Operating Systems 2024/2025

TP Class 02 - Processes

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Some slides based on previous versions from Bruno Cabral, Paulo Marques and Luis Silva.

operating system

noun

the collection of software that directs a computer's operations, controlling and scheduling the execution of other programs, and managing storage, input/output, and communication resources.

Abbreviation: OS

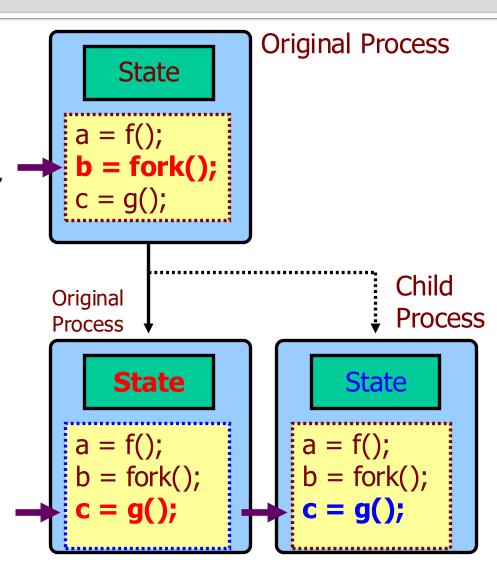
Source: Dictionary.com



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Process Model

- Process creation in Unix is based on spawning child processes which inherit all the characteristics of their fathers
 - Variables, program counter, open files, etc.
 - Spawning a process is done using the fork() system call
- After forking, each process will be executing having different variables and different state.
 - The Program Counter will be pointing to the next instruction
 - Changing a variable in the child program doesn't affect its father (and vice-versa)



Process Management

- Each process has an unique identifier (PID). Each process has a father, which is also identified (PPID).
- pid t getpid(void);
 - Returns the PID of the current process.
- pid_t getppid(void);
 - Returns the PID of the parent process.
- pid_t fork(void);
 - Creates a new process which inherits all its father's state. It returns 0 to the child process and the child's PID to the original process.
- pid_t wait(int* status);
 - Waits until a child process exits. The status of the child is set in status. (status is the return value of the process). Returns -1 in case of error.
- pid t waitpid(pid t who, int* status, int options);
 - Same as wait() but allows to wait for a particular child. By using WNOHANG in options, allows for checking if a child has already exited without blocking. 0 in who means "wait for any child". Returns -1 in case of error.

Using processes to do different things

```
#include <stdio.h>
#include <unistd.h>
                                      The key idea is to create
#include <sys/wait.h>
                                      asymmetry between processes
#include <sys/types.h>
int main()
 pid_t id;
 id = fork();
 if (id == 0)
      printf("[%d] I'm the son!\n", getpid());
      printf("[%d] My parent is: %d\n", getpid(), getppid());
 else
      printf("[%d] I'm the father!\n", getpid());
     wait(NULL);
  return 0;
```

And the result is...

```
pmarques@null:~/IPC$ gcc -Wall hello_world.c -o hello_world
pmarques@null:~/IPC$ ./hello_world
[2190] I'm the father!
[2191] I'm the son!
[2191] My parent is: 2190
pmarques@null:~/IPC$ _
```

But why did we do wait(NULL)?

```
...
printf("[%d] I'm the father!\n", getpid());
wait(NULL);
...
```

wait(NULL) ensures that the father process waits for any of its children, discarding child exit status.

Process Termination in UNIX

- A process is only truly eliminated by the operating system when its father calls wait()/waitpid() on it.
 - This allows the parent check things like the exit code of its son's
- Zombie Process: One that has died and its parent has not acknowledged its death (by calling wait())
 - Be careful with this if your are designing servers. They are eating up resources!!
- Orphan Process: One whose original parent has died. In that case, its parent becomes *init* (process 1). (not always true; in some recent versions of operating systems the parent becomes the nearest sub-reaper process a sub-reaper fulfills the role of *init*(1) for its descendant processes.)

Let's generate some Zombies

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <sys/wait.h>
#include <sys/types.h>
void worker() {
  printf("[%d] Hi, I'm a worker process! Going to die...\n",
         getpid());
int main()
  for (int i=0; i<10; i++) {
    if (fork() == 0) {
      worker();
      exit(0);
  printf("[%d] Big father is sleeping!\n", getpid());
  sleep(10);
  return 0;
```

Let's see adoption

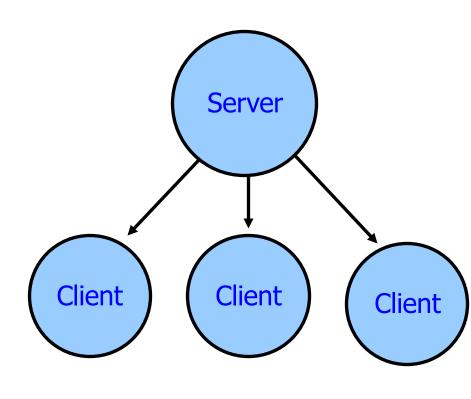
```
#include (...)
void worker()
 sleep(10);
  printf("[%d] Let's see who my dady is: %d\n", getpid(),
         getppid());
int main()
  for (int i=0; i<10; i++)
    if (fork() == 0)
         worker();
         exit(0);
  printf("[%d] Big dady is going away!\n", getpid());
  return 0;
```

How to structure code

```
if ((id - fork()) == 0)
          // Huge amount of code
         else
          // Huge amount of code
                                 Fairly common...
Don't do it!
```

How to structure code

```
void client(int id)
 // Client code
if ((id = fork()) == 0)
  client(id);
  exit(0);
else if (id == -1)
  error();
// Original process code
// ...
```



Note: You still have to consider how to take care of zombies

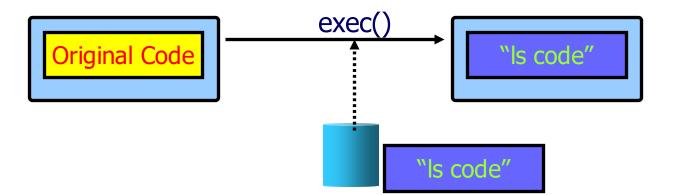
How a process becomes another executable

- Somehow the OS must be able to execute code starting from an executable file
 - e.g. how does the shell (bash) becomes 'ls'?
- int execl(const char *path, const char *arg, ...);
- int execlp(const char *file, const char *arg, ...);
- int execle(const char *path, const char *arg, ..., char *const envp[]);
- int execv(const char *path, char *const argv[]);
- int execvp(const char *file, char *const argv[]);
- "exec family" of functions
 - Allow to substitute the current process executable image by another one. The substitution is complete!
 - The functions that have a 'p' make use of the environment PATH; The functions that have a 'v' make use of a pointer to an array on parameters; The functions that have an 'l' have the parameters passed separated by commas
 - Make sure that the first parameter is the name of the program!

Simple program that lists the files in the current directory

```
pmarques@null:~/IPC$ g++ -Wall list_files.c -o list_files
pmarques@null:~/IPC$ ./list_files
. adopt.c list_files posix_version zombie.c
.. hello_world list_files.c posix_version.c
adopt hello_world.c posix zombie
pmarques@null:~/IPC$ _
```

- Note: A successful exec() never returns
 - The code, the stack, the heap, it's all replaced by the new executable



demo4.c

The corresponding code...

```
#include <stdio.h>
#include <stdib.h>
#include <unistd.h>
#include <sys/types.h>

int main()
{
   if (execlp("ls", "ls", "-a", NULL) == -1)
       perror("Error executing ls: ");
   else
      printf("This cannot happen!\n");

   return 0;
}
```

Using an array
can be more
flexible...
char* ls_param[] = { "ls", "-a", NULL };

if (execvp(ls_param[0], ls_param) == -1)
 perror("Error executing ls: ");

Code to wait for the end of an exec call...

```
makeargv(buf, DELIM, &args);
pid_t child = fork();
if (child == 0)
    execvp(args[0], args);
    perror("Erro ao executar o comando");
    exit(-1);
waitpid(child, NULL, 0);
```

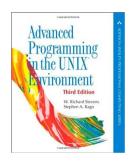
Information about processes in Linux

- Command ps (process status)
 - reports a snapshot of the current processes
 - To see all processes in the OS in full-format listing ps -ef
 - To select all processes from user "user1" ps -ef | grep user1
 - Some process state codes shown in ps
 - Seen when ps is used with a state output modifier
 - E.g.: select all processes and show specific columns
 - ps -eo uid, pid, tty, time, cmd, stat
 - R running
 - T stopped by a job control signal or because it is being traced
 - Z defunct ("zombie")
 - + the process is in the foreground process group.

Class demos included

- Demo01 create process demo01.c
- Demo02 zombies demo02.c
- Demo02 adoption demo03.c
- Demo04 exec

References



Advanced Programming in the UNIX Environment

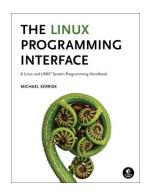
3rd Edition (2013)

W. Richard Stevens, Stephen A. Rago

Addison-Wesley

Chapter 8: Process Control

Where to learn more?



The Linux Programming Interface 2010 Michael Kerrisk No Starch Press

Chapters 24 to 27

INTRODUCTION TO ASSIGNMENT 03-"PROCESSES"

Thank you! Questions?



I keep six honest serving men. They taught me all I knew. Their names are What and Why and When and How and Where and Who.

—Rudyard Kipling