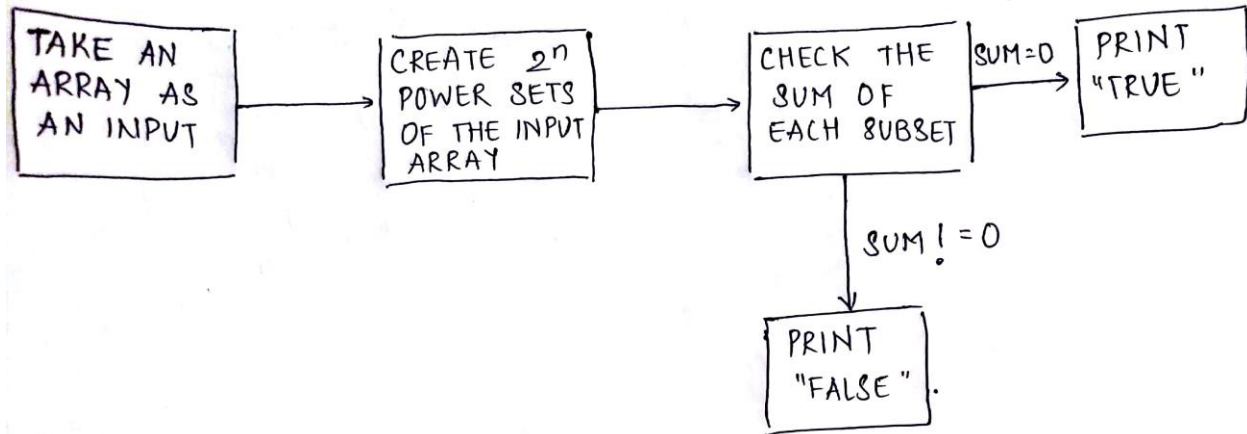


### Report: Assignment 3 –Zero.java



The program selects a bunch of elements from the given set which sums up to zero. To solve this problem, we have written a function `checkZero` which accepts start and end pointers to iterate on a static array. The algorithm starts with making  $2^n$  power sets of the input array. While these combinations are made, we store the sum of every element added to the combination in a variable called as `sum`. Also, we keep a track of all the elements added to combination in an array called as `temp`. Later, if we get a set of elements which sum up to zero, `temp` array is used to show the user which specific elements were used to make sum zero.

### Design changes during implementation

The current algorithm is of  $O(2^n)$ . We tried to reduce the complexity to  $O(2^{n/2})$  by making few modifications to the current algorithm, which were:

- 1) Split the array in mid-way. Say now we have array A and array B.
- 2) Find out all the power set of all the possible sums for each of the element in the arrays (Each will run for  $O(2^{n/2})$ ). Assume we have now 2 new generated arrays, `A_Pow` and `B_Pow`.
- 3) Moreover, as you make these power set, maintain 2 new arrays which will store the combination of the elements used from the set A/B for making the sum. Say its named as `A_Comb` and `B_Comb`. These will be used to show which elements were used to make the sum 0.
- 4) Now sort `A_Pow` and `B_Pow` using Bubble sort and while swapping an element from `A_Pow` or `B_Pow`, swap the respective combination of elements from `A_Comb` or `B_Comb`. Here, Bubble sort will run for  $O(n^2)$ .
- 5) Given the sort list, now we need to check if the element in `A_Pow` and element in the `B_Pow` sums up to 0 in time  $O(2^{n/2})$ .
- 6) To do that, scan the `A_Pow` in decreasing order and `B_Pow` in increasing order.
- 7) Whenever the sum of the current element in the `A_Pow` and the current element in the `B_Pow` is more than 0, the algorithm moves to the next element in the `A_Pow`.

- 8) If it is less than 0, the algorithm moves to the next element in the B\_Pow. If two elements with sum 0 are found, then stops.

The only good thing about this algorithm is that it's  $O(2^{n/2})$  but the overhead of finding 2 power set's, then maintaining an array of combination of those sums, then sorting both the arrays with bubble sort and then iterating the final sorted arrays was too huge. And hence we thought of going with our current approach.