathcal{P}\left(\bigcup_i A_i\right) = \binom{n}{k} 2^{1-\binom{k}{2}}$$

However, since we already have  $\binom{n}{k} 2^{1-\binom{k}{2}} < 1$  as the requirement, then

$$\mathcal{P}(\exists \text{ a monochromatic } K_k) < 1$$

Which means

$$\mathcal{P}(\exists \text{ a coloring with no monochromatic } K_k) > 0$$

Thus there exists a colouring with no monochromatic  $K_k$ . Hence,  $R(k, k) > n$ .  $\square$

By using this theorem, we can show the lower bound of Erdos bound, which is  $R(u, u) > 2^{\frac{u}{2}}$  for  $u \geq 3$ .

**Theorem 15.** For all integers  $u \geq 3$ ,

$$R(u, u) > 2^{\frac{u}{2}}.$$

*Proof.* For  $u \geq 3$ , let  $n = \lfloor 2^{\frac{u}{2}} \rfloor$ , then

$$\binom{n}{u} 2^{1-\binom{u}{2}} \leq \frac{n^u}{u!} 2^{1-\frac{u(u-1)}{2}} \leq \frac{(2^{\frac{u}{2}})^u}{u!} \cdot 2^{1-\frac{u^2}{2}+\frac{u}{2}} = \frac{2^{1+\frac{u}{2}}}{u!}$$

However, for  $u \geq 3$ ,  $\frac{2^{1+\frac{u}{2}}}{u!} < 1$ , and according to above theorem,  $R(u, u) > 2^{\frac{u}{2}}$ .  $\square$

### 5.3.2 Explicit Constructions

Explicit constructions of graphs with no monochromatic  $K_s$  or  $K_t$  provide alternative lower bounds. For example:

- Certain constructions based on finite projective planes or other combinatorial designs provide specific lower bounds for multicolor Ramsey numbers.
- As we show in the proof of  $R(4, 4) = 18$ , we show that 2-coloring of  $K_{17}$  avoids both a red  $K_4$  and a blue  $K_4$ , showing  $R(4, 4) > 17$ .

## 5.4 Asymptotics of Ramsey Numbers

For large  $s$  and  $t$ , the asymptotic behavior of Ramsey numbers becomes a key focus. The following results summarize the current understanding:

- For diagonal Ramsey numbers  $u = s = t$ :

$$2^{\frac{u}{2}} < R(u, u) < 4^{u-1}$$

- For  $R(3, t)$ :

$$R(3, t) \sim \frac{t^2}{\log t}.$$

- For  $R(4, t)$ :

$$R(4, t) = \Omega\left(\frac{t^3}{\log^4(t)}\right)$$

The gap between upper and lower bounds illustrates the challenges inherent in Ramsey theory, motivating further research into both theoretical and computational approaches.

## 5.5 Summary of Bounds

The bounds on Ramsey numbers provide crucial insights into their growth and behavior. Upper bounds demonstrate the inevitability of structure in large graphs, while lower bounds highlight the existence of colorings that avoid certain configurations. Together, these results form a foundation for understanding the complexity and depth of Ramsey theory.

# 6 Multicolor Ramsey Numbers

In addition to the classical two-color Ramsey numbers  $R(s, t)$ , the concept generalizes naturally to more than two colors. Multicolor Ramsey numbers, denoted  $R_k(s_1, s_2, \dots, s_k)$ , extend Ramsey theory to  $k$ -colorings of the edges of complete graphs.

## 6.1 Definitions and Notation

For  $k$ -colorings, the multicolor Ramsey number  $R_k(s_1, s_2, \dots, s_k)$  is defined as the smallest integer  $n$  such that for any edge-coloring of the complete graph  $K_n$  with  $k$  colors, there exists a monochromatic complete subgraph  $K_{s_i}$  in color  $i$  for some  $i \in \{1, 2, \dots, k\}$ . Formally,

$$R_k(s_1, s_2, \dots, s_k) = \min\{n \mid \forall k\text{-colorings of } K_n, \exists \text{ a monochromatic } K_{s_i} \text{ in color } i\}.$$

For the special case where  $s_1 = s_2 = \dots = s_k = s$ , the notation simplifies to  $R_k(s, s, \dots, s)$ , which represents the smallest  $n$  such that every  $k$ -coloring of  $K_n$  contains a monochromatic  $K_s$  in some color.

## 6.2 Key Results in Multicolor Ramsey Numbers

### 6.2.1 Existence of Multicolor Ramsey Numbers

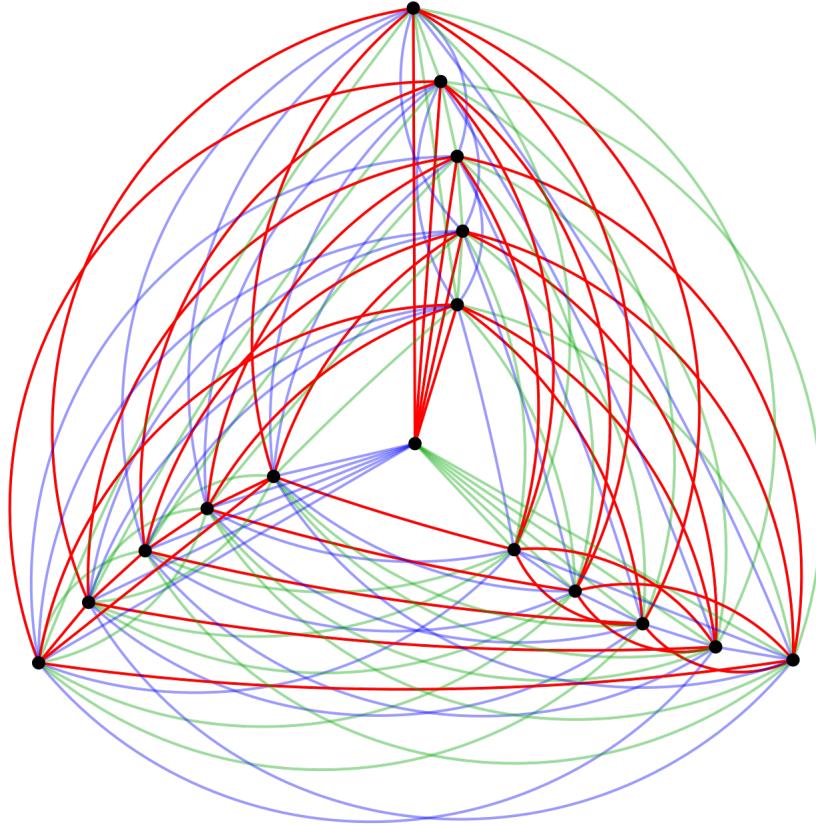
Ramsey's theorem ensures that  $R_k(s_1, s_2, \dots, s_k)$  exists for all  $s_1, s_2, \dots, s_k \geq 1$ . The existence follows from the same principles as the two-color case, extended to  $k$ -colorings.

### 6.2.2 Bounds for Multicolor Ramsey Numbers

The growth of  $R_k(s_1, s_2, \dots, s_k)$  as  $k$  increases is rapid and significantly larger than in the two-color case. Some important bounds are as follows:

- For  $k = 3$  and  $s_1 = s_2 = s_3 = 3$ ,  $R_3(3, 3, 3) = 17$ .

*Proof.* Below here is a proof for  $R_3(3, 3, 3) = 17$ . First, we want to make sure that  $R_3(3, 3, 3) > 16$ , that is true, since here is a counterexample for  $R_3(3, 3, 3) \leq 16$ .



We must show that 17 is a sufficient number. In any chromatic graph with 3 colors and 17 vertices, we choose a vertex  $v$ . At least 6 edges start at  $v$  must all have the same color. If in the subgraph consisting of  $v$  and the 6 other ends of edges there is no triangle of the original color, then all interconnections among the 6 other ends of edges must be of the other two colors. But this is the case  $R(3, 3) = 6$ , and there is an interconnected set of 3 here.  $\square$

- For  $k = 3$  and  $s_1 = s_2 = s_3 = s$ , it is known that:

$$R_k(s, s, \dots, s) \text{ grows exponentially as } k^{s/2}.$$

- General upper bound:

$$R_k(s_1, s_2, \dots, s_k) \leq R_{k-1}(s_1, s_2, \dots, s_{k-1}, R(s_k, s_k)).$$

### 6.2.3 Diagonal Multicolor Ramsey Numbers

For diagonal multicolor Ramsey numbers  $R_k(s, s, \dots, s)$ , asymptotic bounds reveal the rapid growth:

$$k^{s/2} \leq R_k(s, s, \dots, s) \leq k^{s-1}.$$

The gap between the lower and upper bounds remains an active area of research, particularly for large  $k$ .

## 6.3 Examples of Multicolor Ramsey Numbers

Explicit values of multicolor Ramsey numbers are known only for small  $s$  and  $k$ . Some examples include:

- $R_3(3, 3, 3) = 17$ : Any 3-coloring of  $K_{17}$  contains a monochromatic  $K_3$ .
- $R_3(4, 4, 4)$ : Exact value unknown, but bounds suggest  $50 \leq R_3(4, 4, 4) \leq 59$ .
- $R_4(3, 3, 3, 3)$ : Bounds indicate exponential growth as  $k$  increases.

## 6.4 Techniques for Multicolor Ramsey Numbers

Several methods are used to establish bounds or exact values for multicolor Ramsey numbers:

### 6.4.1 Recursive Techniques

For  $R_k(s_1, s_2, \dots, s_k)$ , recursion plays a key role. By extending the two-color recurrence relation:

$$R_k(s_1, s_2, \dots, s_k) \leq R_{k-1}(s_1, s_2, \dots, s_{k-1}, R(s_k, s_k)),$$

we can derive bounds for higher  $k$ -color cases using lower  $k$ -color results.

### 6.4.2 Probabilistic Methods

Probabilistic methods generalize to  $k$ -colorings, demonstrating the existence of colorings that avoid monochromatic cliques of specified sizes. These methods yield lower bounds of the form:

$$R_k(s, s, \dots, s) \geq k^{s/2}.$$

### 6.4.3 Computational Approaches

As in the two-color case, computational methods play a significant role in determining exact values or tighter bounds for multicolor Ramsey numbers. Algorithms based on exhaustive search or heuristic methods have been applied to small cases.

## 6.5 Applications of Multicolor Ramsey Numbers

Multicolor Ramsey numbers have applications in various areas of mathematics and science:

- **Graph Theory:** Understanding multicolor Ramsey numbers provides insights into extremal problems and the structure of large graphs under  $k$ -color constraints.
- **Number Theory:** Multicolor generalizations of van der Waerden's theorem and related results link Ramsey theory to arithmetic progressions.
- **Theoretical Computer Science:** Problems in distributed systems and network design often involve  $k$ -color variants of Ramsey-type problems.

## 6.6 Open Problems in Multicolor Ramsey Numbers

Multicolor Ramsey numbers present several open challenges:

- Determining exact values for small  $k$  and  $s$ , such as  $R_3(4, 4, 4)$ .
- Narrowing the gap between upper and lower bounds for  $R_k(s, s, \dots, s)$  as  $k$  grows.
- Developing efficient algorithms for computing multicolor Ramsey numbers.

## 6.7 Summary

Multicolor Ramsey numbers generalize classical Ramsey numbers to  $k$ -color settings, introducing new challenges and opportunities. The rapid growth of  $R_k(s, s, \dots, s)$  as  $k$  increases highlights the complexity of the problem, while known bounds and exact values provide a foundation for further exploration. The interplay between theoretical, probabilistic, and computational approaches continues to drive progress in this rich area of combinatorics.

# 7 Applications and Generalizations

Ramsey numbers and Ramsey theory, in general, have far-reaching applications across mathematics and science. This section explores these applications, focusing on their implications in graph theory, number theory, and computer science. Additionally, we examine significant generalizations of Ramsey theory to infinite structures and other combinatorial domains.

## 7.1 Applications of Ramsey Numbers

### 7.1.1 Graph Theory

Ramsey numbers are a fundamental concept in graph theory, particularly in the study of extremal problems. Key applications include:

- **Structure in Large Graphs:** Ramsey numbers provide guarantees about the emergence of structure in large graphs, irrespective of how their edges are colored. For example,  $R(3, 3) = 6$  implies that any  $K_6$ , no matter how its edges are 2-colored, contains a monochromatic triangle.
- **Turan's Theorem:** Ramsey numbers connect to Turan's theorem, which provides bounds on the maximum number of edges in a graph that avoids specific subgraphs.
- **Graph Coloring:** Ramsey theory has implications for edge and vertex coloring problems, offering insights into how constraints on coloring force the existence of particular substructures.

### 7.1.2 Number Theory

Ramsey theory extends naturally to number theory, where it guarantees structure within sets of integers. Two prominent examples include:

- **Van der Waerden's Theorem:** For any given positive integers  $r$  and  $k$ , there is some number  $N$  such that if the integers  $\{1, 2, \dots, N\}$  are colored, each with one of  $r$  different colors, then there are at least  $k$  integers in arithmetic progression whose elements are of the same color. The least such  $N$  is the Van der Waerden number  $W(r, k)$ .

*Proof.* We will prove the special case  $W(2, 3) \leq 325$ . Let  $c(n)$  be a coloring of the integers  $\{1, \dots, 325\}$ . We will find three elements of  $\{1, \dots, 325\}$  in arithmetic progression that are the same color.

Divide  $\{1, \dots, 325\}$  into the 65 blocks  $\{1, \dots, 5\}, \{6, \dots, 10\}, \dots, \{321, \dots, 325\}$ , thus each block is of the form  $\{5b + 1, \dots, 5b + 5\}$  for some  $b \in \{0, \dots, 64\}$ . Since each integer is colored either red or blue, each block is colored in one of 32 different ways. By the *pigeonhole principle*, there are two blocks among the first 33 blocks that are colored identically. That is, there are two integers  $b_1$  and  $b_2$ , both in  $\{0, \dots, 32\}$ , such that

$$c(5b_1 + k) = c(5b_2 + k)$$

for all  $k \in \{1, \dots, 5\}$ . Among the three integers  $5b_1 + 1, 5b_1 + 2, 5b_1 + 3$ , there must be at least two that are of the same color. (The *pigeonhole principle* again.) Call these  $5b_1 + a_1$  and  $5b_1 + a_2$ , where the  $a_i$  are in  $\{1, 2, 3\}$  and  $a_1 < a_2$ . Suppose (without loss of generality) that these two integers are both red. (If they are both blue, just exchange red and blue in what follows.)

Let  $a_3 = 2a_2 - a_1$ . If  $5b_1 + a_3$  is red, then we have found our arithmetic progression:  $5b_1 + a_i$  are all red.

Otherwise,  $5b_1 + a_3$  is blue. Since  $a_3 \leq 5$ ,  $5b_1 + a_3$  is in the  $b_1$  block, and since the  $b_2$

block is colored identically,  $5b_2 + a_3$  is also blue.

Now let  $b_3 = 2b_2 - b_1$ . Then  $b_3 \leq 64$ . Consider the integer  $5b_3 + a_3$ , which must be  $\leq 325$ . We will consider its color.

If it is red, then  $5b_1 + a_1$ ,  $5b_2 + a_2$ , and  $5b_3 + a_3$  form a red arithmetic progression. But if it is blue, then  $5b_1 + a_3$ ,  $5b_2 + a_3$ , and  $5b_3 + a_3$  form a blue arithmetic progression. Either way, we are done.  $\square$

However, this bound is extremely bad when comparing to the correct result, that is  $W(2, 3) = 9$ . We will want to prove why  $W(2, 3) = 9$ .

*Proof.*  $W(2, 3) > 8$ , since with 8, we have this construction

1	2	3	4	5	6	7	8
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We now prove that  $W(2, 3) \leq 9$ , we will consider cases with 9 numbers, depending on the color of 4, 5, 6. Obviously, they cannot have the same color, so we divide them into parts, 2 of them will go to part  $A$  and the other go to  $B$ . We will consider other numbers to put in these sets.

- $4, 6 \in A$ . This case is not hard to eliminate, since  $2, 5, 8$  must then be an arithmetic triple in  $B$ : Consider, respectively, the triples  $2, 4, 6$ , and  $4, 5, 6$ , and  $4, 6, 8$ . If we place any of  $2, 5, 8 \in A$ , one of these three arithmetic triples ends up in  $A$ .
- $4, 5 \in A$ . Then  $3, 6 \in B$  (consider, respectively, the triples  $3, 4, 5$  and  $4, 5, 6$ ), so  $9 \in A$  (consider  $3, 6, 9$ ), but then  $1, 7 \in B$  (consider  $1, 5, 9$  and  $5, 7, 9$ ), so  $2, 8 \in A$  (consider  $1, 2, 3$  and  $6, 7, 8$ ). We now see this case cannot be either, because the triple  $2, 5, 8$  is in  $A$ .
- $5, 6 \in A$ . This case is same as Case 2, by symmetry. (In this case,  $4, 7 \in B$ , so  $1 \in A$ , so  $3, 9 \in B$ , so  $2, 8 \in A$ , and we see that the triple  $2, 5, 8$  is in  $A$ .)

Thus,  $W(2, 3) \leq 9$ , and  $W(2, 3) > 8$ , then  $W(2, 3) = 9$ .  $\square$

- **Schur's Theorem:** For any  $k \geq 2$ , there is  $n > 3$  such that for any  $k$ -coloring of  $\{1, 2, \dots, n\}$ , there are three integers  $x, y, z$  of the same color such that  $x + y = z$ .

These results demonstrate the interplay between Ramsey theory and additive number theory, where colorings correspond to partitions of sets of integers.

### 7.1.3 Theoretical Computer Science

Ramsey theory finds significant applications in theoretical computer science, particularly in areas such as:

- **Distributed Systems:** Ramsey-type arguments are used to ensure fault tolerance and consistency in distributed networks, where systems must maintain certain guarantees under adversarial conditions.
- **Complexity Theory:** Ramsey numbers serve as benchmarks for understanding the computational complexity of combinatorial problems. Determining Ramsey numbers is an inherently hard problem, motivating algorithmic advancements.
- **Data Structures:** Ramsey-type results are used in data storage and retrieval problems, ensuring that structured subsets exist for efficient computation.

#### 7.1.4 Physics and Network Theory

In physics, Ramsey theory appears in studies of phase transitions and critical phenomena. For example:

- **Percolation Theory:** Ramsey-theoretic principles help analyze connectivity in random networks.
- **Neural Networks:** The inevitability of structure in large systems informs models of neural activity and pattern recognition.

In network theory, Ramsey numbers provide insights into the robustness of networks under different conditions, such as ensuring reliable connections in communication systems.

## 7.2 Generalizations of Ramsey Theory

### 7.2.1 Infinite Ramsey Theory

The principles of Ramsey theory extend naturally to infinite sets, leading to profound results in logic and set theory. The infinite version of Ramsey's theorem states:

**Theorem 16** (Infinite Ramsey's Theorem). *For any infinite set  $S$  and any finite  $k$ -coloring of  $\binom{S}{2}$ , there exists an infinite subset  $S' \subseteq S$  such that all edges in  $\binom{S'}{2}$  are the same color.*

This result has significant implications in:

- **Set Theory:** Ramsey-type results underpin many partition principles, such as those in the study of ordinal numbers and cardinality.
- **Logic:** Infinite Ramsey theory contributes to the foundations of mathematics, particularly in model theory and proof theory.

### 7.2.2 Hales-Jewett Theorem

The Hales-Jewett theorem generalizes Ramsey theory to higher-dimensional grids, stating that for any  $k$ -coloring of a sufficiently large  $d$ -dimensional grid, there exists a monochromatic combinatorial line. Formally:

**Theorem 17** (Hales-Jewett Theorem). *For all integers  $d, k \geq 1$  and all  $n$ , there exists a  $H(d, k)$  such that any  $k$ -coloring of  $[n]^d$  contains a monochromatic combinatorial line.*

This result has applications in both finite and infinite combinatorics and forms a cornerstone of multidimensional Ramsey theory.

### 7.2.3 Applications to Hypergraphs

Ramsey theory generalizes to hypergraphs, where edges consist of subsets of vertices of size greater than 2. The multicolor Ramsey number  $R_k^{(r)}(s_1, s_2, \dots, s_k)$  for  $r$ -uniform hypergraphs extends the concept of Ramsey numbers to these structures. Applications of hypergraph Ramsey theory include:

- **Design Theory:** Ensuring the existence of combinatorial designs within hypergraphs.
- **Coding Theory:** Analyzing the properties of error-correcting codes.

## 7.3 Summary

The applications and generalizations of Ramsey theory illustrate its profound impact across mathematics, computer science, and physics. From understanding the inevitability of order in finite and infinite settings to providing tools for analyzing complex systems, Ramsey theory offers insights that transcend its combinatorial origins. These results inspire ongoing research into both fundamental problems and practical applications, ensuring that Ramsey theory remains a vibrant and evolving field.

# 8 Open Problems

Despite significant advances in Ramsey theory, many questions about Ramsey numbers remain unresolved. This section highlights key open problems and areas of active research, particularly regarding the determination of exact values, bounding strategies, and computational challenges.

## 8.1 Exact Values of Ramsey Numbers

### 8.1.1 Two-Color Ramsey Numbers

Determining exact values of  $R(s, t)$  for larger  $s$  and  $t$  remains a central challenge in Ramsey theory. While values such as  $R(3, 3) = 6$  and  $R(4, 4) = 18$  are well-established, many questions remain open:

- **Diagonal Ramsey Numbers:** The exact value of  $R(5, 5)$  is unknown but lies between 43 and 46. Improving these bounds is a significant goal.
- **Off-Diagonal Ramsey Numbers:** Determining  $R(3, t)$  for larger  $t$  is an active area of research. For example, while  $R(3, 4) = 9$  is known, exact values for  $R(3, 5)$  and beyond are elusive.

### 8.1.2 Multicolor Ramsey Numbers

For multicolor Ramsey numbers, exact values are known for only a few cases. Open questions include:

- What is the exact value of  $R_3(4, 4, 4)$ ? Current bounds suggest  $50 \leq R_3(4, 4, 4) \leq 59$ .
- How does  $R_k(s, s, \dots, s)$  grow as  $k$  increases? Although bounds such as

$$k^{s/2} \leq R_k(s, s, \dots, s) \leq k^{s-1}$$

exist, closing this gap remains a challenge.

## 8.2 Improving Bounds

While probabilistic and combinatorial methods provide useful bounds for Ramsey numbers, these bounds are often far apart, particularly for large  $s$  and  $t$ . Key questions include:

- Can new probabilistic techniques tighten the lower bounds for  $R(s, s)$ ?
- Can advances in combinatorial constructions or regularity methods refine the upper bounds for  $R(s, t)$ ?
- For multicolor Ramsey numbers, how can recursive relations or computational approaches narrow the gap between bounds?

## 8.3 Asymptotic Behavior

The asymptotic growth of Ramsey numbers is one of the most intriguing aspects of Ramsey theory. Open questions in this area include:

- For diagonal Ramsey numbers, does the true growth rate of  $R(s, s)$  lie closer to  $2^{s/2}$  or  $4^s$ ? Understanding the asymptotics of  $R(s, s)$  is critical to resolving this question.
- How do off-diagonal Ramsey numbers  $R(3, t)$  behave as  $t \rightarrow \infty$ ? Current results suggest  $R(3, t) \sim \frac{t^2}{\log t}$ , but more precise asymptotics are needed.
- What are the asymptotic properties of multicolor Ramsey numbers  $R_k(s, s, \dots, s)$  as  $k$  grows?

## 8.4 Computational Challenges

Determining exact values of Ramsey numbers is computationally intensive, and even small cases require significant resources. Open problems in this area include:

- **Algorithm Development:** Can more efficient algorithms be developed to compute Ramsey numbers or verify their bounds?

- **Computational Complexity:** What is the exact complexity class of computing  $R(s, t)$ ? Determining whether the problem lies in NP or a higher complexity class remains unresolved.
- **Heuristics and Approximations:** Can heuristic approaches or machine learning techniques be applied to predict Ramsey numbers or improve existing bounds?

## 8.5 Generalizations and Extensions

Several generalizations of Ramsey theory pose open questions:

- **Hypergraph Ramsey Numbers:** How can bounds for  $R_k^{(r)}(s, s, \dots, s)$ , the  $k$ -color Ramsey number for  $r$ -uniform hypergraphs, be improved?
- **Infinite Ramsey Theory:** While the infinite version of Ramsey's theorem is well-established, questions about the combinatorial properties of infinite structures remain open.
- **Geometric Ramsey Theory:** How do Ramsey-type results extend to geometric settings, such as points in the plane or higher-dimensional spaces?

## 8.6 Interdisciplinary Connections

Ramsey theory has implications in other disciplines, raising questions such as:

- How can Ramsey-theoretic principles improve fault tolerance and robustness in distributed systems?
- What new connections can be drawn between Ramsey theory and machine learning, particularly in understanding patterns in large datasets?
- Can Ramsey theory be applied to quantum computing, particularly in analyzing entanglement or other structural properties of quantum systems?

## 8.7 Summary

The open problems in Ramsey theory reflect the depth and complexity of this field. From determining exact values and improving bounds to exploring asymptotic behavior and computational challenges, Ramsey theory remains a vibrant area of mathematical research. These open questions not only drive theoretical advancements but also inspire interdisciplinary applications, ensuring that Ramsey theory continues to evolve and expand.

## 9 References

Below here are all the resources I have used for this thesis.

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7. On the Ramsey numbers  $R(3, 8)$  and  $R(3, 9)$ , Charles M. Grinstead, Sam M. Roberts (1981)
8. The Ramsey number  $R(3, t)$  has order of magnitude  $t^2 / \log t$ , J. H. Kim (1995)