




Review for Midterm 2

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Gaussian elimination \rightarrow LU factorization
(with partial pivoting) (P)
- ③ Overdetermined Linear Systems (5 lec.)  $A \vec{x} = \vec{b}$ where $A \in \mathbb{R}^{m \times n}$
(least square)
Normal eqn $A^T A \vec{x} = A^T \vec{b} \rightarrow$ QR factorization

Preliminaries

Two Types of Errors

- absolute error
- relative error

Floating-Point Numbers

- binary scientific notation:

$$\pm \left(1 + \frac{b_1}{2} + \frac{b_2}{2^2} + \cdots + \frac{b_d}{2^d} \right) 2^E,$$

where b_i is 0 or 1 and E is an integer.

- d determines the *resolution*
- the range of E determines the *scope* or *extent*
- IEEE Standard (double-precision; 64 bits)
 - $d = 52$ and $-1022 \leq E \leq 1023$
 - $\boxed{\text{eps}} = 2^{-52} \approx 2 \times 10^{-16}$
 - `realmin`, `realmax`

Floating-Point Numbers (cont')

- Key features
 - On any interval of the form $[2^E, 2^{E+1})$, there are 2^d evenly-spaced f-p numbers.
 - The spacing between two adjacent f-p numbers in $[2^E, 2^{E+1})$ is $2^{E-d} = 2^E \boxed{\text{eps}}$.
 - The gap between 1 and the next f-p number is $\boxed{\text{eps}}$, the machine epsilon.
 - Representation error (in relative sense) is bounded by $\frac{1}{2} \boxed{\text{eps}}$.

Conditioning (of a problem)

- The condition number measures the ratio of error in the result (or output) to error in the data (or input).
- Recall the definition of condition number $\kappa_f(x)$
- A large condition number implies that the error in a result may be much greater than the round-off error used to compute it.
- Catastrophic cancellation is one of the most common sources of loss of precision.

Stability (of an algorithm)

- When an algorithm produces much more error than can be explained by the condition number, the algorithm is unstable.

Square Linear Systems

Polynomial Interpolation

- Polynomial interpolation leads to a square linear system of equations with a Vandermonde matrix.

Gaussian Elimination and (P)LU Factorization

- A triangular linear system is solved by backward substitution or forward elimination.
- A general linear system is solved by Gaussian elimination.
- Gaussian elimination (with partial pivoting) is equivalent to (P)LU factorization.
- Solving a triangular linear system of size $n \times n$ takes $\sim n^2$ flops.
- PLU factorization takes $\sim \frac{2}{3}n^3$ flops.

Norms

A *norm* generalizes the notion of length for vectors and matrices.

- **Vector p -norm**

$$\|\mathbf{v}\|_p = \left(\sum_{i=1}^n |b_i|^p \right)^{1/p}, \quad p \in [1, \infty)$$

and

$$\|\mathbf{v}\|_\infty = \max_i |v_i|$$

- **Matrix p -norm (induced)**

$$\|A\|_p = \max_{\|\mathbf{x}\|_p=1} \|A\mathbf{x}\|_p, \quad p \in [1, \infty]$$

- **Frobenius norm (non-induced)**

$$\|A\|_F = \left(\sum_i \sum_j |a_{i,j}|^2 \right)^{1/2}$$

- **MATLAB:** `norm` can calculate both vector and matrix norms

Row and Column Operations

Various row and column operations can be emulated by matrix multiplications.
("Left-multiplication for row actions, right-multiplication for column actions")

- row/column extraction (unit vector)
- row/column swap (elementary permutation matrix)
- row/column rearrangement (permutation matrix)
- row replacement $R_i \rightarrow R_i + cR_j$ (Gaussian transformation matrix)

Conditioning/Stability

- Partial pivoting is needed for numerical stability.
- The matrix condition number is equal to the condition number of solving a linear system of equations.

Programming Notes

- Built-in functionalities
 - `backslash (\)`
 - `lu`
 - `norm`
 - `cond, condest, linsolve`
- Demonstration/Instructional codes
 - `backsub` and `forelim`
 - `GEnp` and `GEpp`
 - `mylu` and `myplu`

Overdetermined Linear Systems

Polynomial Approximation

- The most common solution to overdetermined systems is obtained by least squares, which minimizes the 2-norm of the residual vector.
- Least squares is used to find fitting functions that depend linearly on the unknown parameters.
- Equivalence of the LLS problem and the normal equation
 - linear algebra proof
 - calculus proof

QR Factorization

- Orthogonal sets of vectors are preferred to nonorthogonal ones in computing. (no catastrophic cancellation)
- Matrices with orthonormal columns and orthogonal matrices enjoy many *nice* analytical properties.
- QR factorization plays a role in LLS similar to that of LU factorization in square linear systems.

$$\bullet \quad Q^{-1} = Q^T \quad \leftarrow \quad Q^T Q = Q Q^T = I$$

$$\bullet \quad \|Q\vec{x}\|_2 = \|\vec{x}\|_2$$

$$\bullet \quad \|Q\|_2 = 1$$

$$\bullet \quad \text{cond}_2(Q) = \kappa_2(Q) = 1$$

$$\kappa_2(Q) = \|Q\|_2 \|Q^{-1}\|_2$$

Two Types of QR Factorization

For $A \in \mathbb{R}^{m \times n}$, $m \geq n$:

- Thick QR factorization: $A = QR$
 - $Q \in \mathbb{R}^{m \times m}$ orthogonal
 - $R \in \mathbb{R}^{m \times n}$ upper triangular
 - obtained by using successive Householder transformation matrices for *triangularization*
- Thin: $A = \hat{Q}\hat{R}$
 - $\hat{Q} \in \mathbb{R}^{m \times n}$ orthonormal columns
 - $\hat{R} \in \mathbb{R}^{n \times n}$ upper triangular
 - obtained by Gram-Schmidt *orthonormalization* procedure

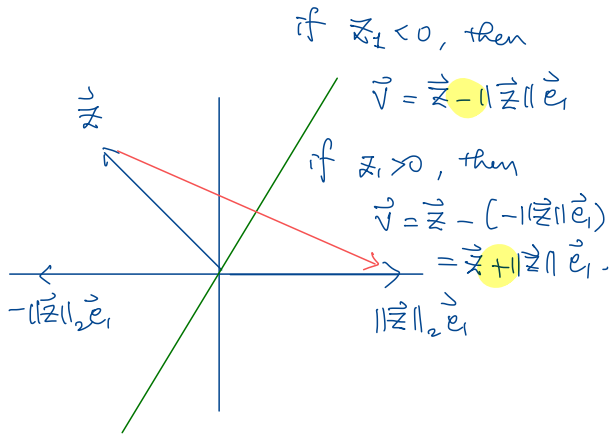
Householder Transformation Matrices

- A Householder transformation matrix H (associated with a vector \mathbf{z}) is a reflection matrix which is

- symmetric,
- orthogonal, and
- transforms \mathbf{z} to $\pm \|\mathbf{z}\|_2 \mathbf{e}_1$.

$$\tilde{\mathbf{v}} = \|\tilde{\mathbf{z}}\|_2 \tilde{\mathbf{e}}_1 - \tilde{\mathbf{z}}$$

$$H = I - 2 \frac{\tilde{\mathbf{v}} \tilde{\mathbf{v}}^T}{\tilde{\mathbf{v}}^T \tilde{\mathbf{v}}}$$



3 (c) Show that $H\vec{z} = \|\vec{z}\|_2 \vec{e}_1$. (Recall: $\vec{v} = \|\vec{z}\|_2 \vec{e}_1 - \frac{\vec{z}}{\|\vec{z}\|_2}$.)

Soln

Note that

$$\begin{aligned}\vec{v}^T \vec{v} &= (\|\vec{z}\|_2 \vec{e}_1 - \vec{z})^T (\|\vec{z}\|_2 \vec{e}_1 - \vec{z}) \\&= (\|\vec{z}\|_2 \vec{e}_1^T - \vec{z}^T) (\|\vec{z}\|_2 \vec{e}_1 - \vec{z}) \\&= \|\vec{z}\|_2^2 \underbrace{\vec{e}_1^T \vec{e}_1}_1 - \|\vec{z}\|_2 \underbrace{\vec{e}_1^T \vec{z}}_{\vec{z}^T \vec{e}_1} \xrightarrow{\text{same}} + \underbrace{\vec{z}^T \vec{z}}_{\|\vec{z}\|_2^2} \\&= \underline{2\|\vec{z}\|_2^2} - 2\|\vec{z}\|_2 \vec{z}^T \vec{e}_1\end{aligned}$$

Scratch

$$H = I - \frac{2\vec{v}\vec{v}^T}{\vec{v}^T \vec{v}}$$

$$H\vec{z} = \left(I - \frac{2\vec{v}\vec{v}^T}{\vec{v}^T \vec{v}}\right) \vec{z}$$

$$= \vec{z} - \frac{2\vec{v}(\vec{v}^T \vec{z})}{\vec{v}^T \vec{v}} \checkmark$$

Programming Notes

- Built-in functionalities
 - backslash (\)
 - qr
- Demonstration/Instructional codes
 - `lsqrfact`: solving least squares using QR
 - `gs`: Gram-Schmidt (for homework)

Visualization of matrix p-norm

- Unit vectors

$$\bullet \quad \|A\|_p := \max_{\vec{x} \neq 0} \frac{\|A\vec{x}\|_p}{\|\vec{x}\|_p} = \max_{\|\vec{x}\|_p=1} \|A\vec{x}\|_p$$

Fitting a circle

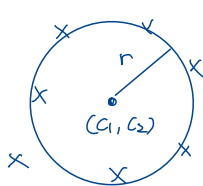
Given data $\{(x_k, y_k) : k=1, \dots, m\}$.

Find c_1 , c_2 , and r as in

$$(x - c_1)^2 + (y - c_2)^2 = r^2$$

$$x^2 - 2xc_1 + c_1^2 + y^2 - 2yc_2 + c_2^2 = r^2$$

$$x^2 + y^2 = 2xc_1 + 2yc_2 + \underbrace{r^2 - c_1^2 - c_2^2}_{c_3}$$



Note $c_3 = r^2 - c_1^2 - c_2^2$
 $\Rightarrow r = \sqrt{c_1^2 + c_2^2 + c_3}$

$$\begin{bmatrix} 2x_1 & 2y_1 & 1 \\ 2x_2 & 2y_2 & 1 \\ \vdots & \vdots & \vdots \\ 2x_m & 2y_m & 1 \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ c_3 \end{bmatrix} = \begin{bmatrix} x_1^2 + y_1^2 \\ x_2^2 + y_2^2 \\ \vdots \\ x_m^2 + y_m^2 \end{bmatrix}$$

Parametric rep'n

$$\begin{cases} x(\theta) = c_1 + r \cos(\theta) \\ y(\theta) = c_2 + r \sin(\theta) \end{cases}$$

for $\theta \in [0, 2\pi]$