

**Pair Sum** ↗ ascending

Return pair in sorted array with target sum.

2	7	11	15
0	1	2	3

target = 913

$O(n^2)$

① Brute  
 ↪ pairs → target (Pair Sum)

```

for (int i = 0; i < n; i++) {
  for (int j = i + 1; j < n; j++) {
    (i, j)
    if (arr[i] + arr[j] == target) {
      ans.pb(i);
      ans.pb(j);
      return ans;
    }
  }
}
  
```

**Print Screen**

code.cpp

```

1 pairSum(vector<int>, int)
2 using namespace std;
3
4
5 vector<int> pairSum(vector<int> nums, int target) {
6     vector<int> ans;
7     int n = nums.size();
8
9     for (int i = 0; i < n; i++) {
10        for (int j = i + 1; j < n; j++) {
11            if (nums[i] + nums[j] == target) {
12                ans.push_back(i);
13                ans.push_back(j);
14                return ans;
15            }
16        }
17    }
18
19    return ans;
20
21
22 int main() {
23     vector<int> nums = {2, 7, 11, 15};
24     int target = 9;
  
```

PORTS PROBLEMS DEBUG CONSOLE OUTPUT TERMINAL

apnacollege@Shradha DSASeries %

**Print Screen**

```

code.cpp
12     ans.push_back(i);
13     ans.push_back(j);
14     return ans;
15 }
16
17 }
18
19 return ans;
20 }
21
22 int main() {
23     vector<int> nums = {2, 7, 11, 15};
24     int target = 13;
25
26     vector<int> ans = pairSum(nums, target);
27     cout << ans[0] << " " << ans[1] << endl;
28     return 0;
29 }
30
31
32
PORTS  PROBLEMS  DEBUG CONSOLE  OUTPUT  TERMINAL
● apnacollege@Shradha DSAseries % g++ -std=c++11 code.cpp && ./a.out
0, 1
● apnacollege@Shradha DSAseries % g++ -std=c++11 code.cpp && ./a.out
0, 2
○ apnacollege@Shradha DSAseries %
Print Screen

```

## Pair Sum

Return pair in sorted array with target sum.

2	7	11	15
0	1	2	3

↑  
i

↑  
j

① PS > tar    j--  
② PS < tar    i++  
③ PS = tar    ans ⇒ (i, j)

Print Screen

tar = 26  
PS = 26

① Brute →  $O(n^2)$   
② Optimized  
2 pointer approach

## Pair Sum

Return pair in sorted array with target sum.

2	7	11	15
0	1	2	3

① PS > tar    j--  
② PS < tar    i++  
③ PS = tar    ans ⇒ (i, j)

Print Screen

tar = 26  
PS = 26

$O(n)$

i = 0, j = n-1  
while (i < j) {  
    PS = arr[i] + arr[j]  
    if (PS > tar) j--  
    else if (PS < tar) i++  
    else return (i, j)  
}



code.cpp

```

1 pairSum(vector<int>, int)
2 #include <vector>
3 using namespace std;
4
5 vector<int> pairSum(vector<int> nums, int target) {
6     vector<int> ans;
7     int n = nums.size();
8
9     int i = 0, j = n-1;
10
11     while(i < j) {
12         int pairSum = nums[i] + nums[j];
13         if(pairSum > target) {
14             j--;
15         } else if(pairSum < target) {
16             i++;
17         } else {
18             ans.push_back(i);
19             ans.push_back(j);
20             return ans;
21         }
22     }
23
24     return ans;
25 }
26
27 int main() {
28     vector<int> nums = {2, 7, 11, 15};
29     int target = 13;
30 }

```

PORTS

PROBLEMS

DEBUG CONSOLE

OUTPUT

TERMINAL

```

apnacollege@Shradha DSAseries % g++ -std=c++11 code.cpp && ./a.out
0, 1
apnacollege@Shradha DSAseries % g++ -std=c++11 code.cpp && ./a.out
0, 2
apnacollege@Shradha DSAseries %

```

Print Screen

code.cpp

```

13 if(pairSum > target) {
14     j--;
15 } else if(pairSum < target) {
16     i++;
17 } else {
18     ans.push_back(i);
19     ans.push_back(j);
20     return ans;
21 }
22
23 return ans;
24 }
25
26
27 int main() {
28     vector<int> nums = {2, 7, 11, 15};
29     int target = 26;
30
31     vector<int> ans = pairSum(nums, target);
32     cout << ans[0] << " " << ans[1] << endl;
33     return 0;
34 }
35
36
37

```

PORTS

PROBLEMS

DEBUG CONSOLE

OUTPUT

TERMINAL

```

apnacollege@Shradha DSAseries % g++ -std=c++11 code.cpp && ./a.out
0, 2
apnacollege@Shradha DSAseries % g++ -std=c++11 code.cpp && ./a.out
2, 3
apnacollege@Shradha DSAseries %

```

Print Screen

Majority Element

$m = 5$   
 $MJ > \frac{n}{2}$  times

$\{1, 2, 2, 1, 1\}$   
 $1 \rightarrow 3 \Rightarrow 3 > \frac{5}{2}$

$MJ > 2x$   
 $3 > 2$

- Brute Force
- Optimize
- Moore's

Print Screen

## Majority Element


$n = 5$   $MJ > 2 \times$

$MJ > \frac{n}{2}$  times

① Brute Force  
② Optimize  
③ Moore's

$2 \rightarrow 2$   
①  $\rightarrow 3 > \frac{n}{2}$  ✓

Print Screen



## Majority Element


$n = 5$   $MJ > 2 \times$

$MJ > \frac{n}{2}$  times

① Brute Force  $O(n^2)$   
② Optimize  
③ Moore's

```
for (int val : nums) {
    freq = 0
    for (int el : nums) {
        if (el == val) {
            freq++
        }
    }
    if (freq > n/2) → MJ
}
```

Print Screen



SDE SHEET - Google Sheets x Majority Element - LeetCode x

leetcode.com/problems/majority-element/

### 169. Majority Element

Solved ✓

Easy Topics Companies

Given an array `nums` of size `n`, return the majority element.

The majority element is the element that appears more than  $\lfloor \frac{n}{2} \rfloor$  times. You may assume that the majority element always exists in the array.

**Example 1:**  
Input: `nums = [3,2,3]`  
Output: `3`

**Example 2:**  
Input: `nums = [2,2,1,1,1,2,2]`  
Output: `2`


**Constraints:**  
• `n == nums.length`

```
class Solution {
public:
    int majorityElement(vector<int>& nums) {
        int n = nums.size();
        for (int val : nums) {
            int freq = 0;
            for (int el : nums) {
                if (el == val) {
                    freq++;
                }
            }
            if (freq > n/2) {
                return val;
            }
        }
    }
};
```

Testcase Test Result

Case 1 Case 2 +

nums = [3,2,3] Print Screen





169. Majority Element

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Input: `nums = [3,2,3]`  
Output: `3`

Example 2:  
Input: `nums = [2,2,1,1,2,2]`  
Output: `2`

Constraints:

- `n == nums.length`

```

4      int n = nums.size();
5
6      for(int val : nums) {
7          int freq = 0;
8
9          for(int el : nums) {
10             if(el == val) {
11                 freq++;
12             }
13         }
14         if(freq > n/2) {
15             return val;
16         }
17     }
18     return -1;
19 }
20

```

Testcase Case 1 Print Screen

num =

## Majority Element

{ 1, 2, 2, 1, 1 }

↓

① [1, 1, 1, 2, 2]  $O(n \log n)$

②  $O(n)$

↓  $n=9$

[ 0, 0, 1, 1, 2, 2, 2, 2, 2 ]

↑  $f = 1, 2, 3 > \frac{n}{2}$

Print Screen

① Brute Force  $O(n^2)$

② Optimize  $\rightarrow$  sorting

③ Moore's

## Majority Element

{ 1, 2, 2, 1, 1 }

↓

① [1, 1, 1, 2, 2]  $O(n \log n)$

②  $freq = 1; ans = nums[0]$

for ( $i = 1; i < n; i++$ ) {

if ( $nums[i] == nums[i-1]$ )

freq++

else

freq = 1; ans = nums[i]

}

Print Screen

① Brute Force  $O(n^2)$

② Optimize  $\rightarrow$  sorting

③ Moore's

## Majority Element

[1, 1, 2, 2, 2]

- ① Brute Force  $O(n^2)$
- ② Optimize  $O(n \log n)$
- ③ Moore's

{ 1, 2, 2, 1, 1 }

②  $freq = 1$ ; ans = nums[0]

for (i = 1; i < n; i++) {  
    if (nums[i] == nums[i-1])  
        freq++;  
    else  
        freq = 1; ans = nums[i];  
}

if (freq > n/2)  
    return ans

Print Screen



Screenshot of the LeetCode problem page for "169. Majority Element". The page shows the problem description, examples, and constraints. The code editor is open, showing a C++ solution that sorts the array and then counts the frequency of each element to find the majority element.

**169. Majority Element**

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**Example 1:**  
Input: `nums = [3,2,3]`  
Output: `3`

**Example 2:**  
Input: `nums = [2,2,1,1,1,2,2]`  
Output: `2`

**Constraints:**

- `n == nums.length`
- `1 <= n <= 5 * 104`
- `-109 <= nums[i] <= 109`

**Code:**

```
1 class Solution {
2 public:
3     int majorityElement(vector<int>& nums) {
4         int n = nums.size();
5
6         //sort
7         sort(nums.begin(), nums.end());
8
9         //freq count
10        int freq = 1, ans = nums[0];
11        for(int i=1; i<n; i++) {
12            if(nums[i] == nums[i-1]) {
13                freq++;
14            } else {
15                freq=1;
16                ans = nums[i];
17            }
18
19            if(freq > n/2) {
20                return ans;
21            }
22        }
23    }
24 }
```

**Testcase:** Case 1 Case 2 +

nums = [3,2,3]

Print Screen



Screenshot of the LeetCode problem page for "169. Majority Element". The page shows the problem description, examples, and constraints. The code editor is open, showing a C++ solution that sorts the array and then counts the frequency of each element to find the majority element.

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**Constraints:**

- `n == nums.length`
- `1 <= n <= 5 * 104`
- `-109 <= nums[i] <= 109`

**Code:**

```
9 //freq count
10 int freq = 1, ans = nums[0];
11 for(int i=1; i<n; i++) {
12     if(nums[i] == nums[i-1]) {
13         freq++;
14     } else {
15         freq=1;
16         ans = nums[i];
17     }
18
19     if(freq > n/2) {
20         return ans;
21     }
22 }
```

**Testcase:** Case 1 Case 2 +

nums = [3,2,3]

Print Screen





169. Majority Element

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Input: `nums = [3,2,3]`  
Output: 3

Example 2:

Input: `nums = [2,2,1,1,1,2,2]`  
Output: 2

Constraints:

- `n == nums.length`
- `1 <= n <= 5 * 104`
- `-109 <= nums[i] <= 109`

```

13         freq++;
14     } else {
15         freq=1;
16         ans = nums[i];
17     }
18
19     if(freq > n/2) {
20         return ans;
21     }
22 }
23
24 return ans;
25

```

Testcase: Case 1 Case 2 +

nums = [3,2,3]

Print Screen

moore's voting algo

## Majority Element

{ 1, 2, 2, 1, 1 }  $> n/2$

freq = 0, ans = 0

```

for(int i=0; i<n; i++) {
    if(freq == 0)
        ans = nums[i];
    if(ans == nums[i]) freq++;
    else freq--;
}
return ans;

```

Print Screen

① Brute Force  $O(n^2)$   
 ② Optimize  $O(n \log n)$   
 ③ Moore's  $O(n)$

same el  $\Rightarrow$  freq++  
 diff el  $\Rightarrow$  freq--

169. Majority Element

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Input: `nums = [2,2,1,1,1,2,2]`  
Output: 2

Constraints:

- `n == nums.length`
- `1 <= n <= 5 * 104`
- `-109 <= nums[i] <= 109`

```

1 class Solution {
2 public:
3     int majorityElement(vector<int>& nums) {
4         int freq = 0, ans = 0;
5
6         for(int i=0; i<nums.size(); i++) {
7             if(freq == 0) {
8                 ans = nums[i];
9             }
10            if(ans == nums[i]) {
11                freq++;
12            } else {
13                freq--;
14            }
15        }
16        return ans;
17    }
18 }

```

Testcase: Accepted Runtime: 5 ms

Print Screen

Case 1 Case 2

## Majority Element

{ 1, 2, 2, 1, 1 }

freq = 0 1 0 0 1  
ans = 0 1 0 1

moore's voting algo

- 1) Brute Force  $O(n^2)$
- 2) Optimize  $O(n \log n)$
- 3) Moore's  $O(n)$

Print Screen



## Majority Element (Variation)

{ 1, 2, 2, 1, 1 }

[1, 2, 3, 4]

① ✓  
②  $ans \rightarrow freq > n/2 \rightarrow ans$   
 $ans \rightarrow -1$

Print Screen



Screenshot of a web browser showing the LeetCode problem page for "169. Majority Element". The page includes the problem description, examples, and a code editor with a C++ solution.

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Input: `nums = [3,2,3]`  
Output: `3`

**Example 2:**  
Input: `nums = [2,2,1,1,1,2,2]`  
Output: `2`

**Constraints:**  
`n == nums.length`

**Code:**

```
C++  
13  
14  
15  
16  
17  
18  
19  
20  
21  
22  
23  
24  
25  
26  
27  
28  
29  
};  
  
freq--;  
}  
}  
  
int count = 0;  
for(int val : nums) {  
    if(val == ans) {  
        count++;  
    }  
}  
  
if(count > n/2) => ans  
else => -1  
  
return ans;
```

**Testcase:** Accepted Runtime: 5 ms

**Print Screen**

