Iteration

Looping solution through successive (sub)cases of a problem

```
Example: Factorials
```

```
N! = N * (N-1) * (N-2) * . . . * 2 * 1
```

Observations:

- Requires temporary variable (prod)
- Requires looping mechanism: loop index counter initialization, completion testing, incrementing
- Loop boundary conditions -- critical

Recursion

Define problem solution in terms of smaller problems of the same type.

```
Example: Factorials
```

$$N! = N * (N-1) * (N-2) * . . . * 2 * 1$$

$$N! = N * (N-1) * (N-2) * . . . * 2 * 1 (N-1)!$$

$$N! = N * (N-1)!$$
 (N > 0)
 $0! = 1$ (N = 0) BASE CASE

Characteristics of a recursive function:

- A recursive function calls itself (one or more times)
- Each recursive call solves an identical, but smaller, problem
- A test for the base case enables the recursive calls to stop
- Eventually, one of the smaller problems must be the base case

Key Concepts for Constructing Recursive Solutions:

- 1. How can you define the problem in terms of a smaller problem of the same type?
- 2. How does each recursive call diminish the size of the problem?
- 3. What instance of the problem can serve as the base case?
- 4. As the problem size diminishes, will you reach this base case?

Advantage of Recursion:

• Often "elegant" or "simple" solution

DisAdvantage of Recursion:

- None in general
- Sometimes execution "efficiency" (vs. "iterative" solution)
- May be "harder" to debug ("box" drawings)

Example: Recursive void function to write a String backwards

H	i	T	h	е	r	е	

```
// Iterative version.
void WriteBackward(stringType S, int Size)
{
    while (Size > 0)
    {
       cout << S[Size-1];
       --Size;
    }
}</pre>
```

Example: Compute terms in Fibonacci sequence

```
1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, ...
```

```
int Rabbit(int N)
// Computes a term in the Fibonacci sequence.
// Precondition: N is a positive integer.
// Postcondition: Returns the Nth Fibonacci number.
   if (N <= 2)
      return 1;
   else
      return Rabbit(N-1) + Rabbit(N-2);
}
int RabbitIterative(int N)
    if (N \le 2)
        return 1;
    // Iterate: ... rab_2 rab_1 rab ....
    int rab_2 = 1;  // 1
int rab_1 = 1;  // 2
    int rab;
    for(int i = 3; i \le N; ++i)
         rab = rab_2 + rab_1; // Sum of prior two iterations
         rab_2 = rab_1;
         rab_1 = rab;
    }
    return(rab);
}
```

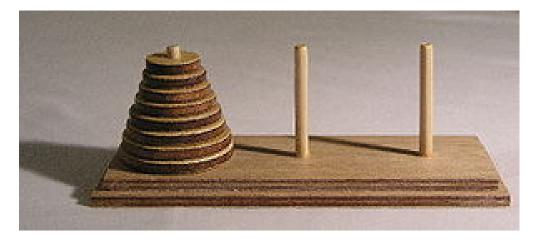
Example: Determine how many ways to choose K items from a collection of N items.

Example: Searching Lists

- Find largest item in arbitrary list
- Find kth smallest item in a list
- Find item in sorted list

```
void BinarySearch(const int A[], int First, int Last,
                  int Value, int& Index)
// Searches the array items A[First] through A[Last]
         for Value by using a binary search.
// Precondition: 0 <= First, Last <= SIZE-1, where</pre>
        SIZE is the maximum size of the array, and
//
        A[First] \leftarrow A[First+1] \leftarrow ... \leftarrow A[Last].
//
// Postcondition: If Value is in the array, Index is
// the index of the array item that equals Value;
//
   otherwise Index is -1.
   if (First > Last)
      Index = -1; // Value not in original array
   else
   {
      // Invariant: If Value is in A,
      // A[First] <= Value <= A[Last]</pre>
      int Mid = (First + Last)/2;
      if (Value == A[Mid])
         Index = Mid; // Value found at A[Mid]
      else if (Value < A[Mid])</pre>
         BinarySearch(A, First, Mid-1, Value, Index);
      else
         BinarySearch(A, Mid+1, Last, Value, Index);
   }
}
```

Example: Towers of Hanoi



```
Goal: Move initial stack of disks to another pole.
Rules:
1) Move 1 disk at a time.
2) Never place a larger disk on top of a smaller disk.
```

```
// Recursive solution is much more obvious than iterative
// How would you solve this problem iteratively???
```