



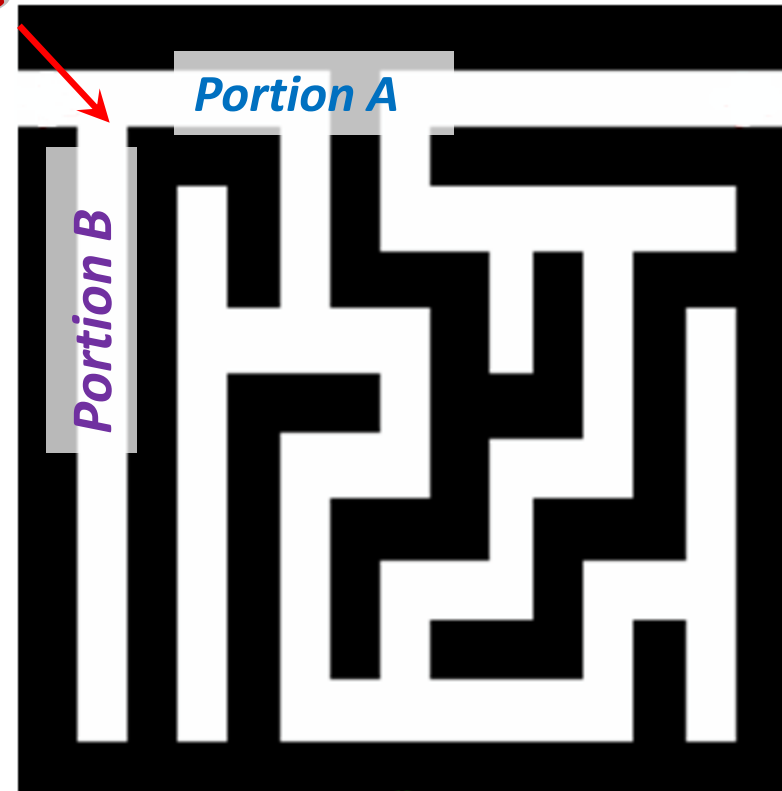
Back tracking Algorithm

Backtracking



- Backtracking is a technique used to solve problems with a large search space, by systematically trying and eliminating possibilities.
- A standard example of backtracking would be going through a maze.
 - At some point in a maze, you might have two options of which direction to go:

Junction

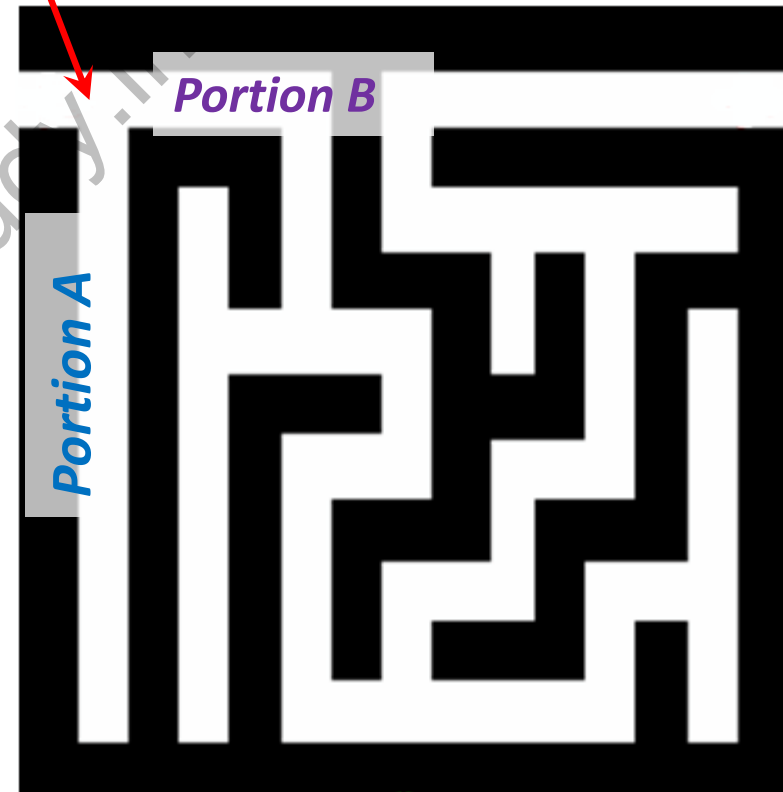


Backtracking



- One strategy would be to try going through **Portion A** of the maze.
 - If you get stuck before you find your way out, then you "**backtrack**" to the junction.
- At this point in time you know that **Portion A** will **NOT** lead you out of the maze,
 - so you then start searching in **Portion B**

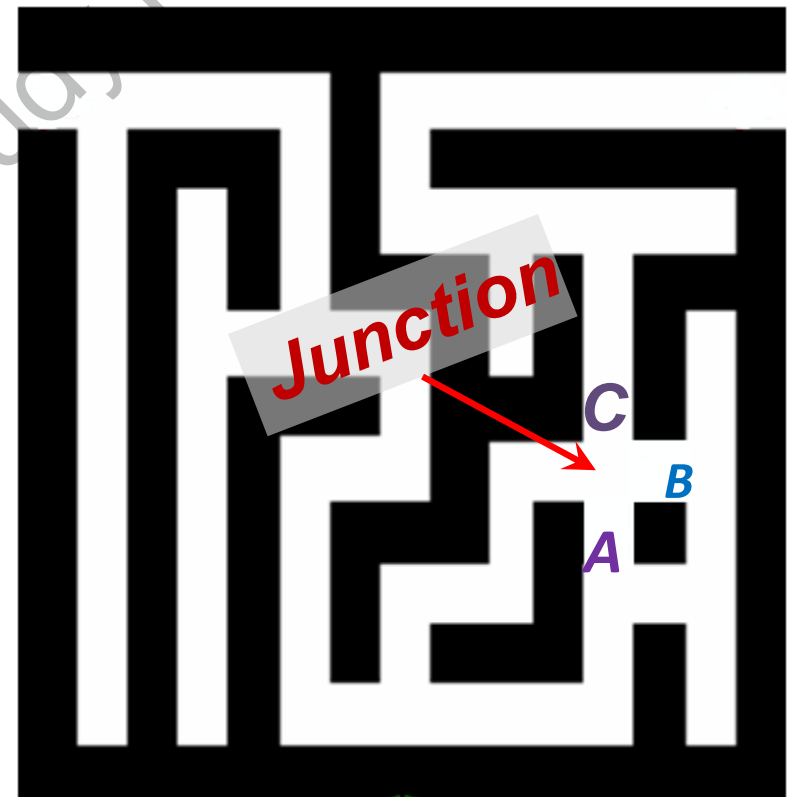
Junction



Backtracking



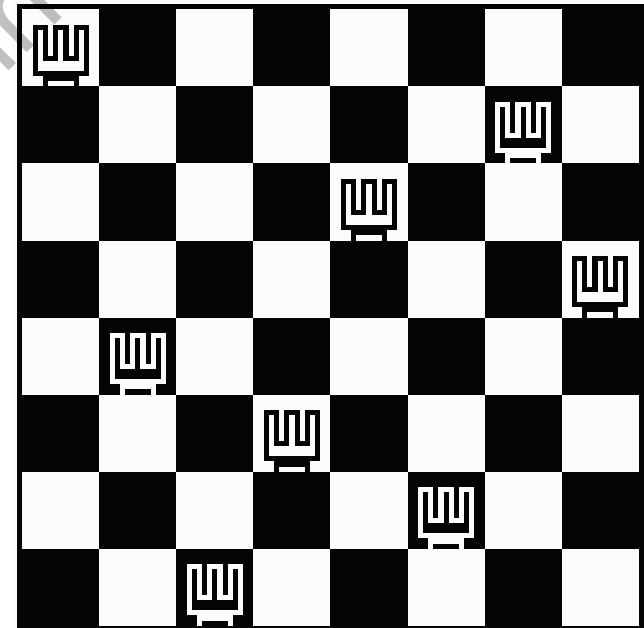
- Clearly, at a single junction you could have even more than 2 choices.
- The backtracking strategy says to try each choice, one after the other,
 - if you ever get stuck, "**backtrack**" to the junction and try the next choice.
- If you try all choices and never found a way out, then there IS no solution to the maze.



Backtracking – Eight Queens Problem



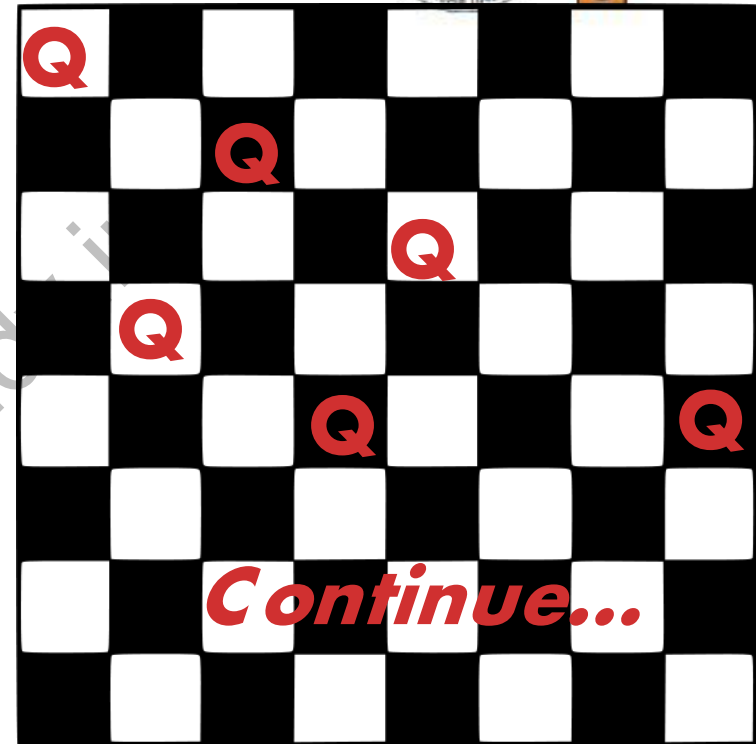
- Find an arrangement of **8** queens on a single chess board such that no two queens are attacking one another.
- In chess, queens can move all the way down any row, column or diagonal (so long as no pieces are in the way).
 - Due to the first two restrictions, it's clear that each row and column of the board will have exactly one queen.



Backtracking – Eight Queens Problem



- The backtracking strategy is as follows:
 - 1) Place a queen on the first available square in row 1.
 - 2) Move onto the next row, placing a queen on the first available square there (that doesn't conflict with the previously placed queens).
 - 3) Continue in this fashion until either:
 - a) you have solved the problem, or
 - b) you get stuck.
 - When you get stuck, remove the queens that got you there, until you get to a row where there is another valid square to try.



Animated Example:

<http://www.hbmeyer.de/backtrack/achtdamen/eight.htm#up>

Backtracking – Eight Queens Problem



- When we carry out backtracking, an easy way to visualize what is going on is a tree that shows all the different possibilities that have been tried.
- On the board we will show a visual representation of solving the 4 Queens problem (placing 4 queens on a 4x4 board where no two attack one another).

Backtracking – Eight Queens Problem



- The neat thing about coding up backtracking, is that it can be done recursively, without having to do all the bookkeeping at once.
 - Instead, the stack or recursive calls does most of the bookkeeping
 - (ie, keeping track of which queens we've placed, and which combinations we've tried so far, etc.)



perm[] - stores a valid permutation of queens from index 0 to location-1.

location – the column we are placing the next queen

usedList[] – keeps track of the rows in which the queens have already been placed.

void solveItRec(int perm[], int location, struct onesquare usedList[]) {

```
if (location == SIZE) {  
    printSol(perm);  
}
```

← Found a solution to the problem, so print it!

```
for (int i=0; i<SIZE; i++) {
```

← Loop through possible rows to place this queen.

```
    if (usedList[i] == false) {
```

← Only try this row if it hasn't been used

```
        if (!conflict(perm, location, i)) {
```

← Check if this position conflicts with any previous queens on the diagonal

```
            perm[location] = i;  
            usedList[i] = true;  
            solveItRec(perm, location+1, usedList);  
            usedList[i] = false;
```

- ←
- 1) mark the queen in this row
 - 2) mark the row as used
 - 3) solve the next column location recursively
 - 4) un-mark the row as used, so we can get ALL possible valid solutions.

```
        }
```

```
    }
```

```
}
```

```
}
```

Backtracking – 8 queens problem - Analysis



- Another possible brute-force algorithm is generate the permutations of the numbers 1 through 8 (of which there are $8! = 40,320$),
 - and uses the elements of each permutation as indices to place a queen on each row.
 - Then it rejects those boards with diagonal attacking positions.
- The backtracking algorithm, is a slight improvement on the permutation method,
 - constructs the search tree by considering one row of the board at a time, eliminating most non-solution board positions at a very early stage in their construction.
 - Because it rejects row and diagonal attacks even on incomplete boards, it examines only 15,720 possible queen placements.
- A further improvement which examines only 5,508 possible queen placements is to combine the permutation based method with the early pruning method:
 - The permutations are generated depth-first, and the search space is pruned if the partial permutation produces a diagonal attack



Thank You !!!

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