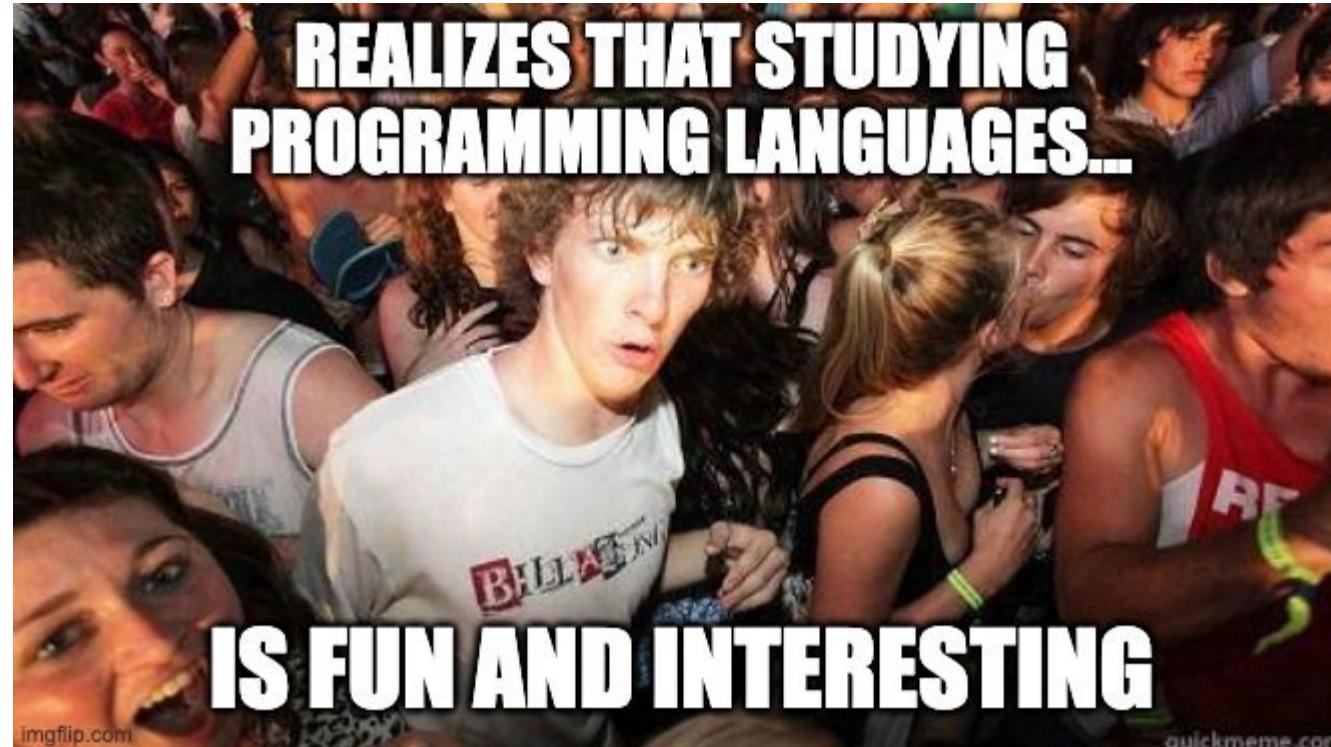


# CS131: Programming Languages

## Fall 2023

Carey Nachenberg

climberkip@gmail.com; ENG VI 299



# Useful Information



Mon: 4-5pm, Wed: 1-2pm  
Location: Eng VI 364



Our Class Hub

<https://bruinlearn.ucla.edu/courses/168646>



Syllabus, Assignments, PPTs, Lectures Notes, Solutions

<https://ucla-cs-131.github.io/fall-23-website>



Your Questions - Our Answers

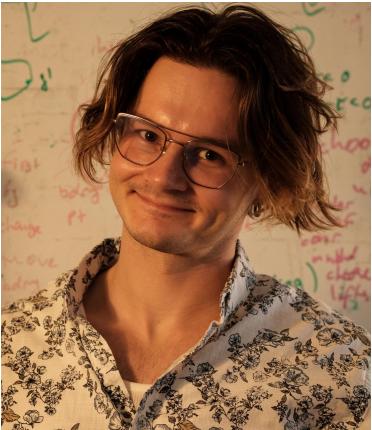
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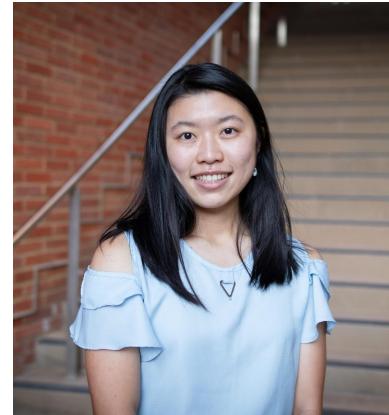
# Meet Our Team



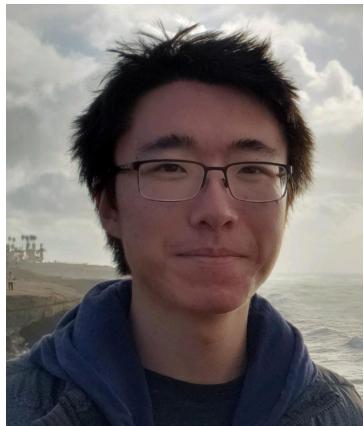
Carey Nachenberg



Andrey Storozhenko  
Discussions 1D and 1G



Bonnie Liu  
Discussions 1B and 1H



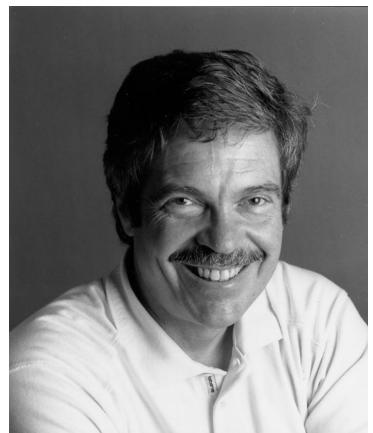
Brian Wang  
Discussions 1C and 1F



Justin Cui  
Discussions 1A and 1E

# Contributors

I couldn't have created this course without the hours of explanations/content/feedback from the following folks:



Alan Kay



Angelina Lue



Devan Dutta



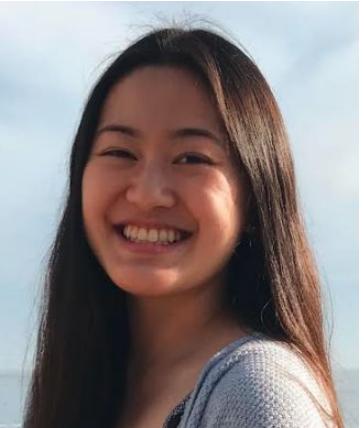
James Wu



Kevin Tan



Kyle Chui



Maggie Li



Matthew Wang



Paul Eggert



Timothy Gu



Todd Millstein

Guess what?

I love Comic Sans sooooo much...

I'm gonna use it for 131 as well! 

OK, just kidding... We'll use Corbel!

# What's CS131 About?



Languages come and go, but they're all built upon a common set of **paradigms** and **building-blocks**.

Our goal in CS131 is to learn these paradigms and building-blocks, understanding the implications of each.



This will enable you to pick up new languages fast, select the best language for a problem, and even invent your own language!



# Your Goal by Week 10?

Be able to explain – in your own words - the major language paradigms and building blocks, and their tradeoffs.

Given any new language\*, be able to grok how it works and write correct, efficient code in a matter of hours.



# Assignments and Exams



10%

30%

25%

35%

CS131 has **homework**, **projects**, a **midterm** and a **final exam**.

Homewo



Our Midterm is on  
Thursday, Nov 7<sup>th</sup> from  
6 to 8pm!

arn.

EFFORT

new language!

your understanding of  
key programming  
language concepts

your ability to analyze  
languages you've  
never seen before

your ability to write  
correct code in our  
focus languages

# Be Wary of Online Resources

There's still a ton of debate about



13 Answers

Sorted by: Highest score (default) ▾

▲ Python is strongly, dynamically typed.

513 • **Strong** typing means that the type of a value doesn't change in unexpected ways. A string containing only digits doesn't magically become a number, as may happen in Perl. Every change of type requires an explicit conversion.

▼

□ Dynamic typing means that runtime objects (values) have a type, as opposed to static

If you want more details, limit searches to  
lecture notes from major universities.



Warning: Stack Overflow, Medium, etc.  
have lots of incorrect explanations!

# Learning Resources



There is no required textbook  
(I couldn't find any good ones!)

But... we'll be posting **slides** and **lecture notes**, and...

## Side-by-side Language Comparisons

Rosetta Code  
<http://rosettacode.org/>

Hyperpolyglot  
<https://hyperpolyglot.org/>

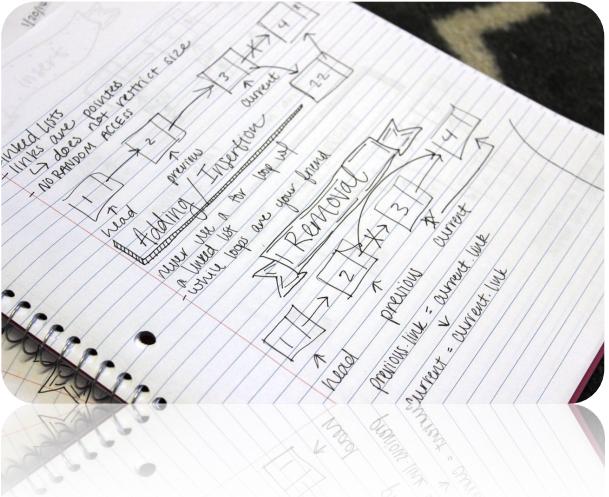
## Gentle Language Introductions

Learn You a Haskell for Great Good  
<http://learnyouahaskell.com/>

W3 Schools Python Intro  
<https://www.w3schools.com/python/>

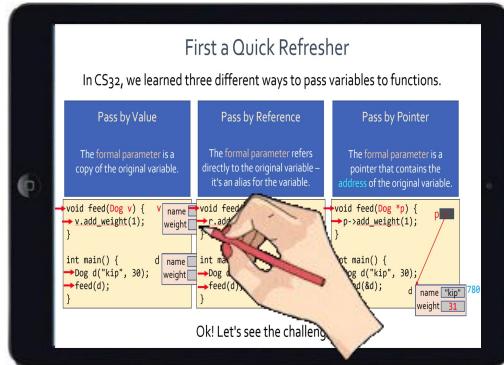
tutorialspoint  
<https://www.tutorialspoint.com/prolog>

# Attention Note Takers!



Many students are used to taking detailed notes during lecture...

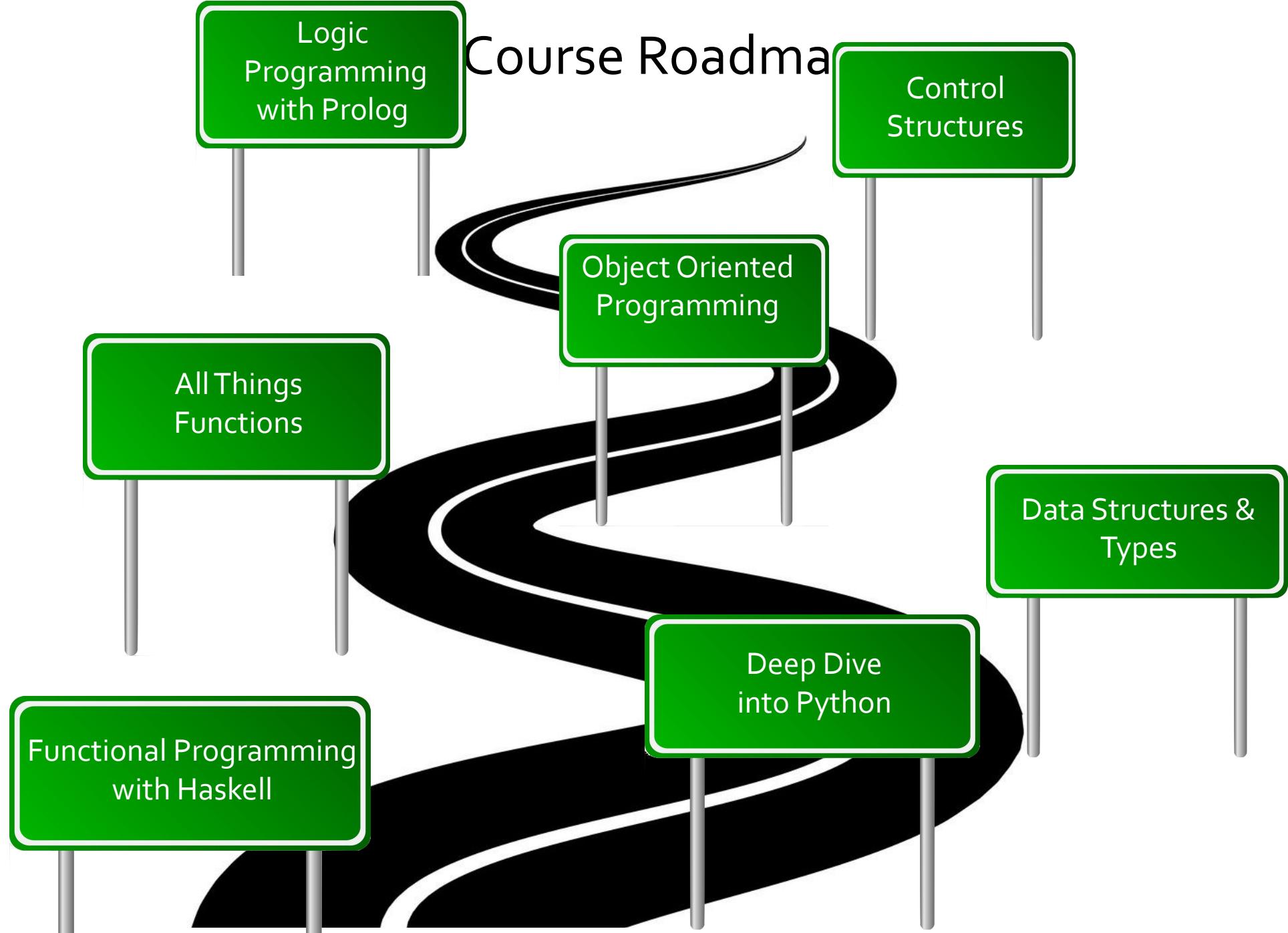
If that's you, you may find it difficult to keep up since I show lots of code examples.



Code examples

Suggestion: Download our lecture notes or slides ahead of time and take notes on them...

# Course Roadmap

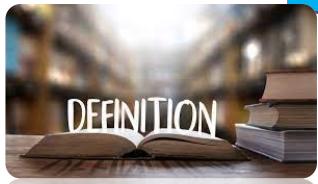


Ok, let's  
go!





# What is a Programming Language?



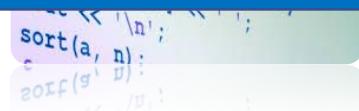
A programming language is a structured system of communication designed to express computation based on some formalism.

A programming language can be used to specify the behavior of a computer.



When you see this bar, it means:

You have 1 minute to discuss with your neighbor and come up with an answer.



In CS131, we'll be focusing on high-level, "universal" languages like C++ and Python – that is, those capable of implementing pretty much any valid program.

# What were the world's first *high-level* programming languages?

First Interpreted Language:

[Shortcode](#)

John Mauchly '49



A program was a set of statements – each a math expression you wanted to compute.

First Compiled Language:

[A-0](#)

Grace Hopper '51



A program was a sequence of subroutine calls with arguments. The compiler converted the program into machine code.

# How many different languages exist?

Wikipedia lists **687** different programming languages!

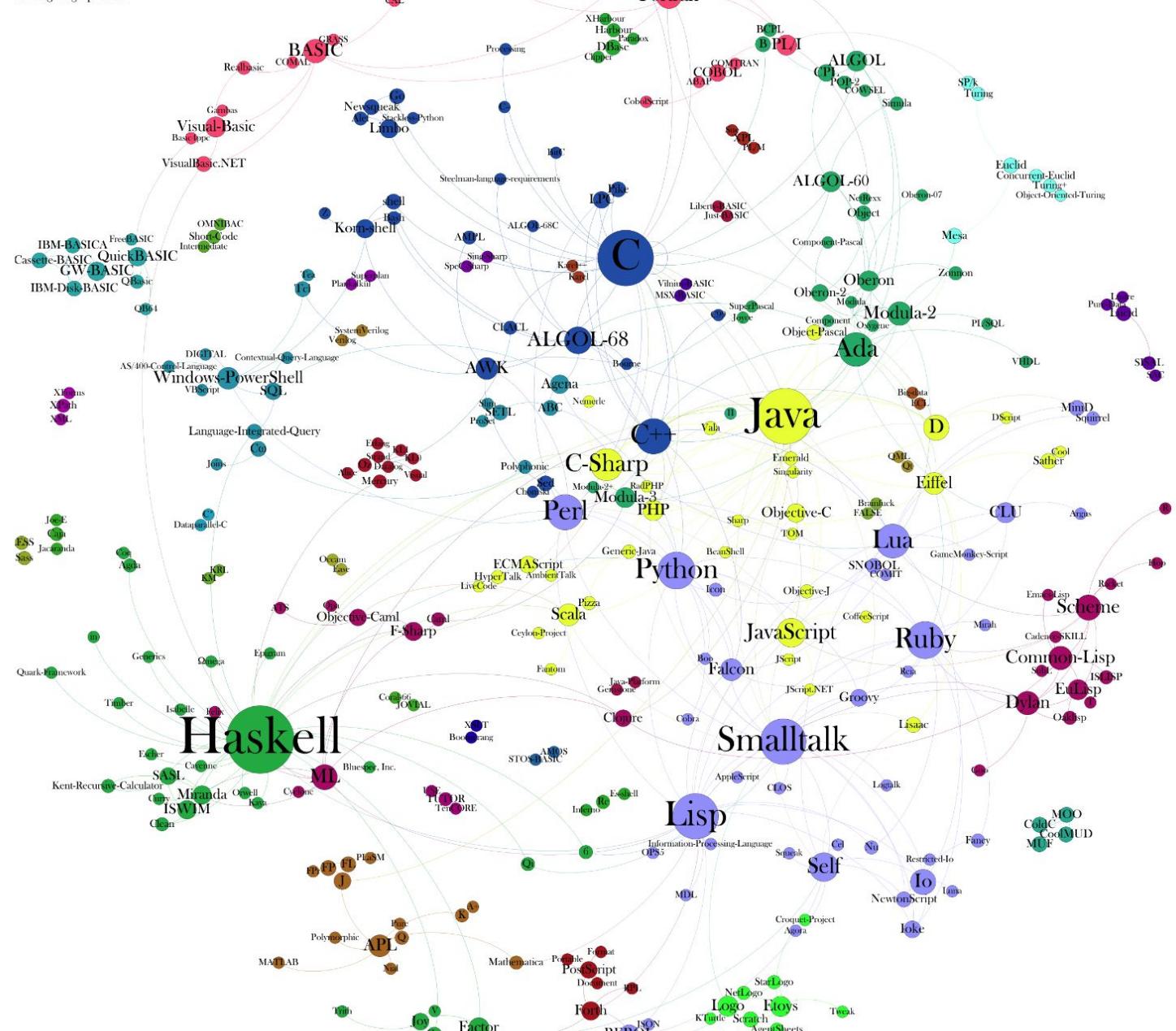
# But there are thousands!

Here's a nice visualization  
showing how  
interconnected they are!

## The Graph Of Programming Languages

By Brendan Griffen

[www.griffsgraphs.com](http://www.griffsgraphs.com)



# Why So Many Languages?

Each language is designed to address a particular category of **problem**, **research area**, or **business domain**.

**JavaScript** was designed for **front-end** use cases

**Simula** was developed to build **simulations**

**SQL** was designed to **query databases**



There are also thousands of **Domain-Specific Languages** (DSLs) that folks have designed to accelerate problem-solving for specific use cases!

How much can a language's design  
impact problem-solving speed?



Let's find out!



# Challenge: How Small is the Smallest Spell Checker?

Given an **array of sentences** and an **array holding an English dictionary**, what's the minimum number of lines of code it would take to build a spell checker?

You can choose any language you like!

```
results = spell_check(s, d);
// results should hold: "admyt" and "lyt"
```

S "I must admyt, it is lit"  
"So LYT... I can't take it!"  
"Lyt; AdMyT it!"

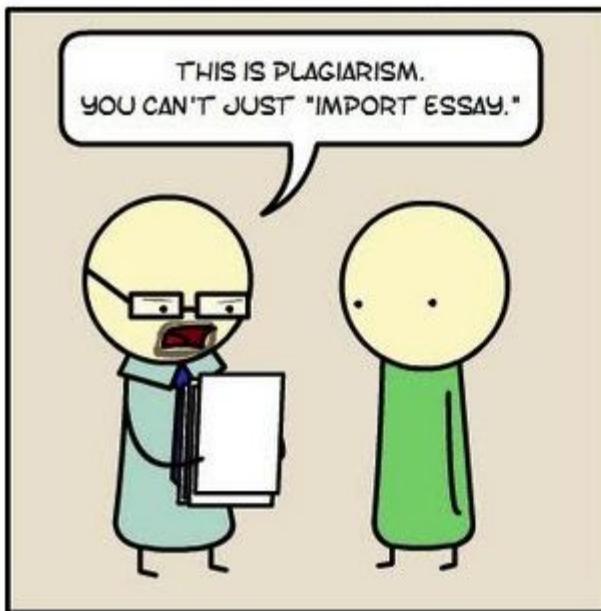
Answer: 3 lines!

```
# Python spell checker using list comprehensions
import re

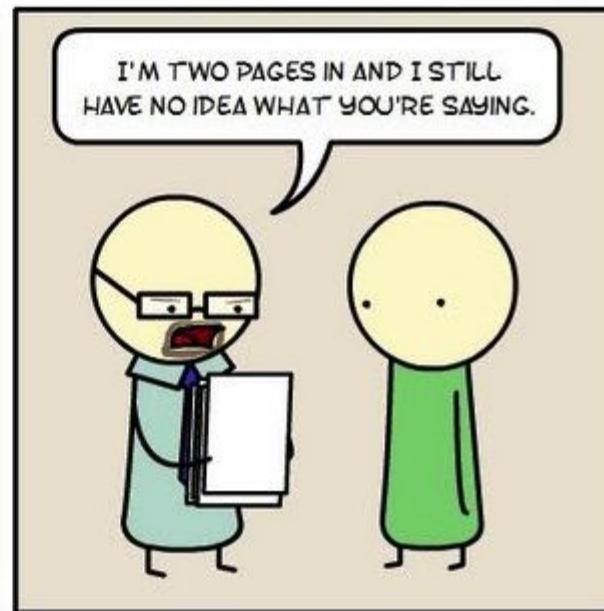
def spell_check(s, d):
    return {w for w in [w for ln in s for w in re.split(r'[.,!;]',ln)]
            if w.lower() not in d and w != ' '}
```

Assume the dictionary is all lower case and in alphabetical order.

## Writing an essay in different languages



PYTHON



JAVA

# Programming and Starbucks?!?!



&



# Programming Languages are Like Starbucks Drinks

Just as we have different **classes** of drinks...

Coffee Drinks



Hot Chocolate



Teas

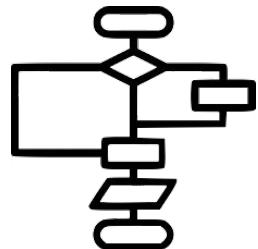


Ciders



There are different programming language **paradigms**:

Imperative



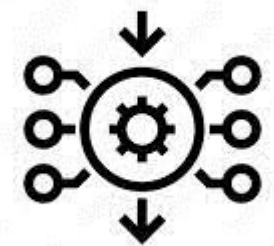
Object Oriented



Functional



Logic



# The Major

This quarter, we're going to

And just like coffee drinks, many languages are hybrids!

C

## Imperative Languages

What you're used to - programs are made up of statements, loops, and mutable variables.

```
int a = 0;  
while (a++ < 100)  
    cout << a << endl;
```

C++

## OOP Languages

Programs are organized into classes and objects which send messages, e.g., obj.msg(x), to each other to get things done.

```
class Dog { ... };
```

Haskell

## Functional Languages

Programs are made up of math-like functions. No statements or iteration, only expressions, functions, constants and recursion!

```
f(n) = if n>0 then n*f(n-1)  
       else 1
```

Scala

## Logic Languages

Programs define a set of facts, e.g.: like(dogs, meat), category(beef, meat).

Programs also define rules:  
likes(A,C): if like(A,B) and category(C, B)

You then query: likes(dogs, beef)?

Smalltalk

Prolog

# Programming Languages are Like Starbucks Drinks



# Coffee Choices

**Coffee:** light, medium, dark-roast

Milk: 2%, skim, whole, soy, almond, oat

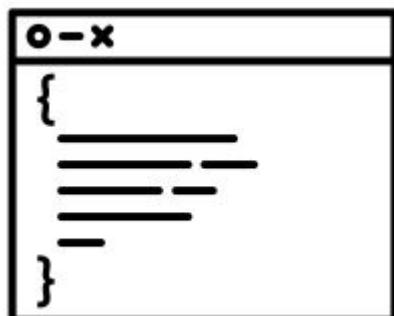
**Flavoring:** chocolate, caramel, vanilla

**Sweetener:** sugar, Splenda, stevia

Starbucks lets you  
customize coffee across  
key dimensions...

For instance, C++ made the following design choices...

But Python made a totally different set of choices!



Type checking: static typing, dynamic typing

Parameter passing: by-value, by-reference,  
by-pointer, by-object reference, by-name

## Scoping: lexical scoping, dynamic scoping

**Memory management:** manual, automatic

# Language Dimensions: More Examples

A

Pro: Detects typing bugs at compile time, when they're easy to fix  
Con: More verbose code

```
int a; // variable declared first!  
a = 5; // then it can be assigned
```

vs.

A var

Pro: Faster dev, simpler code  
Con: Less readable code, may detect errors only at runtime

```
a = 5 # 1st assignm't creates var
```

The language  
considers

Pro: Makes for simpler code  
Con: May hide conversion bugs

```
int a = 5;  
double b = a; // OK!
```

S.

The language  
forwards

Pro: Detects bugs more easily  
Con: More verbose code

```
int a = 5;  
double b = (double)a; // OK
```

The language

Pro: Catches out-of-bounds bugs  
Con: Runtime checks slow execution!

```
int arr[] = new int[10];  
  
cout << arr[10]; // Blocked!
```

v3.

The la

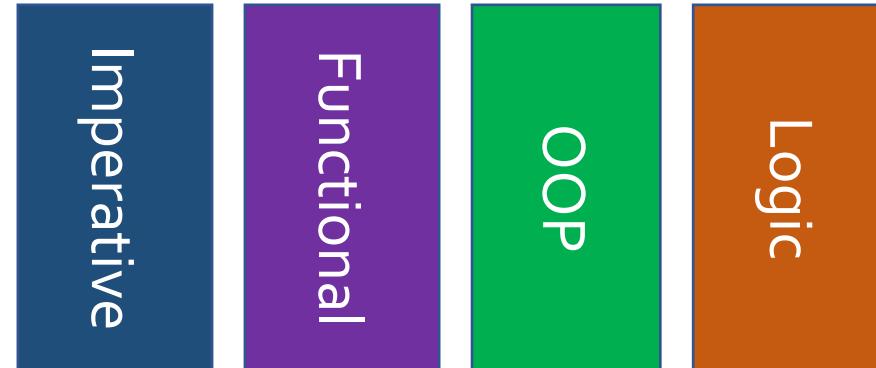
Pro: Faster  
Con: Permits nasty memory bugs!

```
int arr[10];
```

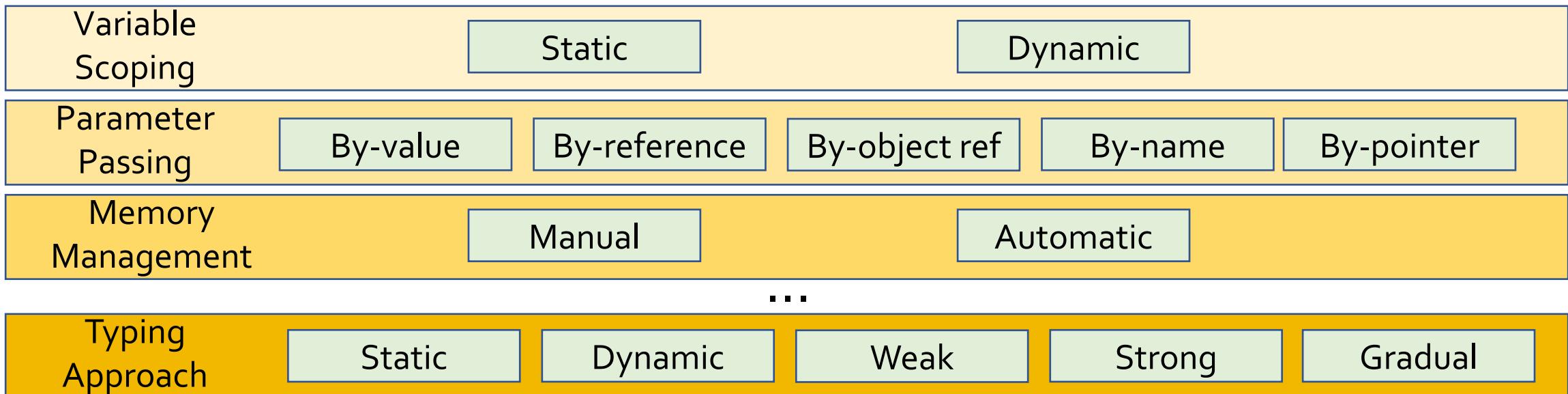
```
cout << arr[10]; // Allowed! Eek!
```

# Summarizing Our Learning Approach

We'll deep-dive into the major paradigms



We'll go wide and survey the spectrum of design options





An important skill is the ability to logically deduce which design choices each language has made.

So throughout the quarter, we'll be training ourselves to do this with challenges.



Let's see an example.

# First a Quick Refresher

In CS32, we learned three different ways to pass variables to functions.

## Pass by Value

The formal parameter is a copy of the original variable.

```
void feed(Dog v) {    v
    v.add_weight(1);
}

int main() {
    Dog d("kip", 30);
    feed(d);
}
```

## Pass by Reference

The formal parameter refers directly to the original variable – it's an alias for the variable.

```
void feed(Dog &r) {
    r.add_weight(1);
}

int main() {
    Dog d("kip", 30);
    feed(d);
}
```

## Pass by Pointer

The formal parameter is a pointer that contains the address of the original variable.

```
void feed(Dog *p) {
    p->add_weight(1);
}

int main() {
    Dog d("kip", 30);
    feed(&d);
}
```

p

d  
name "kip"  
weight 31

780

Ok! Let's see the challenge!

# Classify That Language

Clearly, `change_major()`

Reassigning `s` only changes the local pointer, not the original object!

new object?

`major('Econ')`

es life choices, so starts from  
new('Carence', 'Music')

Let's see how it works!

```
# main program
student1 = Student.new(
  life_changes(student1)
  student1.print_my_detail()
```

changes

`s`

Tech  
by c

So, this language must NOT be passing arguments by reference, or `student1`'s name *and* major would have changed.

passing by pointer in C++:

By value

By p

This changes `student1`'s major.

`student1`

So, this language must NOT be passing arguments by value, or `student1`'s major would still be CompSci.

This is Ruby!

That didn't change the original variable's value!

print the following:

Carey's major is Econ.

# We'll Deep Dive Into Three Languages

Of the thousands of languages, this quarter we'll deep-dive into **Haskell**, **Prolog**, and **Python**.



While you're familiar with **imperative programming** and **OOP** (via C++), we need to introduce you to the other major language paradigms.

**Haskell** is one of the simplest and easiest-to-grok **functional languages**.

**Prolog** is one of the few examples of **logic languages**.

**Python** is a **dynamically-typed language** (and it's super useful).

Next, let's learn what it means to specify  
the details of a programming language!



Some language specs, like the one for C++, are written in plain English!!!

# eeded to specify the details of a programming Language?

## Patterns



A language area

Other

And some languages have NO written spec at all (like Perl v1-v4)!

Whatever the compiler/interpreter does is canonical!

### 8.5.2.2 Increment

- 1 The operand of prefix `++x` shall be an address. The result is the updated value. `++x` is equivalent to `x = x + 1`. For information on compound increments, see 8.5.1.3.
- 2 The operand of prefix `--x` shall be the properties of its decrementation, see 8.5.1.3.

Most modern languages use an encoding called Extended Backus–Naur Form (EBNF) to define the language's syntax.

$$\begin{array}{c} \textit{longcon} = \textit{strid}_1 \dots \textit{strid}_n \\ \hline E, v' \vdash \textit{atpat} \end{array}$$

$$\begin{array}{c} \textit{longcon} = \textit{strid}_1 \dots \textit{strid}_n \\ \hline E, v' \vdash \textit{atpat} \end{array}$$

$$E(\textit{longexco})$$

$$E, v \vdash \textit{longexco}$$

$$\begin{array}{c} E(\textit{longexcon}) = \textit{en} \quad v \notin \{\textit{en}\} \times \text{Val} \\ \hline E, v \vdash \textit{longexcon} \textit{atpat} \Rightarrow \text{FAIL} \end{array} \quad (157)$$

$$\begin{array}{c} s(a) = v \quad s, E, v \vdash \textit{atpat} \Rightarrow \text{VE/FAIL}, s \\ \hline s, E, a \vdash \textit{ref} \textit{atpat} \Rightarrow \text{VE/FAIL}, s \end{array} \quad (158)$$

$$\begin{array}{c} E, v \vdash \textit{pat} \Rightarrow \text{VE/FAIL} \\ \hline E, v \vdash \textit{var as pat} \Rightarrow \{\textit{var} \mapsto v\} + \text{VE/FAIL} \end{array} \quad (159)$$

The operand of prefix `--` and `++`: For postfix increment and

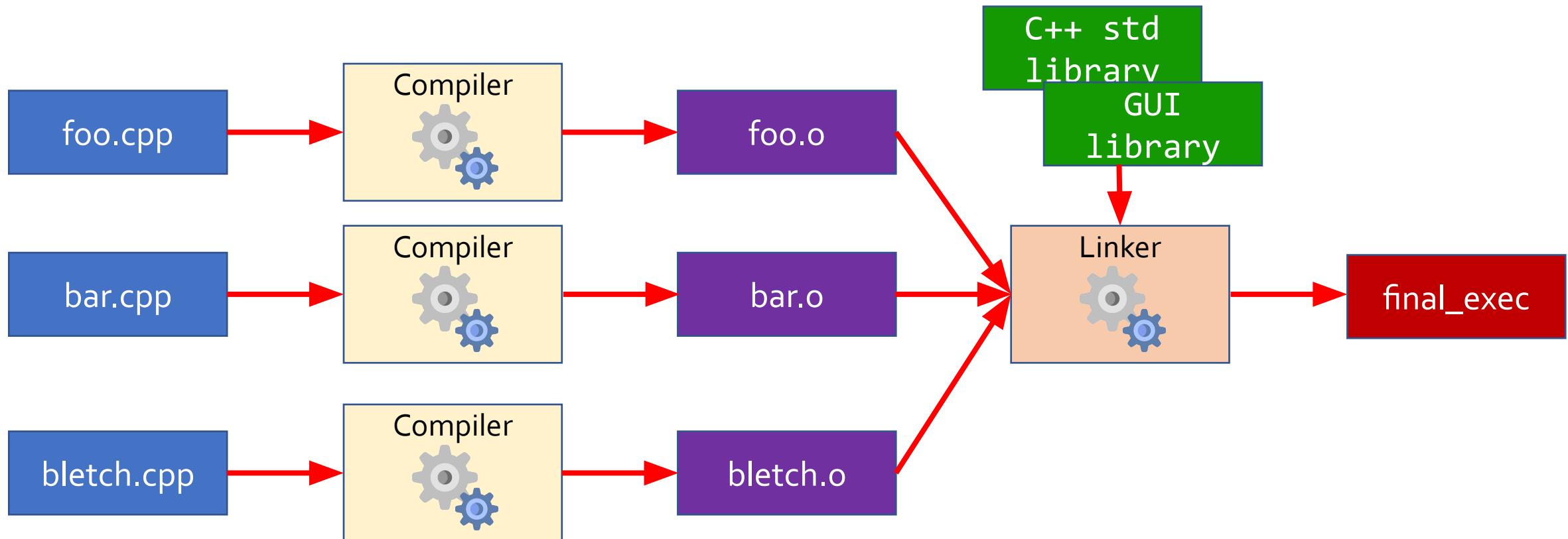


# Our Tools

# What's a Compiler? What's a Linker?

A **compiler** is a program that translates **program source** into **object modules** (which are either machine language, or bytecode targeted at an interpreter).

A **linker** is a program that combines **multiple object modules** and **libraries** into a single **executable file** or **library**.



file.cpp

```
while(i<5)  
    x = x + i*3;
```

e.g., that x and i are both the tokens and compiler + and

Bytecode is a binary encoding – like a machine language, but not for a real CPU.

It is processed by an interpreter or just-in-time (JIT) compiler.



Machine  
or byte

55 48 8  
00 00 6

The output is an annotated tree with type and other info added to the nodes of the tree.

# How Does a Compiler Work?

If the syntax is valid, the parser converts the tokens into a representation of the program's structure.

The IR generator produces an abstract representation of the program – it's totally independent of the original language.

The representation could be a tree structure, or CPU-agnostic instructions.

```
load i  
load_constant 5  
less_than  
jump_if_zero L1  
...
```

The  
IR

IF

WHILE

ASSIGN

ADD

MULT

VARIABLE

INTCONST

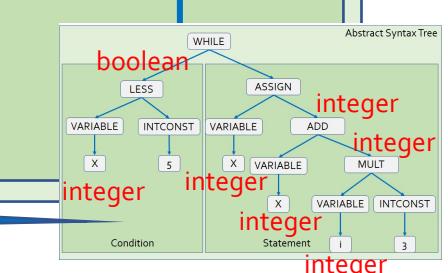
i

Intermediate  
representation

Intermediate  
Representation  
Generator

Statement

Condition



func: type func\_name '(' params ')' '{' { type var\_name ';' } { stmt } '}'  
stmt: 'if' '(' expr ')' stmt [ 'else' stmt ]  
[ 'while' '(' expr ')' stmt ]  
[ assg ]' stmt

[ or } ]' ';' pr | ...

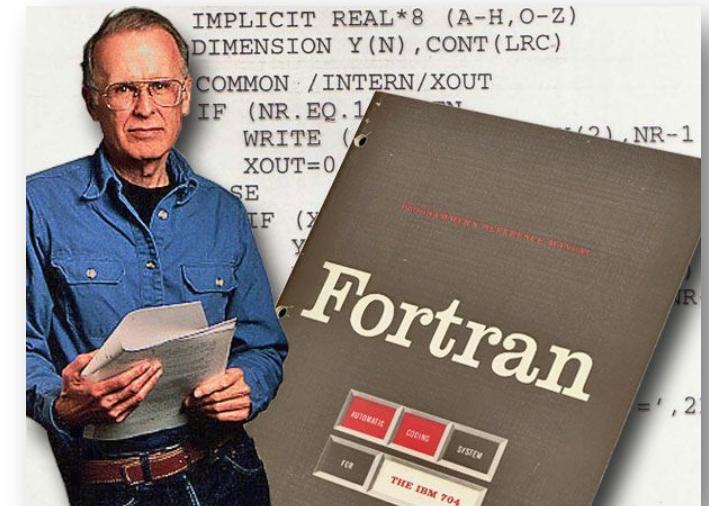
# For Funzies



What was the world's first mainstream compiler (for Fortran) written in?

Answer: Assembly language.

Makes sense, right?  
What else would it be written in?



John Backus

# For Funzies, Part 2



What language are most C++ compilers written in?

Answer: C++?!?!? Through a process called "**bootstrapping**"

```
// C++ compiler in C++!
class Compiler {
public:
    Compiler();
    void compile();
};
```

c++-compiler.cpp

to-c-translat  
or.c

```
// C++ compiler in C
struct Compiler
{
    ...
};
void init(Compiler& c)
void compile(Compiler& c);
```

c++-to-c-translat  
or.exe

Now we have a  
working C++  
compiler!



o  
wor  
c++compiler.exe

First you build a C++ to C  
translator... in some existing  
language like C.

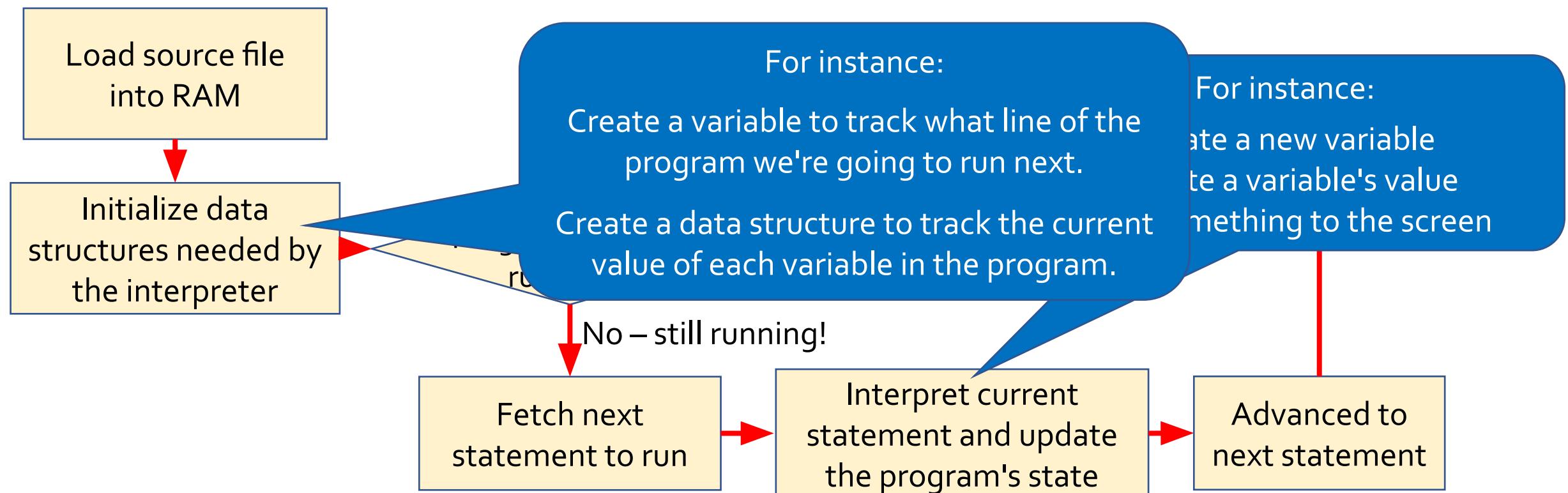
c++-compiler.cpp

c++compiler.exe

# What's an Interpreter?

An **interpreter** is a program that **directly executes program statements**, without requiring them to first be compiled into machine language.

How does an interpreter work?



# What's an Interpreter?

An **interpreter** is a program that **directly executes program statements**, without requiring them to first be compiled into machine language.

Ok, let's build an interpreter for a simple programming language!

Consider the following simple program in our new language which:

initializes a variable called **x** to **10**,  
adds **20** to the variable,  
prints the value of **x**,  
and finally **terminates**.

```
set x 10
add x 20
print x
end
```

# An Interpreter In One Slide!

We'll need to track where we are in our program.

```
// specify the code to interpret
vector<string> program = {"set x 10", "add x 20", "print x", "end"};
int cur_line = 0;
bool terminated = false;

void interpret(const vector<string> &program) {
    for (int i = 0; i < program.size(); ++i) {
        cout << "Line " << i << ": " << program[i] << endl;
        if (program[i] == "end") {
            terminated = true;
        } else if (program[i] == "set") {
            variable_to_value["x"] = stoi(program[i].substr(4));
        } else if (program[i] == "add") {
            variable_to_value[program[i].substr(4)] += stoi(program[i].substr(8));
        } else if (program[i] == "print") {
            cout << variable_to_value[program[i].substr(4)] << endl;
        }
        ++cur_line;
    }
}
```

We also need to keep track which variables have been defined so far.

Let's use a map to associate each variable name, e.g., x, with its value.

We'll start executing at line 0.

which causes execution to terminate.

stoi() converts a string to an integer.

```
if (tokens[0] == "end") {
    terminated = true;
} else if (tokens[0] == "set") {
    variable_to_value[tokens[1]] = stoi(tokens[2]);
} else if (tokens[0] == "add") {
    variable_to_value[tokens[1]] += stoi(tokens[2]);
} else if (tokens[0] == "print") {
    cout << variable_to_value[tokens[1]] << endl;
}
++cur_line;
```

variable's value

Finally, we advance to the next line in our program.

First we'll parse our statement into separate tokens:  
"set x 10" → {"set", "x", "10"}

while our program hasn't terminated itself, keep executing the next statement.

Let's assume we have a function that can split a string up into its parts.

Here's how we'll use it...

... see the details!

... add", we update the value by adding the specified amount.

... map and print it.

```
int main() {
    vector<string> program = {"set x 10", "add x 20", "print x", "end"};
    interpret(program); // 30
}
```

# Interpreters are Soooo Cool

that you'll be building your own



Our four class projects require you to build an interpreter for a new programming language!

And as we learn new language concepts, you'll add them to your interpreter!

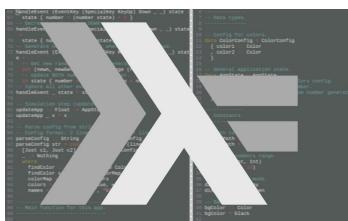


Giving you a deep,  
hands-on understanding.

# A Few Final Things...



Your first homework will be posted ***Wednesday of this week*** on Gradescope.



Until then, please install the Haskell compiler and Python 3.x interpreter.

# A Few Final Things...

CS131 is a notoriously difficult class.



So ask questions...

Stop me if something's unclear...



Don't be afraid to ask for help.

Our goal is not just to cover the material, but for you to deeply understand it and have fun in the process.

# Onward to Functional Programming With Haskell!

```
56 handleEvent (EventKey (SpecialKey KeyUp) Down _) state =  
57     state { number = (number state) - 1 }  
58 -- Decrease number.  
59 handleEvent (EventKey (SpecialKey KeyDown) Down _) state =  
60     state { number = (number state) + 1 }  
61 -- Increase number.  
62 handleEvent (EventKey (SpecialKey KeyLeft) Down _) state =  
63     state { number = (number state) - 1 }  
64 -- Generate new random number.  
65 handleEvent (EventKey (SpecialKey KeyRight) Down _) state =  
66     state { number = (number state) + 1 }  
67 -- Generate new random number.  
68 handleEvent (EventKey (SpecialKey KeyGenerate) Down _) state =  
69     let (new, newGen) = randomRange (number state)  
70     state { number = new, generator = newGen }  
71 -- Ignore all other events.  
72 handleEvent _ state = state  
73  
74 -- Simulation step (update application state).  
75 updateApp :: Float -> AppState  
76 updateApp _ x = x  
77  
78 -- Parse config from string.  
79 -- Config format: 2 lines.  
80 -- Line 1: color1, color2  
81 parseConfig :: String -> ColorConfig  
82 parseConfig str = case lines str of  
83     [Just c1, Just c2] -> ColorConfig { color1 = c1, color2 = c2 }  
84     _ -> Nothing  
85     where  
86         findColor :: String -> Maybe Color  
87         findColor s = case words s of  
88             [name, hex] -> Just $ Color { name = name, hex = hex }  
89             _ -> Nothing  
90         colorMap = Map.fromList [ ("red", "#ff0000"),  
91             ("green", "#00ff00"), ("blue", "#0000ff") ]  
92         colors = Map.toList colorMap  
93         names = Map.keys colorMap  
94  
95 -- Main function for this app.  
96 main = do  
97     configStr <- getLine "Enter config file path:  
98     configPath <- readLine configStr  
99     config <- parseConfig configPath  
100    putStrLn $ "Color 1: " ++ show (color1 config)  
101    putStrLn $ "Color 2: " ++ show (color2 config)  
102  
103 -- Data types.  
104  
105 -- Config for colors.  
106 data ColorConfig = ColorConfig  
107     { color1 :: Color  
108     , color2 :: Color  
109     }  
110  
111 -- General application state.  
112 data AppState = AppState  
113     { number :: Int,  
114     generator :: RandomGen  
115     }  
116  
117 -- Constants.  
118  
119 -- Path to config file.  
120 configPath :: Path  
121 configPath = "config.txt"  
122  
123 -- Random number range.  
124 randomRange :: (Int, Int)  
125 randomRange = (1, 100)  
126  
127 -- Colors.  
128 colors :: Map String Color  
129  
130 -- Background color.  
131 bgColor :: Color  
132 bgColor = black
```

# Appendix

# yntax or Semantic Error?

Which errors will be detected by the **parser** vs. the **semantic analyzer**?

Detected by the  
Parser

Detected by the  
Semantic Analyzer



```
int
cou
} int burp = 5
cout << burp
```

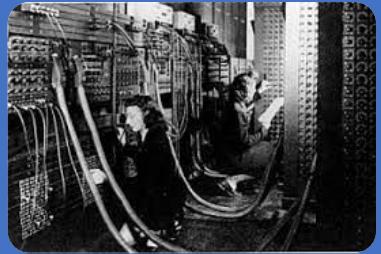
# An Early History of Modern Languages

## Machine Language

A program is a sequence of numbers which represent instructions, e.g., (0x6b=multiply); the instructions are directly interpreted by the CPU.

```
55 48 89 e5 c7 45 fc  
05 00 00 00 6b 45 fc  
03 89 45 f8
```

ENIAC '40s



They're hard to read!



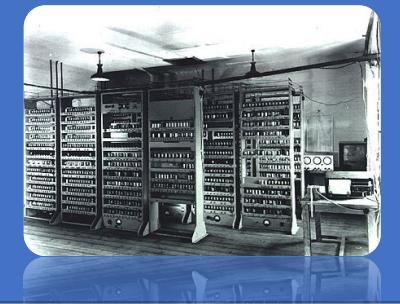
Development is super slow

## Assembly Language

A program is a sequence of barely human-readable operations and operands ("add eax, 5") which have a 1-to-1 correspondence with machine language instructions.

```
pushq %rbp  
movq %rsp, %rbp  
movl $5, -4(%rbp)  
imull $3, -4(%rbp), %eax  
movl %eax, -8(%rbp)
```

ESDAC '49



They're not portable to other CPUs!

machine and assembly languages?



They're error-prone!



There's no error checking!

# What's a Transpiler?

A **transpiler** is a program that **converts source code** from one language into another **with the same level of abstraction**.



Why/when would we want to do that??

Enable migration to a "safer" version of a language,  
yet ensure compatibility with existing platforms



TypeScript is like a strict  
version of JavaScript.  
It greatly reduces bugs.



Support a new language's syntax before you  
have a full compiler for it

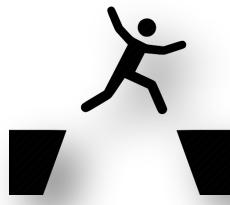


Compiler

Stroustrup's first C++ "compiler"  
was really a transpiler that  
generated C code!

OK, now let's see a transpiler in one slide!

# Language Choices: Your Turn



What other big language design choices can you think of?

This quarter, we're going to learn all the major language design choices and options for each!

# What's CS131 About?

To do this, we'll:

*Go Deep*

Dive deep into each language paradigm:  
Imperative, Functional, OOP, Logic

Go wide, surveying the common building blocks  
that every language is based on

*Go Wide*



Be hands-on – you'll be implementing your own  
programming language from scratch this quarter!

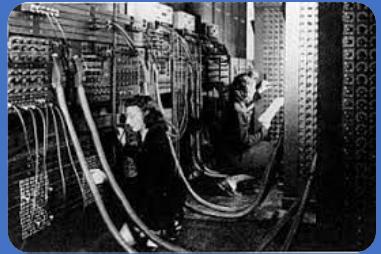
# But First... A Throwback to CS33!

## Machine Language

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ENIAC '40s



They're hard to read!



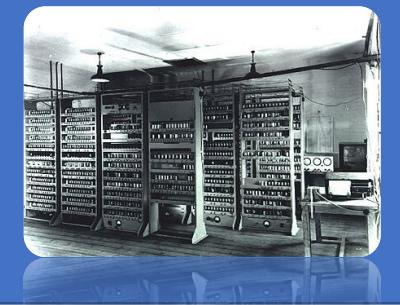
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ESDAC '49



They're not portable to other CPUs!

machine and assembly languages?



They're error-prone!



There's no error checking!



# What are the benefits of a high-level language over assembly/machine language?



Programs are easier to understand!



They're portable to many CPUs!



Development is faster



There's more error checking!

# Does a Compiler \u26a1

e.g., that **x** and **i** are both the and comp + and

If the syntax is valid, the parser converts the tokens into a tree that represents program's operations.

file.cpp

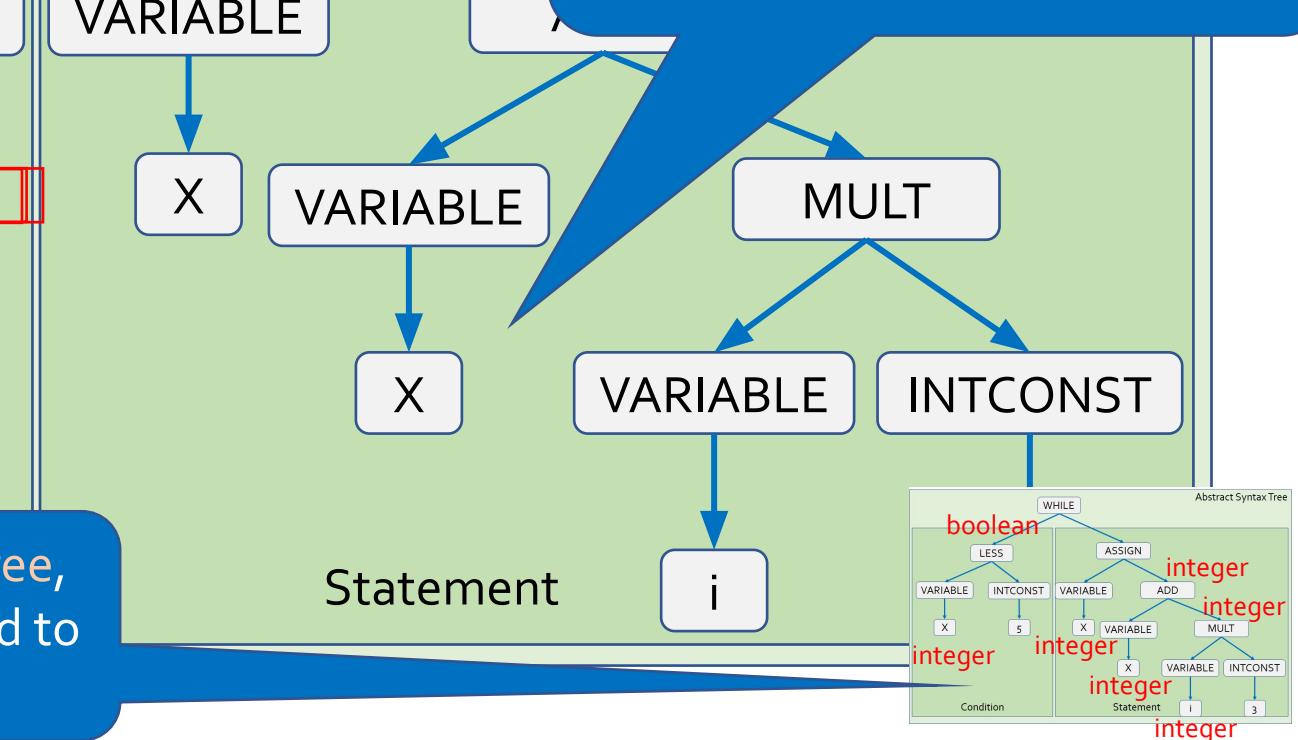
```
while(i<5)  
    x = x + i*3;
```

Bytecode is a binary encoding – like a machine language, but not for a real CPU.

It is processed by an interpreter or just-in-time (JIT) compiler.



The output is an annotated tree, with type and other info added to the nodes of the tree.



```
func: type func_name '(' params ')'
      '{' { type var_name ';' } { stmt } '}'  
stmt: 'if' '(' expr ')' stmt [ 'else' stmt ]
      | 'while' '(' expr ')' stmt
      | 'for' '(' [ assg ] ';' [ expr ] ';' [ assg ] ')' stmt
      | 'return' [ expr ] ';'  
      | func_name '(' [expr { ',' expr } ] ')' ';'  
      | '{' { stmt } '}'  
expr: const | variable | expr '<' expr | ...
```

# bytecode output.

# Another Example: Sudoku Solver

How many LoC would it take to build a Sudoku solver in **Python**?

```
% Prolog solution for Sudoku in ~15 lines!
sudoku(Rows) :-
    length(Rows, 9), maplist(same_length(Rows), Rows),
    append(Rows, Vs), Vs ins 1..9,
    maplist(all_different, Rows),
    transpose(Rows, Columns),
    maplist(all_different, Columns),
    Rows = [As,Bs,Cs,Ds,Es,Fs,Gs,Hs,Is],
    blocks(As, Bs, Cs),
    blocks(Ds, Es, Fs),
    blocks(Gs, Hs, Is).

blocks([], [], []).
blocks([N1,N2,N3|Ns1], [N4,N5,N6|Ns2], [N7,N8,N9|Ns3]) :-
    all_different([N1,N2,N3,N4,N5,N6,N7,N8,N9]),
    blocks(Ns1, Ns2, Ns3).
```

Answer: ~45 lines

Ok, here's the 15-line solution in **Prolog**!

Clearly, a language can have a massive impact on problem-solving efficiency!

```
class Student:  
    def __init__(self, name, debt):  
        self.name = name  
        self.debt = debt
```

This updates the debt of our original object, through the pointer!

This updates the debt of our original object, through the pointer!

```
→ def go_to uc/a(self, qtrs):
```

Clearly grow up() successfully  
Reassigning s only changes  
the local pointer, not the  
original object!

# original object!

Let's see how  
it works!

# # main program

```
student1 = Student("Carey")
grow_up(student1)
student1.my_debt();
```

student1

So, this language must NOT be passing arguments by value, or student1's debt would stay at zero.

# guage

The code to the left...  
prints the following:

Carey has \$12000 in debt.

"p

Technically, in Python we call this "pass by object reference," but it's just like passing by pointer in C++!

By Vale

# ~~By reference?~~

# By pointer?

# This is Py

# This is Python!

# The "Right" Answer

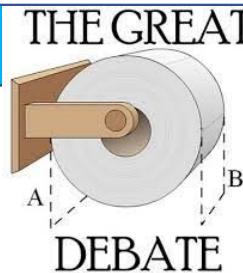
What's better? An  $O(N \cdot \log_2 N)$  sorting algorithm or a  $O(N^2)$  version?

What's better? Placing a return type **before** a function's name/parameters or **after**? Why?

```
// C++  
double sqrt(double val)  
{ ... }
```

```
// Swift  
func sqrt(val: Double) -> Double  
{ ... }
```

Guess what? There's no  
right answer...



just arguments for/  
against each approach!

As we compare programming languages, there often won't be a "right" answer.

So what's important is being able to back up our answer with arguments.