

## Euclid's algorithm

Consider two natural numbers  $a$  and  $b$ . Usually  $a$  is allowed to be an integer (i.e., it can be negative), but to keep things simple here we will say that  $a, b \in \mathbb{N}$ ,  $a$  and  $b$  are examples of counting or natural numbers, also known as the positive integers.

We can find their *greatest common divisor*, written  $(a, b)$ . First we write the unique prime factorization of  $a$  and  $b$ :

$$\begin{aligned} 180 &= 2 \times 2 \times 3 \times 3 \times 5 \\ 140 &= 2 \times 2 \times 5 \times 7 \\ \gcd(140, 180) &= 2 \times 2 \times 5 = 20 \end{aligned}$$

Pick out the common factors and the  $\gcd(a, b)$  will be their product. (We will develop a theorem on unique prime factorization in another chapter).

However, it is important that we do not need to actually factor  $a$  and  $b$ , as we'll see.

The algorithm works like this. Find integers  $r \geq 0$  and  $q > 0$  such that

$$a = b \cdot q + r$$

- If  $r = 0$  we are done:  $b$  divides  $a$  equally. Otherwise
  - switch  $a = b$  and  $b = r$  and repeat.

Then  $b$  is the gcd of the original  $a$  and  $b$ .

In our example

$$180 = 140 \times 1 + 40$$

$$140 = 40 \times 3 + 20$$

$$40 = 20 \times 2 + 0$$

$$\text{gcd} = 20$$

*Proof.*

Let  $n = a + b$  and suppose that  $p$  evenly divides  $a$  and  $b$ , that is,  $p$  is a common factor of both.

Then  $a = px$  and  $b = py$  so

$$n = a + b = px + py = p(x + y)$$

$p$  evenly divides  $a + b$ .

More important for us,  $p$

$$a - b = px - py = p(x - y)$$

$p$  evenly divides  $a - b$ .

□

To speed things up, we find the largest multiple of  $b$ ,  $mb$  such that

$$mb < a < (m + 1)b$$

And then we repeat, finding the common factor of  $b$  and  $a - mb$ .

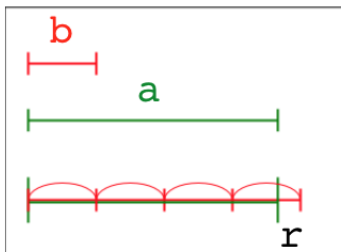
**longer proof**

Here is the reason this works. First, we can always find  $q$  and  $r$  such that

$$a = b \cdot q + r$$

where  $0 \leq r < b$  (since if  $r = b$ , then  $a = b \cdot (q + 1) + 0$ ).

This is a version of the Archimedean property for positive integers.



It may be paraphrased by saying

given a bathtub full of water and a teaspoon, it is possible to empty the bathtub.

Either  $a = b \cdot q$  and we are done or:

$$b \cdot q < a < b \cdot q + b$$

So then

$$a - bq > 0$$

$$a - bq < b$$

With  $r = a - bq$ , we obtain  $0 < r < b$ .

Let  $u$  be the largest integer that divides both  $a$  and  $b$  (the greatest common divisor)

$$a = su$$

$$b = tu$$

Then

$$su = q \cdot tu + r$$

$$r = su - q \cdot tu$$

$$r = u(s - q \cdot t)$$

So  $u$  divides  $r$ .

Hence every common divisor of  $a$  and  $b$  is also a divisor of  $b$  and  $r$ .

#### recursive program

Here are two examples of programs in different styles that implement the algorithm (with no error checking):

```
def gcd(a,b):
    r = a % b
    if r == 0:
        return b
    return gcd(b,r)
```

```
def gcd(a,b):
    r = a % b
    while r != 0:
        a,b = b,r
        r = a % b
    return b
```

The first version is *recursive*, it may call itself. The second uses a **while** loop to accomplish the same thing.

We're using the built-in mod function `%` from Python, but could do something like this:

```
def mod(a,b):
    if a == 0 or b == 0:
        raise ValueError
    if a == b:
        return 0
    if a < b:
        a,b = b,a
```

```
c = b + b
if c > a:
    return a - b
next = c + b

while True:
    if next > a:
        return a - c
    c = next
    next = c + b
```