

1. In Lecture 8 (using a somewhat different approach than in the reading), we derived an expression for the signed distance  $r = d(\mathbf{H}, \underline{x})$  between an arbitrary point  $\underline{x}$  and a hyperplane  $\mathbf{H}$  given by  $g(\underline{x}) = w_0 + \underline{w}^T \underline{x} = 0$ , all in *nonaugmented* feature space. This question explores this topic further.

- (a) Derive, an expression for the signed distance  $r = d(\mathbf{H}, \underline{x}^{(+)})$  between an arbitrary point  $\underline{x}^{(+)}$  and a hyperplane  $\mathbf{H}$ :  $g(\underline{x}^{(+)}) = \underline{w}^{(+T)} \underline{x}^{(+)} = 0$  in *augmented* feature space. Set up the sign of your distance so that  $\underline{w}^{(+)}$  points to the positive-distance side of  $\mathbf{H}$ .
- (b) In *weight* space, using augmented quantities, derive an expression for the signed distance  $r_{\underline{w}}$  between an arbitrary point  $\underline{w}^{(+)}$  and a hyperplane  $g(\underline{w}^{(+)}) = \underline{w}^{(+T)} \underline{x}^{(+)} = 0$ . Set up the sign of your distance so that  $\underline{x}^{(+)}$  points to the positive-distance side of the hyperplane.

2. For a 2-class learning problem with one feature, you are given four training data points (in augmented space):

$$\underline{x}_1^{(1)} = (1, -2); \underline{x}_2^{(1)} = (1, 1); \underline{x}_1^{(2)} = (1, 3); \underline{x}_2^{(2)} = (1, 5)$$

- (a) Plot the data points in 2D feature space. Draw a linear decision boundary  $\mathbf{H}$  that correctly classifies them, showing which side is positive. Give the weight vector  $\underline{w}$  that corresponds to your  $\mathbf{H}$ .
- (b) Plot the data points  $\underline{x}_n$  as lines in 2D weight space, showing the positive side ( $g_{\underline{x}_n}(\underline{w}) > 0$ ) of each with a small arrow. Also show the side of each line that corresponds to correct classification with shading or cross-hatching (similar to plot in Lecture 6 notes, p. 9). Show the final solution region.
- (c) Plot the weight vector  $\underline{w}$  of  $\mathbf{H}$  from part (a) as a point in weight space. Is  $\underline{w}$  in the solution region?

For parts (d)-(f) below, you will use the same dataset of points  $\underline{x}_n$  given above, but you will first reflect them and use the reflected data points  $z_n \underline{x}_n$  for your plots.

- (d) Plot the reflected data points  $z_n \underline{x}_n$  in 2D feature space. Draw the linear decision boundary  $\mathbf{H}$  from your part (a) and its weight vector. Does  $\mathbf{H}$  correctly classify all the data points?

**Hint:** write down the general condition for correct classification of reflected data points.

- (e) Plot the data points  $z_n \underline{x}_n$  as lines in 2D weight space, showing the positive side of each with a small arrow. Also show the side of each line that corresponds to correct classification with shading or cross-hatching. (**Same hint** applies as for part (d).) Show the final solution region.
- (f) Also, plot the weight vector  $\underline{w}$  of  $\mathbf{H}$  from part (a) as a point in your weight space plot of part (e). Is  $\underline{w}$  in the solution region of part (e)?

3. This problem uses the notation we used in Lecture 5, and  $m$  is a positive integer. For the following computational complexity:

$$p(m) = 2m + 1000$$

- (a) Is  $p(m) = O(m)$ ?

If yes, prove your answer by letting  $a=1$ , and solve for what  $m_0$  we have  $m \geq p(m) \quad \forall m \geq m_0$ . If you need a larger  $a$ , then state what value of  $a$  will work. Find the smallest such integer  $m_0$  for the value of  $a$  you used.

If no, justify why not.

- (b) Is  $p(m) = \Omega(m)$ ?

If yes, prove your answer by letting  $b=1$ , and solve for what  $m_1$  we have  $m \leq p(m) \quad \forall m \geq m_1$ . If you need a smaller  $b$ , then state what value of  $b$  will work. Find the smallest such integer  $m_1$  for the value of  $b$  you used.

If no, justify why not.

- (c) Is  $p(m) = \Theta(m)$ ?

Justify your answer.

4. This problem also uses the notation of Lecture 5.

- (a) Suppose we have a function  $p(m)$  that can be expressed as:

$$p(m) = p_1(m) + p_2(m) + p_3(m)$$

and we have:

$$p_k(m) = O(q_k(m)), \quad k = 1, 2, 3 \tag{i}$$

Prove that:

$$p(m) = O(q_1(m) + q_2(m) + q_3(m)) \tag{ii}$$

**Hints:** (i) Use the definition of big-O.

- (ii) If you find the problem statement unclear or confusing, try looking at the example in the appendix below.
- (b) Is a similar statement to (a) true for  $\Omega(\cdot)$ ? (That is, if you replace each  $O(\cdot)$  in part (a) with  $\Omega(\cdot)$ , would the last equation be true?) Justify your answer.
5. Consider the following computational complexity:
- $$p(m) = m^2 \log_2 m + \frac{2m^3}{\log_2 m} + (\log_2 m)^3$$
- (a) Is  $p(m) = O(m^3)$ ? If yes, justify it by showing it satisfies the definition. If no, justify why not.
- Hint:**  $p(m)$  is the sum of 3 terms. Try setting  $a=1$ , and check each term individually to see if it is  $O(m^3)$  (if it is, then give the value of each  $m_0$ ). Then use the result of Problem 4(a).
- (b) Is  $p(m) = \Omega(m^3)$ ? Justify your answer.
6. For a nearest-means classifier, with  $C = 2$  classes (held constant),  $\frac{1}{2}N$  data points in each class (assume  $N$  is even), and  $D$  features, find the computational time complexity and space complexity, for the operations given in (a) and (b) below.

For all parts of this problem, assume a completely serial computer. Your answer may be in the form of a big- $O$  upper bound; express the big- $O$  bound in simplest form, without making it looser (higher) than necessary. As part of your answer to each part, please:

- (i) Give a brief statement of the algorithm you are analyzing
- (ii) Show your reasoning in calculating the complexity
- (a) Computing each mean vector from the training data (the training phase)
- (b) Classifying  $M$  data points, given the mean vectors (the classification phase)
- (c), (d) Repeat parts (a) and (b), except for a  $C$ -class problem, in which  $C$  is a variable.

**Hint for (d):** you may use without proof that finding the minimum of  $C$  values that are unsorted can be done in  $O(C)$  time and  $O(1)$  space.

**Appendix** (example) is on the next page.

#### Appendix - Example (relates to Problem 4)

Suppose we want to find and prove the (tightest) asymptotic upper bound of  $p(m)$ , with:

$$p(m) = 3m^3 + 100m^2 \log_2 m + 0.1(2^m)$$

Applying the definition directly to  $p(m)$  (especially to prove your bound, including finding  $m_0$ ) might be difficult. Instead, you could use the result of Problem 4a, to apply the big-O bound to each term independently:

$$3m^3 = O(m^3)$$

$$100m^2 \log_2 m = O(m^2 \log m)$$

$$0.1(2^m) = O(2^m)$$

Then using Problem 4a equation (ii), we can conclude:

$$\begin{aligned} 3m^3 + 100m^2 \log_2 m + 0.1(2^m) &= O(m^3 + m^2 \log m + 2^m) \\ &= O(2^m) \end{aligned}$$