

Common Data Structure Operations

Data Structure	Time Complexity								Space Complexity
	Average				Worst				Worst
	Access	Search	Insertion	Deletion	Access	Search	Insertion	Deletion	
<u>Array</u>	$\theta(1)$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$
<u>Stack</u>	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$
<u>Queue</u>	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$
<u>Singly-Linked List</u>	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$
<u>Doubly-Linked List</u>	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$	$\theta(n)$	$\theta(1)$	$\theta(1)$	$\theta(n)$
<u>Skip List</u>	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n \log(n))$
<u>Hash Table</u>	N/A	$\theta(1)$	$\theta(1)$	$\theta(1)$	N/A	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$
<u>Binary Search Tree</u>	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$
<u>Cartesian Tree</u>	N/A	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	N/A	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$
<u>B-Tree</u>	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$
<u>Red-Black Tree</u>	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$
<u>Splay Tree</u>	N/A	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	N/A	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$
<u>AVL Tree</u>	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$
<u>KD Tree</u>	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(\log(n))$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$	$\theta(n)$

Loops:

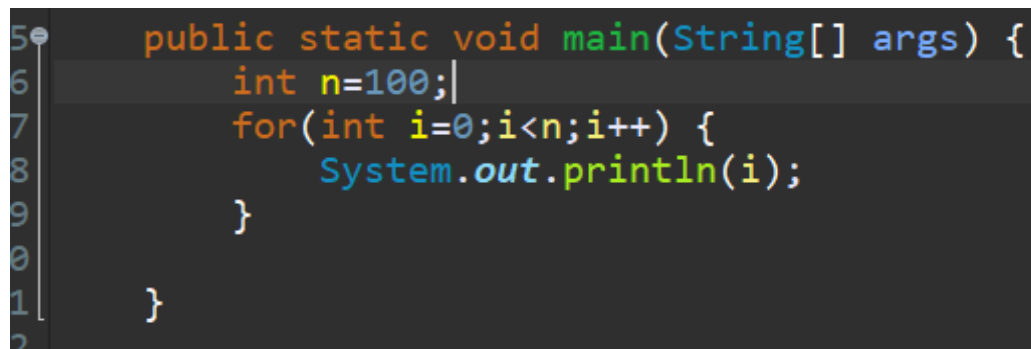
1. Simple Loop: $O(n)$

```
for (i = 1 to n){           // n  
  
    x = y + z;              // constant time  
  
}
```

Time complexity: $O(n)$

Note: constant time can be neglected

Example: $O(n)$

A screenshot of a code editor with a dark background. It shows a Java main method with a loop. The code is: public static void main(String[] args) { int n=100; for(int i=0;i<n;i++) { System.out.println(i); } }. The lines are numbered 5 through 11 on the left margin.

```
5 public static void main(String[] args) {  
6     int n=100;  
7     for(int i=0;i<n;i++) {  
8         System.out.println(i);  
9     }  
10  
11 }
```

2. Nested Loop: $O(n^2)$

```
for (i = 1 to n){           // n
```

```
for (j = 1 to n){           // n
x = y + z;                 // constant time
}}
```

Time complexity: $O(n^2)$

Example 1: $O(n^2)$

```
4
5 public static void main(String[] args) {
6     int n=100;
7     for(int i=0;i<n;i++) {
8         for(int j=0;j<n;j++) {
9             System.out.println(i);
10        }
11    }
12 }
```

Example 2: $O(n^2)$ Consecutive Statements.

Here the time complexity of the first loop is $O(n)$ and the nested loop is $O(n^2)$. so we will take whichever is higher into the consideration.

```

4
5 public static void main(String[] args) {
6     int n=100;
7     for(int m=0;m<n;m++) {
8         System.out.println(m);
9     }
10    for(int i=0;i<n;i++) {
11        for(int j=0;j<n;j++) {
12            System.out.println(i);
13        }
14    }
15 }
16 }

```

Example 3: $O(n)$ with if-else loop.

The time complexity of the if statement is $O(1)$ and else is $O(n)$. as $O(n) > O(1)$ time complexity of this program is $O(n)$

```

4
5 public static void main(String[] args) {
6     int n = 100;
7     int length = 10;
8     if (length < 10) {
9         System.out.println("not enough length");
10    } else {
11
12        for (int m = 0; m < n; m++) {
13            System.out.println(m);
14        }
15    }
16 }
17 }
18 }

```

Example 4: $O(\log n)$ logarithmic complexity

```
5 public static void main(String[] args) {  
6     int n = 100;  
7     for (int m = n; m > 1; m = m / 2) {  
8         System.out.println(m);  
9     }  
10 }  
11
```

+-----+-----+

Iteration	Value of M
-----------	------------

+-----+-----+

1	n
---	---

2	$n/2$
---	-------

3	$n/2^2$
---	---------

4	$n/2^3$
---	---------

| . | . |

| . | . |

| k | $n/2^k$ |

+-----+-----+

loop will run k times so that $m > 1$.

$m > 1$; putting value of m for as $n/2^k$

$$n/2^k > 1$$

$$2^k > n$$

$$k > \log n$$

Example 5: $O(\sqrt{n})$

```

5 public static void main(String[] args) {
6     int n=100;
7     int i=1,s=1;
8     while(s<=n) {
9         i++;
0         s=s+i;
1         System.out.println("*");
2     }
3 }
4
5 }
6

```

```

+-----+-----+-----+

```

Iteration	value of i	value of s=s+i
-----------	------------	----------------

```

+-----+-----+-----+

```

1	1	1+1
---	---	-----

2	2	1+1+2
---	---	-------

3	3	1+1+2+3
---	---	---------

4	4	1+1+2+3+4
---	---	-----------

| . | . | . |

| . | . | . |

| k | k | 1+1+2+3+4+..+k |

+-----+-----+-----+

number of time the loop will run is k

so $1+2+3+..+k > n$ (that's when loop breaks)

$$k(k+1)/2 > n$$

$$k^2 + k > n$$

$$k^2 > n \text{ (ignoring } k \text{ as } k^2 > k)$$

$$k > \sqrt{n}$$

so time complexity of program is \sqrt{n}

Example 6: $O(\sqrt{n})$

```
4
5 public static void main(String[] args) {
6     int n=100;
7     for(int i=1;i*i<=n;i++) {
8         System.out.println("*");
9     }
10 }
11
```

let's say loop runs k times

so $k^2 > n$ (to break the condition)

$$k > \sqrt{n}$$

so time complexity is \sqrt{n}

Example 7: $O(n^2 \log n)$

```

4
5 public static void main(String[] args) {
6     int n = 100;
7     for (int i = n / 2; i <= n; i++) {
8         for (int j = 1; j + n / 2 < n; j++) {
9             for (int k = 1; k <= n; k = k * 2) {
10                 System.out.println("*");
11             }
12         }
13     }
14 }
15
16 }
17

```

first loop will run $n/2$ times

The second loop will run $n/2$ times as $j+n/2 < n$ is the condition. so $n/2$ is already added to the iterator so only $n/2$ times loop will run.

third loop : $2^k > n$ (iteration wise $2^0, 2^1, 2^2, \dots, 2^k$)

$k > \log n$

so time complexity is $n/2 * n/2 * \log n$

so $n^2 \log n$ is the time complexity.

Example 8: $O(n \log^2 n)$

```

public static void main(String[] args) {
    int n = 100;
    for (int i = n / 2; i <= n; i++) {
        for (int j = 1; j < n; j=j*2) {
            for (int k = 1; k <= n; k = k * 2) {
                System.out.println("*");
            }
        }
    }
}

```

first loop will run $n/2$ times

second and third loop as per above example will run $\log n$ times

so time complexity = $n/2 * \log n * \log n = O(n \log^2 n)$

Example 9: $O(n)$ with break statement

```

6      public static void main(String[] args) {
7          int n = 100;
8          for (int i = 1; i <= n; i++) {
9              for (int j = 1; j < n; j++) {
10                 System.out.println("*");
11                 break;
12             }
13         }
14     }

```

The first loop will run N times, the second will break out after every first iteration. so it will run 1 time

so time complexity is $O(n)$ instead of $O(n^2)$

Example 10: $O(n \log n)$

```

6      public static void main(String[] args) {
7          int n = 100;
8          for (int i = 1; i <= n; i++) {
9              for (int j = 1; j < n; j=j+i) {
10                 System.out.println("*");
11                 break;
12             }
13         }
14     }
15

```

i=1	i=2	i=3	i=n
j=j+i=j+1	j=j+i=j+2	j=j+i=j+3		
j=1	j=1	j=1		
j=2	j=3	j=4		
j=3	j=5	j=7		
j=4	j=7	j=10		
.				
.				
n	n/2	n/3	1

Value of j for
particular value of i



so $n + n/2 + n/3 + \dots + 1$ times total loop will run

$$n(1 + 1/2 + 1/3 + 1/4) = n \log n$$

so time complexity is $n \log n$.

Example 11: $O(n^2)$

```
5 public static void main(String[] args) {  
6     int n = 100;  
7     for (int i = 1; i <= n/3; i++) {  
8         for (int j = 1; j < n; j=j+4) {  
9             System.out.println("*");  
10            break;  
11        }  
12    }  
13 }  
14 }
```

outer loop will run $n/3$ times

inner loop will run $n/4$ times

so total time complexity is $(n/3) * (n/4) = n^2/12 = O(n^2)$

Example 12: $O(\log^2 n)$

```

6      public static void main(String[] args) {
7          int n = 100;
8          int i = 1, j = 0;
9          while (i < n) {
10             j = n;
11             while (j > 0) {
12                 j = j / 2;
13             }
14             i = i * 2;
15         }
16     }
17 }
18

```

outer loop will run $\log n$ times [$2^k > n$]

inner loop will run $\log n$ times [$n/2^k > 1$]

so total time complexity = $\log n * \log n = \log^2 n$

Example 13: $O(n^5)$

```

6      public static void main(String[] args) {
7          int n = 100;
8          for (int i = 0; i < n; i++) {
9              for (int j = 1; j < i * i; j++) {
10                 if (i % j == 0) {
11                     for (int k = 0; k < j; k++) {
12                         System.out.println("*");
13                     }
14                 }
15             }
16         }
17     }
18 }

```

first loop will run n times

second loop will run n^2 times [$1, 2^2, 3^2, \dots, n^2$]

The third loop will run n times. approx as there is no clear pattern.

total time complexity = $n * n * n^2 = O(n^4)$

Example 14: $O(n)$

```
int count = 0;
for (int i = N; i > 0; i /= 2) {
    for (int j = 0; j < i; j++) {
        count += 1;
    }
}
```

loop will run $N + N/2 + N/4 + \dots + N/2^n$

$= N(1 + 1/2 + 1/4 + \dots + 1/2^n)$

$$=N(1+1)$$

$$=2N = O(n)$$

sum of the series for reference

$$[S_n = 1/2 + 1/4 + \dots + 1/2^n$$

$$2S_n = 2/2 + 2/4 + 2/8 + \dots + 2/2^n$$

$$2S_n = 1 + 1/2 + 1/4 + \dots + 1/2^{n-1}$$

$$2S_n = 1 + S_n - 1/2^n$$

$$S_n = 1 - 1/2^n$$

$$S_n = 1 \text{ (when } n = \infty \text{)}$$
