

String Pattern Matching

Reference:

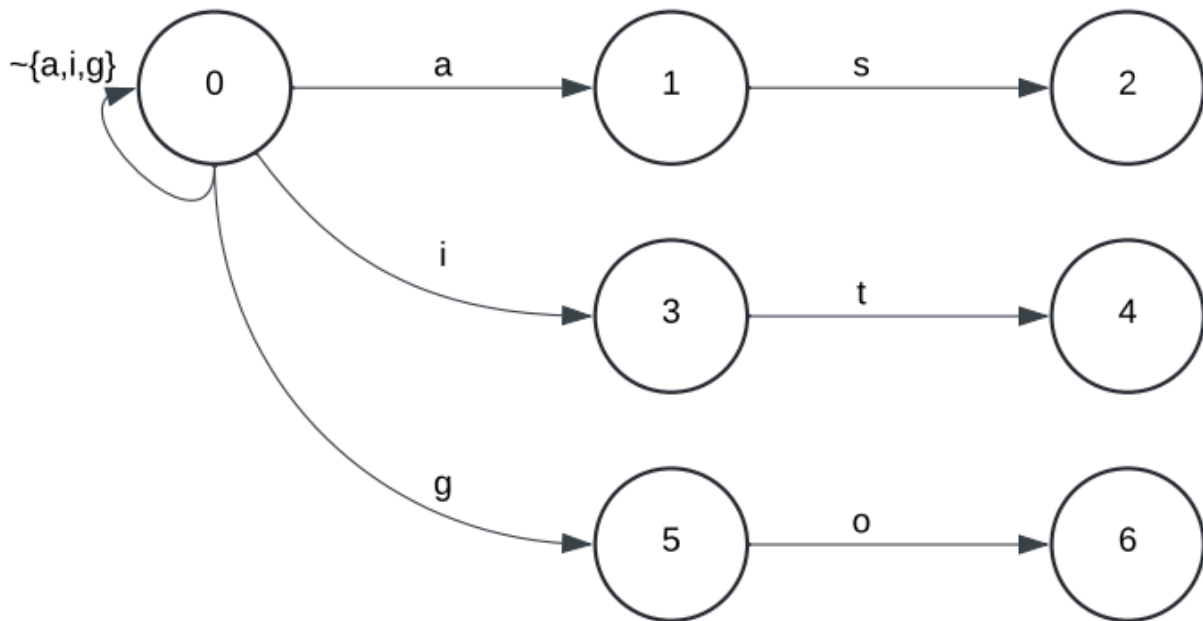
Aho, A. V., & Corasick, M. J. (1975). Efficient string matching: An aid to bibliographic search. *Communications of the ACM*, 18(6), 333–340. <https://doi.org/10.1145/360825.360855>

1) Chosen a simple set of 3 strings : {"as", "it", "go"}

Pattern Matching Machine consists of Goto, failure and output functions as shown below.

Fig1: Pattern Matching Machine

(a) Goto Function



(b) Failure Function

i	1	2	3	4	5	6
f(i)	0	0	0	0	0	0

(c) Output Function

i	output(i)
2	as
4	it
6	go

Computation of goto function:

All our pattern matching machines have the property that $g(0, \sim r) \neq \text{fail}$ for all input symbols r . We shall see that this property of the goto function on state 0 ensures that one input symbol will be processed by the machine in every machine cycle.

goto(0, 'a') = 1, goto(1, 's') = 2,
 goto(0, 'i') = 3, goto(3, 't') = 4,
 goto(0, 'g') = 5, goto(5, 'o') = 6

Computation of failure function:

The failure function is constructed from the goto function. Let us define the depth of a state s in the goto graph as the length of the shortest path from the start state to s . Thus in above Figure 1 (a), the start state is of depth 0, states 1, 3 and 5 are of depth 1, states 2, 4, and 6 are of depth 2.

Computed the failure function for all states of depth 1, then for all states of depth 2 until the failure function has been computed for all states (except state 0 for which the failure function is not defined).

The algorithm to compute the failure function f at a state is conceptually quite simple. We make $f(s) = 0$ for all states of depth 1. Now suppose f has been computed for all states of depth less than d . The failure function for the states of depth d is computed from the failure function for the states of depth less than d . The states of depth d can be determined from the nonfail values of the goto function of the states of depth $d-1$. Specifically, to compute the failure function for the states of depth d , we consider each state r of depth $d-1$ and perform the following actions.

1. If $g(r, a) = \text{fail}$ for all a , do nothing.
2. Otherwise, for each symbol a such that $g(r, a) = s$, do the following:
 - (b) Execute the statement $\text{state} = f(\text{state})$ zero or more times, until a value for state is obtained such that $g(\text{state}, a) \neq \text{fail}$. (Note that since $g(0, a) \neq \text{fail}$ for all a , such a state will always be found.)
 - (c) Set $f(s) = g(\text{state}, a)$.

For the states of depth 1:

failure function for 1, 3 and 5 which are of the states of depth 1.

$f(1) = f(3) = f(5) = 0$ since depth is 1

For the states of depth 2:

failure function for 2, 4 and 6 which are of the states of depth 2.

$f(2)$:

Set $\text{state} = f(1)$

$f(1) = 0$ since $g(0, 's') = \text{fail}$

Hence, $f(2) = 0$

f(4):
 Set state=f(3)
 As f(3)=0
 Then g(0,'t')= fail
 Hence, f(4) = 0.

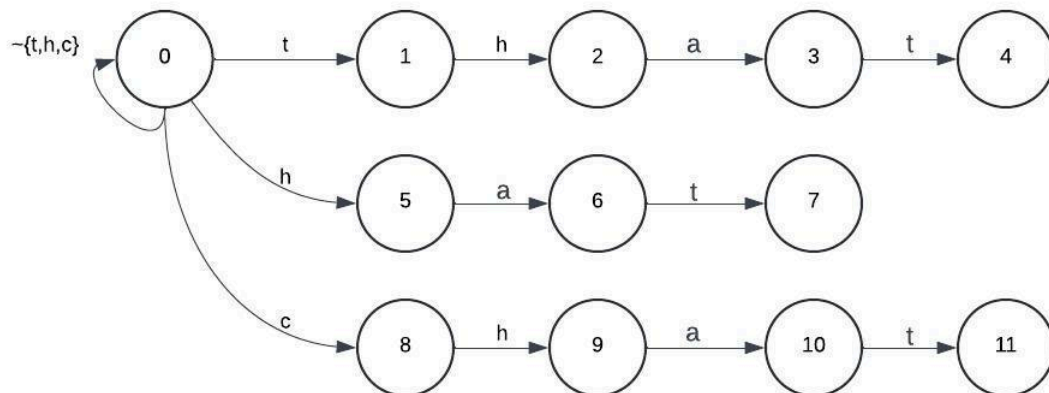
f(6):
 Set state = f(5)
 As f(5)= 0
 Then g(0, 'o') =fail
 Hence, f(6) = 0.

Hence,
 The output “as” is associated with state 2
 The output “it” is associated with state 4
 The output “go” is associated with state 6

The above set of strings resulted in the failure states of 0. Considering the below set of strings to illustrate non-zero failure states of output.

{“that”, “hat”, “chat”}

(a) Goto Function



(b) Failure Function

i	1	2	3	4	5	6	7	8	9	10	11
f(i)	0	5	6	7	0	0	1	0	5	6	7

(c) Output Function

i1	output(i)
4	{that}
7	{hat}
11	{chat, hat}

Computation of goto function:

goto(0, 't') = 1, goto(1, 'h') = 2, goto(2, 'a') = 3, goto(3, 't') = 4
goto(0, 'h') = 5, goto(5, 'a') = 6, goto(6, 't') = 7
goto(0, 'c') = 8, goto(8, 'h') = 9, goto(9, 'a') = 10, goto(10, 't') = 11

Computation of failure function:

Using the failure function algorithm:

For all the input symbols that can be reached directly through root ,we set $f(s)=0$

$f(1) = 0$, $f(5) = 0$, and $f(8) = 0$

$f(2)$:

state = $f(1) = 0$.

goto(0, 'h') = 5, set $f(2) = 5$

Hence, $f(2) = 5$

$f(3)$:

state = $f(2) = 5$.

goto(5, 'a') = 6, set $f(3) = 6$.

Hence, $f(3)=6$

$f(4)$:

state = $f(3) = 6$.

As goto(6, 't') = 7, set $f(4) = 7$

Hence, $f(4)=7$

$f(6)$:

state = $f(5) = 0$.

Since goto(0, 'a') = fail

Hence, $f(6) = 0$

$f(7)$:

state = $f(6) = 0$.

Since goto(0, 't') = 1, set $f(7) = 1$.

Hence, $f(7) = 1$.

$f(9)$:

$\text{state} = f(8) = 0$.

Since $\text{goto}(0, 'h') = 5$, set $f(9) = 5$.

Hence, $f(9) = 5$.

$f(10)$:

$\text{state} = f(9) = 5$.

As $\text{goto}(5, 'a') = 6$, set $f(10) = 6$

Hence, $f(10) = 6$

$f(11)$:

$\text{state} = f(10) = 6$.

As $\text{goto}(6, 't') = 7$, set $f(11) = 7$

Hence, $f(11) = 7$