

✧ Discrete Log (x):

✧ Naive :  $O(p)$

✧ Block size : b

$$\rightarrow a^x \equiv n \pmod{p} \rightarrow x \in [0, p-1] \Rightarrow 1 \leq i \leq \left\lceil \frac{p-1}{b} \right\rceil$$

$$x = ib - j$$

$$\rightarrow i = \left\lceil \frac{x}{b} \right\rceil \rightarrow j = \left\lceil \frac{x}{b} \right\rceil b - x = (b - (x \% b)) \% b \Rightarrow 0 \leq j \leq b-1$$

$$\rightarrow a^{ib} \equiv na^j \pmod{p}$$

→  $P/b$        $b$       values

- Sort these Values
- Iterate & do Bin Search

✧ Giant Step →  $i$

✧ Baby Step →  $j$

$$\star T.C. \rightarrow O\left(\frac{P}{b}\right) + O(b) + O\left(\frac{P}{b} \log \frac{P}{b} + b \log \frac{P}{b}\right)$$

$$\rightarrow O\left(\frac{P}{b} \log \frac{P}{b} + b \log \frac{P}{b}\right)$$

$$\rightarrow \text{Use Hash - Map} \rightarrow O\left(\frac{P}{b} + b\right)$$

$$\Rightarrow b = \sqrt{p} \Rightarrow \underline{\underline{T.C. \rightarrow O(\sqrt{p})}}$$

- ✧  $b$  : block size
- ✧  $n$  : argument
- ✧  $a$  : base
- ✧  $p$  : prime
- ✧  $j$  : Base power
- ✧  $x$  : Exponent

✧ CODE SNIPPET

```
int discrete_log(int arg, int base, int p)
{
    int block = sqrt(p - 1);
    vector<pii> arg_base_power(block); → (naj, j)

    int base_power = 1; → aj
    for (int j = 0; j < block; j++)
    {
        arg_base_power[j] = mp((1ll * arg * base_power) % p, j);
        base_power = (1ll * base * base_power) % p;
    }

    sort(all(arg_base_power));
    int base_power_block = base_power; → ab

    int exponent = -1;
    for (int i = 1; i <= (((p - 1) + block - 1) / block); i++)
    {
        auto temp = lb(all(arg_base_power), mp(base_power_block, 0));
        if (temp != arg_base_power.end()
            && temp->first == base_power_block)
        {
            exponent = (i * block - temp->second) % (p - 1);
            break;
        }
        base_power_block = (1ll * base_power * base_power_block) % p; → aib
    }

    return exponent;
}
```