

# FIT9133 Semester 2 2019

## Programming Foundations in Python

### Week 12:

### Divide-and-Conquer and Recursion

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- Module 6 is aimed to introduce you with:
  - Concepts of **divide-and-conquer**
  - Definitions of **recursion**
    - Base case
    - Recursive case
    - Convergence
  - **Iterative** solutions versus **recursive** solutions
  - Recursive sorting algorithms:
    - **Merge Sort**
    - **Quick Sort**

- There are no interviews for this week.
- Please check the consultation table on Moodle
  - Some of our booked consultation rooms were taken by other units during swotvac.
- SETU is open, please fill it in.
  - You can win gift cards by submitting the SETU.

- Upon completing this module, you should be able to:
  - Recognise computational problems that are solvable using the divide-and-conquer or recursive approach
  - Implement a recursive solution for a specific problem
  - Contrast the two advanced sorting algorithms that adopt the recursive approach (Merge Sort and Quick Sort)



# Concepts of Divide-and-Conquer

- Divide-and-Conquer:
  - Solving a complex problem by **breaking it into smaller manageable sub-problems**
  - Sub-problems can then be solved in a similar way (with the same solution)
  - **Sub-solutions are then combined to produce the final solution for the original problem**
- Basic example: **Binary Search**
  - “Repeatedly” divides the (sorted) list into two sublists until the target item is found

# Binary Search: “Iterative” Approach

```
def binary_search(the_list, target_item):  
    low = 0  
    high = len(the_list)-1  
  
    # repeatedly divide the list into two halves  
    # as long as the target item is not found  
    while low <= high:  
  
        # find the mid position  
        mid = (low + high) // 2  
  
        if the_list[mid] == target_item:  
            return True  
        elif target_item < the_list[mid]:  
            high = mid - 1 # search lower half  
        else:  
            low = mid + 1  # search upper half  
  
    # the list cannot be further divided  
    # the target is not found  
    return False
```

# Binary Search: “Recursive” Approach

```
def rec_binary_search(the_list, target_item):
    # repeatedly divide the list into two halves
    # as long as the target item is not found

    # the list cannot be further divided i.e. item is not found
    if len(the_list) == 0:
        return False
    else:
        # find the mid position
        mid = len(the_list) // 2

        # check if target item is equal to middle item
        if the_list[mid] == target_item:
            return True

        # check if target item is less than middle item
        # search lower half
        elif target_item < the_list[mid]:
            smaller_list = the_list[:mid]
            return rec_binary_search(smaller_list, target_item)

        # check if target item is greater than middle item
        # search upper half
        else:
            smaller_list = the_list[mid+1:]
            return rec_binary_search(smaller_list, target_item)
```

**Recursive Function:**  
A function that calls itself repeatedly.



# Concepts of Recursion

- Recursion:

- A **divide-and-conquer** approach for solving computational problems
- Each problem is “recursively” **decomposed into sub-problems** (which have the same properties the original problem but smaller in size)
- When the sub-problems have reached **the simplest form**, i.e. **a known solution can be defined**
- The **known solutions of these sub-problems are then recomposed** together to produce the solution of the original problem

**Cannot be  
further  
decomposed**

Solution is  
known

- Three key requirements:
  - **Base case**: The recursive function must have a base case (i.e. the simplest form)
  - **Convergence**: The recursive function must be able to decompose the original problem into sub-problems; and must be converging towards the base case
  - **Recursive case**: The recursive function must call itself recursively to solve the sub-problems

## Example: Recursive Addition

- How could we solve an *addition* problem recursively?

```
5 + 1 =  
(4 + 1) + 1 =  
((3 + 1) + 1) + 1 =  
(((2 + 1) + 1) + 1) + 1 =  
((((1 + 1) + 1) + 1) + 1) + 1 = 6
```

```
def recursive_addition(a, b):  
    if a == 0:  
        return b  
    return recursive_addition(a-1, b) + b
```

Partial solution

```
def recursive_addition(a, b):  
    if a == 0:  
        return b  
    return recursive_addition(a-1, b+1)
```

Complete solution

# Iterative Solution vs Recursive Solution

- How could we solve a **factorial** problem (**n!**) *iteratively*?
- E.g.  $5! = 5 * 4 * 3 * 2 * 1$

```
def iterative_factorial(n):  
    factorial = 1  
  
    while n > 0:  
        factorial *= n  
        n -= 1  
  
    return factorial
```

# Iterative Solution vs Recursive Solution

- How could we solve a **factorial** problem (**n!**) *recursively*?
- E.g.  $5! = 5 * 4!$

$$5! = 5 * (4 * 3!)$$

$$5! = 5 * (4 * (3 * 2!))$$

$$5! = 5 * (4 * (3 * (2 * 1!)))$$

$$5! = 5 * (4 * (3 * (2 * (1))))$$

```
def recursive_factorial(n):
```

```
    if n == 1:  
        return 1
```

```
    return n * recursive_factorial(n-1)
```

base case

convergence

recursive case

# Review Questions: Part 1

What does the given recursive function do? (Assume that a and b are positive integers.)

```
def mystery_func1(a, b):  
    if a == 0:  
        return 0  
    return b + mystery_func1(a-1, b)
```

- A. Addition
- B. Subtraction
- C. Multiplication
- D. Division
- E. None of the above



What does the given recursive function do? (Assume that n is an positive integer.)

```
def mystery_func2(n):  
    if n == 1:  
        return 1  
    return n + mystery_func2(n-1)
```

- A. Adding up from 1 to n
- B. Subtracting 1 from n
- C. Multiplying from 1 to n
- D. Dividing n with 1
- E. None of the above

What is the result of the given recursive function? (Assume  $a = 2$  and  $b = 3$ .)

```
def mystery_func3(a, b):  
    if b == 0:  
        return 1  
    if b % 2 == 0:  
        return mystery_func3(a*a, b//2)  
    return mystery_func3(a*a, b//2) * a
```

- A. 5
- B. 6
- C. 8
- D. 16
- E. None of the above

What is the result of the given recursive function? (Assume  $a = \text{"1234"}$ .)

```
def mystery_func4(a):  
    if len(a) == 1:  
        return a  
    return mystery_func4(a[1:]) + a[0]
```

- A. "1234"
- B. "4321"
- C. "4"
- D. "1"
- E. None of the above

How many recursive calls and base cases are there for the given function?

```
def mystery_func5(n):  
    if n == 0:  
        return 0  
    if n == 1:  
        return 1  
    return mystery_func5(n-1) + mystery_func5(n-2)
```

- A. 1 and 1
- B. 1 and 2
- C. 2 and 1
- D. 2 and 2
- E. None of the above

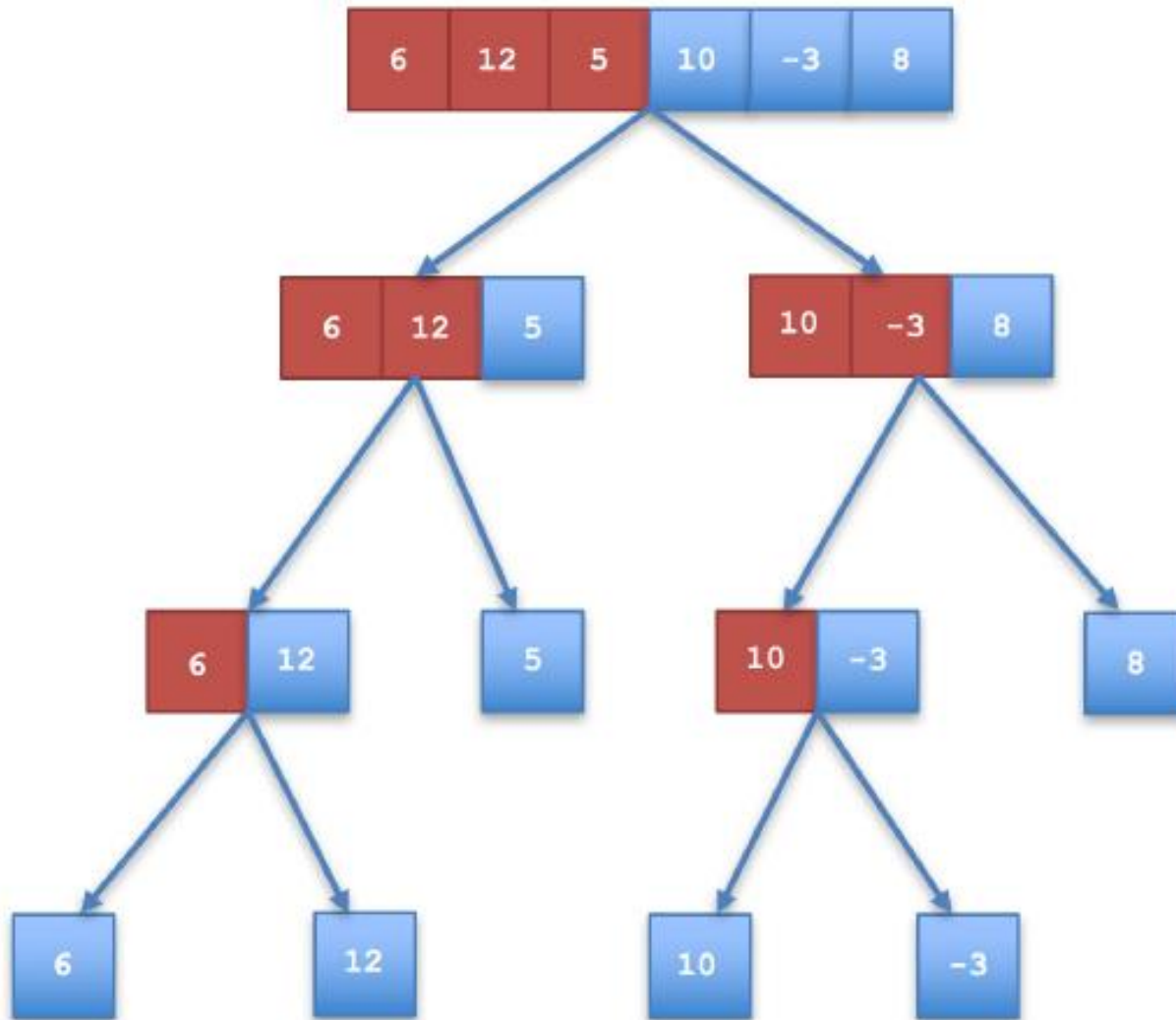


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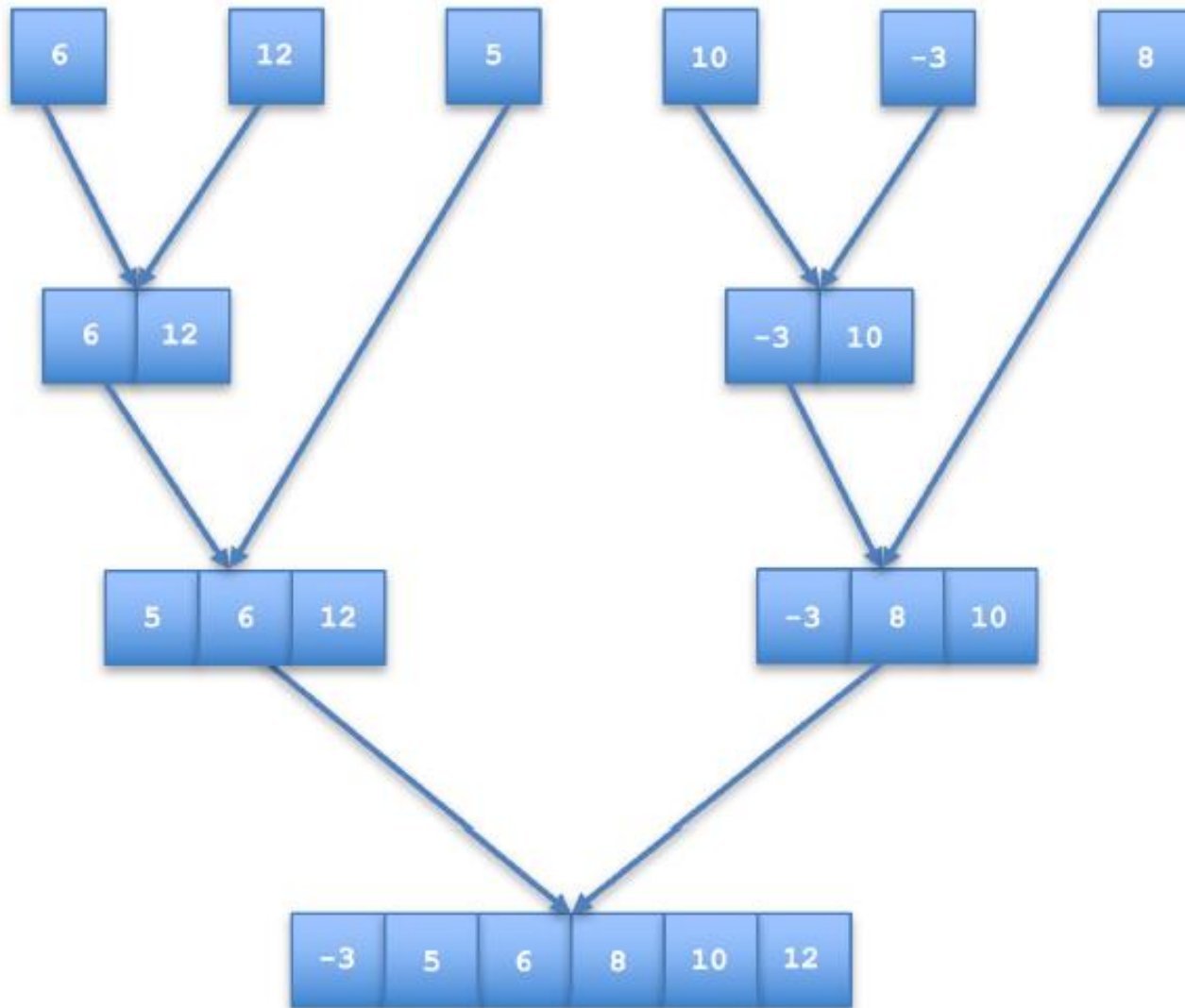
# Recursive Sorting: Merge Sort

- Basic ideas:
  - **Splits** an unsorted list into two halves “recursively” until there is only one element left in each sublist
  - Sublists are then **sorted and merged** until the complete sorted list is obtained
- Divide and conquer:
  - **Base case**: A sublist with length of one (considered *sorted*)
  - **Divide**: Recursively identify the middle point of a list and divide into two halves
  - **Conquer**: Sort the smaller sublists
  - **Combine**: Merge the sorted sublist into one complete list

# Merge Sort: Splitting



# Merge Sort: Merging





# Merge Sort: Implementation

```
def merge_sort(the_list):  
    # obtain the length of the list  
    n = len(the_list)  
  
    if n > 1: # check for base case  
        # find the middle of the list  
        mid = n // 2  
  
        # based upon the middle index, two sublists are then created  
        # from the 0 index until the mid-1 index  
        left_sublist = the_list[:mid]  
        # from the mid index until the n - 1 index  
        right_sublist = the_list[mid:]  
        print("Splitting: " + str(left_sublist) + " and " + str(right_sublist))  
  
        # merge sort is called again on the two new created sublists  
        merge_sort(left_sublist)  
        merge_sort(right_sublist)
```

# Merge Sort: Implementation (continue)

```
# sort and merge
print("Merging " + str(left_sublist) + " with " + str(right_sublist))
i = 0 # index for left sublist
j = 0 # index for right sublist
k = 0 # index for main list

while i < len(left_sublist) and j < len(right_sublist):
    if left_sublist[i] <= right_sublist[j]:
        the_list[k] = left_sublist[i]
        i += 1
    else:
        the_list[k] = right_sublist[j]
        j += 1
    k += 1

# insert the remaining elements into main list
while i < len(left_sublist):
    the_list[k] = left_sublist[i]
    i += 1
    k += 1

while j < len(right_sublist):
    the_list[k] = right_sublist[j]
    j += 1
    k += 1

print("After merging the list is: " + str(the_list))
```



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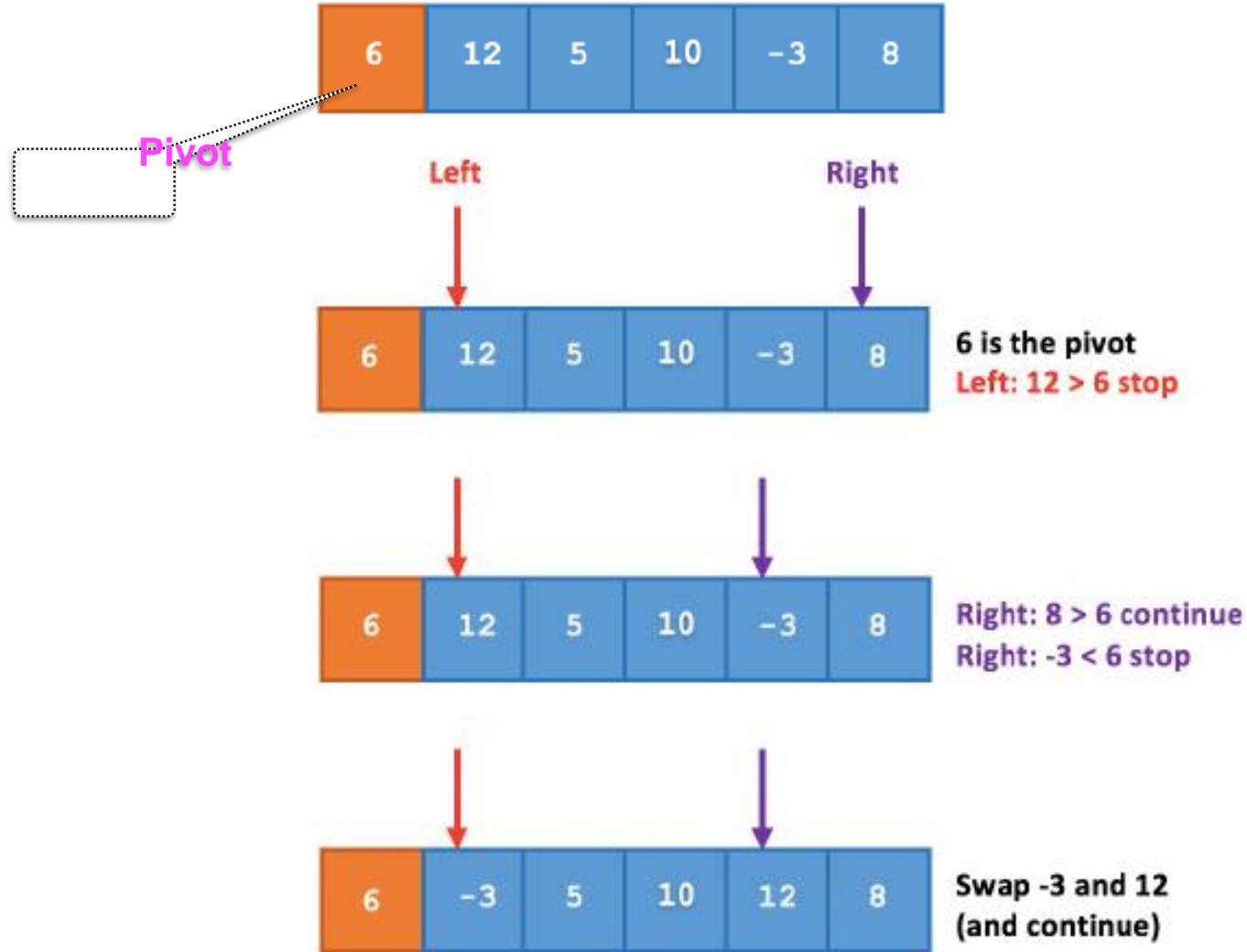
# Recursive Sorting: Quick Sort

# Quick Sort

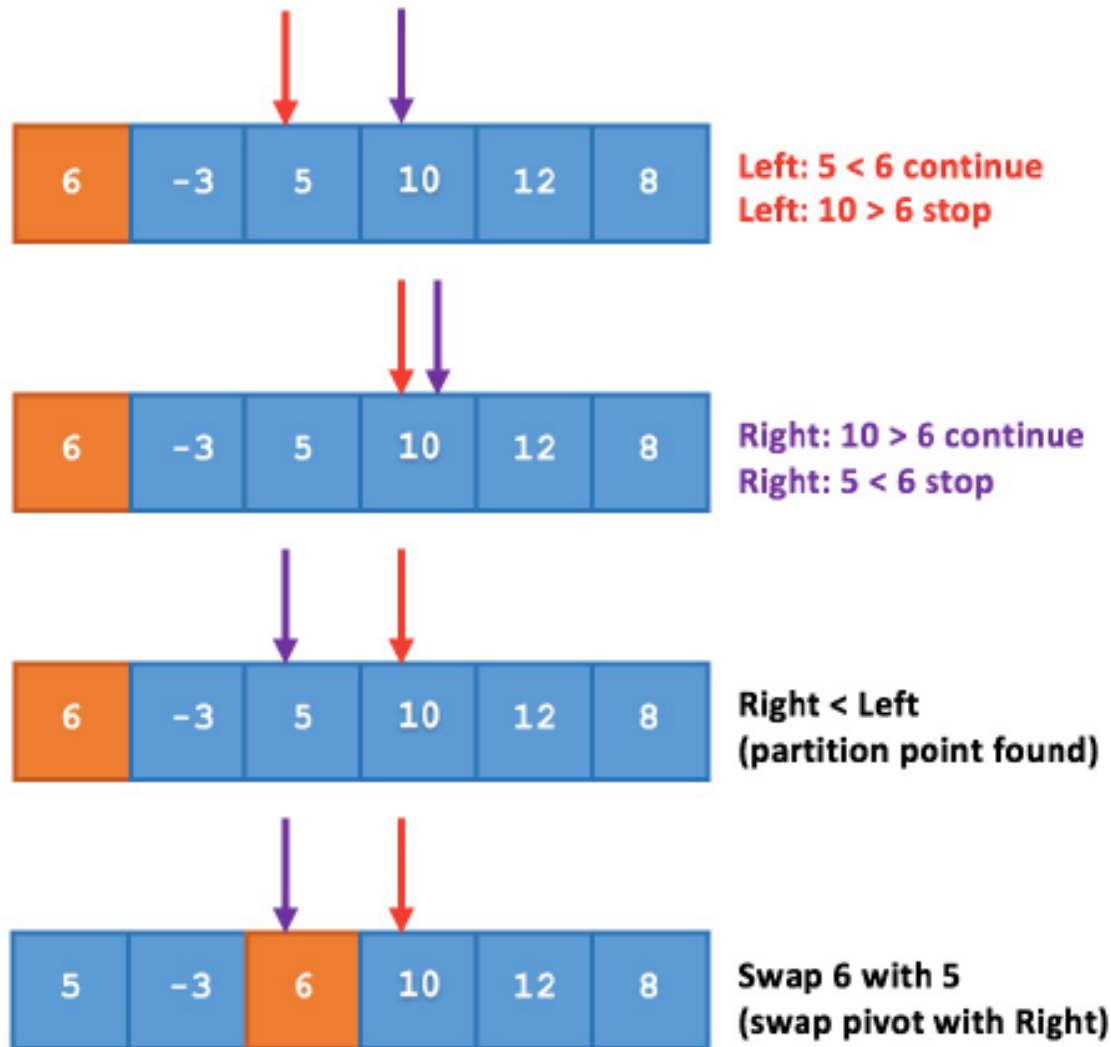
- Basic ideas:
  - Similarly to Merge Sort
  - Major computation is performed in “partitioning” (dividing the list into two partitions)
- Divide and conquer:
  - **Divide**: Select a “pivot” to serve as the partition point
    - Elements smaller than the pivot are relocated to the left of the pivot
    - Elements greater than the pivot are relocated to the right
  - **Conquer**: Recursively partition the sublists based on the pivot chosen for each sublist
  - **Combine**: No computation needed
  - **Base case**: A sublist with length of one (considered *sorted*) or with zero length

first element as  
“pivot”

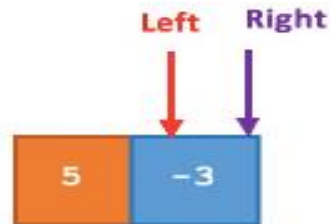
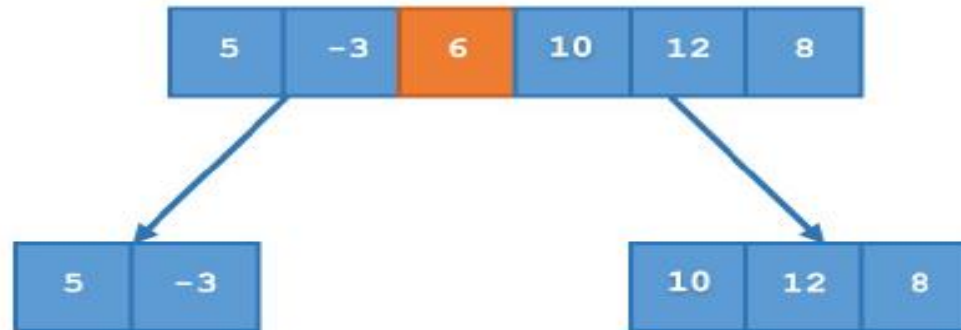
# Quick Sort: Partitioning



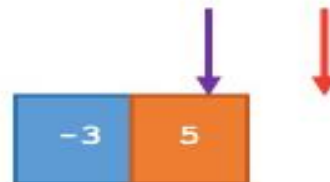
## Quick Sort: Partitioning (continue)



## Quick Sort: Partitioning (continue)



5 is the pivot for the sub-list [5, -3]  
Left:  $-3 < 5$  continue  
Left > Right  
(the end of the sub-list reached)



Swap 5 with -3  
(exchange pivot with Right)

At this stage, another two sub-lists will be created: [-3] and [ ]. Since these are the base case, the partitioning process terminates here. Note that the original list is considered partially sorted at this point. The other sub-list [10, 12, 8] will be sorted using the same partitioning procedure until the original list is completely sorted.

# Quick Sort: Implementation

```
def quick_sort(the_list):  
    # pass the indices of first and last elements of the list  
    first = 0  
    last = len(the_list) - 1  
    quick_sort_aux(the_list, first, last)  
  
def quick_sort_aux(the_list, first, last):  
    # if it is not the base case  
    if first < last:  
        # find the partition point  
        partition_point = partitioning(the_list, first, last)  
  
        print("partition at index: ", partition_point)  
        print(the_list)  
        print ("after partitioning: " + str(the_list[first:partition_point]) +  
              " and " + str(the_list[partition_point+1:last+1]))  
  
        # call the quick sort function again on the new sublists  
        quick_sort_aux(the_list, first, partition_point - 1)  
        quick_sort_aux(the_list, partition_point + 1, last)
```



# Quick Sort: Implementation (continue)

```
def partitioning(the_list, first, last):
    # take first element of the list is the pivot
    pivot_value = the_list[first]
    print("pivot: ", pivot_value)

    # these two indices will help us in locating the index point
    # where the list will be partitioned
    left_index = first + 1
    right_index = last

    complete = False

    while not complete:
        # start with the left index and keep on incrementing it
        # until a value greater than the pivot value is found
        while left_index <= right_index and
            the_list[left_index] <= pivot_value:
            left_index += 1

        # now look for element from the right of the list
        # which is smaller than the pivot value
        while right_index >= left_index and
            the_list[right_index] >= pivot_value:
            right_index -= 1

        # check whether left and right indices have crossed each other
        # if that is the case exit the while loop
        if right_index < left_index:
            complete = True
        else:
            # otherwise swap the two elements
            the_list[left_index], the_list[right_index]
            = the_list[right_index], the_list[left_index]

    # swap the pivot element with the element of the right index
    the_list[first], the_list[right_index]
    = the_list[right_index], the_list[first]

    # return right index which is the partition point
    return right_index
```

# Review Questions: Part 2

Given a list of integers below after the first partitioning of running Quick Sort, which of the integers could likely to be the *pivot*?

2, 5, 3, 7, 10, 12, 9

- A. 7
- B. 10
- C. Either 7 or 10
- D. Neither 7 nor 10

Given a list of integers below to be sorted using Quick Sort. How would the two partitions be if '4' is chosen as the *pivot*?

4, 5, 8, 10, 12, 9, 11

- A. Two almost even partitions
- B. Left partition with 0 elements
- C. Right partition with 0 elements
- D. Not sure

- We have discussed:
  - Concepts of recursion
  - Iterative solutions vs recursive solutions
  - Recursive sorting algorithms (MergeSort, QuickSort)