SQLI

R & G Chapter 5



SQL Roots



- Developed @IBM Research in the 1970s
 - System R project
 - Vs. Berkeley's Quel language (Ingres project)
- Commercialized/Popularized in the 1980s
 - "Intergalactic Dataspeak"
 - IBM beaten to market by a startup called Oracle

SQL's Persistence



- Over 40 years old!
- Questioned repeatedly
 - 90's: Object-Oriented DBMS (OQL, etc.)
 - 2000's: XML (Xquery, Xpath, XSLT)
 - 2010's: NoSQL & MapReduce
- SQL keeps re-emerging as the standard
 - Even Hadoop, Spark etc. mostly used via SQL
 - May not be perfect, but it is useful

SQL Pros and Cons

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- Declarative!
 - Say what you want, not how to get it
- Implemented widely
 - With varying levels of efficiency, completeness
- Constrained
 - Not targeted at Turing-complete tasks
- General-purpose and feature-rich
 - many years of added features
 - extensible: callouts to other languages, data sources

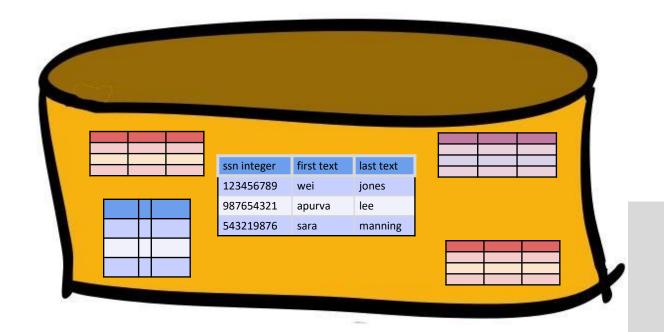
Content Break



Relational Terminology



Database: Set of named Relations



Relational Terminology, Pt 2.

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- Database: Set of named Relations
- Relation (Table):
 - Schema: description ("metadata")
 - Instance: set of data satisfying the schema

ssn integer	first text	last text
123456789	wei	jones
987654321	apurva	lee
543219876	sara	manning

Relational Terminology, Pt. 3

- Database: Set of named Relations
- **Relation** (Table):
 - Schema: description ("metadata")
 - Instance: set of data satisfying the schema
- Attribute (Column, Field)



first text

wei

apurva

sara

Relational Terminology, Pt. 4



- Database: Set of named Relations
- Relation (Table):
 - Schema: description ("metadata")
 - Instance: set of data satisfying the schema
- Attribute (Column, Field)
- *Tuple* (Record, Row)

543219876	sara	manning
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Relational Tables

- itelational lables
- Schema is fixed:
 - unique attribute names, *atomic* types
 - folks (ssn integer, first text, last text)
- Instance can change often
 - a multiset of "rows" ("tuples")

```
{(123456789, 'wei', 'jones'),
(987654321, 'apurva', 'lee'),
(543219876, 'sara', 'manning'),
(987654321, 'apurva', 'lee')}
```



Quick Check 1

• Why is this not a relation?

num integer	street text	zip integer	
84	Maple Ave	54704	
22	High	Street	76425
75	Hearst Ave	94720	

Quick Check 2

• Why is this not a relation?

num integer	street text	num integer
84	Maple Ave	54704
22	High Street	76425
75	Hearst Ave	94720

Quick Check 3

• Why is this not a relation?

first text	last text	addr address
wei	jones	(84, 'Maple', 54704)
apurva	lee	(22, 'High', 76425)
sara	manning	(75, 'Hearst', 94720)

Content Break 2



SQL Language



- Two sublanguages:
 - DDL Data Definition Language
 - Define and modify schema
 - DML Data Manipulation Language
 - Queries can be written intuitively.
- RDBMS responsible for efficient evaluation.
 - Choose and run algorithms for declarative queries
 - Choice of algorithm must not affect query answer.

Example Database



Sailors

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

Boats

<u>bid</u>	bname	color
101	Nina	red
102	Pinta	blue
103	Santa Maria	red

Reserves

<u>sid</u>	bid	day
1	102	9/12/2015
2	102	9/13/2015

The SQL DDL: Sailors



CREATE TABLE Sailors (
sid INTEGER,
sname CHAR(20),
rating INTEGER,
age FLOAT

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

The SQL DDL: Sailors, Pt. 2



CREATE TABLE Sailors (
sid INTEGER,
sname CHAR(20),
rating INTEGER,
age FLOAT
PRIMARY KEY (sid));

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

The SQL DDL: Primary Keys



CREATE TABLE Sailors (sid INTEGER, sname CHAR(20), rating INTEGER, age FLOAT

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

- Primary Key column(s)
 - Provides a unique "lookup key" for the relation
 - Cannot have any duplicate values
 - Can be made up of >1 column
 - E.g. (firstname, lastname)

The SQL DDL: Boats



CREATE TABLE Sailors (
sid INTEGER,
sname CHAR(20),
rating INTEGER,
age FLOAT

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

CREATE TABLE Boats (
bid INTEGER,
bname CHAR (20),
color CHAR(10),
PRIMARY KEY (bid));

<u>bid</u>	bname	color
101	Nina	red
102	Pinta	blue
103	Santa Maria	red

The SQL DDL: Reserves



CREATE TABLE Sailors (sid INTEGER, sname CHAR(20), rating INTEGER, age FLOAT

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

CREATE TABLE Boats (bid INTEGER, bname CHAR (20), color CHAR(10), PRIMARY KEY (bid));

CREATE TABLE Reserves (
sid INTEGER,
bid INTEGER,
day DATE,
PRIMARY KEY (sid, bid, day);

<u>bid</u>	bname	color
101	Nina	red
102	Pinta	blue
103	Santa Maria	red

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

The SQL DDL: Reserves Pt. 2



CREATE TABLE Sailors (sid INTEGER, sname CHAR(20), rating INTEGER, age FLOAT

<u>sid</u> rating age sname Fred 22 39 Jim 3 8 27 Nancy

CREATE TABLE Boats (
bid INTEGER,
bname CHAR (20),
color CHAR(10),
PRIMARY KEY (bid));

CREATE TABLE Reserves (
sid INTEGER,
bid INTEGER,
day DATE,
PRIMARY KEY (sid, bid, day),
FOREIGN KEY (sid) REFERENCES Sailors,

<u>bid</u>	bname	color
101	Nina	red
102	Pinta	blue
103	Santa Maria	red

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

The SQL DDL: Foreign Keys



CREATE TABLE Sailors (sid INTEGER, sname CHAR(20), rating INTEGER, age FLOAT

CREATE TABLE Boats (
bid INTEGER,
bname CHAR (20),
color CHAR(10),
PRIMARY KEY (bid));

CREATE TABLE Reserves (
sid INTEGER,
bid INTEGER,
day DATE,
PRIMARY KEY (sid. bid

day DATE,
PRIMARY KEY (sid, bid, day),
FOREIGN KEY (sid) REFERENCES Sailors,
FOREIGN KEY (bid) REFERENCES Boats);

		<u> </u>	
<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

4	<u>bid</u>	bname	color
	101	Nina	red
	102	Pinta	blue
	103	Santa Maria	red

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

The SQL DDL: Foreign Keys Pt. 2

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- Foreign key references a table
 - Via the primary key of that table
- Need not share the name of the referenced primary key

sid	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

CREATE TABLE Reserves (
sid INTEGER,	
bid INTEGER,	
day DATE,	
PRIMARY KEY (sid, bid, day), FOREIGN KEY (sid) REFERENCES	
FOREIGN KEY (sid) REFERENCES	3
Sailors,	
FOREIGN KEY (bid)	
FOREIGN KEY (bid) REFERENCES Boats);	

4	<u>bid</u>	bname	color
	101	Nina	red
$/\! $	102	Pinta	blue
	103	Santa Maria	red

<u>sid</u>	<u>bid</u>	<u>day</u>	
1	102	9/12	
2	102	9/13	

The SQL DML



Find all 27-year-old sailors:

SELECT *
FROM Sailors AS S
WHERE S.age=27;

• To find just names and rating, replace the first line to:

SELECT S.sname, S.rating

Sailors

<u>sid</u>	sname	rating	age
1	Fred	7	22
2	Jim	2	39
3	Nancy	8	27

Content Break 3



Basic Single-Table Queries



SELECT [DISTINCT] < column expression list>
 FROM < single table>
 [WHERE < predicate>]

- Simplest version is straightforward
 - Produce all tuples in the table that satisfy the predicate
 - Output the expressions in the SELECT list
 - Expression can be a column reference, or an arithmetic expression over column refs

SELECT DISTINCT



SELECT DISTINCT S.name, S.gpa **FROM** students S **WHERE** S.dept = 'CS'

- DISTINCT specifies removal of duplicate rows before output
- Can refer to the students table as "S", this is called an alias

ORDER BY



SELECT S.name, S.gpa, S.age*2 AS a2
 FROM Students S
 WHERE S.dept = 'CS'
 ORDER BY S.gpa, S.name, a2;

- ORDER BY clause specifies output to be sorted
 - Lexicographic ordering
- Obviously must refer to columns in the output
 - Note the AS clause for naming output columns!

ORDER BY, Pt. 2



SELECT S.name, S.gpa, S.age*2 AS a2
 FROM Students S
 WHERE S.dept = 'CS'
 ORDER BY S.gpa DESC, S.name ASC, a2;

- Ascending order by default, but can be overridden
 - DESC flag for descending, ASC for ascending
 - Can mix and match, lexicographically

LIMIT



SELECT S.name, S.gpa, S.age*2 AS a2
 FROM Students S
 WHERE S.dept = 'CS'
 ORDER BY S.gpa DESC, S.name ASC, a2;
 LIMIT 3;

- Only produces the first <integer> output rows
- Typically used with ORDER BY
 - Otherwise the output is non-deterministic
 - Not a "pure" declarative construct in that case output set depends on algorithm for query processing

Aggregates



SELECT [DISTINCT] AVG(S.gpa)
 FROM Students S
 WHERE S.dept = 'CS'

- Before producing output, compute a summary (a.k.a. an aggregate) of some arithmetic expression
- Produces 1 row of output
 - with one column in this case
- Other aggregates: SUM, COUNT, MAX, MIN

GROUP BY



SELECT [DISTINCT] **AVG**(S.gpa), S.dept **FROM** Students S **GROUP BY** S.dept

- Partition table into groups with same GROUP BY column values
 - Can group by a list of columns
- Produce an aggregate result per group
 - Cardinality of output = # of distinct group values
- Note: can put grouping columns in SELECT list

HAVING



SELECT [DISTINCT] AVG(S.gpa), S.dept FROM Students S GROUP BY S.dept HAVING COUNT(*) > 2

- The HAVING predicate filters groups
- HAVING is applied after grouping and aggregation
 - Hence can contain anything that could go in the SELECT list
 - I.e. aggs or GROUP BY columns
- HAVING can only be used in aggregate queries
- It's an optional clause

Putting it all together

SELECT S.dept, AVG(S.gpa), COUNT(*)
FROM Students S
WHERE S.gender = 'F'
GROUP BY S.dept
HAVING COUNT(*) >= 2
ORDER BY S.dept;



Content Break 4



DISTINCT Aggregates

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Are these the same or different?

SELECT COUNT(DISTINCT S.name)
FROM Students S
WHERE S.dept = 'CS';

SELECT DISTINCT COUNT(S.name) FROM Students S WHERE S.dept = 'CS';

What Is This Asking For?

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SELECT S.name, AVG(S.gpa) FROM Students S GROUP BY S.dept;

Content Break 5



SQL DML: General Single-Table Queries



SELECT [DISTINCT] < column expression list>
 FROM < single table>
 [WHERE < predicate>]
 [GROUP BY < column list>
 [HAVING < predicate>]]
 [ORDER BY < column list>]
 [LIMIT < integer>];

Summary



- Relational model has well-defined query semantics
- Modern SQL extends "pure" relational model (some extra goodies for duplicate row, non-atomic types... more in next lecture)
- Typically, many ways to write a query
 - DBMS figures out a fast way to execute a query, regardless of how it is written.