

Galois field 2^4

Let's take a quick look at $\text{GF}(2^4)$ before moving on to the main event ($\text{GF}(2^8)$).

These are the polynomial equivalents of the binary numbers with 4 digits, all except 0000. There are 15 elements.

Degree 0:

$$1$$

Degree 1:

$$x \quad x + 1$$

Degree 2:

$$x^2 \quad x^2 + 1 \quad x^2 + x \quad x^2 + x + 1$$

The new ones are degree 3. Write them as binary equivalents:

$$1000 \quad 1001 \quad 1010 \quad 1011$$

$$1100 \quad 1101 \quad 1110 \quad 1111$$

irreducible polynomial

We need an irreducible polynomial. In $\text{GF}(2^3)$ we could choose from $x^3 + x + 1$ and $x^3 + x^2 + 1$.

Now we need a polynomial of degree 4, which may be obtained either by multiplying a degree 3 times a degree 1, or by multiplying two of degree 2.

Start with the degree 3's and multiply them by either 10 or 11:

```
10000  10010  10100  10110
11000  11010  11100  11110
11000  11011  11110  11101
10100  10111  10010  10001
```

Degree 2's as binary equivalents:

```
100  101  110  111
```

And all multiplied together:

```
10000  10100  11000  11100
10001  11110  11011
10100  10011
10101
```

Write them all without the leading 1 (grouping by the next digit) to see the pattern:

```
10 + 000 001 010 011 100 101 110 111
11 + 000 --- 010 011 100 --- 110 ---
```

I find only 13, so 3 are missing. These are

```
11001 11101 11111
```

We shall choose $11001 = x^4 + x^3 + 1$.

generator

I tried **0x03** and found it is not a generator for this field. Rather than guess again, I wrote a short script that calls the `gmultiply` routine. It's modified to use the irreducible polynomial from above ($11001 =$ decimal 25) and do the mod operation if $n > 15$. I'll put the code at the end

The output is

```
2  4  8  9 11 15  7 14  5 10 13  3  6 12  1  2  4
3  5 15  8  1  3  5 15  8  1  3  5 15  8  1  3  5
4  9 15 14 10  3 12  2  8 11  7  5 13  6  1  4  9
  5  8  3 15  1  5  8  3 15  1  5  8  3 15  1  5  8
..
15  3  8  5  1 15  3  8  5  1 15  3  8  5  1 15  3
```

Unexpectedly, 2 and 4 are generators but 3 and 5 are not. So

```
0010
0010
----
0100 => 0100
0010
----
1000 => 1000
0010
----
10000
11001 mod
-----
1001 => 1001
0010
```

```

-----
10010
11001 mod
-----

```

1011 => 1011

..

I won't show the whole thing.

I got

```

0100 1000 1001 1011 1111 0111 1110 0101
1010 1101 0011 0110 1100 0001 0010

```

and compare that with

```

 2  4  8  9 11 15  7 14  5 10 13  3  6 12  1  2

```

That's a match. With the powers, we can compute a table of logarithms.

Here they are: the elements are in the first row and their logarithms below.

```

 2  4  8  9 11 15  7 14  5 10 13  3  6 12  1  2
 1  2  3  4  5  6  7  8  9 10 11 12 13 14 15 16

```

According to this, 7 and 14 should be multiplicative inverses because their logarithms add to 15, whose anti-logarithm is 1.

Try it:

```

0111 = 7
1110 = 14
-----
111
111
111

```

```

-----
101010
11001 mod
-----
11000
11001 mod
-----
1

```

This is pretty tedious. I wrote code to carry out the procedure, and check the above multiplicative inverses. They are all correct.

The last thing is to try an example of the extended Euclidean algorithm. Find the multiplicative inverse of 5.

a	b	q	bq	r
11001	101	110	11110	111
101	111	1	111	10
111	10	11	110	1

```

-----
111 = 11001 - (110)101
10 = 101 - (1)111
1 = 111 - (11)10
  = 111 - (11)[101 - (1)111]
  = (10)111 - (11)101
  = (10)[11001 - (110)0101] - (11)101
  = (1100)101 + (11)101
  = (1111)101
  = (15)(5)

```

which is correct.

Here is the code I used. First is the current version of gmultiply:

```
def gmod(n,N):
```

```

D = { 3:19, 4:25, 8:283 }
P = D[N]
b = len(bin(P))
while True:
    if n < N**2:
        return n
    n = n ^ (P << (len(bin(n)) - b))

def gmultiply(a,b,N=8): # default N = 8
    s = bin(b)[2:][::-1]
    r = 0
    for c in s:
        if c == '1':
            r = r ^ a
        a = a << 1
    return gmod(r,N)

```