Euclidean algorithm

Goal: find the gcd (greatest common divisor) of two integers a and b.

Example:

```
421 = 111 \times 3 + 88
111 = 88 \times 1 + 23
88 = 23 \times 3 + 19
23 = 19 \times 1 + 4
19 = 4 \times 4 + 3
4 = 3 \times 1 + 1
3 = 1 \times 3 + 0
```

The last non-zero remainder is 1. This is gcd(421,111). Hint: 421 is on this <u>list</u>.

Three similar Python implementations:

```
if b > a:
    a,b = b,a

def gcd(a,b):
    while True:
        m = a % b
        if m == 0:
            return b
        a,b = b,m

print gcd(421,111) # 1
print gcd(60,24) # 12
print gcd(11838*2888, 99991987*2888) # 2888
```

The while True isn't strictly necessary:

```
# requires a > b
def gcd(a,b):
    m = a % b
    while m != 0:
        a,b = b,m
        m = a % b
    return b
```

Recursive version:

```
def gcd(a,b):
    m = a % b
    if m == 0:
        return b
    return gcd(b,m)
```

Explanation

Suppose we have two integers a and b and we do m = a mod b and m is not zero.

The idea is that if a and b have a common divisor, which we seek, then so does m.

The mod operation can be expressed as

```
m = a - nb
```

where nb < a but (n+1)b > a. (If there were an integer n so that nb = a, then we would get zero remainder and return b as the result).

Suppose a and b have a common factor f. We can factor f from each term of the previous equation:

```
m = a - nb

m = f(a/f - nb/f)
```

leaving integer terms inside the parentheses. But clearly [m] is also evently divided by [f].

Thus the problem is reduced, because now we can just find gcd(b,m), because b and m also have the common factor f, and the same logic applies.

Finding the multiplicative inverse

Suppose we know e and want to find d such that $ed \mod p = 1$.

If there exists such a d then

```
\begin{array}{l} \text{ed mod p} = 1 \\ \text{ed mod p} + \text{np mod p} = 1 \end{array}
```

and because $np \mod p = 0$:

$$(ed + np) \mod p = 1$$

We will use this fact in a bit.

link

working through Euclid's theorem

Consider gcd(81,57)

$$81 = 1(57) + 24$$

 $57 = 2(24) + 9$
 $24 = 2(9) + 6$
 $9 = 1(6) + 3$
 $6 = 2(3) + 0$

So gcd(81,57) = 3. What we do next is to find integers (one negative) such that

```
p(a) + s(b) = 3

p(81) + s(57) = 3
```

Rearrange the next to last line (line no. 4) from the gcd calculation:

```
3 = 9 - 1(6)
```

Substitute for 6 from the line no. 3

$$6 = 24 - 2(9)$$
 $3 = 9 - 1[24 - 2(9)]$
 $= 3(9) - 1(24)$

Substitute for 9 from line no. 2

Substitute for 24 from line no. 1

$$24 = 81 - 1(57)$$

$$3 = 3(57) - 7(24)$$

$$= 3(57) - 7[81 - 1(57)]$$

$$= 10(57) - 7(81)$$

$$= -7(81) + 10(57)$$

Thus, we've shown that

$$3 = p(a) + s(b)$$

where p = -7 and s = 10.

But this means that

$$3 = 10(57) - 7(81)$$

and what this means is that $10(57) = 3 \mod 81$.

Paraphrasing the <u>link</u>:

We want to do arithmetic modulo n, and in particular, for division we need to find the inverse of integers mod n. For large numbers, this turns out to be a difficult task (and not always possible).

It is known that a number x has an inverse mod n (i.e., a number y so that $xy = 1 \mod n$) if and only if

```
gcd(x, n) = 1.
```

Our example above had gcd = 3, but we are really interested in cases where gcd = 1 because then we are guaranteed that an inverse exists. The simplest way to arrange this is the choose n prime, because then every integer less than n has gcd = 1 with n.

The following simple Python function finds the inverse:

```
# requires a > b
def eea(a,b):
    s, t = 1, 0
    u, v = 0, 1
    while b != 0:
        q = a / b
        a, b = b, a % b

    tmp = s, t
    s = u - (q * s)
    t = v - (q * t)
    (u,v) = tmp
    return u
```

I combine this with a brute-force approach to finding the inverse:

```
# requires a > b
def caveman(m,p):
    n = 2
    while m*n % p != 1:
        n += 1
    return n
```

and use the two functions like so

```
p = 127
for i in range(2,p):
    r = my_eea(p,i)
    if r < 0: r += p
    print i, caveman(i,p), r</pre>
```

Output

```
> python euclidean.py
2 64 64
3 85 85
4 32 32
5 51 51
6 106 106
7 109 109
8 16 16
...
```

The two methods agree and we may also check by doing (for example)

```
5 * 51 = 255, 127 * 2 = 254
6 * 106 = 636, 127 * 5 = 635
```

For each pair we have that $pr \mod n = 1$.