## Homework 5

**2.1** Consider the following grammar over the alphabet  $\{+, \times, (,), a\}$ .

$$\begin{array}{c} E \rightarrow E + T \mid T \\ T \rightarrow T \times F \mid F \\ F \rightarrow (E) \mid \mathbf{a} \end{array}$$

Give parse trees and derivations for each string.

- a. ab. a+ac. a+a+ad. ((a))
- 2.6 Give context-free grammars generating the following languages.

**d.** 
$$\{x_1 \# x_2 \# \cdots \# x_k | k \ge 1, \text{ each } x_i \in \{a, b\}^*, \text{ and for some } i \text{ and } j, x_i = x_j^{\mathcal{R}}\}$$

2.9 Give a context-free grammar that generates the language

$$A = \{ \mathbf{a}^i \mathbf{b}^j \mathbf{c}^k | i = j \text{ or } j = k \text{ where } i, j, k > 0 \}.$$

Is your grammar ambiguous? Why or why not?

\*2.19 Let CFG G be the following grammar.

$$S 
ightarrow aSb \mid bY \mid Ya$$
  
 $Y 
ightarrow bY \mid aY \mid arepsilon$ 

Let L(G) denote the language of G. Give a simple description of L(G) in English. Use that description to give a CFG for the complement of L(G).

\*2.27 Let  $G = (V, \Sigma, R, \langle STMT \rangle)$  be the following grammar.

$$\langle \text{STMT} \rangle \rightarrow \langle \text{ASSIGN} \rangle \mid \langle \text{IF-THEN} \rangle \mid \langle \text{IF-THEN-ELSE} \rangle \\ \langle \text{IF-THEN} \rangle \rightarrow \text{if condition then } \langle \text{STMT} \rangle \\ \langle \text{IF-THEN-ELSE} \rangle \rightarrow \text{if condition then } \langle \text{STMT} \rangle \text{ else } \langle \text{STMT} \rangle \\ \langle \text{ASSIGN} \rangle \rightarrow \text{a:=1} \\ \Sigma = \{ \text{if, condition, then, else, a:=1} \} \\ V = \{ \langle \text{STMT} \rangle, \langle \text{IF-THEN} \rangle, \langle \text{IF-THEN-ELSE} \rangle, \langle \text{ASSIGN} \rangle \}$$

 ${\cal G}$  is a natural-looking grammar for a fragment of a programming language, but  ${\cal G}$  is ambiguous.

- **a.** Show that G is ambiguous.
- b. Give a new unambiguous grammar for the same language.
- \*2.28 Give unambiguous CFGs for the following languages.
  - **c.**  $\{w | \text{ the number of a's is at least the number of b's in } w\}$
- 7 Prove that every regular language has a context-free grammar. *Hint:* given a DFA, explain how to transform it into an equivalent grammar.