

Interpretations & Anabelian Geometry

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(preliminary report)

GOAL:

To show how

$$G \sim (G, O^\Delta)$$
$$(G, O^+)$$
$$(G, O^{+\mu})$$

can be formalized.

G_{group}

Classical group language

G_{group}

infinitary group language

G^{Γ}

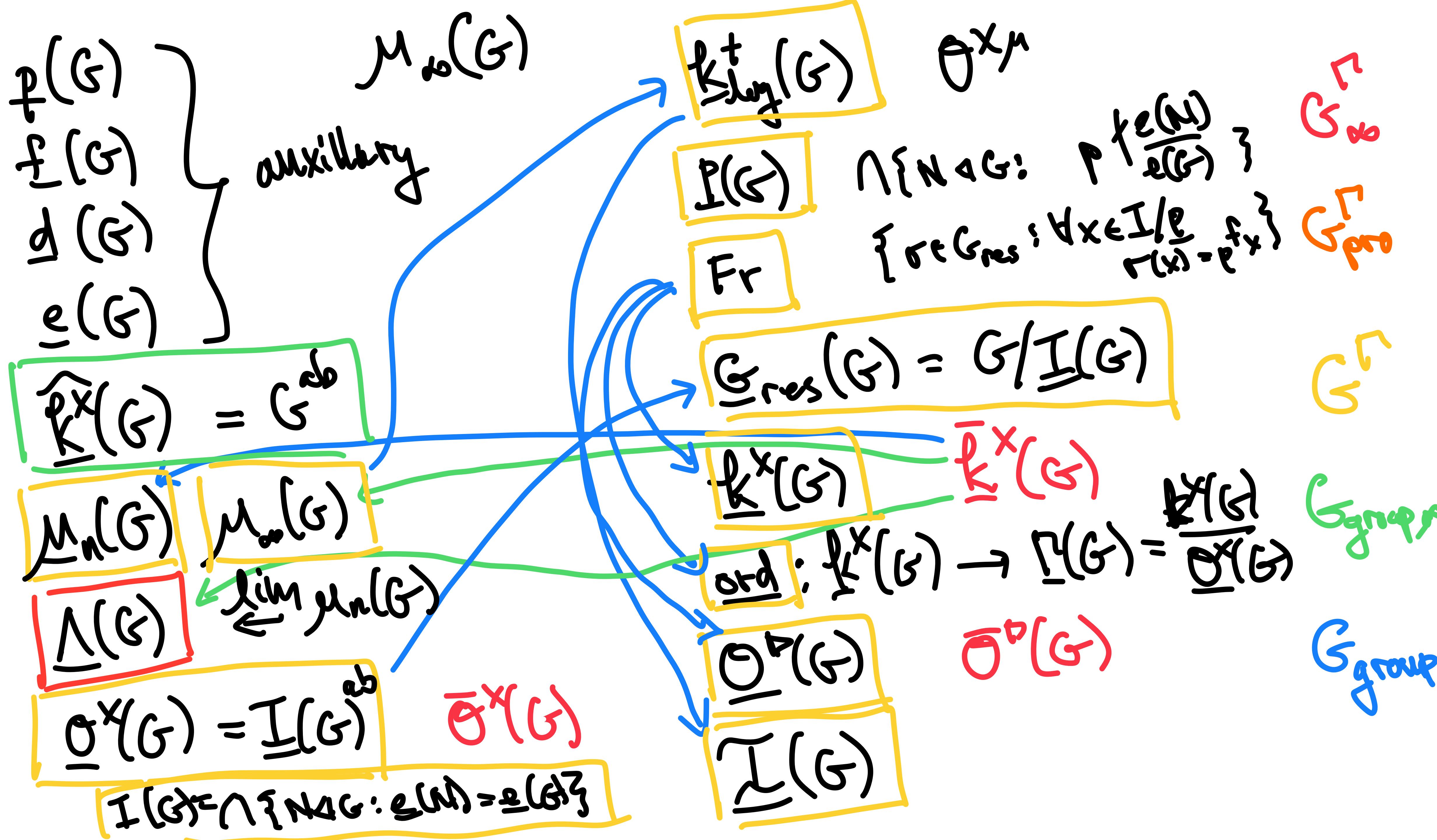
Γ -invariant sets $\Gamma = \text{Aut}(G)$

G_{per}

PO- Γ -invariant sets

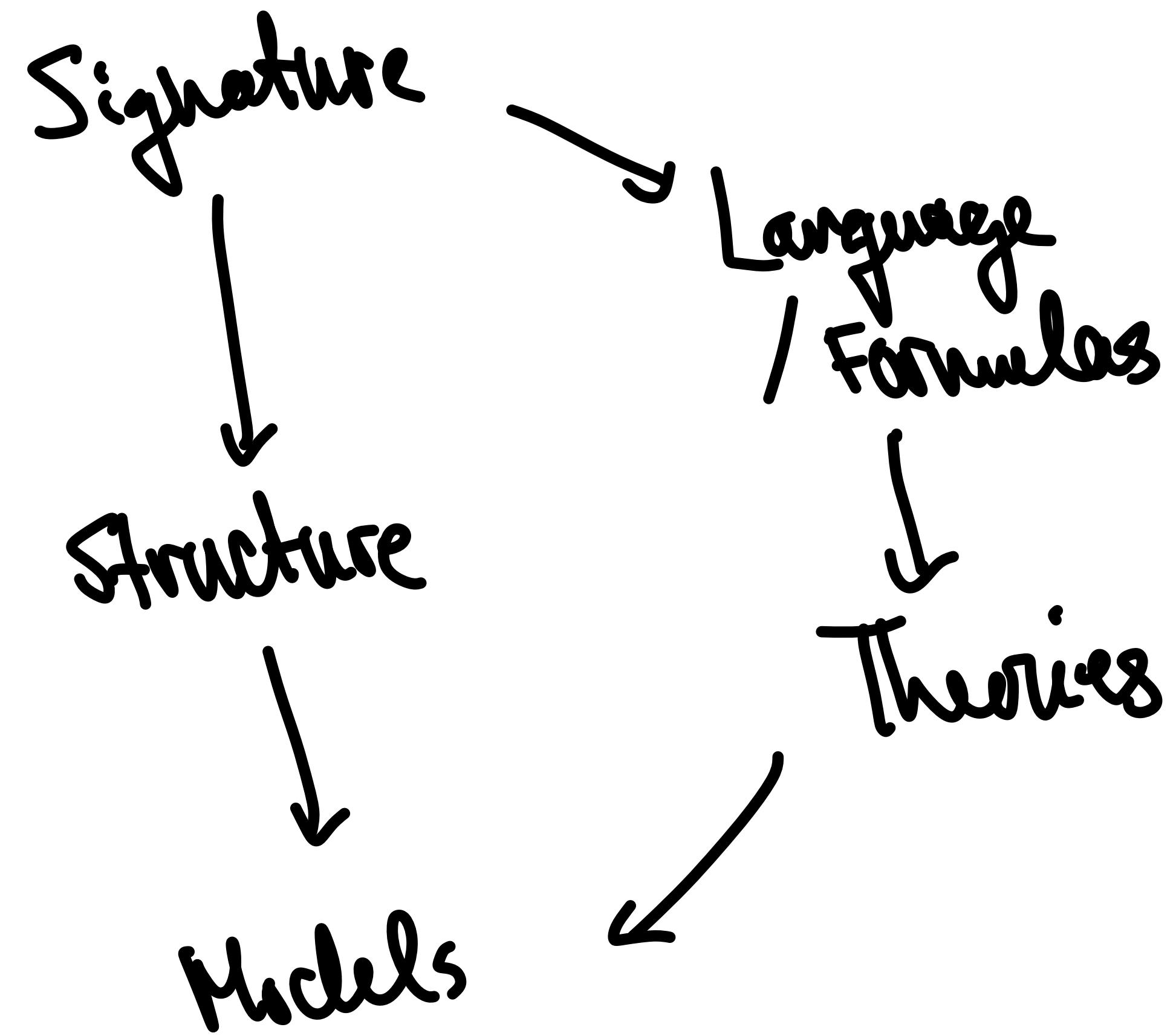
G_{∞}^{Γ}

Γ -invariant sets w/ disjoint unions



Crash Course In Model Theory

Crash Course In Model Theory



Axioms of Groups

First Order Formulas

Crash Course In Model Theory

- 0) ~~$\exists! e \in G, \forall x \in G (x * e = e * x = x)$~~
- 1) $\forall x \in G, \exists y \in G (x * y = y * x = e)$
- 2) $\forall x, y, z \in G ((x * y) * z = x * (y * z))$

Signature \rightarrow constant symbol

Axioms of Groups

First Order Formulas

From the signature we generate a language

Signature

$(G, *, e)$

constant symbol

function symbol

$* : G \times G \rightarrow G$

- 0) $\exists ! e \in G, \forall x \in G (x * e = e \wedge e * x = x)$
- 1) $\forall x \in G, \exists y \in G (x * y = y * x = e)$
- 2) $\forall x, y, z \in G ((x * y) * z = x * (y * z))$

Axioms of Groups

$(G, *, e)$

Another Example: Fields

$(F, +, *, 0, 1)$

- 0) $\exists ! e \in G, \forall x \in G (x * e = e \text{ and } e * x = x)$
- 1) $\forall x \in G, \exists y \in G (x * y = y * x = e)$
- 2) $\forall x, y, z \in G ((x * y) * z = x * (y * z))$

Axioms of Groups

$(G, *, e)$

i) $\forall x \in G, \exists y \in G (x * y = y * x = e)$

ii) $\forall x \in G (S(x) * x = x * S(x) = e)$

New Function Symbol; $S: G \rightarrow G$
 $S(x) = x^{-1}$

Axioms of Groups

Expansion

Old Signature
 $(G, *, e)$

New Signature
 $(G, *, S, e)$

Reduct

UPSHOT: You can have
multiple ways of
talking about the
same object.

Crash Course In Model Theory

Crash Course In Model Theory

σ -rings $\leadsto L(\sigma_{\text{rings}})$

σ -groups $\leadsto L(\sigma_{\text{groups}})$

σ -ordered
abelian
groups $= (\Gamma, +, 0, \leq) \leadsto L(\sigma_{\text{ordered}}_{\text{abelian}}_{\text{groups}})$

σ -valued
fields $= ((K, +, -, \cdot, 0, 1), (\Gamma \cup \{\infty\}, +, -, 0, \infty), \text{val})$
 $\text{val}: K^\times \rightarrow \Gamma \cup \{\infty\}$

$\sigma_{\text{ordered abelian groups}} = (\Gamma, +, 0, \leq) \rightsquigarrow L(\sigma_{\text{ordered abelian groups}})$
 $\sigma_{\text{valued fields}} = ((K, +, -, \cdot, 0, 1), (\Gamma \cup \{\infty\}, +, -, 0, \infty), \text{val})$

two "sorts"

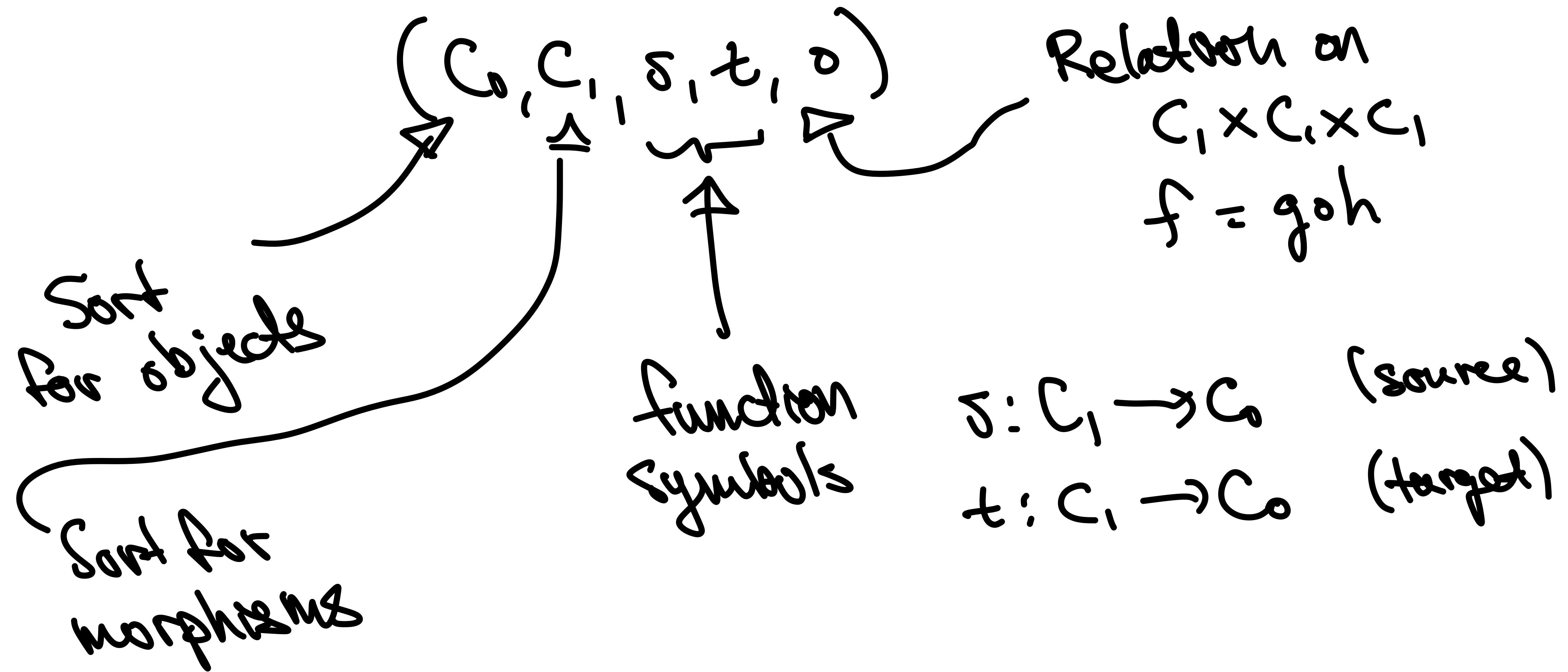
definable sets:
 $X \subseteq K^n \times \Gamma^m$

$\sigma_{\text{valued fields}} = ((K, +, -, \cdot, 0, 1), (\Gamma \cup \{\alpha\}, +, -, 0, \infty), \text{val})$

Similar to Valued Fields : $(G, 0^\nabla)$ or $(G, 0^x)$
seems multi-sorted.

Crash Course In Model Theory

Signature for Categories:



Relation Principle:

Functions, constants, and sorts can be encoded as relations.

UPSHOT:

When making abstract definitions we only need to worry about how it behaves with respect to relations.

Definable Sets

Definable Sets

Fix a signature σ

Fix a σ -structure M

Formula $\phi(x_1, \dots, x_n)$ with n free variables,

$$M \mapsto \{ (a_1, \dots, a_n) \in M^n : \phi(a_1, \dots, a_n) \text{ is TRUE} \}$$

Definable Set

Definable Set

EXAMPLE

$M \mapsto \{(a_1, \dots, a_n) \in M^n : \phi(a_1, \dots, a_n) \text{ is TRUE}\}$

$\sigma \leftarrow$ Language of Fields

$\gamma \leftarrow$ An actual field

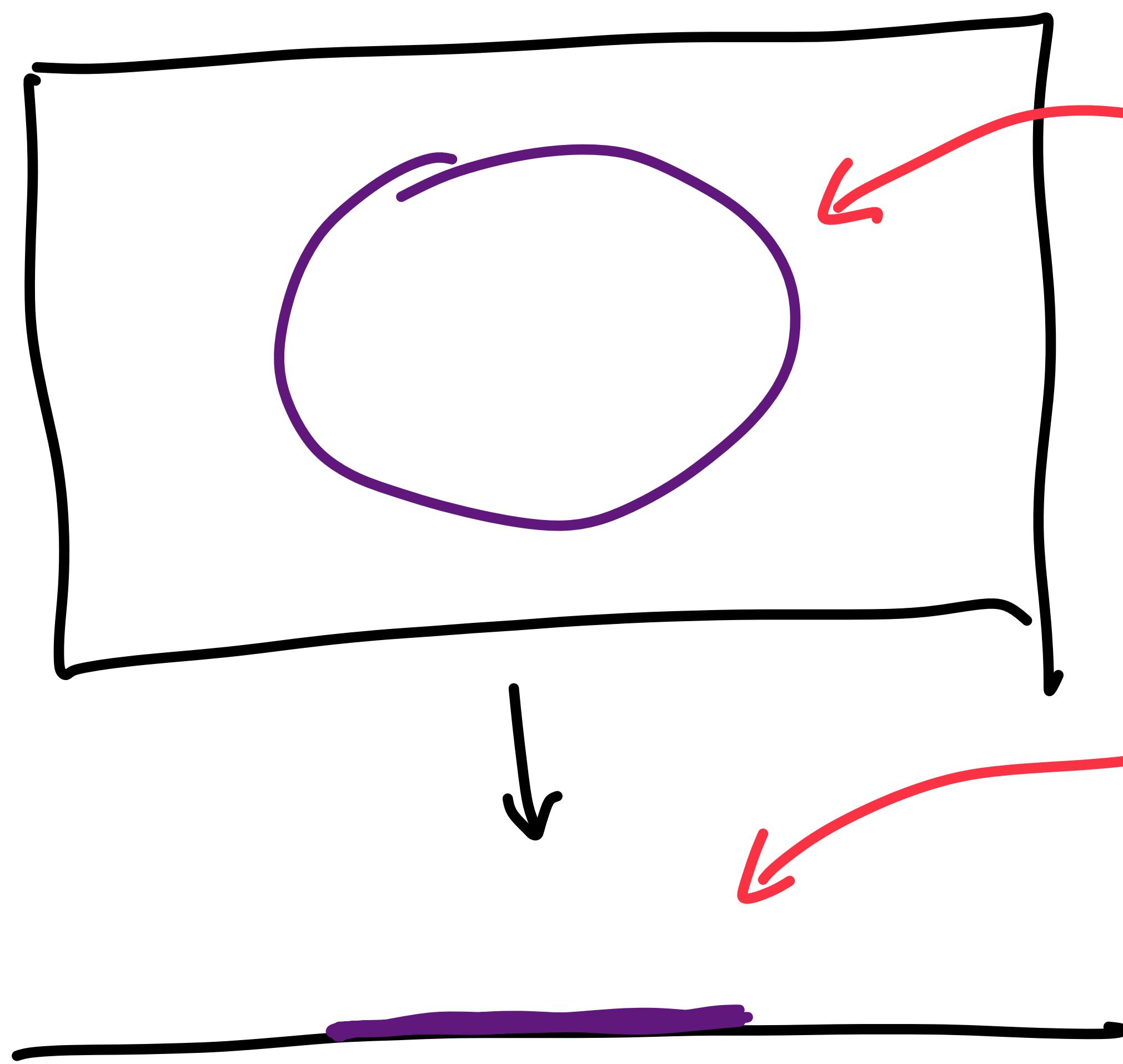
$\phi(x_1, \dots, x_n) \leftarrow \exists y (x^2 + y^2 = 1) = \phi(x)$

Definable Sets

Definable Sets

EXAMPLE

M An actual field
 $M \mapsto \{(a_1, \dots, a_n) \in M^n : \phi(a_1, \dots, a_n) \text{ is TRUE}\}$



$$\begin{aligned} & \{(x, y) \in M^2 : x^2 + y^2 = 1\} \\ & \{x \in M : \exists y (x^2 + y^2 = 1)\} \\ & [\exists y (x^2 + y^2 = 1)] = \phi(x) \end{aligned}$$

Definable Sets

TWO PERSPECTIVES

1) Functor of Points

$$M \mapsto \{ (a_1, \dots, a_n) \in M^n : \phi(a_1, \dots, a_n) \text{ is TRUE} \}$$

2) Subset $X \subseteq M^n$ for fixed M .

$$\{ (a_1, \dots, a_n) \in M^n : \phi(a_1, \dots, a_n) \text{ is TRUE} \}$$

Definable Sets | $\text{Def}_M = \left(\begin{array}{l} \text{Category of Definable} \\ \text{sets in a structure } M \end{array} \right)$

Objects: Definable sets $X \subseteq M^n$

Morphisms: Definable morphisms

$$f: X \rightarrow Y$$

Graph is Definable



[Interpretations]

Interpretations

M, N structures.

An interpretation of N in M is

$$I: X \rightarrow N$$

$\cap I$

M^M

Think:

$M = \mathbb{R}$

$N = \mathbb{C}$

such that the inverse image of every definable set is definable.



such that the inverse image of every definable set is definable.

$$Y \subseteq N^n, \quad I^{-1}(Y) \stackrel{\text{def}}{=} (I \times \cdots \times I)^{-1}(Y)$$

n-times

Interpretations

M, N structures.

An interpretation of N in M is

$$I: X \longrightarrow N$$

\cap
 M^m

such that the inverse image of every definable set is definable.

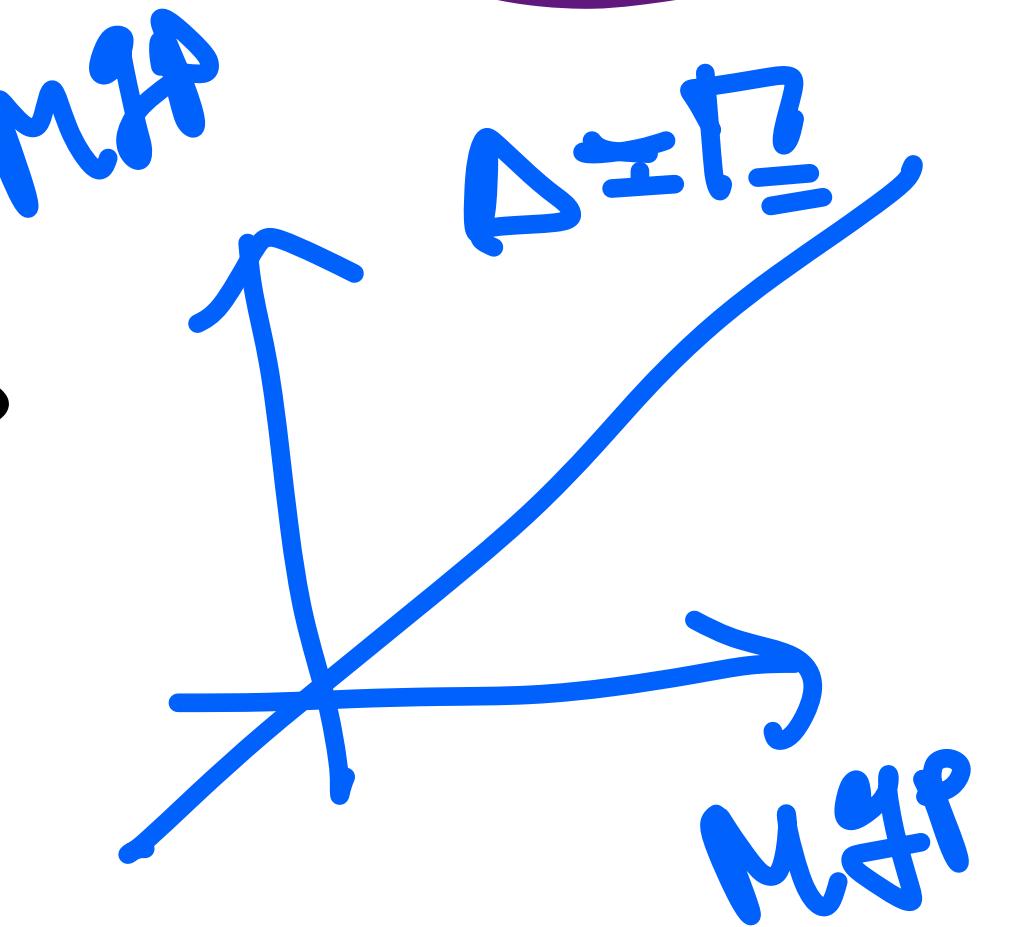
$$I : M^q \longrightarrow M^{qp}$$

$$I(a,b) = a - b.$$

Definable
set for M₀P

\vdash = \equiv = (graph of equality in Mag)

$\subset M^{gp} \times M^{gp}$



$$I(a_1, b_1) = I(a_2, b_2) \iff \begin{cases} a_1 - b_1 = a_2 - b_2 \\ a_1 + b_2 = a_2 + b_1 \end{cases}$$

$$I^{-1}(n) = \{(a_1, b_1, a_2, b_2) \in M^4 : a_1 + b_2 = a_2 + b_1\}$$

Definable
Set on Target

$\Gamma =$

$$I(a_1, b_1) = I(a_2, b_2)$$

Inverse
Image

$$a_1 + b_2 = a_2 + b_1$$

Γ_+

$$\begin{aligned} I(a_3, b_3) &= I(a_1, b_1) \\ &+ I(a_2, b_2) \end{aligned}$$

$$\begin{aligned} a_3 + b_1 + b_2 \\ = a_1 + a_2 + b_3 \end{aligned}$$

EXAMPLE

Lemma: To check that

$$I: X \longrightarrow N$$

$$\begin{matrix} \cap \\ M^m \end{matrix}$$

is an interpretation it suffices to check the
graph of all the relations.

EXAMPLE (G, \ast, e) interprets (G, \ast, \leq, e)

Proof. \ast is id
we just check $\leq \subseteq G \times G$

$$\begin{aligned} \leq &= \{(x, y) \in G^2 : y = s(x)\} \\ &= \{(x, y) \in G^2 : x \ast y = e\} \text{, 4 --- definable in } (G, \ast, e), / \end{aligned}$$

EXAMPLE:

G group

$N \triangleleft G$ normal subgroup.

(G, \star, e, \in_N)

new relation
symbol

G with " \in_N " interprets G/N .

$G \rightarrow G/N$,

$g \mapsto [g]$

$[g] = gN$.

proof idea

$[g_1] = [g_2] \iff$

$g_1 N = g_2 N \iff$

$g_1^{-1} g_2 \in N$.

Definable in
enhanced signature.

R, C structures.

$I: R \ni f \rightarrow C$ interp.
 R_n

Notation: $C(R) = I^{-1}(C)$
Structure constructed
from R .

Notation:

Write $C \leq R \iff R$ interprets C ,

R, C structures.

$I: X \rightarrow C$ interp.

$$R \stackrel{I}{\hookrightarrow} R^n$$

Notation: $C(R) = I^{-1}(C)$
Structure constructed
from R .

Notation:
Write $C \leq R \Leftrightarrow R$ interprets C ,

Composition: $R_1 \leq R_2 \text{ & } R_2 \leq R_3$ implies $R_1 \leq R_3$
 \downarrow $R_3 \vdash R_1(R_2(R))$ w/ the interp.

Defn. A history of computations is a sequence
 $R_1 \leq R_2 \leq \dots \leq R_n$

Representation: $\text{Aut}(R) \xrightarrow{\rho_I} \text{Aut}(C)$ as above.

Defn. I is multi radical $\Leftrightarrow \rho_I$ surjective.

R, C structures.

$I: X \rightarrow C$ interp.

$$R^n$$

Notation: $C(R) = I^{-1}(C)$
Structure constructed
from R .

Composition: $R_1 \leq R_2 \wedge R_2 \leq R_3$ implies $R_1 \leq R_3$

$\rightarrow R_3 \mapsto R_1(R_2(R_3))$ is the interp.

Representations: $\text{Aut}(R) \xrightarrow{\rho_I} \text{Aut}(C)$ as above.

Notation:
Write $C \leq R \iff R$ interprets C ,

Defn. A history of computations is a sequence
 $R_1 \leq R_2 \leq \dots \leq R_n$

Defn. I is multi-radical $\iff \rho_I$ surjective.

Functionality: The interpretation I induces a functor $\text{Def}_C \rightarrow \text{Def}_R$ in the category of definable sets

Functionality: The interpretation I induces a functor $\text{Def}_C \rightarrow \text{Def}_F$ in the category of definable sets

- One can characterize such functors
- One can prove converges!

Functionality: The interpretation I induces a functor $\text{Def}_C \rightarrow \text{Def}_R$ in the category of definable sets.

Theorem:

Any Boolean logical Functor $\text{Def}_C \rightarrow \text{Def}_R$ is induced by an interpretation.

Defn. C and R are bi-interpretable iff $C \leq R$ and $R \leq C$ such that the associated functors induce an equivalence of categories.*

Remark: A formal justification for replacing Frobeniuss with (G_v, O^σ) could be achieved by specifying structures & proving bi-interpretability.

Pre-Topos Completion

Given interpretation of A in B :

$$X \xrightarrow{I} A$$
$$\bigcap_{B^n}$$

→ expansion of B : (B, A, r) relation on $B^n \times A$
new sort

$$r(b, a) \Leftrightarrow I(b) = a.$$

UPSHOT: $\text{Aut}((BA)) = \text{Aut}(B) \times \text{Aut}(A)$

$$\text{Aut}((B, A, r)) = \text{Aut}(B)$$



"on/off switch" for interpretation,

Pre-Topos Completion

Given interpretation of A in B :

$$X \xrightarrow{I} A$$
$$\begin{matrix} C \\ \cap \\ B^n \end{matrix}$$

→ expansion of B : (B, A, r)
new sort

relation on $B^n \times A$
 $r(b, a) \Leftrightarrow I(b) = a.$

(Pre-Topos)
Completion ↪ (Add a sort for every interpretable
equivalence relation.)

Applications to IUT

Fundamental Groups

Object	Description	Interpretation
Π	$\cong \pi_1(Z, \bar{z})$	Given
$\mathbf{p}(G)$	$\text{char}(\mathcal{O}_K/\mathfrak{m}_K)$	the unique prime l such that for all other primes $l' \neq l$ we have $\dim_l(\Pi) - \dim_{l'}(\Pi) > 0$.
$\mathbf{d}(G)$	$[K : \mathbb{Q}_p]$	$\dim_{\mathbf{p}(\Pi)}(\Pi) - \dim_l(\Pi)$ for any prime $l \neq \mathbf{p}(\Pi)$
$\Delta(\Pi)$	$\cong \pi_1(Z_{\bar{K}})$	as in Table 5 or $\Delta(\Pi) = \bigcap \{\Pi_0 \subset \Pi \text{ clopen} : \dim_{\mathbf{p}(\Pi)}(\Pi) - \dim_l(\Pi_0) = \mathbf{d}(\Pi)[\Pi : \Pi_0]\}$ where l is any prime not equal to $\mathbf{p}(\Pi)$.
$\mathbf{G}(\Pi)$	$\cong G_K$	as in Table 5
D_I for $I \in \text{Cusp}(\Pi)$	$\cong D_{\bar{z}/z}, I_{\bar{z}/z}$ inertial group of a cusp	$N_{\Pi}(I)$
$\overline{\Delta}(\Pi)$	$\cong (\pi_1(\overline{Z}_{\bar{K}}))$	$\Delta/J(\Pi)$ where $J(\Pi) \subset \Delta(\Pi)$ smallest open normal containing I for $I \in \text{Cusp}(\Pi)$
$\Lambda(\Pi)$	$\cong \widehat{\mathbb{Z}}(1)$ as G_K -module.	$H^2(\overline{\Delta}, \widehat{\mathbb{Z}})^* = \text{Hom}(H^2(\overline{\Delta}(\Pi), \widehat{\mathbb{Z}}), \widehat{\mathbb{Z}})$
$ \Pi $ (SBT)	$\cong Z $	approximation by elements of $\text{NF}(\Pi)$ [Moc05 Lemma 3.1.i.iv]

$\text{div}(\mathcal{O}^\times(\Pi))$	$\cong \text{div}(\mathcal{O}(Z)^\times)$	$\ker(\alpha) \cap \deg^{-1}(0)$ where $\alpha: \mathbb{Z}^{\oplus \text{Cusp}(\Pi)} \rightarrow H^1(\mathbf{G}(\Pi), \overline{\Delta}(\Pi)^{\text{ab}})$ given by $\sum_{I \in \text{Cusp}(\Pi)} n_I I \mapsto \sum_{I \in \text{Cusp}(\Pi)} n_I s_{D_I}^{\text{ab}}$. See prose at beginning of this appendix.
$\mathcal{O}^\times(\Pi)$	$\cong \mathcal{O}^\times(Z)$	$p^{-1}(\text{div}(\mathcal{O}(\Pi)^\times))$
$\mathcal{O}^\times(\Pi \setminus S)$ (SBT)	$\cong \mathcal{O}^\times(Z \setminus S)$	$\mathcal{O}^\times(\pi_1(\Pi \setminus S))$
$\mathbf{K}^\times(\Pi)$ (SBT)	$\cong \kappa(Z)^\times$	$\lim_{\leftarrow} \mathcal{O}^\times(\Pi \setminus S)$
$\mathbf{k}^\times(\Pi)$	$\cong K^\times$	$\mathcal{O}^\times(\overline{\Pi})$
$\mathbf{K}_0^{\text{geom} \times}(\Pi)$ (SBT)	$\cong \kappa((Z_0)_{\overline{K}_0})^\times$	$\{\eta \in \mathbf{K}^\times(\Pi) : \exists n \in \mathbb{N}, \exists D \in \Pi _{\text{NF}} \quad \eta^n _D = 1\}$ see [Moc08 1.8.i]
$\overline{\mathbf{k}}_0^\times(\Pi)$ (SBT)	$\cong \overline{K}_0^\times$	image of $\mathbf{K}_0^{\text{geom} \times}$ and $ \Pi _{\text{NF}}$ under evaluation; compose with synchronization. [Moc08 1.8.ii]
$\text{ord}_I: \mathbf{K}^\times(\Pi) \rightarrow \mathbb{Z}, \quad I \in \text{Cusp}(S)$ $S \in \text{Open}_{\text{NF}}(\Pi)$ (SBT)	$\text{ord}_s: \kappa(Z)^\times \rightarrow \mathbb{Z}, \quad s \in Z_0 $	$\kappa_f _I \in H^1(I, \Lambda(\Pi))$
$\text{Div}(\Pi _{\text{NF}})$ (SBT)	$\cong \text{Div}((\overline{Z}_0)_{\overline{Q}})$	$\mathbb{Z}^{\oplus Z _{\text{NF}}}$
$H^0(\mathcal{O}_\Pi(D)) \quad D \in \text{Div}(\Pi _{\text{NF}})$ (SBT)	$\cong H^0(\overline{Z}, \mathcal{O}_Z(D)) = \{f \in \kappa((Z_0)_{\overline{Q}})^\times : \{f \in \mathbf{K}_0^{\text{geom}}(\overline{\Pi}) : \text{div}(f) + D \geq 0\} \text{ here } D \in \text{Div}(\overline{Z} _{\text{NF}})$	$\{f \in \mathbf{K}_0^{\text{geom}}(\overline{\Pi}) : \text{div}(f) + D \geq 0\}$ for $D \in \text{Div}((\overline{Z}_0)_{\overline{Q}})$
$\overline{\mathbf{k}}_0(\Pi)^{\text{Kum}}$ (SBT)	$\cong K_0$ as a field	Uchida trick/fundamental theorem of projective geometry (see §???)
$\overline{\mathbf{k}}_0^{\text{geom}}(\Pi)^{\text{Kum}}$ (SBT)	$\cong \kappa(Z_{\overline{K}_0})$ as a field	$\overline{\mathbf{k}}_0^{\text{geom}}(\Pi)^{\text{Kum}} = \overline{\mathbf{k}}_0^{\text{geom} \times}(\Pi)^{\text{Kum}} \cup \{0\}$ and the field structure is induced by the injection into $\bigoplus_{x \in \text{NF}(\Pi)} \kappa(x)$



Cyclotomic Synchronizations

Object	Description	Interpretation
MT	monotheta environment	given
G	$\cong G_K$	given
\overline{M}	$\cong \mathcal{O}_{\overline{K}}^\times$	given
Π	$\cong \pi_1(Z)$, Z hyperbolic curve	given
I	$\cong \widehat{\mathbb{Z}}(1)$, inertia subgroup of $\pi_1(\mathbb{Z}_{\bar{K}})$	given
$\text{sync}_{\overline{G}}^{\overline{M}}: \Lambda(\overline{M}) \rightarrow \Lambda(G)$	Brauer synchronization	$\text{inv}_G \circ H^2(G, \text{sync}_{\overline{G}}^{\overline{M}}) = \text{inv}_{(\overline{M}, G)}$
$\text{sync}_G^{\Pi}: \Lambda(\Pi) \rightarrow \Lambda(G)$	bilinear synchronization	The unique element of $\text{Hom}(\Lambda(\Pi), \Lambda(G)) \cap P$ where P is the positive rational structure. [Moc07a]
$\text{sync}_I^{\Pi}: \Lambda(\Pi) \rightarrow I$	cuspidal synchronization	$d_2^{1,0}(\text{id}_I)$; This is the map on the second page of the spectral sequence associated to the exact sequence $1 \rightarrow I \rightarrow \Delta^{\text{cc}}(X) \rightarrow \overline{\Delta}(X) \rightarrow 1$. One computes $H^0(\overline{\Delta}(X), H^1(I, I)) = \text{Hom}(I, I)$ and $H^2(\overline{\Delta}(\Pi), H^0(I, I)) = \text{Hom}(\Lambda(\Pi), I)$.
$\text{sync}_{\text{int}}^{\text{ext}}: \Lambda(\text{MT})^{\text{ext}} \rightarrow \Lambda(\text{MT})^{\text{int}}$	monotheta cyclotomic synchronization	$s - s^{\text{taut}}$

Table 4: A table of cyclotomic synchronizations.

Galois Groups

Object	Description	Interpretation
G	$\cong G_K$ as topological groups	given
$\mathbf{p}(G)$	$\text{char}(k)$	$l = \mathbf{p}(G) \iff l$ prime and $\log_l(\#\mathbf{G}^{\text{ab/tors}} / l\mathbf{G}^{\text{ab/tors}}) \geq 2$
$\mathbf{f}(G)$	$[k : \mathbb{F}_p]$	$\log_{\mathbf{p}(G)}(1 + \#((\mathbf{G}^{\text{ab/tors}})^{\mathbf{p}(G)})$)
$\mathbf{d}(G)$	$[K : \mathbb{Q}_p]$	$\log_{\mathbf{p}(G)}(\#\mathbf{G}^{\text{ab/tors}} / l\mathbf{G}^{\text{ab/tors}}) - 1$ (for any $l \neq \mathbf{p}(G)$)
$\mathbf{e}(G)$	Inertia degree	$\mathbf{d}(G)/\mathbf{f}(G)$
$\widehat{\mathbf{k}}^\times(G)^{\text{LCFT}}$	$\varprojlim K^\times / K_i^{\times n}$	G^{ab}
$\widehat{\mathbf{k}}^\times(G)^{\text{LCFT}}$	$\varprojlim \varprojlim K_i^\times / K_i^{\times n}$	$\varprojlim_{G_0 \subset G} \widehat{\mathbf{k}}^\times(G_0)$, this limit varies over open subgroups and the transition maps are given by the transfer maps
$\mu_n(G)$	$\cong \mu_n(\overline{K})$ as G_K -modules	n -torsion of $\widehat{\mathbf{k}}^\times(G)$
$\Lambda(G)$	$\cong \widehat{\mathbb{Z}}(1)$ as G_K -modules	$\varprojlim \mu_n(G)$
$\mathcal{O}^\times(G)$	$\cong \mathcal{O}_K^\times$	$\ker(G \rightarrow \mathbf{G}_{\text{res}}(G))$
$\overline{\mathcal{O}}^\times(G)$	$\cong \mathcal{O}_{\overline{K}}^\times$	$\varprojlim_{G_0 \subset G} \mathcal{O}^\times(G)$
$\mathbf{k}_{\log}^+(G)$	$\cong K_{\log}^+$ as topological abelian groups	$\mathcal{O}^\times(G) \otimes_{\mathbb{Z}} \mathbb{Q}$ (same as monoid perfection)

Applications to IUT

$[K : \mathbb{Q}_p]_{\text{tors}}$

finite extn
of \mathbb{Q}_p

$G = (\text{Absolute Galois Group } \mathfrak{g} \text{ } K)$

Galois groups

Object	Description	Interpretation
G	$\cong G_K$ as topological groups	given
$\mathbf{p}(G)$	$\text{char}(k)$	$l = \mathbf{p}(G) \iff l$ prime and $\log_l(\#G^{\text{ab/tors}}/lG^{\text{ab/tors}}) \geq 2$
$\mathbf{f}(G)$	$[k : \mathbb{F}_p]$	$\log_{\mathbf{p}(G)}(1 + \#((G^{\text{ab/tors}})^{\mathbf{p}(G)}))$
$\mathbf{d}(G)$	$[K : \mathbb{Q}_p]$	$\log_{\mathbf{p}(G)}(\#G^{\text{ab/tors}}/lG^{\text{ab/tors}}) - 1$ (for any $l \neq \mathbf{p}(G)$)
$\mathbf{e}(G)$	Inertia degree	$\mathbf{d}(G)/\mathbf{f}(G)$
$\widehat{\mathbf{k}^\times}(G)^{\text{LCFT}}$	$\varprojlim K^\times/K^{\times n}$	G^{ab}
$\widehat{\mathbf{k}^\times}(G)^{\text{LCFT}}$	$\varinjlim \varprojlim K_i^\times/K_i^{\times n}$	$\varinjlim_{G_0 \subset G} \widehat{\mathbf{k}^\times}(G_0)$, this limit varies over open subgroups and the transition maps are given by the transfer maps
$\mu_n(G)$	$\cong \mu_n(\overline{K})$ as G_K -modules	n -torsion of $\widehat{\mathbf{k}^\times}(G)$
$\Lambda(G)$	$\cong \widehat{\mathbb{Z}}(1)$ as G_K -modules	$\varprojlim \mu_n(G)$
$\mathcal{O}^\times(G)$	$\cong \mathcal{O}_K^\times$	$\ker(G \rightarrow \mathbf{G}_{\text{res}}(G))$
$\overline{\mathcal{O}}^\times(G)$	$\cong \mathcal{O}_{\overline{K}}^\times$	$\varinjlim_{G_0 \subset G} \mathcal{O}^\times(G)$
$\mathbf{k}_{\log}^+(G)$	$\cong K_{\log}^+$ as topological abelian groups	$\mathcal{O}^\times(G) \otimes_{\mathbb{Z}} \mathbb{Q}$ (same as monoid perfection)

$\underline{P}(G)$ $\underline{f}(G)$ $\underline{d}(G)$ $\underline{e}(G)$

$$\boxed{\underline{k}^X(G) = G^{ab}}$$

 $M_n(G)$ $M_\infty(G)$ $\Lambda(G)$

$$\underline{\Omega}^X(G) = \underline{I}(G^{ab})$$

$$I(G) = \cap \{N_G : e(M) = \underline{e}(G)\}$$

$[G, G]$
definable

 $\underline{k}_{\log}^+(G)$ $P(G)$ F_r Ω^{X_μ}

$$\{r \in G_m : \forall x \in I \mid \frac{r}{r(x)} - \underline{r}_x \}$$

 $G_{res}(G)$ $G/I(G)$ $\underline{k}^X(G)$ ord $\underline{k}^X(G) \rightarrow \underline{\Gamma}(G) = \frac{\underline{k}(G)}{\underline{\Omega}^X(G)}$ $\underline{\Omega}^\Delta(G)$ $\underline{I}(G)$ $\bar{k}^X(G)$

Galois groups

Interpretation of Torsion

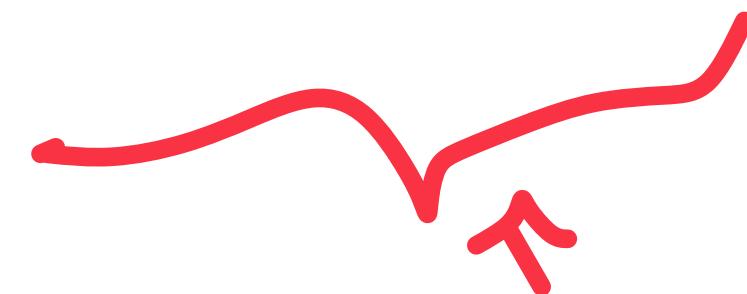
A abelian group

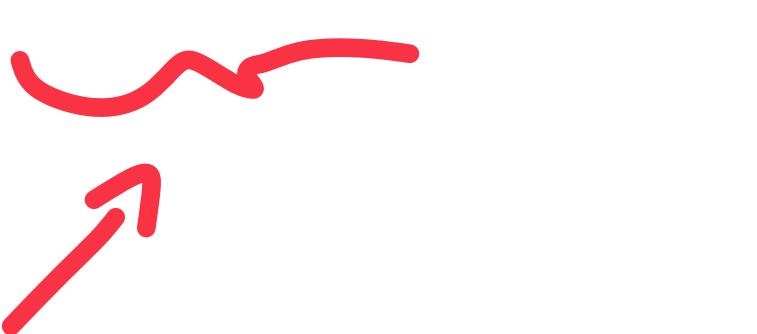
$$a \in A \text{ torsion} \iff \exists n \in \mathbb{N}, \quad na = 0$$

↑
Not in signature

A abelian group

Interpretation of Torsion

$$\exists n \in \mathbb{N}$$


$$n a = 0$$




$$n a = 0$$



$$\underbrace{a + a + \dots + a}_{n - \text{times}} = 0$$

A abelian group

Interpretation of Torsion

$$\exists_{n \in \mathbb{N}} \phi_n(a)$$

$$\begin{aligned} & \text{at } \underbrace{a + a + \dots + a}_{n\text{-times}} = 0 \\ & n a = 0 \end{aligned}$$


$$\exists_{n \in \mathbb{N}} \phi_n(a)$$

$$\phi_1(a) \vee \phi_2(a) \vee \phi_3(a) \vee \dots$$

A abelian group

Interpretation of Torsion

$$\exists_{n \in \mathbb{N}} \quad \exists_{n \in \mathbb{N}}, \phi_n(a)$$

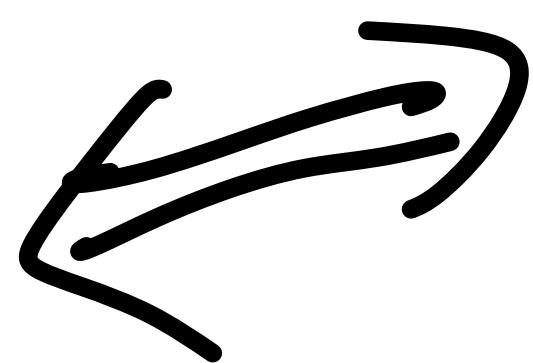
$$\phi_1(a) \vee \phi_2(a) \vee \phi_3(a) \vee \dots$$

$$\begin{aligned} & n a = 0 \\ & \underbrace{a + a + \dots + a}_{n\text{-times}} = 0 \\ & \phi_n(a) \end{aligned}$$

A abelian group

Interpretation of Torsion

$$\exists n \in \mathbb{N} \quad na = 0$$



$$\phi_1(a) \vee \phi_2(a) \vee \phi_3(a) \vee \dots$$

$L_{\omega_1, \omega}$

First Order
Formula!

(Commutator subgroups can be handled similarly)

Interpretation of Torsion

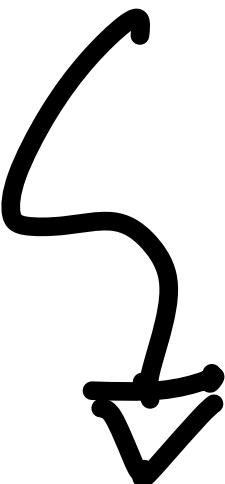
UPSHOT: We will need infinitary logic

First Order Categorical Logic
Model-Theoretical Methods in the Theory
of Topoi and Related Categories

Michael Makkai and Gonzalo E. Reyes

Things Get Worse

G



$$f(G) = \left\{ \begin{array}{l} \text{unique } l \in \mathbb{N} \text{ such that} \\ 1) l \text{ prime} \\ 2) \log_l \#(G^{\text{abtors}}) \geq 2 \end{array} \right\}$$

OK

Things Get Worse

Annotations:

- Red circle around "unique $l \in \mathbb{N}$ such that".
- Red circle around "prime".
- Red circle around "2) \log_l ".
- Green oval around "#(G^{abtors})".
- Green oval around " \log_l ".
- Red circle around " ≥ 2 ".

G

Things Get Worse

$$f(G) = (\ell \text{ prime}) \wedge (\log(\#\text{Gtors}/\ell \text{ Gtors}) \geq 2)$$

$$\underline{f}(G) \approx \log_{f(G)}(1 + \#(\text{Gtors})^{f(G)})$$

$$\underline{d}(G) = \log_{f(G)}\left(\frac{\#\text{Gtors}}{\ell \text{ Gtors}} - 1\right)$$

$$\underline{e}(G) = \underline{d}(G)/\underline{f}(G)$$

Things Get Worse

$$\begin{aligned}
 P(G) &= (\text{l prime}) \wedge \left(\log_{\underline{f}(G)} (\# G^{\text{ab/factors}} / \# G^{\text{ab/tors}}) \leq 2 \right) \\
 \underline{f}(G) &\approx \log_{P(G)} \left(1 + \#(G^{\text{ab/tors}}) f(G) \right) \\
 \underline{d}(G) &= \log_{P(G)} \left(\#(G^{\text{ab/factors}} / \# G^{\text{ab/tors}}) - 1 \right) \\
 e(G) &= \underline{d}(G) / \underline{f}(G)
 \end{aligned}$$

$$\begin{aligned}
 G \rightsquigarrow \underline{\mathcal{I}}(G) &= \bigcup \{ N \triangle G : e(N) = e(G) \} \\
 \underline{\mathcal{G}}_{\text{reg}}(G) &= G \setminus \underline{\mathcal{I}}(G) \\
 \underline{\mathcal{O}}^+(G) &= \varinjlim \underline{\mathcal{O}}^+(G_0)
 \end{aligned}$$

Things Get Worse

$$\underline{I}(G) = \bigcap \{ N \triangleleft G : \underline{\epsilon}(N) = \underline{\epsilon}(G) \}$$

$$G_{\text{res}}(G) = G / \underline{I}(G)$$

$$\bar{\underline{\theta}}^x(G) = \varinjlim \underline{\theta}^x(G_0)$$

$$\begin{aligned} p(G) &= (\text{l prime}) \& \left(\log_l (\# G^{\text{ab/fors}} / l G^{\text{ab/fors}}) \leq 2 \right) \\ f(G) &\approx \log_{p(G)} \left(1 + \#(G^{\text{ab/fors}})^{p(G)} \right) \\ d(G) &= \log_{p(G)} \left(\#(G^{\text{ab/fors}} / l G^{\text{ab/fors}}) - 1 \right) \\ \underline{\epsilon}(G) &= \underline{d}(G) / \underline{f}(G) \end{aligned}$$

Even using infinitary logic it is unclear how to
formalize!

- topology
- direct limits
- weird constants

|Return To Structures|

G ↗

Return To Structures

JOURNAL OF THE
AMERICAN MATHEMATICAL SOCIETY
Volume 9, Number 1, January 1996

ZARISKI GEOMETRIES

EHUD HRUSHOVSKI AND BORIS ZILBER

1. INTRODUCTION

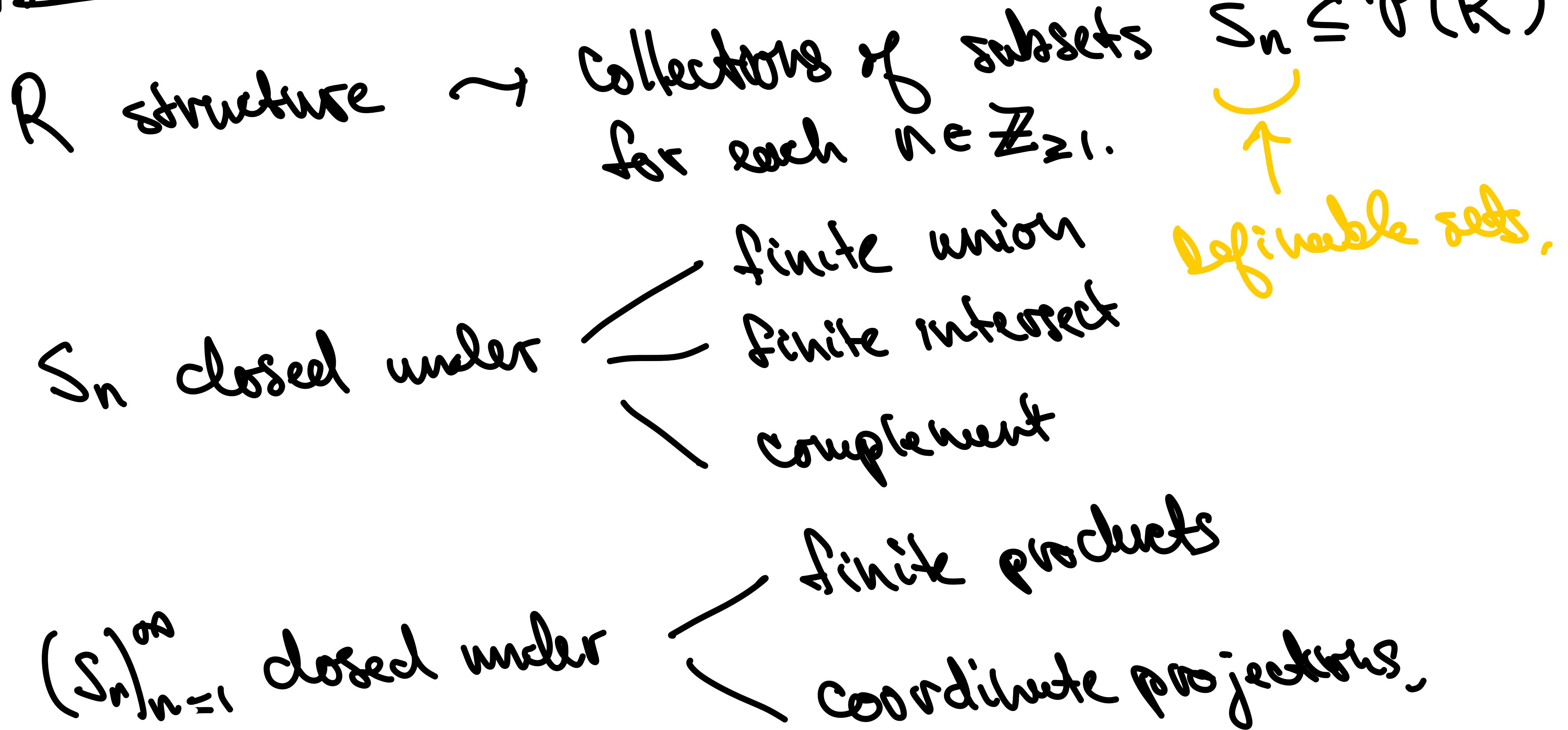
Let k be an algebraically closed field. The set of ordered n -tuples from k is viewed as an n -dimensional space; a subset described by the vanishing of a polynomial, or a family of polynomials, is called an *algebraic set*, or a *Zariski closed set*. Algebraic

LECTURES ON THE AX–SCHANUEL CONJECTURE

BENJAMIN BAKKER AND JACOB TSIMERMAN

ABSTRACT. Functional transcendence results have in the last decade found a number of important applications to the algebraic and arithmetic geometry of varieties X admitting flat or hyperbolic uniformizations: Pila and Zannier’s new proof of the Manin–Mumford conjecture, the proof of the André–Oort conjecture for A_g , and the generic Shafarevich conjecture for hypersurfaces

Return To Structures



Defn An abstract structure is

Return To Structures

collections of subsets $S_n \subseteq \mathcal{P}(R^n)$
for each $n \in \mathbb{Z}_{\geq 1}$. such that :

- 1) S_n closed under
 - finite union
 - finite intersect
 - complement
- 2) $(S_n)_{n=1}^{\infty}$ closed under
 - finite products
 - coordinate projections,

Return To Structures

classical
First order
Structure

equivalent

Defn An abstract structure is
collections of subsets $S_n \subseteq P(\mathbb{R}^n)$
for each $n \in \mathbb{Z}_{\geq 1}$. such that:

- 1) S_n closed under
 - finite union
 - finite intersect
 - complement
- 2) $(S_n)_{n=1}^{\infty}$ closed under
 - finite products
 - coordinate projections,

G

Return To Structures

let G be a set and let Γ
 $\subseteq \text{Perm}(G)$ be a subgroup.

Defn. The Γ -structure is the
Abstract structure (S_n) where
 $S_n \subseteq \mathcal{P}(G^n)$ is the collection of
 Γ -invariant sets.

Defn An abstract structure is
collections of subsets $S_n \subseteq \mathcal{P}(R^n)$
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- 1) S_n closed under
 - finite union
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- 2) $(S_n)_{n=1}^{\infty}$ closed under
 - finite products
 - coordinate projections,

G^Γ

let G be a set and let $\Gamma \subseteq \text{Perm}(G)$ be a subgroup.

Defn. The Γ -structure is the Abstract structure (S_n) where $S_n \subseteq \mathcal{P}(G^n)$ is the collection of Γ -invariant sets.

Return To Structures

Application: Take $G = \text{Gal}(\mathbb{R}/\mathbb{K})$, $[\mathbb{K}:\mathbb{Q}_p] < \infty$
Take $\Gamma = \text{Aut}(G)$, (topological)

Defn An abstract structure is collections of subsets $S_n \subseteq \mathcal{P}(R^n)$ for each $n \in \mathbb{Z}_{\geq 1}$. such that:
1) S_n closed under finite union, finite intersect, complement
2) $(S_n)_{n=1}^{\infty}$ closed under finite products, coordinate projections,

Application: Take $G = \text{Gal}(\mathbb{R}/K)$, $[K:\mathbb{Q}_p] < \infty$
Take $\Gamma = \text{Aut}(G)$, (topological)

G^{Γ}

G_{group}

(usual group language) \subseteq (infinitary group language) \subseteq (language of Γ -invariant sets)

$$\begin{aligned} p(G) \\ f(G) \\ d(G) \\ e(G) \end{aligned}$$

$$\widehat{\underline{K}}^x(G) = G^{ab}$$

$$\begin{aligned} M_n(G) & \quad M_\infty(G) \\ \Lambda(G) & \quad \lim_{n \rightarrow \infty} M_n(G) \end{aligned}$$

$$\underline{O}^X(G) = \underline{I}(G)$$

$$\cap \{NAG : e(M) = \underline{e}(G)\}$$

$[G, G]$
definable

$$\begin{aligned} \underline{K}_{\log}^+(G) \\ P(G) \end{aligned}$$

$$P(G)$$

$$Fr$$

$$\left\{ \sigma \in G_{\text{res}} : \forall x \in I \mid \frac{\sigma(x)}{r(x)} = \frac{p(x)}{f(x)} \right\}$$

G_{group}

$$\begin{aligned} G_{\text{res}}(G) & \quad G/I(G) \end{aligned}$$

$$\underline{K}^x(G)$$

$$\underline{I}^x(G)$$

$$\underline{O}^\Delta(G)$$

$$\underline{I}(G)$$

$$\widehat{\underline{K}}^x(G)$$

$$\underline{I}^x(G) \rightarrow \underline{P}(G) = \frac{\underline{K}(G)}{\underline{O}^x(G)}$$

G^{\square}

G_{group}

$p(G)$
 $f(G)$
 $d(G)$
 $e(G)$

auxillary

$[G, G]$
definable

$\widehat{\exp}(G) = G^{ab}$

$\mu_n(G)$
 $\mu_\infty(G)$

$\lim \mu_n(G)$

$\Lambda(G)$

$\widehat{o}(G) = \widehat{I}(G^{ab})$

$I(G) = \bigcap_{n=1}^{\infty} N(G : e^n(G)) = \underline{e}(G)$

$k^t_{\mathrm{dmg}}(G)$
 $P(G)$
 Fr
 $G_{\mathrm{reg}}(G)$
 $k^x[G]$
 $\underline{\mathrm{ord}}$
 $O^\Delta(G)$
 $\pi(G)$

$\alpha x \mu$

$\{ \text{sets } G_x : \forall x \in I / e \mid f(x) = e^x \}$ group

$G/I(G)$

group

Group A

$\widehat{f}_k^*(g)$

$$\widehat{\text{rk}}^X(G) = G^{ab}$$

$$\lim_{n \rightarrow \infty} \mu_n(G) = \mu(G)$$

$$\underline{\alpha}(G) = \underline{I}(G)$$

$$I(G) = \bigcap \{ N \in \mathcal{G} : e(N) = \emptyset \}$$

$$I(G) = \cap \{NAG : e(A) = \underline{\epsilon}(G)\}$$

Γ -invariant collection

Definition: A Γ -invariant collection is some
such that \mathcal{A} itself is invariant under Γ .

$$\mathcal{A} \subseteq \mathcal{P}(G^n)$$

Lemma: If \mathcal{A} is Γ -invariant then

- 1) $\bigcap_{A \in \mathcal{A}} A \in G^\Gamma$ -definable
- 2) $\bigcup_{A \in \mathcal{A}} A \in G^\Gamma$ -definable.

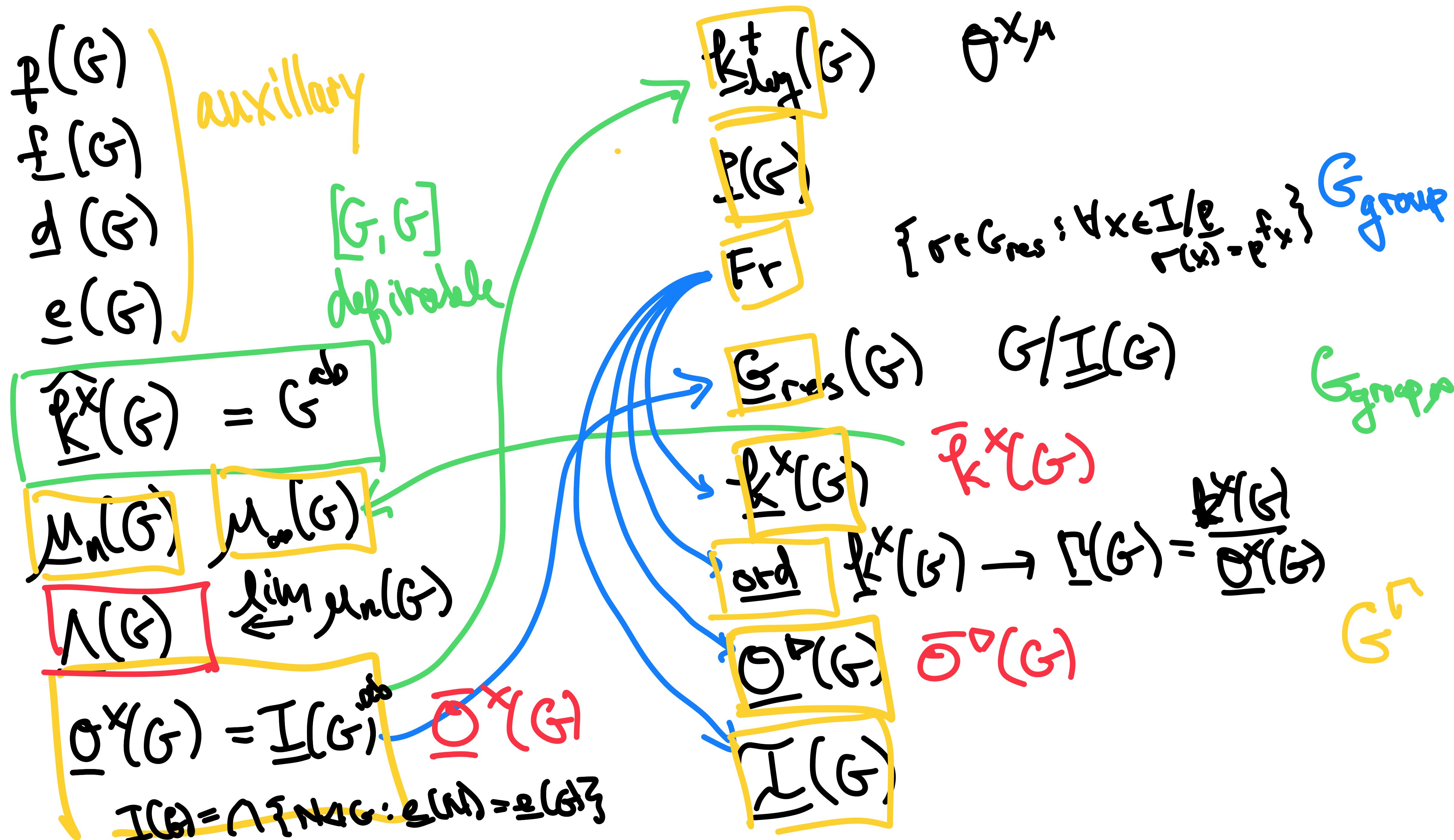


$$I(G) = \bigcap \{NAG : e(N) = \underline{\omega}(G)\}$$

Definition: A Γ -invariant collection is some

$$\mathcal{A} \subseteq \mathcal{P}(G^n)$$

such that \mathcal{A} itself is invariant under Γ .



G_∞) New "fragment"
for dealing
with frags

$$\underline{\Omega}^X(G) = \varinjlim \underline{\Omega}^X(G_\alpha)$$

$$\left(\coprod_{G_\alpha \subseteq G} \underline{\Omega}^X(G_\alpha) \right)$$

an infinite
disjoint union!

WANT: $\coprod_{G_\alpha \subseteq G} \underline{\Omega}^X(G_\alpha) \xrightarrow{I} \varinjlim \underline{\Omega}^X(G_\alpha)$

\mathcal{M} a structure.

Definable sets: $S_1 \subseteq \mathcal{P}(\mathcal{M})$, $S_2 \subseteq \mathcal{P}(\mathcal{M}^2)$, $S_3 \subseteq \mathcal{P}(\mathcal{M}^3)$, ...
(classical)

In powers of \mathcal{M} .

Definable sets:
(infinitary)

$S_\alpha \subseteq \mathcal{P}(\mathcal{M}^\alpha)$

$\gamma < \kappa$
ordinal

G_∞

\mathcal{M} a structure.

Definable sets:

(infinitary)

$$S_\alpha \subseteq \wp(\mathcal{M}^\alpha)$$

$$\alpha < \kappa$$

ordinal

Definable sets:

(what we need)

$$S_\kappa \subseteq \wp(F_\kappa(\mathcal{M}))$$

$v \in V$ a collection

where

contains $\prod_{\gamma < \kappa} M$

Defn We define the collection \mathcal{G}^{Γ} as the set of
 $(S_v)_{v \in V}$, where $S_v \subseteq P(F_v(G))$ or the collection
of Γ -invariant sets.

G_{group}

Classical group language

G_{group}

infinitary group language

G^{Γ}

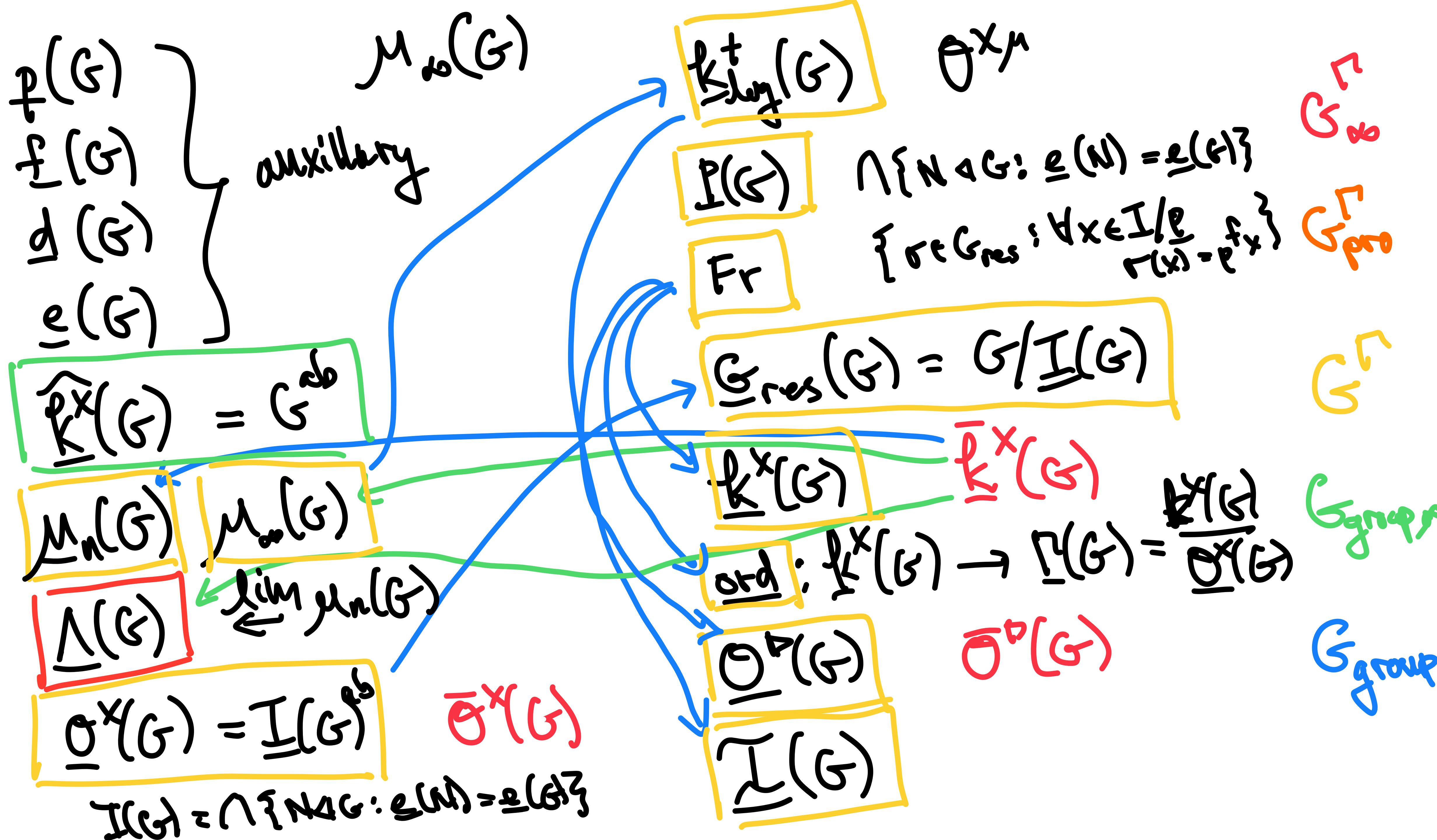
Γ -invariant sets $\Gamma = \text{Aut}(G)$

G_{per}

PO- Γ -invariant sets

G_{∞}^{Γ}

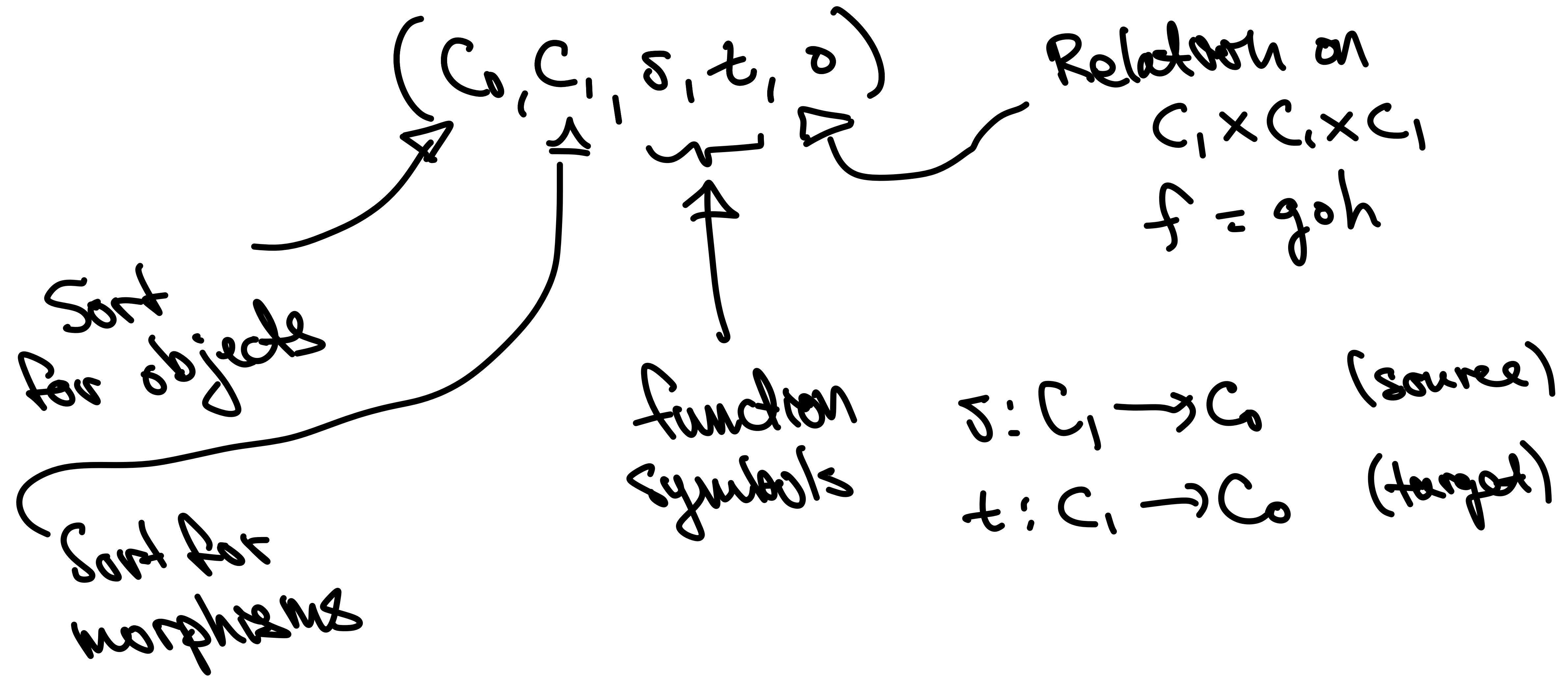
Γ -invariant sets w/ disjoint unions



THANK You!

BONUS: Species &
Mutations.

Signature for Categories:



Interpreted Categories

$$(C_0, C_1, \sigma, \tau, \delta) = C$$

$$I_0: X_0 \rightarrow C_0, \quad I_1: X_1 \rightarrow C_1 \quad \left. \begin{array}{l} \text{Interpretation} \\ \text{of } C \text{ in } M \end{array} \right\}$$

\cap

M^{n_0}

\cap

M^{n_1}

(Graph of Source/Target) $\Gamma_s, \Gamma_t \subseteq C_1 \times C_0$

(Graph of Composition) $\Gamma_o \subseteq C_1 \times C_0 \times C_1$

$$I_0: X_0 \rightarrow C_0$$
$$M^{n_0}$$

$$I_1: X_1 \rightarrow C_1$$
$$M^{n_1}$$

$$\Gamma_x \cap \Gamma_s \subseteq C_1 \times C_0$$

$$X_0 = \{ \vec{m} \in M^{n_0} : \Psi_0(\vec{m}) \}$$

$$X_1 = \{ \vec{m} \in M^{n_1} : \Psi_1(\vec{m}) \}$$

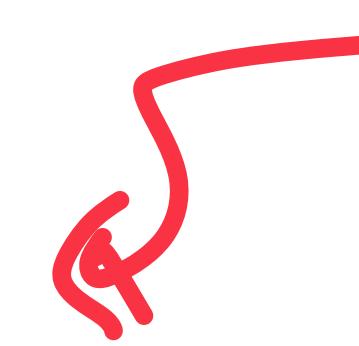
$$\Gamma_0 \subseteq C_1 \times C_1 \times C_1$$

$$\Gamma_s \subseteq C_1 \times C_0 \rightsquigarrow I^{-1}(\Gamma_s) = X_s \subseteq X_1 \times X_0$$

$$\Gamma_x \subseteq C_1 \times C_0$$

given by formula Φ_s

$$\Gamma_0 \subseteq C_1 \times C_1 \times C_1 \rightsquigarrow I^+(\Gamma_0) \subseteq X_1 \times X_1 \times X_1$$

$f \circ g = h$  unravel this definition

$$X_0 = \{ \vec{m} \in M^{n_0} : \Phi_0(\vec{m}) \} \quad X_1 = \{ \vec{m} \in M^{n_1} : \Phi_1(\vec{m}) \}$$

Definition 3.1.

(i) A 0-species \mathfrak{S}_0 is a collection of conditions given by a *set-theoretic formula*

$$\Phi_0(\mathfrak{E})$$

involving an ordered collection $\mathfrak{E} = (\mathfrak{E}_1, \dots, \mathfrak{E}_{n_0})$ of sets $\mathfrak{E}_1, \dots, \mathfrak{E}_{n_0}$ [which we think of as “indeterminates”], for some integer $n_0 \geq 1$; in this situation, we shall refer to \mathfrak{E} as a *collection of species-data* for \mathfrak{S}_0 . If \mathfrak{S}_0 is a 0-species given by a set-theoretic formula $\Phi_0(\mathfrak{E})$, then a 0-specimen of \mathfrak{S}_0 is a *specific* ordered collection of n_0 sets $E = (E_1, \dots, E_{n_0})$ in some *specific* ZFC-model that satisfies $\Phi_0(E)$. If E is a 0-specimen of a 0-species \mathfrak{S}_0 , then we shall write $E \in \mathfrak{S}_0$. If, moreover, it holds, in any ZFC-model, that the 0-specimens of \mathfrak{S}_0 form a *set*, then we shall refer to \mathfrak{S}_0 as *0-small*.

(ii) Let \mathfrak{S}_0 be a 0-species. Then a 1-species \mathfrak{S}_1 *acting on* \mathfrak{S}_0 is a collection of

(iii) A *species* \mathfrak{S} is defined to be a pair consisting of a 0-species \mathfrak{S}_0 and a 1-species \mathfrak{S}_1 acting on \mathfrak{S}_0 . Fix a species $\mathfrak{S} = (\mathfrak{S}_0, \mathfrak{S}_1)$. Let $i \in \{0, 1\}$. Then we shall refer to an i -specimen of \mathfrak{S}_i as an i -specimen of \mathfrak{S} . We shall refer to a 0-specimen (respectively, 1-specimen) of \mathfrak{S} as a *species-object* (respectively, a *species-morphism*) of \mathfrak{S} . We shall say that \mathfrak{S} is *i-small* if \mathfrak{S}_i is *i-small*. We shall refer to a species-morphism $F : E \rightarrow E'$ as a *species-isomorphism* if there exists a species-morphism $F' : E' \rightarrow E$ such that the composites $F \circ F'$, $F' \circ F$ are *identity* species-morphisms; in this situation, we shall say that E , E' are *species-isomorphic*. [Thus, one verifies immediately that *composites of species-isomorphisms* are species-isomorphisms.] We shall refer to a species-isomorphism whose domain and codomain are equal as a *species-automorphism*. We shall refer to as *model-free* [cf. Remark 3.1.1 below] an i -specimen of \mathfrak{S} equipped with a description via a *set-theoretic formula* that is “*independent* of the ZFC-model in which it is given” in the sense that for any pair of universes V_1, V_2 of some ZFC-model such that $V_1 \in V_2$, the set-theoretic formula determines the *same* i -specimen of \mathfrak{S} , whether interpreted relative to the ZFC-model determined by V_1 or the ZFC-model determined by V_2 .

$\mathfrak{F} : \mathfrak{E} \rightarrow \mathfrak{E}'$. If, in some ZFC-model, $E, E' \in \mathfrak{S}_0$, and F is a *specific* ordered collection of n_1 sets that satisfies the condition $\Phi_1(E, E', F)$, then we shall refer to the data (E, E', F) as a *1-specimen* of \mathfrak{S}_1 and write $(E, E', F) \in \mathfrak{S}_1$; alternatively, we shall denote a 1-specimen (E, E', F) via the notation $F : E \rightarrow E'$ and refer to E (respectively, E') as the *domain* (respectively, *codomain*) of $F : E \rightarrow E'$.

(b) $\Phi_{1 \circ 1}$ is a set-theoretic formula

$$\Phi_{1 \circ 1}(\mathfrak{E}, \mathfrak{E}', \mathfrak{E}'', \mathfrak{F}, \mathfrak{F}', \mathfrak{F}'')$$

involving three collections of species-data $\mathfrak{F} : \mathfrak{E} \rightarrow \mathfrak{E}'$, $\mathfrak{F}' : \mathfrak{E}' \rightarrow \mathfrak{E}''$, $\mathfrak{F}'' : \mathfrak{E} \rightarrow \mathfrak{E}''$ for \mathfrak{S}_1 [i.e., the conditions $\Phi_0(\mathfrak{E})$; $\Phi_0(\mathfrak{E}')$; $\Phi_0(\mathfrak{E}'')$; $\Phi_1(\mathfrak{E}, \mathfrak{E}', \mathfrak{F})$; $\Phi_1(\mathfrak{E}', \mathfrak{E}'', \mathfrak{F}')$; $\Phi_1(\mathfrak{E}, \mathfrak{E}'', \mathfrak{F}'')$ hold]; in this situation, we shall refer to \mathfrak{F}'' as a *composite of \mathfrak{F} with \mathfrak{F}'* and write $\mathfrak{F}'' = \mathfrak{F}' \circ \mathfrak{F}$ [which is, *a priori*, an abuse of notation, since there may exist *many* composites of \mathfrak{F} with \mathfrak{F}' — cf. (c) below]; we shall use similar terminology and notation for 1-specimens in specific ZFC-models.

- c) Given a pair of 1-specimens $F : E \rightarrow E'$, $F' : E' \rightarrow E''$ of \mathfrak{S}_1 in some ZFC-model, there *exists a unique composite* $F'' : E \rightarrow E''$ of F with F' in the given ZFC-model.
- d) Composition of 1-specimens $F : E \rightarrow E'$, $F' : E' \rightarrow E''$, $F'' : E'' \rightarrow E'''$ of \mathfrak{S}_1 in a ZFC-model is *associative*.
- e) For any 0-specimen E of \mathfrak{S}_0 in a ZFC-model, there exists a [necessarily unique] 1-specimen $F : E \rightarrow E$ of \mathfrak{S}_1 [in the given ZFC-model] — which we shall refer to as the *identity 1-specimen* id_E of E — such that for any 1-specimens $F' : E' \rightarrow E$, $F'' : E \rightarrow E''$ of \mathfrak{S}_1 [in the given ZFC-model] we have $F \circ F' = F'$, $F'' \circ F = F''$.

Definition 3.3. Let $\mathfrak{S} = (\mathfrak{S}_0, \mathfrak{S}_1)$; $\underline{\mathfrak{S}} = (\underline{\mathfrak{S}}_0, \underline{\mathfrak{S}}_1)$ be species.

(i) A mutation $\mathfrak{M} : \mathfrak{S} \rightsquigarrow \underline{\mathfrak{S}}$ is defined to be a collection of set-theoretic formulas Ψ_0, Ψ_1 satisfying the following properties:

(a) Ψ_0 is a set-theoretic formula

$$\Psi_0(\mathfrak{E}, \underline{\mathfrak{E}})$$

involving a collection of species-data \mathfrak{E} for \mathfrak{S}_0 and a collection of species-data $\underline{\mathfrak{E}}$ for $\underline{\mathfrak{S}}_0$; in this situation, we shall write $\mathfrak{M}(\mathfrak{E})$ for $\underline{\mathfrak{E}}$. Moreover, if, in some ZFC-model, $E \in \mathfrak{S}_0$, then we require that there exist a unique $\underline{E} \in \underline{\mathfrak{S}}_0$ such that $\Psi_0(E, \underline{E})$ holds; in this situation, we shall write $\mathfrak{M}(E)$ for E .

(b) Ψ_1 is a set-theoretic formula

$$\Psi_1(\mathfrak{E}, \mathfrak{E}', \mathfrak{F}, \underline{\mathfrak{F}})$$

(iv) Let $\vec{\Gamma}$ be an oriented graph, i.e., a graph Γ , which we shall refer to as the underlying graph of $\vec{\Gamma}$, equipped with the additional data of a total ordering, for each edge e of Γ , on the set [of cardinality 2] of branches of e [cf., e.g., [AbsTopIII], §0]. Then we define a mutation-history $\mathfrak{H} = (\vec{\Gamma}, \mathfrak{S}^*, \mathfrak{M}^*)$ [indexed by $\vec{\Gamma}$] to be a collection of data as follows:

- (a) for each vertex v of $\vec{\Gamma}$, a species \mathfrak{S}^v ;
- (b) for each edge e of $\vec{\Gamma}$, running from a vertex v_1 to a vertex v_2 , a mutation $\mathfrak{M}^e : \mathfrak{S}^{v_1} \rightsquigarrow \mathfrak{S}^{v_2}$.

In this situation, we shall refer to the vertices, edges, and branches of $\vec{\Gamma}$ as vertices, edges, and branches of \mathfrak{H} . Thus, the notion of a “mutation-history” may be thought of as a species-theoretic version of the notion of a “diagram of categories” given in [AbsTopIII], Definition 3.5, (i).

(ii) Let $\mathfrak{M}, \mathfrak{M}' : \mathfrak{S} \rightsquigarrow \underline{\mathfrak{S}}$ be mutations. Then a morphism of mutations $\mathfrak{Z} : \mathfrak{M} \rightarrow \mathfrak{M}'$ is defined to be a set-theoretic formula Ξ satisfying the following properties:

(a) Ξ is a set-theoretic formula

$$\Xi(\mathfrak{E}, \underline{\mathfrak{F}})$$

involving a collection of species-data \mathfrak{E} for \mathfrak{S}_0 and a collection of species-data $\underline{\mathfrak{F}} : \mathfrak{M}(\mathfrak{E}) \rightarrow \mathfrak{M}'(\mathfrak{E})$ for \mathfrak{S}_1 ; in this situation, we shall write $\mathfrak{Z}(\mathfrak{E})$ for $\underline{\mathfrak{F}}$. Moreover, if, in some ZFC-model, $E \in \mathfrak{S}_0$, then we require that there exist a unique $\underline{F} \in \underline{\mathfrak{S}}_1$ such that $\Xi(E, \underline{F})$ holds; in this situation, we shall write $\mathfrak{Z}(E)$ for \underline{F} .

(b) Suppose, in some ZFC-model, that $F : E_1 \rightarrow E_2$ is a species-morphism of \mathfrak{S} . Then one has an equality of composite species-morphisms $\mathfrak{M}'(F) \circ \mathfrak{Z}(E_1) = \mathfrak{Z}(E_2) \circ \mathfrak{M}(F) : \mathfrak{M}(E_1) \rightarrow \mathfrak{M}'(E_2)$. In particular, if one fixes a ZFC-model, then a morphism of mutations $\mathfrak{M} \rightarrow \mathfrak{M}'$ determines a natural transformation between the functors determined by $\mathfrak{M}, \mathfrak{M}'$ in the ZFC-model — cf. (i).

INTER-UNIVERSAL TEICHMÜLLER THEORY IV

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involving a collection of species-data $\mathfrak{F} : \mathfrak{E} \rightarrow \mathfrak{E}'$ for \mathfrak{S}_1 and a collection of species-data $\underline{\mathfrak{F}} : \underline{\mathfrak{E}} \rightarrow \underline{\mathfrak{E}'}$ for $\underline{\mathfrak{S}}_1$, where $\underline{\mathfrak{E}} = \mathfrak{M}(\mathfrak{E})$, $\underline{\mathfrak{E}'} = \mathfrak{M}'(\mathfrak{E}')$; in this situation, we shall write $\mathfrak{M}(\mathfrak{F})$ for $\underline{\mathfrak{F}}$. Moreover, if, in some ZFC-model, $(F : E \rightarrow E') \in \mathfrak{S}_1$, then we require that there exist a unique $(\underline{F} : \underline{E} \rightarrow \underline{E}') \in \underline{\mathfrak{S}}_1$ such that $\Psi_0(E, E', F, \underline{F})$ holds; in this situation, we shall write $\mathfrak{M}(F)$ for \underline{F} . Finally, we require that the assignment $F \mapsto \mathfrak{M}(F)$ be compatible with composites and map identity species-morphisms of \mathfrak{S} to identity species-morphisms of $\underline{\mathfrak{S}}$. In particular, if one fixes a ZFC-model, then \mathfrak{M} determines a functor from the category determined by \mathfrak{S} in the given ZFC-model to the category determined by $\underline{\mathfrak{S}}$ in the given ZFC-model.

BONUS: Infinitary
languages.

The language $L_{k,\lambda}(z)$

for ordinals

- FORMULAS:
- $\exists(x :: i < \gamma) , \gamma \leq k$
 - $\bigwedge e_\alpha(\vec{x}) , \beta \leq \gamma$ [same for \vee]
 $d \leq \beta$
 - Formulas have less than k many free variables (projections of k -many vars)
 - Subsets of M^d for $d \leq k$
 - Intersection of λ many sets

BONKS: LasCar ^Types

- Types) . Let A be a structure .
- The type of $\vec{a} \in A^n$ is the collection $\phi \in L_A$ with one free variable such that $\phi(\vec{a})$,
 - $S_n(A) = (\text{Stone-space of types})$

Example, F field
 \bar{F} alg closed field.

$A = (\bar{F}, +, \cdot, 0, 1, (a)_{a \in F})$

constant symbol
 for every $a \in F$

type $b \in \bar{F}$
 ↑ determined by
 minimal poly
 $m_b(x) \in F[x]$.
 points of same type are
 Galois conjugates.

Types

Example, F field
 \bar{F} alg closed field.

$A = (\bar{F}, +, \cdot, 0, 1, (a)_{a \in F})$
constant symbol!
for every $a \in F$

$\cdot S_n(A) = (\text{stone-space of types}) = \text{Spec}(F[x_1, \dots, x_n])$

B structure.

[Types]

A substructure of B .

$$\text{Aut}(B/A) = \{\sigma \in \text{Aut}(B) : \forall a \in A, \sigma(a) = a\},$$

Defn. The Loskar types (of $b \in B''$) are the set

$$B''/\text{Aut}(B/A)$$

Example: $G = \text{Gal}(\bar{K}/K)$, $(K:\mathbb{Q}_p) \subset_B$, $H = \bar{\theta}^\sigma, \bar{\theta}^x, \bar{\theta}^{xm}$

$$\text{Aut}(G, H) \rightarrow \text{Aut}(G), \quad B = (G, H), \quad A = G$$

Loskar types then are pairs (g, m) where m is up to indeterminacy.

BONUS: QUESTIONS

Question: What is the interpretation-theoretic content of the mathematical theory of Frobenioids?

THE GEOMETRY OF FROBENIoids I: THE GENERAL THEORY

SHINICHI MOCHIZUKI

June 2008

ABSTRACT. We develop the theory of *Frobenioids*, which may be regarded as a category-theoretic abstraction of the theory of divisors and line bundles on models of finite separable extensions of a given function field or number field. This sort of abstraction is analogous to the role of *Galois categories* in Galois theory or *monoids* in the geometry of log schemes. This abstract category-theoretic framework preserves many of the important features of the classical theory of divisors and line bundles on models of finite separable extensions of a function field or number field such as the *global degree* of an arithmetic line bundle over a number field, but also exhibits interesting new phenomena, such as a “*Frobenius endomorphism*” of the Frobenioid associated to a number field.

Introduction

- §0. Notations and Conventions
- §1. Definitions and First Properties
- §2. Frobenius Functors
- §3. Category-theoreticity of the Base and Frobenius Degree
- §4. Category-theoreticity of the Divisor Monoid
- §5. Model Frobenioids
- §6. Some Motivating Examples

Appendix: Slim Exponentiation

Index

Question: What do Leibniz types have to do with indeterminacies?

Question: Can you find a classical first
order example of non-anabelian
transport?

Question:

Why is it the case that when the absolute Grothendieck conjecture holds for π_X then π_X interprets a field? Is this necessary?

BONUS: Complete
theories.

• Structures are theories over the empty set.

$\phi \models M$ for all $M \in \text{Str}_\sigma$

$\Rightarrow \text{Def}(T) = \text{Def}(\phi) \Rightarrow$ functions $\text{Def}_\sigma(\phi) \rightarrow \text{sets}$
are structures.

• $T = \text{Th}(M) = \left(\begin{array}{l} \text{All Formulas} \\ \text{That Are True} \\ \text{For } M. \end{array} \right)$

Bonus: Functional Algorithms

INTER-UNIVERSAL TEICHMÜLLER THEORY II: HODGE-ARAKELOV-THEORETIC EVALUATION

SHINICHI MOCHIZUKI

December 2020

ABSTRACT. In the present paper, which is the second in a series of four papers, we study the **Kummer theory** surrounding the Hodge-Arakelov-theoretic evaluation — i.e., evaluation in the style of the **scheme-theoretic Hodge-Arakelov theory** established by the author in previous papers — of the [reciprocal of the l -th root of the] **theta function** at **l -torsion points** [strictly speaking, shifted by a suitable 2-torsion point], for $l \geq 5$ a prime number. In the first paper of the series, we studied “*miniature models of conventional scheme theory*”, which we referred to as $\Theta^{\pm\text{ell}}\text{NF-Hodge theaters}$, that were associated to certain data, called *initial Θ -data*, that includes an *elliptic curve* E_F over a *number field* F , together with a *prime number* $l \geq 5$. The underlying Θ -Hodge theaters of these $\Theta^{\pm\text{ell}}\text{NF-Hodge theaters}$ were *glued* to one another by means of “ Θ -links”, that identify the [reciprocal of the l -th root of the] **theta function** at primes of bad reduction of E_F in one $\Theta^{\pm\text{ell}}\text{NF-Hodge theater}$ with [2 l -th roots of] the *q-parameter* at primes of bad reduction of E_F in another $\Theta^{\pm\text{ell}}\text{NF-Hodge theater}$. The theory developed in the present paper allows one to construct certain new versions of this “ Θ -link”. One such new version is the $\Theta_{\text{gau}}^{\times\mu}$ -

Example 1.7. Radial and Coric Data I: Generalities.

(i) In the following discussion, we would like to consider a certain “*type of mathematical data*”, which we shall refer to as **radial data**. This notion of a “type of mathematical data” may be *formalized* — cf. [IUTchIV], §3, for more details. From the point of view of the present discussion, one may think of a “type of mathematical data” as the input or output data of a “**functorial algorithm**” [cf. the discussion of [IUTchI], Remark 3.2.1]. At a more concrete level, we shall assume that this “type of mathematical data” gives rise to a *category*

$$\mathcal{R}$$

— i.e., each of whose *objects* is a specific collection of radial data, and each of whose *morphisms* is an isomorphism. In the following discussion, we shall also consider another “type of mathematical data”, which we shall refer to as **coric data**. Write

$$\mathcal{C}$$

for the category obtained by considering specific collections of coric data and isomorphisms of collections of coric data. In addition, we shall assume that we are given a *functorial algorithm* — which we shall refer to as **radial** — whose *input data* consists of a collection of radial data, and whose *output data* consists of a collection of coric data. Thus, this functorial algorithm gives rise to a *functor* $\Phi : \mathcal{R} \rightarrow \mathcal{C}$. In the following discussion, we shall assume that this functor is *essentially surjective*. We shall refer to the category \mathcal{R} and the functor Φ as *radial* and to the category \mathcal{C} as *coric*. Finally, if I is some *nonempty index set*, then we shall often consider collections

$$\{\Phi_i : \mathcal{R}_i \rightarrow \mathcal{C}\}_{i \in I}$$

of copies of Φ and \mathcal{R} , such that the various copies of Φ have the *same codomain* \mathcal{C} — cf. Fig. 1.1 below. Thus, one may think of each \mathcal{R}_i as the category of radial data *equipped with a label* $i \in I$, and isomorphisms of such data.

(ii) We shall refer to a triple $(\mathcal{R}, \mathcal{C}, \Phi : \mathcal{R} \rightarrow \mathcal{C})$ [or to the triple consisting of

BONUS : Pre-topics Completion

- Finite limits
 - Stable finite gaps
 - Stable images
 - stable division,
- proto pos



* $T = Th(M)$, $\text{Def}(T) \approx \text{Def } M$

* $\text{Def}(T) = \text{Def}(T^{\text{eq}})$

\downarrow

Skolem Construction.

topos theory

- Abstract \mathcal{B} -interpretations, give an equal valence of categories between pre-deps completions.
- $\text{Def}(T) \neq \text{Def}(T^g)$ in general.

$$I_1: X_1 \rightarrow N$$

\cap
 M_{α_1}

$$I_2: X_2 \rightarrow N$$

\cap
 M_{α_2}

$$I_1 \sim I_2 \iff Eq(I_1, I_2) = \{ (\vec{m}_1, \vec{m}_2) \in X_1 \times X_2 : I(\vec{m}_1) \leq I(\vec{m}_2) \}$$

definable.

A Bi-interpretation is I_1, I_2 such that

$I_1, I_2 \sim id$

$I_2, I_1 \sim id.$
