

# Redes de Comunicação 2024/2025

## TP08

### Delay, loss and throughput in packet-switched networks (part 2)

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## TP08: Delay, loss and throughput (part 2)

### *Overview:*

- Transmission of data
- Maximum data rate
- Nyquist formula
- Shannon formula

# Transmission of data

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- Data transmission involves converting bits in electric, electromagnetic or optical signals
- Signals may suffer from:
  - Attenuation: reduction of the amplitude of the signals
  - Delay distortion: different speeds for different frequencies
- The transmission of binary data involves:
  - Modulation (AM, FM, PM, mixing)
  - Power of the information signal vs Power of the noise
  - A propagation delay related with the transmission medium

# Maximum Data Rate (channel capacity)

- Data rate governs the speed of data transmission
- Depends upon 3 factors:
  - Bandwidth available
  - Number or levels in digital signal (resulting from modulation)
  - Quality of the channel – the level of noise
- Two theoretical formulas were developed to calculate the data rate in digital communication systems:
  - One by Nyquist, considering a noiseless channel
  - Other by Shannon, for a noisy channel

# Nyquist formula

- Assumes a noise-free channel
- Defines the theoretical maximum bit rate, from:
  - Bandwidth available (in Hz)
  - Number or levels in the digital signal (resulting from modulation)

## *Factor of “2” in the formula*

arises from the need to sample the signal at least twice per cycle (according to the Nyquist sampling theorem) in order to accurately reconstruct the original signal.

$$C = 2 B \log_2 M \quad \text{bps}$$

C – Maximum Bit Rate

B – Bandwidth

M – Levels of signalling

- Bandwidth is fixed, so data rate is directly proportional to the number of signal levels

# Nyquist formula (exercises)

- Consider a noiseless channel with a bandwidth of 3000 Hz transmitting a signal with two signal levels. What can be the maximum (theoretical) bit rate?

$$\text{BitRate} = 2 * 3000 * \log_2(2) = 6000\text{bps}$$

- We need to send 265 kbps over a noiseless channel with a bandwidth of 20 kHz. How many signal levels do we need?

$$265000 = 2 * 20000 * \log_2(L)$$

$$\log_2(L) = 6.625$$

$$L = 2^{6.625} = 98.7 \text{ levels}$$

*Increasing the number of signal levels...*

...also increases the complexity of the transmission and reception processes, and may require higher signal-to-noise ratios for reliable communication.

# Shannon formula

- Channels in reality are always noisy
- Defines the theoretical highest data rate for a noisy channel, from:
  - Bandwidth available (in Hz)
  - The signal-to-noise (S/N) ratio
  - Bandwidth is fixed, so data rate is directly proportional to S/N

$$C = B \log_2 (1 + S/N) \quad \text{bps}$$

S – Signal power (Watt)

N – Noise Power (Watt)

or...

$$C = B \log_2 (1 + 10^{(SNR(dB)/10)})$$

If S/N ratio in dB

- SNR may be expressed in dB (decibels), given by:

$$SNR(dB) = 10 * \log_{10} (S/N)$$

Example: For a S/N of 1000:  $10 * \log_{10} (1000) = 30 \text{ dB}$

# Shannon formula (exercises)

- A telephone line normally has a bandwidth of 3000 Hz (300 to 3300 Hz) assigned for data communication. The SNR is usually 3162. What will be the capacity for this channel?

$$C = 3000 * \log_2(1 + \text{SNR}) = 3000 * 11.62 = 34860 \text{ bps}$$

- The SNR is often given in decibels. Assume that SNR(dB) is 36 and the channel bandwidth is 2 MHz. Calculate the theoretical channel capacity.

$$\text{SNR(dB)} = 10 * \log_{10}(\text{SNR})$$

$$\text{SNR} = 10^{(\text{SNR(dB)})/10}$$

$$\text{SNR} = 10^{3.6} = 3981$$

$$\text{Hence, } C = 2 * 10^6 * \log_2(3982) = 24 \text{ MHz}$$



# TP08: Summary

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*What we have covered here?*

- Transmission of data
- Attenuation and distortion
- Nyquist and Shannon formulas
- Usage examples (exercises)