Reliable and Automatic Composition of Language Extensions to C

Making language extension practical with AbleC

Ted Kaminski, Lucas Kramer, Travis Carlson, and Eric Van Wyk (University of Minnesota)

October 2017

What do we mean by language extensions?

```
cilk int count matches (Tree *t) {
typedef datatype Tree Tree;
                                       match ( t ) {
                                         Fork(t1.t2.str): {
datatype Tree {
  Fork ( Tree*, Tree*, const char* );
                                           int res_t, res_t1, res_t2;
  Leaf ( const char* );
                                           spawn res t1 = count matches( t1 );
};
                                           spawn res t2 = count matches( t2 );
                                           res t = (str = ~/foo[0-9]+/) ? 1 : 0:
                                           svnc:
                                           cilk return res t1 + res t2 + res t :
                                         Leaf(str): { return (str =~ /foo[0-9]+/) ? 1 : 0; };
```

Language extension research program goals

Expression problem¹ criteria, as applied to AST:

- 1. Introduce both new syntax and new analysis
- 2. Static checking of this representation
- 3. Without modifying the original code
- 4. Separate compilation

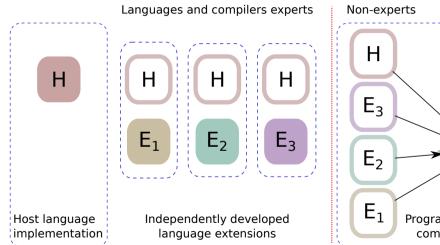
Independent extensibility²:

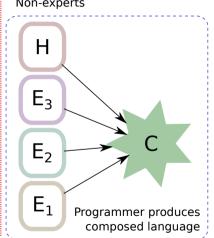
5. "Unordered," composable extensions

Our criteria:

6. Automatic composition, no glue code

Library model of language extension





Distinguishing characteristics

Many approaches to extensible languages.

- ExtendJ (using JastAdd)
- SujarJ (using Spoofax)
- Wyvern/VerseML
- mbeddr

- ▶ XTC
- ➤ XoC
- object algebras
 - and others

What distinguishes this work is how two questions are answered:

- 1. Who composes the language features?
- 2. How expressive are the supported features?

Success! Composition without conflict.

- ► Copper & Silver: tools for creating extensible compilers
- ► AbleC: our application to C

The general operation of AbleC

```
pts.xc
edu:umn:cs:melt:IndData
                                   cpp
     org:bar:cilk
     com:foo:RegEx
                                 pts.i
                                my-ableC
        Silver
edu:umn:cs:melt:ableC
                                 pts.c
                                  gcc
                                 a.out
```

 $\begin{cases} -s canning \\ -p ar sing \\ -AST \ construction \\ -type \ checking \\ -optimization \\ -C \ code \ generation \end{cases}$

Previous work

Allowing expressive extensions:

Context-aware scanning: safe syntactic overlap between extensions

Forwarding: tool for solving the expression problem

Modular analyses ensure reliable composition:

Modular determinism analysis: no unexpected syntactic conflicts Modular well-definedness analysis: well-defined attribute grammar

Concurrent work

Coherent non-interference: extensions behave as specified

This talk

- Modular analyses impose restrictions
- Previous work, let's skip over that
- What kinds of extensions can we build?
- What kinds of extensions can't we build for plain C?
- What host language modifications allow for more kinds of extensions?

Extensions to AbleC

```
cilk int count matches (Tree *t) {
typedef datatype Tree Tree;
                                       match ( t ) {
                                         Fork(t1,t2,str): {
datatype Tree {
  Fork ( Tree*, Tree*, const char* );
                                           int res_t, res_t1, res_t2;
  Leaf ( const char* );
                                           spawn res t1 = count matches( t1 );
};
                                           spawn res t2 = count matches( t2 );
                                           res t = (str = ~/foo[0-9]+/) ? 1 : 0:
                                           svnc:
                                           cilk return res t1 + res t2 + res t :
                                         }:
                                         Leaf(str): { return (str =~ /foo[0-9]+/) ? 1 : 0; };
```

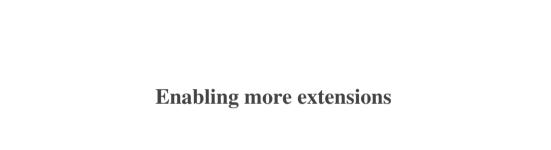
Extensions to AbleC

```
transform {
  for (unsigned i : m, unsigned j : n) {
   c[i][i] = 0:
   for (unsigned k : p) {
     c[i][j] += a[i][k] * b[k][j];
} bv {
  split i into (unsigned i outer,
                unsigned i inner: (m - 1) / NUM THREADS + 1);
  parallelize i_outer into (NUM_THREADS) threads;
  tile i_inner, j into (TILE_DIM, TILE_DIM);
  split k into (unsigned k outer,
                unsigned k_unroll : UNROLL_SIZE,
                unsigned k vector : VECTOR_SIZE);
 unroll k unroll:
 vectorize k vector;
```

Extensions to AbleC

Other extensions

- ► Sqlite: describing schemas, writing queries, LINQ-like
- ► Tensor/matrix/vector: more scientific computing applications
- ► **Go concurrency**: other parallel computing models



GCC extensions

- Something we already have to support
- ▶ One in particular especially useful: statement-expressions
- General problem with stratified grammars
- ► Transformations are local

```
1 + ({ foo x; f(&x); x.val; })
```

Operator overloading

- Cannot be implemented as an extension
 - ► (Changes meaning of host productions)
- Quite useful in enabling extensions, though
- ► An alternative non-syntactic "hook" from host into extension

```
matrix(A * x + V)
VS
A * x + V
```

Lifting declarations

- ► Extension translation ("forwarding") is local
- ► Sometimes need non-local transformations

```
lambda (int x) \rightarrow (*z = x * y + *z)
```

- ► Lift out a function declaration
- Lift out a closure type declaration

Type qualifiers

- General problem with annotation-driven analysis
 - ▶ Need to have an equivalent host-language tree
- ► Previous work³ has extended C with generic type qualifiers
 - ► Changes meaning of host again, not an extension in our sense
- ► New qualifiers as language extensions (See our GPCE paper this year!)

```
typedef datatype Expr Expr;
datatype Expr {
   Add (Expr * nonnull, Expr * nonnull);
   Mul (Expr * nonnull, Expr * nonnull);
   Const (int);
};
```

Enabling more extensions

Restrictions on extensions are host-language-relative

GCC extensions escaping the confines of expressions

Operator overloading additional options for "hooking into" extensions

Lifting declarations escaping the confines of local scope

Type qualifiers some kinds of annotation-driven analysis

Host language type system is important

Summary

- Extended notion of expression problem
- Reliable & automatic composition of language extensions
- Breadth of extensions possible
- Language extension impacts language design
 - Clear line between extension and modification
 - We can get experience with extensions before standardization

Thanks!

Get in touch:

- Eric Van Wyk <evw@cs.umn.edu>
- ► Ted Kaminski <tedinski@cs.umn.edu>

Check things out:

- ► melt.cs.umn.edu
- ► github.com/melt-umn



