A APPENDIX

A.1 Hausdorff Distance-Aware Progressive Polyhedron Simplification

${\bf Algorithm~1:} \ {\bf Progressive~Polyhedron~Simplification}$

```
Input : Polyhedron P
    Output: Multi-LOD Representation P'
 _{1}P'\leftarrow P;
 <sup>2</sup> for i^{th} facet f'_i in P' and f_i in P do
    f_i'.assoc\_facets \leftarrow \{f_i\};
 4 while Lowest LOD not reached do
         for each vertex v in P' do
            set v as unprocessed;
 6
 7
         while P' has unprocessed vertex v do
              V_c \leftarrow all vertices in P' connected to v;
 8
              for each v' in V_c + v do
                   set v' as processed;
10
              f' \leftarrow facet constructed with vertices in V_c;
11
12
              f'.interior \leftarrow v.is\_protruding;
              f'.exterior \leftarrow v.is\_recessing;
13
              for each facet f around vertex v do
14
                   f'.interior \leftarrow f'.interior&f.interior;
15
                   f'.exterior \leftarrow f'.exterior & f.exterior;
16
                   for each facet fa in f.assoc_facets do
17
 18
                         f'.assoc\_facets \leftarrow f'.assoc\_facets + f_a;
                        f_a.assoc\_facets \leftarrow f_a.assoc\_facets + f';
 19
                   remove f from P';
20
              if f'.interior = f alse && f'.exterior = f alse then
21
                  determine state of f' evaluating f'. assoc\_facets;
22
              add f' to P';
23
              f'.removed\_vertex \leftarrow v;
24
25
         for each facet f' in P' do
              for each point p in sample (f', \epsilon) do
26
                   d \leftarrow distance(p, f'.assoc\_facets) + \frac{\epsilon}{\sqrt{2}}
27
                    f'.hdist \leftarrow max(f'.hdist, d);
28
                   P'.hdist \leftarrow max(P'.hdist, d);
29
         for each facet f in P do
30
              for each point p in sample (f, \epsilon) do
31
                    f' \leftarrow \text{closest facet in } f.assoc\_facets \text{ to } p;
32
                   d \leftarrow distance(p, f') + \frac{\epsilon}{\sqrt{2}};
33
                    f'.phdist \leftarrow max(f'.phdist, d);
34
                   P'.phdist \leftarrow max(P'.phdist, d);
35
         for each facet f' in P' do
36
              if f'.removed\_vertex \neq null then
37
               encoding f'.removed_vertex;
38
              \begin{aligned} &\text{encoding } f'.interior?0: \lceil \frac{f'.hdist*255}{P'.hdist} \rceil; \\ &\text{encoding } f'.exterior?0: \lceil \frac{f'.phdist*255}{P'.phdist} \rceil; \end{aligned}
40
41 encoding the remaining primitives on P' as LOD 0;
```

A.2 Performance Analysis of Computing with Facet-Association and BVH-tree

Consider the calculation of Hausdorff distances and proxy Hausdorff distances for the facets of a low-LOD polyhedron P' in relation to the original polyhedron *P*. During the Hausdorff distance calculation, we construct a BVH tree comprising |P| facets, and we sample $SR \times |P'|$ points from the surface of P'. Each BVH tree lookup operation consumes time equivalent to $t \times \log |P|$, with t representing a constant coefficient associated with tree lookups in a given execution environment. Consequently, the time required for calculating the Hausdorff distance is proportional to $SR \times |P'| \times t \times \log |P|$. Similarly, the time spent on computing the proxy Hausdorff distance is proportional to $SR \times |P| \times t \times \log |P'|$. Importantly, it is evident that |P'| < |P|, resulting in the inequality $SR \times |P'| \times t \times \log |P| < SR \times |P| \times t \times \log |P'|$. Therefore, calculating the proxy Hausdorff distance demands more time than computing the Hausdorff distance. Additionally, as more rounds of simplification are performed, the value of |P'| decreases, resulting in reduced computation time for low-LOD polyhedrons.

With the FA method, each sampled point on facet f merely needs to compute its distance from the facets associated with f. The time required for calculating the Hausdorff distance and proxy Hausdorff distance can be estimated as $SR \times |P'| \times t \times r' \times |P|$ and $SR \times |P| \times t \times r \times |P'|$, respectively. Here, r' represents the average portion of facets in P associated with one facet in P', and r represents the average portion of facets in P' associated with one facet in P. Remarkably, as the facets in low-LOD polyhedrons and the original polyhedrons are associated region by region, it can be approximated that r is roughly equal to r'. Consequently, calculating the Hausdorff and proxy Hausdorff distances consumes approximately the same amount of time. Moreover, as the simplification process progresses, the values of both r and r' increase, while the value of |P'| decreases. Consequently, the computation time does not exhibit a strictly decreasing trend with varying LODs.