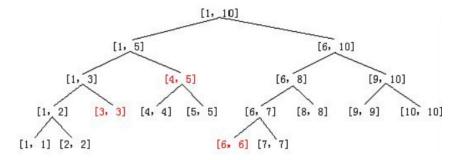
## Cover

Input file: standard input
Output file: standard output

Time limit: 1 second Memory limit: 256 megabytes

Segment tree is a kind of binary tree, and each of its node can cover an interval. A node in the segment tree is a leaf node if and only if it simply covers one index, which means its interval only includes an integer. To a non-leaf node covering interval [a,b], its left child's interval is  $[a,\lfloor\frac{a+b}{2}\rfloor]$ , and its right child's interval is  $[\lfloor\frac{a+b}{2}\rfloor+1,b]$ . So we can determine the structure of one segment tree uniquely when we know the coverage interval of its root node [1,n].

It can be proved that to a given interval [l, r], there is only one **gorgeous-cover** method, we note it as w[l, r], which represents the minimal number of intervals needed to cover [l, r].



For example, in this [1,10] segment tree, w[3,6]=3. The selected intervals have been painted red ([3,3],[4,5],[6,6]).

Given n, l, r, k, you need to find the shortest interval [l', r'] in the [1, n] segment tree which satisfies  $l' \le l \le r \le r'$  and  $w[l', r'] \le k$ . It's guaranteed that  $k \le w[l, r]$ .

## Input

The first line contains an integer  $T(T \le 10^5)$ , the number of the test cases.

Each of the next T lines contains 4 integers  $n, l, r, k (1 \le l \le r \le n \le 10^{18}, 1 \le k \le w[l, r])$ .

## Output

To each of the test cases, output one integer in a line, the length of the shortest interval. Note that the length of interval [l, r] is r - l.

## Example

standard input	standard output
2	9
10 3 6 1	5
10 2 6 3	