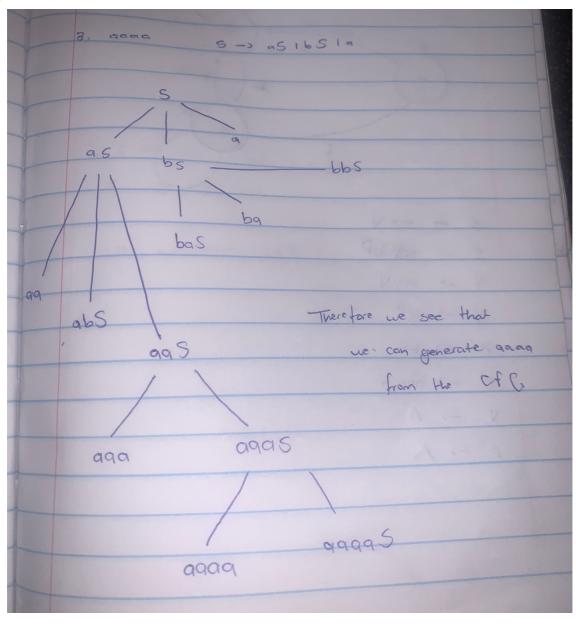
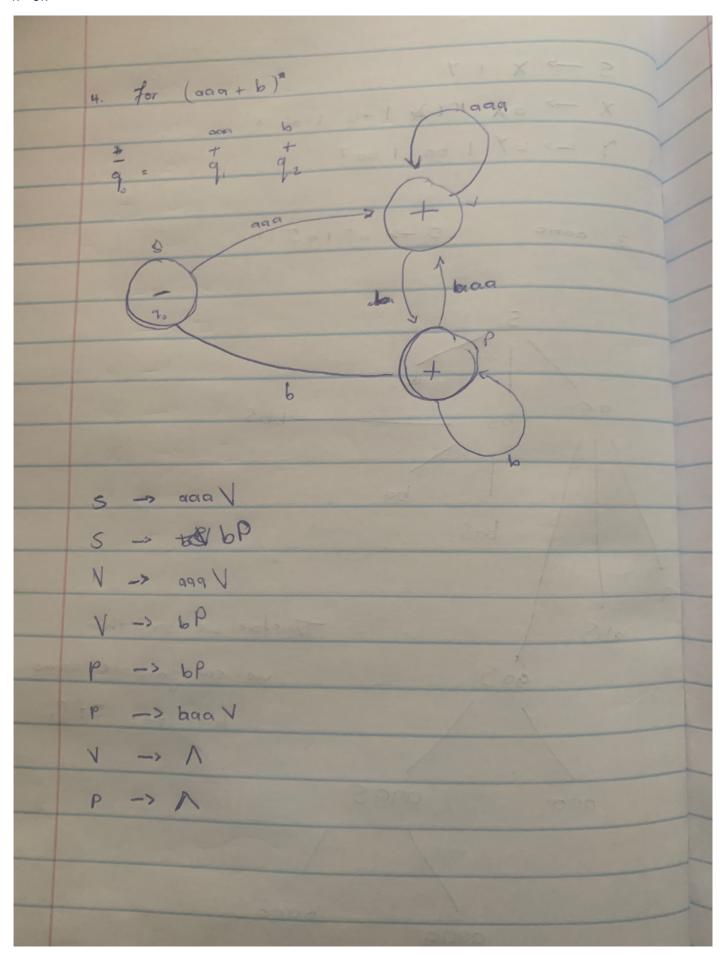
- 1. (b + a)*a(b + a)*a(a + b)*This is the language with atleast 2 a's
- 2. S -> X | Y X -> aX | ab | abX Y -> bY | ba | baY
- 3. ok



4. ok



Since only S -> aX and X -> aX contain the X that leads to Λ then only X -> a

therefore

Defines the same CFG

$$Z \rightarrow a \mid b \mid baY \mid \Lambda$$

Split

Remove the **∧**

Reformulated

S -> aXY
S -> Z
X -> bbb
$X \rightarrow bX$
X -> aZ
Y -> aS
Y -> ∧
Z -> a
Z -> b
Z - > baY
Z -> \

S -> aXY S -> Z X -> bbb X -> bX X -> aZ Y -> aS Y -> aX Y -> ba Z -> a Z -> b

S -> Z
X -> ZZZ X -> ZX
X -> ZZ Y -> ZS
Y -> ZX Y -> ZZ
Z -> ZZY Z -> a Z -> b

Chomsky Normal Form

X -> ZR

R -> ZZ

X -> ZX

X -> ZZ

Y -> ZS

Y -> ZZ

Y -> ZX

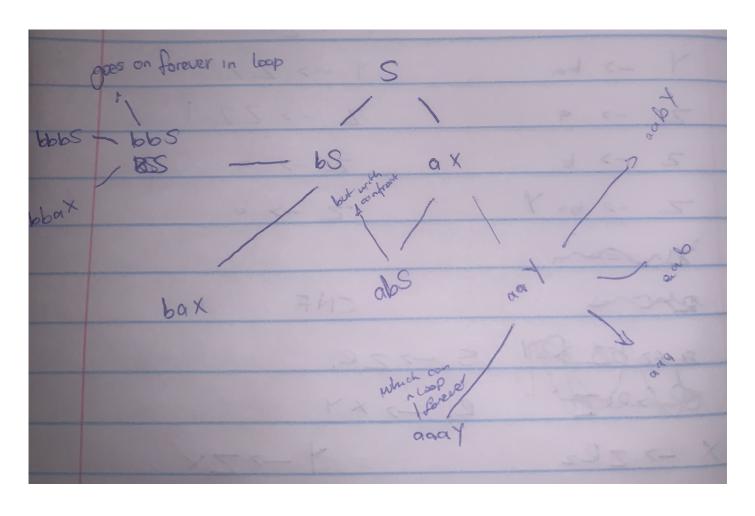
Z -> ZR

R -> ZY

Z -> a

Z -> b

7. S -> bS | ax X -> bS | aY Y -> aY | bY | a | b



b*a(b*a)*a(a + b)*The language of the first 2 b's having an a after them. It also has at least 2 a's.

