Algorithm simplification for Heisenberg group case

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1 Introduction and Background

Following [BBPT, Algorithm 1] we implement a Python program to implement said algorithm in the setting of a Heisenberg substitution. We explain the variation from Algorithm 1 in our python program implementation. Specifically, the use of certain functions in $Aux_Check.py$ and $Aux_0.py$; and deviation from the algorithm in lines 10-11 and 14-18.

Let us recall shortly the setting from [BBPT, Section 3.3]. The action in the Heisenberg group $G := H_3(\mathbb{R})$, is given by

$$(x, y, z) \cdot (a, b, c) := (x + a, y + b, z + c + \frac{1}{2}(xb - ay))$$
 for any $(x, y, z), (a, b, c) \in H_3(\mathbb{R})$.

The metric d_H is left invariant metric on $H_3(\mathbb{R})$ satisfying $d_H((x,y,z),(a,b,c)) = \|(x,y,z)^{-1}\cdot(a,b,c)\|_{CK}$, with $\|\cdot\|_{CK}$ being the Cygan-Koranyi norm on $H_3(\mathbb{R})$ given by

$$\|(x,y,z)\|_{CK} := \sqrt[4]{(x^2+y^2)^2+z^2}.$$

An underlying dilation family on $H_3(\mathbb{R})$ is given by $(D_{\lambda})_{\lambda>0}$ where

$$D_{\lambda}: H_3(\mathbb{R}) \to H_3(\mathbb{R}), \quad D_{\lambda}(x, y, z) := (\lambda x, \lambda y, \lambda^2 z),$$

The lattice in $H_3(\mathbb{R})$ we consider is $\Gamma := H_3(2\mathbb{Z})$, which is invariant under D_{λ} for any $\lambda \in \mathbb{N}$. $V := [-1,1)^3$ defines a fundamental domain for Γ in G. Altogether $\mathcal{D}_H = (G, d_H, (D_{\lambda})_{\lambda > 0}, \Gamma, V)$ is a dilation datum.

We are interested in checking whether inclusions of the form

$$xF_1 \subseteq D^N(\gamma)V(N, F_2) \tag{1.1}$$

hold, for $x, \gamma \in \Gamma$, finite nonempty sets $F_1, F_2 \subseteq \Gamma$ and $N \in \mathbb{N}$.

2 Simplifications for the Heisenberg setting

We make some basic observations regarding the Heisenberg substitution case. First, it is easily verifiable that $V(1) = [-\lambda_0, \lambda_0)^2 \times [-\lambda_0^2, \lambda_0^2)$. We now turn to establish some estimates regarding V(n).

Lemma 2.1. Let $n \in \mathbb{N}$ and $\lambda_0 \geq 3$. Then th following estimates hold. Let $\pi_X : H_3(\mathbb{R}) \to \mathbb{R}$, $\pi_Y : H_3(\mathbb{R}) \to \mathbb{R}$ and $\pi_Z : H_3(\mathbb{R}) \to \mathbb{R}$ be projections of the first, second and third coordinates accordingly.

• $\pi_X(V(n)) = \pi_Y(V(n))$ and

$$\pi_X\big(V(n)\big) = \begin{cases} \left[-\lambda_0^n, \lambda_0^n\right), & \text{if } \lambda_0 \notin 2\mathbb{N} \\ \left[-\lambda_0^n - \sum_{k=1}^{n-1} \lambda_0^k, \lambda_0^n - \sum_{k=1}^{n-1} \lambda_0^k\right), & \text{if } \lambda_0 \in 2\mathbb{N} \end{cases}$$
(2.1)

• For all $(x,y) \in \pi_X(V(n)) \times \pi_Y(V(n))$, the set $\{z \in \mathbb{R} : (x,y,z) \in V(n)\}$ is a connected set of length $2\lambda_0^{2n}$.

One can verify the previous lemma inductively, which is left to the interested reader.

2.1 Shifting sets with respect to the Heisenberg action

Looking at Algorithm 1, we must verify inclusion (1.1) for $|D^N[V] \cap \Gamma|$. We need to check for each such x, at most $|B(e, R_{F_2} + C_+) \cap \Gamma|$ values of γ_x for (1.1). For each such γ_x , we must also compute $D^N(\gamma_x)V(N, F_2)$. This last computation demands naively |V(N)| many steps. This implies that a naive application of algorithm has a run time of the order λ_0^{8N} .

Luckily, our sets are of rectangular form. i.e., there are $a_1, a_2, c_1, c_2 \in 2\mathbb{Z}$ such that F is of the form

$$F = \left([a_1, a_2)^2 \times [c_1, c_2) \right) \cap \Gamma.$$

It can be also represented as a union of "discrete" intervals,

$$F = \bigsqcup_{m_1, m_2 \in 2\mathbb{Z}, a_1 < m_i < a_2} \{ (m_1, m_2) \} \times ([c_1, c_2) \cap 2\mathbb{Z}).$$

In general, sets of the form

$$F = \bigcup_{(m_1, m_2) \in W} \{ (m_1, m_2) \} \times \left([c_1^{(m_1, m_2)}, c_2^{(m_1, m_2)}) \cap 2\mathbb{Z} \right)$$
 (2.2)

where $W \subseteq (2\mathbb{Z})^2$, will be called a union of discrete intervals, or UDI in short.

Lemma 2.2. Let F be a UDI, $\gamma \in \Gamma$ and $N \in \mathbb{N}$. Then, both γF and $V(N, F) \cap \Gamma$ are UDIs.

Proof. Note first that the map $\eta \mapsto \gamma \cdot \gamma$ is a bijection for any γ . Note also that

$$V(N+1,F) = D\big[V(N,F) \cap \Gamma\big] \cdot V(1).$$

It holds for distinct $\eta_1, \eta_2 \in \Gamma$ that $D(\eta_1)V(1) \cap D(\eta_2)V(1) = \emptyset$. It then follows that is enough to verify the claim for sets of the form $F = \{(m_1, m_2)\} \times ([c_1, c_2) \cap 2\mathbb{Z})$. For such sets, computing γF gives that $\gamma F = \{(\gamma_1 + m_1, \gamma_2 + m_2)\} \times (\gamma_3 + \Delta(\gamma_1, \gamma_2; m_1, m_2) + [c_1, c_2) \cap 2\mathbb{Z})$, where $\Delta(\gamma_1, \gamma_2; m_1, m_2) := \frac{1}{2}(\gamma_1 m_2 - \gamma_2 m_1)$. By an inductive argument, it is enough to show that $V(1, F) \cap \Gamma$ is a UDI. Note that

$$D[F] = \{(\lambda_0 m_1, \lambda_0 m_2)\} \times ([\lambda_0^2 c_1, \lambda_0^2 c_2 - 2\lambda_0^2] \cap 2\lambda_0^2 \mathbb{Z}),$$

which in turn implies for any $w_1, w_2 \in [-\lambda_0, \lambda_0) \cap 2\mathbb{Z}$ we have

$$D[F] \cdot \Big(\{(w_1, w_2)\} \times [-\lambda_0^2, \lambda_0^2) \Big) = \Big\{ (\lambda_0 m_1 + w_1, \lambda_0 m_2 + w_2) \Big\} \times \Big(\lambda_0 \Delta(m_1, m_2; w_1, w_2) + \left[\lambda_0^2(c_1 - 1), \lambda_0^2(c_2 - 1)\right] \Big).$$

We deduce that

$$V(1,F) = \bigcup_{w_1,w_2 \in [-\lambda_0,\lambda_0) \cap 2\mathbb{Z}} \left\{ (\lambda_0 m_1 + w_1, \lambda_0 m_2 + w_2) \right\} \times \left(\lambda_0 \Delta(m_1, m_2; w_1, w_2) + \left[\lambda_0^2(c_1 - 1), \lambda_0^2(c_2 - 1) \right) \right),$$

We therefore prefer to operate on sets more efficiently by saving them as UDIs. The computations of the sets V(N, F) for UDIs is implemented in the functions ShiftFaces, DilatEdge and FundemFaceIter in Aux 0.py.

2.2 Searching for γ_x for x

In algorithm 1 line 11, we wish to compute a search radius to find possible γ_x for each x, for a susupected testing domain $F \in \Gamma$. Note first, that as long as we find a $\gamma_x = (\gamma_1, \gamma_2, \gamma_3)$ satisfying inclusion (1.1) for each x, it does not matter how it was found. Therefore, it is enough that given a list of γ_x 's they satisfy the inclusion $xT_0 \subseteq D^n(\gamma)V(n,T)$ for all $x \in D^n[V]$ to show T is a testing domain.

However, given that the dilation with respect to the XY coordinates behaves as in the case of Euclidean block substitutions, we may compute efficiently what $\gamma_1, \gamma_2 \in 2\mathbb{Z}$ may be. Recall that all our suspected testing domains are discrete boxes of the form

$$F = \{-2a_1, -2a_1 + 2, ..., 2a_2 - 2\}^2 \times ([-c_1, c_2) \cap 2\mathbb{Z})$$

Since $V(n,F) = \bigsqcup_{\gamma \in F_0} D^n(\gamma)V(n)$, it follows from equation (2.1) that

$$\pi_X\big(V(n,F)\big) = \begin{cases} \left[-(2a_1+1)\lambda_0^n, (2a_2-1)\lambda_0^n \right), & \text{if } \lambda_0 \notin 2\mathbb{N} \\ \left[-(2a_1+1)\lambda_0^n - \sum_{k=1}^{n-1} \lambda_0^k, (2a_2-1)\lambda_0^n - \sum_{k=1}^{n-1} \lambda_0^k \right), & \text{if } \lambda_0 \in 2\mathbb{N} \end{cases}$$

We denote $W_N(a_1, a_2)$ As our suspected testing domains are mostly of the form

$$F_0 = \{-2, 0\}^2 \times ([-c_1, c_2) \cap 2\mathbb{Z}),$$

we can be more explicit in our estimates by plugging in $a_1 = 1$ and $a_2 = 1$. This gives us a set $\mathcal{N}_{XY}(x, N; F, F_0) \in (2\mathbb{Z})^2$ such that $(\gamma_1, \gamma_2) \in \mathcal{N}_{XY}(x, N; F, F_0)$ if and only if

$$x_i + \Pi_X[F] \subseteq \lambda_0^N \gamma_i + \pi_X(V(N, F_0))$$
 for $i \in \{1, 2\}$.

A direct computation gives us that $|\mathcal{N}_{XY}(x, N; F, F_0)| \leq 4$. It then remains to find whether there exists values of $\gamma_3 \in 2\mathbb{Z}$ satisfying inclusion (1.1).

The set F attains $(a_1 + a_2)^2$ many distinct XY-values. Let us denote

$$I_{(m_1,m_2)}(x;F) := \{ z \in \mathbb{R} : (x_1 + m_1, x_2 + m_2, z) \in xF \} \quad \text{for any} \quad m_1, m_2 \in \left[-2a_1, 2(a_2 - 1) \right) \cap 2\mathbb{Z}.$$

Let us denote

$$J_{(k_1,k_2)}(\gamma,N;F_0) := \{ z \in \mathbb{R} : (k_1,k_1,z) \in D^N(\gamma)V(N,F_0) \}$$

for all (k_1, k_2) which are projections of the set $D^N(\gamma)V(N, F_0)$. We can compute these sets by the function FindCorrLst in $Aux_Check.py$. If $I_{(m_1,m_2)}(x;F) = [\ell_{(m_1,m_2)},u_{(m_1,m_2)})$ and $J_{(k_1,k_2)}(\gamma,N;F_0) = [\tilde{\ell}_{(k_1,k_2)}^{(\gamma)},\tilde{u}_{(k_1,k_2)}^{(\gamma)})$, then $xF \subseteq D^N(\gamma)V(N,F_0)$ if and only if for every (m_1,m_2) we have that

$$u_{(m_1, m_2)} \le \tilde{u}_{(k_1, k_2)}^{(\gamma)}$$
 and $\ell_{(m_1, m_2)} \ge \tilde{\ell}_{(k_1, k_2)}^{(\gamma)}$

when $x_1 + m_1 = k_1$ and $x_2 + m_2 = k_2$ for all (m_1, m_2) . Alternatively, we may ask whether there exists an integer $s \in 2\mathbb{Z}$ such that

$$u_{(m_1,m_2)} \le \tilde{u}_{(k_1,k_2)}^{(\gamma_1,\gamma_2,0)} + s \cdot \lambda_0^{2N} \quad \text{and} \quad \ell_{(m_1,m_2)} \ge \tilde{\ell}_{(k_1,k_2)}^{(\gamma_1,\gamma_2,0)} + s \cdot \lambda_0^{2N}$$
 (2.3)

for all $m_1, m_2 \in [-2a_1, 2(a_2 - 1)) \cap 2\mathbb{Z}$. Each such pair of inequalities give certain bounds on s, which must be satisfied simultaneously for some $s \in 2\mathbb{Z}$. These bounds are computed in the function PossShift in $Aux_Check.py$

References

[BBPT] Ram Band, Siegfried Beckus, Felix Pogorzelski, and Lior Tenenbaum. Periodic approximation of substitution subshifts, In preparation.