

Assignment 3

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May 28, 2018

1 Task 1

1.1 Type *word*

This definition of word is basically from requirement of assignment 3. We say two words v, w that v is an absolute prefix of w as $v < w$ which is defined as $v \leq w \wedge v \neq w$.

1.2 abstract Data Type *Dict*

According to the specified problem statement in the assignment, we could describe the syntactic data type *Dict* as below. The encapsulated state is a dictionary word set W .

$$Dict = (W = \phi, \left(\begin{array}{l} \text{proc } addword^{Dict}(\text{word } w) \cdot b, W : [\text{TRUE}, b = b_0 \wedge W = W_0 \cup \{w\}] \\ \text{func } checkword^{Dict}(\text{word } w) : \mathbb{B} \cdot \text{var } b \cdot b, W : [\text{TRUE}, b = (w \in W_0)]; \text{return } b \\ \text{proc } delword^{Dict}(\text{word } w) \cdot b, W : [w \in W, b = b_0 \wedge W = W_0 \setminus \{w\}] \end{array} \right))$$

2 Task 2

Now we would like to refine *Dict* to a second data type *DictA* where we replace W with a trie t , the corresponding trie domain $D = \mathbf{dom}(t)$. It represents the set of all tries according to the domain. We shall use this definition later in our refinement.

2.1 Datat Invariant

$$\forall w \in \mathbf{dom}(t), t(w) = 1, w' \leq w (w' \in \mathbf{dom}(t))$$

2.2 Data Type Refinement

This suggests we should first build up a inductively defined predicate to ensure the provable relations between $DictA$ and $Dict$.

$$r = (W = \{w \in \mathbf{dom}(t) | t(w) = 1\})$$

which we can translate into a function from concrete to abstract values:

$$f(t) = \{w \in \mathbf{dom}(t) | t(w) = 1\}$$

With that in mind we propose the initialisation predicate $init^{DictA} = (\mathbf{dom}(t) = \{\epsilon\} \wedge f(t) = \phi)$ and operations given as follows.

```

proc  $addword^{DictA}(\mathbf{word} \ w) \cdot b, t :$ 
  [  $\text{TRUE}, b = b_0 \wedge t = t_0 \cup \{w \mapsto 1\} \cup \{w' < w \wedge w' \notin \mathbf{dom}(t) | w' \mapsto 0\}$  ]
func  $checkword^{DictA}(\mathbf{word} \ w) \mathbb{B} \cdot \mathbf{var} \ b \cdot b, t :$ 
  [  $\text{TRUE}, b = (w \in \mathbf{dom}(t) \wedge t(w) = 1)$  ]; return  $b$ 
proc  $delword^{DictA}(\mathbf{word} \ w) \cdot b, t :$ 
  [  $\text{TRUE}, b = b_0 \wedge (w \notin \mathbf{dom}(t) \vee t := t : w \mapsto 0)$  ]

```

2.3 Proof of Refinement

We need start the proof with the initialisation:

$$\begin{aligned}
& init^{DictA} \Rightarrow init^{Dict} [f(t)/_W] \\
& \Leftrightarrow \langle \text{Definition of } init^{DictA} \text{ and } init^{Dict} \rangle \\
& \forall w \in \mathbf{dom}(t) (t(w) = 0) \Rightarrow W = \phi
\end{aligned}$$

Since all our precondition of concrete is trivial which all of them are TRUE, we don't need to proof the condition (3_f). But condition (4_f) must be checked for all three operations. For the *addword* we proof:

$$\begin{aligned}
& pre_{addword^{Dict}} [f(t_0)/_W] \wedge post_{addword^{DictA}} \\
& \Leftrightarrow \langle \text{Definition of } addword^{Dict} \text{ and } addword^{DictA} \rangle \\
& \text{TRUE} [f(t_0)/_W] \wedge b = b_0 \wedge t = t_0 \cup \{w \mapsto 1\} \cup \{w' < w \wedge w' \notin \mathbf{dom}(t) | w' \mapsto 0\} \\
& \Rightarrow \langle \text{Definition of } f \rangle \\
& f(t) = f(t_0 \cup \{w \mapsto 1\} \cup \{w' < w \wedge w' \notin \mathbf{dom}(t) | w' \mapsto 0\}) \\
& \Leftrightarrow \langle \text{Logic} \rangle \\
& f(t) = f(t_0) \cup \{w\} \\
& \Leftrightarrow \langle \text{Definition of } addword^{Dict} \text{ and } addword^{DictA} \rangle \\
& post_{addword^{Dict}} [f(t_0), f(t)/_{W_0, W}]
\end{aligned}$$

For the *checkword* we proof:

$$\begin{aligned}
& pre_{checkword}^{Dict} [f(t_0)/W] \wedge post_{checkword}^{DictA} \\
\Leftrightarrow & \langle \text{Definition of } checkword^{DictA} \text{ and } checkword^{Dict} \rangle \\
& TRUE [f(t_0)/W] \wedge b = (w \in \mathbf{dom}(t) \wedge t = t_0) \\
\Rightarrow & \langle \text{Definition of } f \rangle \\
& b = (w \in f(t)) \wedge t = t_0 \\
\Leftrightarrow & \langle \text{Definition of } checkword^{DictA} \text{ and } checkword^{Dict} \rangle \\
& post_{checkword}^{Dict} [f(t_0), f(t)/W_0, W]
\end{aligned}$$

For the *delword* we proof:

$$\begin{aligned}
& pre_{delword}^{Dict} [f(t_0)/W] \wedge post_{delword}^{DictA} \\
\Leftrightarrow & \langle \text{Definition of } delword^{DictA} \text{ and } delword^{Dict} \rangle \\
& w \in f(t_0) \wedge b = b_0 \wedge (w \notin \mathbf{dom}(t) \vee t := t : w \mapsto 0) \\
\Rightarrow & \langle \text{Definition of } f \rangle \\
& f(t) = f(t_0) \setminus \{w\} \\
\Leftrightarrow & \langle \text{Definition of } delword^{DictA} \text{ and } delword^{Dict} \rangle \\
& post_{delword}^{Dict} [f(t_0), f(t)/W_0, W]
\end{aligned}$$

3 Task 3

We derive our code in to fours part by *init* and its operations.

3.1 init

From the spec we have:

$$\begin{aligned}
& \mathbf{dom}(t) = \{\epsilon\} \wedge f(t) = \phi \\
\sqsubseteq & \langle \text{ass} \rangle \\
& t := \{\epsilon \mapsto 0\}
\end{aligned}$$

3.2 addword

From the spec we have:

$$\begin{aligned}
& \mathbf{proc} \text{ addword}^{DictA}(\mathbf{word} \ w). \\
& \quad \sqcup b, t : [TRUE, b = b_0 \wedge t = t_0 \cup \{w \mapsto 1\} \cup \{w' < w \wedge w' \notin \mathbf{dom}(t) | w' \mapsto 0\}] \neg(A1) \\
(A1) \sqsubseteq & \langle \text{c-frame} \rangle \\
& t : [TRUE, t = t_0 \cup \{w \mapsto 1\} \cup \{w' < w \wedge w' \notin \mathbf{dom}(t) | w' \mapsto 0\}]
\end{aligned}$$

3.3 checkword

From the spec we have:

```
func checkwordDictA(word w)  
   $\sqsubseteq \mathbb{B} \cdot \mathbf{var} \ b \cdot b, t : [\text{TRUE}, b = (w \in \mathbf{dom}(t) \wedge t = t_0)]; \mathbf{return} \ b \sqsubseteq_{(C1)}$ 
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```
 $(C1) \sqsubseteq \langle \text{c-frame} \rangle$   
 $\mathbb{B} \cdot \mathbf{var} \ b \cdot b : [\text{TRUE}, b = (w \in \mathbf{dom}(t) \wedge t = t_0)]; \mathbf{return} \ b$ 
```

3.4 delword

From the spec we have:

```
proc delwordDictA(word w)  
   $\sqsubseteq \cdot b, t : [\text{TRUE}, b = b_0 \wedge (w \notin \mathbf{dom}(t) \vee t := t : w \mapsto 0)] \sqsubseteq_{(D1)}$ 
```