

# Notes on “A Practical Cryptanalysis of the Algebraic Eraser”

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Let *KEY* denote the paper “Key Agreement, the Algebraic Eraser<sup>TM</sup>, and Lightweight Cryptography”.

### 1. Notation

- $\mathbb{F}$  is a finite field of small order.  $S_n$  is the symmetric group on  $n$  elements.  $N = \text{GL}_n(\mathbb{F})$  is the general linear group over  $\mathbb{F}$ .
- $M$  is a subgroup of  $\text{GL}_n(\mathbb{F}(t_1, \dots, t_n))$ , where the  $t_i$ 's are algebraically independent commuting indeterminates, such that  $M$  is contained in the subgroup of  $\text{GL}_n(\mathbb{F}(t_1, \dots, t_n))$  of matrices whose determinant can be written as  $at$  for some non-zero element  $a \in \mathbb{F}$  and some, possibly empty, word  $\mathbf{t}$  in the elements  $t_i$  and their inverses.
- $\overline{M}$  is the subgroup of  $\text{GL}_n(\mathbb{F}(t_1, \dots, t_n))$  generated by permuting the indeterminants of  $M$  in all possible ways.
  - $S_n$  acts on  $\overline{M}$  by permuting the indeterminants  $t_i$ .
- $\varphi : \overline{M} \rightarrow \text{GL}_n(\mathbb{F}(t_1, \dots, t_n))$  is the evaluation homomorphism, where the evaluation elements are the non-zero elements  $\tau_1, \dots, \tau_n \in \mathbb{F}$ .
  - This function is denoted as  $\Pi : M \rightarrow N$  in *KEY*.
- The semidirect product  $\overline{M} \rtimes S_n$  has the operation defined as

$$(a, g)(b, h) = (a^g b, gh)$$

for all  $(a, g), (b, h) \in \overline{M} \rtimes S_n$  where  $^g a$  is the left action of  $g \in S_n$  on  $a \in \overline{M}$ .

- In the CBKAP,  $C$  and  $D$  (i.e.  $C = D$ ) are the subgroups of  $\text{GL}_n(\mathbb{F})$  consisting of all invertible matrices of the form  $l_0 + l_1 \kappa + \dots + l_r \kappa^r$  where  $\kappa$  is a fixed matrix,  $l_i \in \mathbb{F}$ , and  $r \geq 0$ .
  - In this paper,  $C$  and  $D$  are only assumed to be subgroups of  $\text{GL}_n(\mathbb{F})$  that commute element wise:  $cd = dc$  for all  $c \in C$  and  $d \in D$ ... Well actually, in Section 4 they assume that  $C = \text{Alg}^*(C)$ , a more general assumption.
  - $C$  and  $D$  are denoted as  $N_A$  and  $N_B$  in *KEY*.
- $\Omega = \text{GL}_n(\mathbb{F}) \times S_n$ . This is denoted as  $N \times S$  in *KEY* section 4.
- $\widehat{S_n} = \overline{M} \rtimes S_n$ . This is denoted as  $M \rtimes S$  in *KEY* section 4.
- $*$  :  $\Omega \times \widehat{S_n} \rightarrow \Omega$  is defined as

$$(s, g) * (b, h) = (s\varphi(^g b), gh)$$

for all  $(s, g) \in \Omega$  and  $(b, h) \in \widehat{S}_n$ . This is the Algebraic Eraser function defined in KEY, so  $*$ -commuting makes sense.

- $\bullet : \text{GL}_n(\mathbb{F}) \times \Omega \rightarrow \Omega$  is defined as

$$x \bullet (s, g) = (xs, g)$$

for all  $x \in \text{GL}_n(\mathbb{F})$  and all  $(s, g) \in \Omega$ .

- Let  $A$  and  $B$  be subgroups of  $\widehat{S}_n$  that  $*$ -commute: for all  $(a, g) \in A$ ,  $(b, h) \in B$ , and  $\omega = (x_\omega, s_\omega) \in \Omega$ ,

$$(\omega*(a, g))*(b, h) = (x_\omega\varphi^{(s_\omega a)}, s_\omega g)*(b, h) = (x_\omega\varphi^{(s_\omega b)}, s_\omega h)*(a, g) = (\omega*(b, h))*(a, g)$$

## 2. Proposed Attack

- For an arbitrary group  $H \subseteq M_n(\mathbb{F}_q)$ , we write  $\text{Alg}(H)$  for the set of all  $\mathbb{F}_q$ -linear combinations of matrices in  $H$ . We write  $\text{Alg}^*(H)$  for the set of all invertible matrices in  $\text{Alg}(H)$ .
- Let  $C = \text{Alg}^*(C)$ . Let  $\kappa_1, \kappa_2, \dots, \kappa_r \in C$  be a basis for  $\text{Alg}(C)$ .
- Let  $P \leq A$  be defined as (the pure subgroup)

$$P = \{(\alpha, g) \in A : g = e\}.$$

- Let  $Q \leq N$  be  $Q = \text{Alg}^*(\varphi(P))$ . This implies that for all  $\alpha' \in Q$ ,

$$\alpha' = \sum_{i=1}^k l_i \varphi(\alpha_i)$$

where  $k \geq 0$ ,  $l_i \in \mathbb{F}_q$ , and  $(\alpha_i, e) \in P$ .

- Eve finds elements  $\tilde{c} \in C$ ,  $\alpha' \in Q$  and  $(\tilde{a}, g) \in \widehat{G}$  such that

$$(p, g) = \tilde{c} \bullet (\alpha', e) * (\tilde{a}, g)$$

and elements  $(\alpha_i, e) \in P$  and  $l_i \in \mathbb{F}_q$  such that

$$\sum_{i=1}^k l_i \varphi(\alpha_i) = \alpha'.$$

If Eve finds this it is GAME OVER FOR ALICE AND BOB.