A Framework for Managing the Imperfect Modularity of Variability Implementations

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ABSTRACT

In many industrial settings, the common and varying features of related software-intensive systems, as their reusable units, are likely to be implemented by a combined set of traditional techniques. Features do not align perfectly well with the used language constructs, *e.g.*, classes, thus hindering the management of implemented variability. Herein, we provide a detailed framework to capture, model, and trace this imperfectly modular variability in terms of variation points with variants. We describe an implementation of this framework, as a [..1] domain specific language, and report on its application on four [..2] case studies and usage for variability management, showing its feasibility.

1. Introduction

Software product line (SPL) engineering is the common methodological process for developing together a set of related software-intensive systems. The process is intended to achieve mass customization with methodological reuse of the software systems seen as products, leading to lower cost and time-to-market and improved quality. At the domain level, [...¹] the considered systems are commonly described in terms of their common and varying features, as their reusable units, in a *feature model* (FM) [38, 39, 20]. Then, in a forward engineering approach, their features are realized in different software assets, including reusable code assets. Consequently, the objective is that various systems can be systematically derived from a common and managed set of software assets.

[..2] In many industrial settings, it is often the case that variability of code assets is typically realized by a diverse and combined set of traditional techniques, such as inheritance, design patterns, overloading, or generic types. However, their used code units, such as classes, or methods in object-oriented programming, do not align well with the domain features[..3]. This hinders the management of the variability of code assets, which [..4] is extensive in realistic SPL settings.

[...⁵] For example, let us consider the system of JavaGeom [48], as a realistic feature-rich system. It is an open-source geometry library for Java that is easily understandable. Although not presented [...⁶] as an SPL, it is architected around well-identified features[...⁷], being the size of medium-size SPLs. In Figure 1 is given an excerpt of its [...⁸] feature model, which consists of five features. The abstract feature [79] [...⁹] JavaGeomSPL does not require an

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¹removed: the variability of considered systems is commonly

²removed: Although variability is largely studied in the context of software product lines and product families [20], most modern software intensive systems, being organized as a software product line or not, are variability-intensive [??]. In systems that are not part of a product line, that is, variability-rich systems, the implemented variability in their code assets is neither explicit nor documented. This is because the

³removed: or, simply put, neither features nor variation points with variants (their definition is given in Section 2) are by-product of traditional techniques [17].

⁴removed: seems to be extensive in real variability-rich systems [? 78].

⁵removed: For example, let us consider the system of JavaGeom [48], an example of a real variability-rich system, which is not organized as a software product line but is variability-intensive, based on its documentation, and is used in a previous study [78].

⁶removed: and organized

⁷removed: Then, compared to similar real variability-rich systems that are of over 100K LoC and thousands of potential *vp*-s with variants [? 78], JavaGeom has over 35K LoC, being the size of medium-size variability-rich systems.

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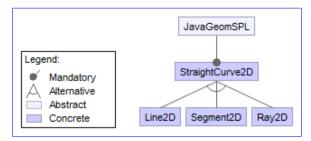


Figure 1: An excerpt of the feature model (FM) of JavaGeom variability-rich system (The FM is created using the FeatureIDE tool [51])

```
1
      /* Class level variation point, vp_AbsLine2D */
                                                                 /* First variant, v_Line2D, of vp_AbsLine2D */
                                                          21
     public abstract class AbstractLine2D extends
2
                                                                public class StraightLine2D extends
                                                          22

→ AbstractSmoothCurve2D implements

→ AbstractLine2D implements SmoothContour2D,

         SmoothOrientedCurve2D, LinearElement2D {
                                                                    Cloneable, CircleLine2D {
                                                          23
                                                                         /* Omitted code */
          /* Method level variation point.
4
                                                                }
                                                          24
             vp_distance */
          /* First variant of vp_distance */
         public double distance(Point2D p) {
                                                          25
                                                                 /* Second variant, v_Segment2D, of vp_AbsLine2D
              return distance(p.x(), p.y());
                                                                 c, */
                                                                public class LineSegment2D extends
                                                          26
                                                                    AbstractLine2D implements Cloneable,
10
          /* Second variant of vp_distance */
                                                                    CirculinearElement2D {
11
         public double distance(double x, double y)
                                                                         /* Omitted code */
                                                          27
12
                                                          28
              Point2D proj = projectedPoint(x, y);
13
              if (contains(proj))
14
15
                  return proj.distance(x, y);
                                                                 /* Third variant, v_Ray2D, of vp_AbsLine2D */
                                                          29
              /* Omitted code */
16
                                                                public class Ray2D extends AbstractLine2D
                                                          30
              return dist;
17

→ implements Cloneable {
18
                                                          31
                                                                         /* Omitted code */
              /* Omitted code */
19
                                                                }
                                                          32
     }
```

Listing 1: An excerpt of variability implementation in the JavaGeom variability-rich system considered as a software product line

implementation, it is used to represent conceptually the JavaGeom [...¹⁰] product line. As for the other four concrete features, StraightCurve2D, Line2D, Segment2D, and Ray2D, a form of their implementation is given in Listing 1, using Java language. At first sight, the mapping between features in Figure 1 and used code units, classes, in Listing 1 seems to be straightforward, but it is more complicated. [...¹¹] The feature names and class names are mostly different, for instance, the feature name Line2D is different from its implementation class StraightLine2D. Besides, the classes extend or implement other classes and interfaces, showing more dependencies at the code level than that features have at the domain level. Then, the [...¹²] variability implemented by method overloading, distance() in lines 4 - 18, is not modeled in the FM. This single example, where only the techniques of inheritance and overloading are used, shows some difficulties in mapping features at the [...¹³] reusable code assets. In real variability-rich systems, the problem is exacerbated as features are often not implemented by a single class or method. This is due partly to the many design choices that influence the usage of different implementation techniques to structure the variability implementations, and partly to the crosscutting nature of features with code units [42].

To overcome such variability management difficulties, existing approaches use annotations, often preprocessor directives [49], or have proposed to put into separate modules, for example, into aspects, feature modules, or delta

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<sup>11</sup>removed: First, the
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¹³removed: code assets of a given variability-rich system.

modules [6], all lines of code that belong to each specific domain feature [7, 71]. In addition, different tools are proposed, for example, pure::variants [73], code tagging [34], or those that propose to use colors, at the representation layer of code, for distinguishing the lines of code that implement each feature [40]. Analysing which technique, tool, or combination of them [56] fits best to implement or manage some variability is hard to conclude as none of them is yet a preferred one [6], while the SPL application domains are themselves largely diverse. But, except for the tool-based approaches, most of the proposed solutions affect the initial decomposition or design of reusable code assets, such as their object-oriented design [7], and require a massive code refactoring[...¹⁴]. As a consequence, when an object-oriented design is refactored, for example, to a feature-oriented design [6], the logical relation among code units is usually lost at the code level. For instance, the alternative relation among the three child features in Figure 1 may lose, which relation is realized by subclasses through the technique of inheritance in Listing 1.

Facing these issues, our work takes the assumption that one can keep unchanged the initial design of [...¹⁵] reusable code assets and manage their variability in terms of *variation points* with *variants* (their definition is given in Section 2), as initial variability abstractions [36]. To the best of our knowledge, no variability management approach is currently addressing the early steps of capturing and modeling the variability of code assets[...¹⁶] [45, 5, 13, 14, 63, 68], especially not when a subset of variability implementation techniques are used in combination [36, 23].

Towards an approach that meets [..¹⁷]out assumption and goal on managing variability implementations in code assets[..¹⁸], we proposed, in a previous work [77], a three step framework to manage such variability when a combined set of traditional techniques is used (Section 4). Such techniques expose a form of what we define as *imperfectly modular variability* (Section 3), as domain features do not perfectly align with the used code units. [..¹⁹]

The framework is devised by analysing up to ten common variability implementation techniques and by studying the related literature. Its main contributions are as follows.

- 1. The framework provides the means to make explicit the variability of code assets by capturing the used variability implementation techniques in terms of variation points (*vp*-s) with variants, as variability abstractions.
- 2. It then makes possible to use the captured *vp*-s with variants to model the variability of code assets in a fragmented way into so-called *technical variability models* (Section 4.2.2). Towards this, five types of *vp*-s are identified and suggested for use, whereas what is a fragment aims at being flexible. It can be a package, file, or class with variability that worth to be modeled separately.
- 3. The framework also introduces the possibility to establish n–to–m trace links between the modeled variability of code assets and their respective features in some feature model at the domain level.

Herein, we extend this work in several aspects:

- First, we [..²⁰] detail the concept of *imperfectly modular variability* further and determine its induced challenges (Section 3), which motivate the whole presented work.
- We present the three steps of our framework with all necessary details, formal definitions [...²¹] and related illustrations (Section 4), mainly by using a running [...²²] case study introduced in Figure 1 and Listing 1.
- We present the implementation of the framework as a small [..²³][..²⁴]domain specific language (DSL) in Scala, publicly available. Its technical foundations are also given in Section 5.
- We report on a series of experiments on the application of this DSL to four [...²⁵] case studies (Section 6). This

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¹⁹removed: The framework is devised by analysing up to ten common variability implementation techniques used in variability-rich systems, which are revealed in a recent catalog of 21 traditional and emerging techniques [76] (*cf.* Section 2.2.3). This catalog is at first based on 7 previous frameworks, taxonomies, and catalogs [6, 29, 57, 60, 62, 61, 23]. For the purpose of our framework, we analysed the characteristics properties of these techniques, such as binding time.

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evaluation was conducted around four research questions to study the encountered implementation techniques, types of vp-s, [...²⁶] multiplicity of trace links, and to evaluate the overhead in modeling with the DSL.

- We also evaluate the capacity of usage of the proposed framework for variability management by adapting a previous work on early consistency checking [75] and showing how the [..²⁷] framework enables to integrate such [..²⁸] checking method in one of the [..²⁹] case studies (Section 7).

Threats to the validity of the framework and its implementation are discussed in Section 8. We also put our contribution into perspective with some related work, including industrial approaches for managing variability at the implementation level (Section 9). Finally, Section 10 concludes this article by summarizing the contributions and discussing future work.

2. Background

[..30] In an SPL, the key aspect of variability management is the ability to trace the specified variability of software systems to its realization in reusable code assets.

2.1. Variability modeling

Since its first presentation by Kang et al. [38], feature modeling has progressively been widely adopted as a means for scoping and modeling the variability of software systems within an SPL in terms of features [26]. A feature model (FM) is a tree structure of features, consisting of *mandatory*, *optional*, *or*, and/or *alternative* logical relations between features with their cross-tree constraints, *implies* and/or *excludes*, that are expressed in propositional logic. In Figure 1 is given an example of an FM, for the JavaGeom [..31] product line. Its main feature, [..32] JavaGeomSPL, represents conceptually the JavaGeom domain. It has a mandatory feature, StraightCurve2D, with three alternative features, Line2D, Segment2D, and Ray2D. While an FM can have cross-tree constraints between features, there is none in this example. Semantically, an FM represents all the valid software systems (*i.e.*, the feature configurations) within an SPL[..33].

2.2. Variability realization

 $[..^{34}]$

2.2.1. Reusable code assets

[..35] In realistic SPL settings, despite the programming paradigm (*e.g.*, object-oriented, or functional), the implementation of reusable code assets complies, mostly, to a commonality and variability approach [23]. Specifically, a domain is decomposed into subdomains[..36], then within each subdomain [..37] the commonality is factorized from the variability. Thus, the reusable code assets consist of three parts: the core, commonalities, and variabilities. The core part is what remains of the system in the absence of any particular feature [80], that is, the assets that are included in any software system. A commonality is a common part of the related variant parts within a subdomain. After the commonality is factorized from the variability and implemented, it becomes part of the core (*i.e.*, it is buried in the core [23]), except when it represents some optional variability. On the other hand, the variant parts are used to distinguish the software systems in the domain. A subdomain can have more than one common and variant part. The

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³⁰removed: Some of the key aspects of the variability management of a system, being it organized as an SPL or not, is the ability to model the realized variability of code assets and then possibly to trace it to its specified variability. In the following, we provide a background on variability modeling, realization, and traceability.

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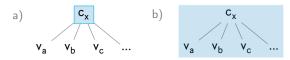


Figure 2: The variation point (vp) concept as: a) a variable point (i.e., the c_x), b) a variable part (i.e., the c_x with variants). The c_x is the common part for variants v_a , v_b , and v_c .

core with the commonalities and variabilities of all subdomains constitute the wholeness of reusable code assets in [...³⁸] an SPL.

2.2.2. Commonality and variability abstractions

While features are commonly used [..39] as problem-oriented abstractions, the commonalities and variabilities in reusable code assets are usually abstracted in terms of variation points (*vp*-s) with variants, respectively, as solution-oriented abstractions [25, 26, 64]. Moreover, unlike features that are mere names, *vp*-s with variants are related to concrete elements in reusable code assets. Originally, "*a variation point identifies one or more locations at which the variation will occur*" [36], whereas the way that a variation point is going to vary is expressed by its variants.

In the literature, there are few different understandings of a vp [74, pg. 16]. Mostly, a vp is recognized (i) as an abstraction of a location or point that varies in reusable code assets (cf. Figure 2a) [36], (ii) as a variable part (cf. Figure 2b) [9], or (iii) as an abstraction of a variability implementation technique [20, pg. 48]. But, these all have a common meaning. Specifically, the *variable part* is like an organizing container [9]. It contains the *location* where some variability happens in reusable code assets (i.e., the variation point), the variants, and the used *technique* to implement them. In this work we use the same meaning of vp-s with variants.

2.2.3. Variability implementation techniques

The used techniques for implementing variability in an [..40] SPL are known as variability implementation techniques [28, 76]. They came from several programming paradigms and are supported by different constructs in different programming languages, which in turn offer different properties for [..41] vp-s with variants. Examples of such techniques are inheritance, preprocessor directives, feature modules, or some design patterns. For example, in Listing 1 is shown the usage of two variability implementation techniques, inheritance and overloading. They implement two variation points (see the comments in Listing 1), lines 1 - 20, and lines 4 - 18, respectively, which have different properties. Both of them have different granularity, of *class*, and *method* levels, and during the product derivation, they will be resolved at different times, *runtime* and *compile time*, respectively.

Because of the large diversity of techniques, an important issue for variability implementation is the ability to evaluate and choose a technique that will fulfill best some given variability requirements. Therefore, for various techniques, several evaluation schemas have emerged, mainly from academia, in the form of frameworks, taxonomies, and catalogs. They all evaluate a different subset of techniques regarding different properties, some of which overlap. For instance, Svahnberg et al. [72] propose to group and evaluate techniques by 2 main properties, according to the size of software entities with variability or variable (components, frameworks, and lines of code), and their latest binding time. Then, the others provide some more concrete and enriched sets of techniques and evaluation properties. For instance, Patzke and Muthig [60] evaluate 8 techniques with up to 6 properties, Patzke and Muthig [61] evaluate 1 technique based on 10 properties, Muthig and Patzke [57] evaluate 8 techniques with up to 8 properties, Gacek and Anastasopoules [29] evaluate 7 techniques with up to 7 properties, Patzke et al. [62] evaluate 7 techniques with up to 14 properties, or Apel et al. [6] evaluate 10 techniques with up to 15 properties.

A more systematic evaluation is provided by a recent catalog [76] that provides an updated set of 21 variability implementation techniques. They are categorized by their 24 different properties that they may support or belong to, such as the logical relation that a technique provides among variants of a *vp*, their granularity, binding time, traceability, language-based, tool-based, or in which programming paradigm it belongs to. As a main property, all 21 techniques

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³⁹removed: to model the domain variability of an SPL or variability-rich system (*i.e.*, as problem-oriented abstractions), the commonalities and variabilities in code assets are usually abstracted in terms of variation points (*vp*-s) with variants, respectively (*i.e.*, as solution-oriented abstractions for the realized variability of code assets) [25, 26, 64].

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are categorized into traditional or emerging techniques:

- Traditional techniques are those that have emerged and evolved separately and quite before the emergence of the SPL paradigm. They encompass methods that are used for single system development but provide the necessary mechanisms to be good candidates for SPL engineering. Although, using these techniques feature does not have a first-class representation in implementation. Such techniques are inheritance, overloading, or design patterns. For instance, the traditional technique of overloading provides an *alternative* logical relation between variants of a *vp*, but it cannot be used to provide an *or* or *optional* logical relation between them [29, 60, 76].
- Emerging techniques have appeared with the SPL engineering advances. Here, the concept of feature is a first-class citizen at the code level. Such techniques are frames [10], feature modules [6], or delta modules [66]. For instance, the emerging technique of feature modules provides good support for realizing the *or* logical relation between feature modules in code assets, but it hardly supports the realization of *alternative* or *optional* logical relation between them [6, 76].

Instead of eliminating this diversity, for example, by proposing to implement the whole variability only at the class level, or by mixing two languages, such as the case of C with #ifdef-s, or proposing to use another code unit, such as feature module, that crosscuts with the prior design of code, in this work we strive to [..42] capture and model the different properties of vp-s with variants which are implemented by traditional techniques. This enables to aim at variability management in variability-rich systems implemented with such [..43] techniques. Apart from this, in Section 4.1.2 we detail four characteristic properties of vp-s with variants that are essential for managing the variability in code assets, such as logical relation, binding time, evolution, and granularity. It is important to emphasize that, in principle, each vp is associated with a single variability implementation technique, whereas the same technique can be used to implement several vp-s within [..44] an SPL domain.

2.3. Variability traceability

The CoEST ⁴⁵ defines a trace (noun 2) as "a specified triple of elements comprising: a source artifact, a target artifact, and a trace link associating the two artifacts", and the traceability as "the potential for the traces to be established and used" [21].

The concept of traceability is already used for different reasons in *single* software engineering, mostly to trace requirements. But, in [..⁴⁶] SPL engineering the relevant entities that are required to be traced and the semantics of trace links are different. Specifically, an explicit capturing of variability information and their modeling of dependencies and relationships separate from other development artifacts is needed [14]. This is well known as variability traceability, which is specific to SPL engineering [4, 5]. It includes the ability to trace the specified variability in problem space to the realized variability in solution space, for example, the ability to trace features in Figure 1, as source artifacts, to their respective vp-s with variants in implementation in Listing 1, as target artifacts.

According to the generic traceability model [33], trace links are basically *established* to meet a specific *usage* purpose and need to be *maintained* during [..⁴⁷] an SPL domain engineering process [63]. Moreover, variability trace links can be established for different reasons and by different stakeholders [22, 5]. For example, they can be used for (semi)automating different processes during the application engineering [63], namely for resolving the variability during product derivation, evolving, checking consistency, addressing, or comprehending variability. In this work, we focus on the ability to establish these variability trace links and use them for checking the consistency of variability. For example, after tracing features in Figure 1 to their respective *vp*-s with variants in Listing 1, trace links can be used to check whether each specified feature in the FM is addressed at [..⁴⁸] code level, or whether the *alternative* relation between three features, Line2D, Segment2D, and Ray2D, is properly considered during their implementation in code.

 $^{^{42}}$ removed: model the implemented variability of variability-rich systems by capturing and modeling different properties of vp-s with variants which are implemented by traditional techniques. Specifically, considering the scope of our study, we consider the ten object-oriented traditional techniques and software design patterns provided in the catalog, which are commonly used to implement variability-rich systems. The other techniques, such as preprocessors or feature modules aim at a methodological development of related software systems as an SPL, therefore they are outside the scope of this study.

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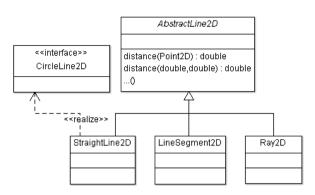


Figure 3: The detailed design of the given excerpt of *JavaGeom* variability-rich system in Listing 1

 $\begin{aligned} & StraightCurve2D \longmapsto \textit{AbstractLine2D} \\ & Line2D \longmapsto \textit{StraightLine2D} \\ & Segment2D \longmapsto \textit{LineSegment2D} \\ & Ray2D \longmapsto \textit{Ray2D} \end{aligned}$

Note: Along this work, the non-italic names stand for features, the italic names stand for code unit, whereas a prefix "vp_" and "v_" is used for variation points and variants. These prefixes are reintroduced in Section 4.1

Figure 4: The mapping of domain features in Figure 1 to the code units in Figure 3. The \longmapsto has the "implemented by" meaning

3. Imperfectly modular variability

Variability management would be simplified in case that each domain feature was implemented in a modular way, by a single language code unit [6], as the mapping of domain features to their implementation would be straightforward. But, this is clearly not the case [..⁴⁹] in realistic SPL settings, with the current language constructs and design methods, where variability is implemented by using a combined set of traditional techniques, such as inheritance, overloading, generic types, or design patterns. These techniques do not enable [..⁵⁰] to completely align the domain features to the places where they are implemented, leading to what we define as *imperfectly modular variability* at the implementation level. On the contrary, a perfect alignment of a domain feature to its implementation would make traces trivial to be established.

To illustrate what imperfectly modular variability means, let us consider the given features in Figure 1 for JavaGeom [...⁵¹]PL and the excerpts of their respective implementation as in Listing 1. Focusing on the implementation techniques in Listing 1, the design of which is shown in Figure 3, the abstract class, AbstractLine2D, is a vp and its three subclasses, StraightLine2D, LineSegment2D, and Ray2D, represent its variants that are created by specializing its implementation. In this way, features in the FM in Figure 1 seem to have a direct and modular mapping in implementation, as is shown in Figure 4, where \longmapsto has the *implemented by* meaning.

Actually, this perfect form of modularity hardly exists. The variation point <code>AbstractLine2D</code> and its variants have a lot of internal variabilities. For instance, the method <code>distance()</code> in <code>AbstractLine2D</code> is another <code>vp</code> with two alternative variants implemented using the technique of overloading (<code>cf. Listing 1</code> in lines 6 - 9 and 11 - 18). The variants of <code>distance()</code> can be considered as implementations in finer grained modules. They represent some technical and nested variability, which could be specified or not in the FM, but still need to be documented, traced, and managed (<code>e.g.</code>, in order to be resolved). Then, the feature <code>Line2D</code> uses several <code>vp-s</code>, such as <code>AbstractLine2D</code>, <code>CircleLine2D</code>, the variant <code>StraightLine2D</code> (<code>cf. Figure 1</code> and <code>Figure 3</code>), plus their finer grained <code>vp-s</code> mentioned just before.

Imperfect modularity comes from the fact that a feature is a domain concept and its refinement in reusable code assets is a set of *vp*-s with variants, even if they are modular, meaning that it may not have a direct and single mapping. Also, while it would be preferable to use a single variability implementation technique to implement [...⁵²] an SPL, none of the existing techniques is yet "a preferred one" [6], which could cover all kinds and properties of variability that may appear in different domains. In our illustrative example, the traditional techniques of inheritance and overloading are used together, which induces a *class* and *method* granularity of variability with *runtime* and *compile time* binding, respectively. In these techniques, the conceptual features do not have a first-class representation in implementation. Still, a degree of modularization of variability can be achieved when used with variability management or traceability in mind. Therefore, we propose the following definition.

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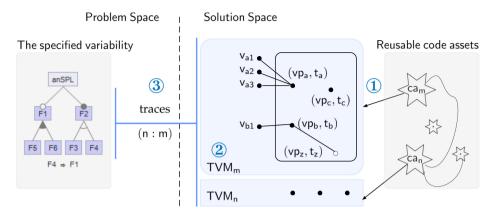


Figure 5: A three step framework for variability management of reusable code assets. TVM_m stands for technical variability model of the reusable code asset ca_m , with vp-s $\{vp_a, vp_b, ...\}$ and their respective variants $\{v_{a1}, v_{a2}, ...\}$ that are realized by different techniques $\{t_a, t_b, ...\}$. Whereas, $\{f_1, f_2, ...\}$ are features in the FM.

Definition 1. An imperfectly modular variability in implementation occurs when some variability is realized in a methodological way with several traditional variability implementation techniques used in combination. The code is not necessarily shaped in terms of domain features, but still, it is designed with the variability traceability in mind.

Our motivation in this work is the ability to manage this imperfectly modular variability[..⁵³]. Towards that, we determine the three following challenges.

Challenge₁: The diverse *vp*-s with variants of [..⁵⁴] reusable code assets, realized by different traditional techniques, need to be captured and modeled in some way, in order to trace them to domain features.

Challenge₂: The way variability is modeled at the implementation level and traced to the domain features should be well-suited for variability management, namely for checking their consistency.

Challenge₃: The feasibility of an approach addressing $Challenge_1$ and $Challenge_2$ should be demonstrated by a tool support.

The $Challenge_1$ is actually motivated by the $Challenge_2$. The importance of capturing and modeling the different types of vp-s and their implementation techniques becomes especially important during the usage of trace links. For instance, when trace links are exploited to resolve variability in [..⁵⁵] an SPL, knowing the binding time of variants in a vp is necessary. Then, the logical relation between variants in a vp is also required to check the consistency between the variability at the specification and implementation levels. In addition, knowing whether a vp is open for adding new variants or unimplemented is needed during variability evolution. All this indicates that considering the variability implementation techniques during the variability modeling of code assets, as challenged in $Challenge_1$, has a strong impact on the multiple usages of trace links, that is, in variability management ($Challenge_2$).

[..⁵⁶] In this article, we propose a tooled framework to capture the imperfectly modular variability at the implementation level, in a forward engineering process, so to be able to exploit the captured trace links, *e.g.*, for consistency checking.

4. A three step framework

To address $Challenge_1$, we propose a three step framework for managing the imperfectly modular variability of [...⁵⁷] reusable code assets. It supports the following tasks, which are also depicted in Figure 5:

⁵³removed: of variability-rich systems.

⁵⁴removed: code assets of variability-rich systems

⁵⁵removed: a variability-rich system,

⁵⁶removed: In this article, we propose a tooled framework to capture, model, and trace the imperfectly modular variability at the implementation level of a variability-rich system, in a forward engineering process, so to be able to exploit the captured trace links, *e.g.*, for consistency checking.

⁵⁷removed: code assets of a given variability-rich system

- 1. Capturing the imperfectly modular variability in terms of vp-s with variants, as abstract concepts.
- 2. Modeling this variability in terms of *vp*-s with variants while keeping the consistency with their implementation in reusable code assets.
- 3. Establishing the trace links between the specified and implemented variabilities.

The first two steps are mostly based on a conducted study of 21 variability implementation techniques regarding their 24 characteristic properties [76][...⁵⁸], which was briefly discussed in Section 2.2.3. We use this large set of summarized characteristic properties to understand what characterizes the variability in implementation[...⁵⁹], and what should be considered as abstractions so to model them properly in terms of variation points with variants. [...⁶⁰] Thus, we first [...⁶¹] analyse what characterizes only traditional techniques [...⁶²] (*cf.* Section 2.2.3), which lead to the form of [...⁶³] imperfectly modular of variability. Then, [...⁶⁴] in addition to the established trace links of variability during the third step, we identified that 4 from the 21 characteristic properties are crucial for a variability management framework, namely for consistency checking, by which we would address *Challenge*₂. Those four properties are the [...⁶⁵] logical relation between variants of a vp, their [...⁶⁶] binding time, evolution, and granularity. The third step, related to trace links, is based on some other existing frameworks on variability traceability [4, 5].

[..⁶⁷] In the following, we give a detailed and formal representation of each of these three steps of the framework, which are [..⁶⁸] associated with illustrative examples.

4.1. Capturing the variability of reusable code assets

As previously discussed, [..⁶⁹] the variabilities in reusable code assets that [..⁷⁰] have a form of imperfect modularity usually [..⁷¹] consist of a common part [..⁷²][..⁷³] and a variable part. It may happen that a whole code asset is also a variable asset, such as for example a source file, a package, or a class[..⁷⁴][..⁷⁵]. The variable part consists of [..⁷⁶] technique with a mechanism for [..⁷⁷] creating the variants, a way for resolving the variants, and the variants themselves (*cf.* Figure 2b). Toward addressing the first step of the framework, we naturally abstract them by using the concepts of variation points (vp-s) with variants. [..⁷⁸] The variable part can have several vp-s and their variants, which can have dependencies with the vp-s and variants in the same or in other assets. Therefore,

Definition 2. Capturing the variability of reusable code assets means to abstract the [...⁷⁹] place where the variability happens, namely [...⁸⁰] a variable asset element in a reusable code asset, or a whole variable code asset, and to represent [...⁸¹] it by a variation point (vp) or variant concept.

Capturing a vp with its variants is specific to the used technique to implement variability. Therefore, we propose the following formalization. Let vp_x be a specific variation point and v_{xy} one of its variants, where $x, y \in \mathbb{N}$. We

```
<sup>58</sup>removed: . This study
<sup>59</sup>removed: and how it can be captured and modeled
<sup>60</sup>removed: Notably
61 removed: analysed
62 removed: , being the focus of our study
<sup>63</sup>removed: imperfect modularity
<sup>64</sup>removed: we noticed
65 removed: logical relation
<sup>66</sup>removed: binding time, evolution, and granularity. On the other hand, the third stepregarding the trace of variabilities
<sup>67</sup>removed: The second or third step begins only after the completion of its preceding step.
<sup>68</sup>removed: illustrated by
<sup>69</sup>removed: a variable part in
70removed: has
<sup>71</sup>removed: consists
<sup>72</sup>removed: with its variants (
<sup>73</sup>removed: Figure 2b).
<sup>74</sup>removed: , is a common part or varies (
<sup>75</sup>removed: is a variant). In effect, the
<sup>76</sup>removed: an implementation technique that denotes
<sup>77</sup>removed: factorizing the commonality among related variants, for
<sup>78</sup>removed: In general, it is expected that some targeted code assets of a system
<sup>79</sup>removed: places
<sup>80</sup>removed: the common code assets, or their common elements, with their varying code assets or their varying elements,
<sup>81</sup>removed: them by variation points (vp-s) with variants concepts, respectively
```

assume that the vp_x is implemented by a single traditional technique t_x . The set \mathbb{T} of possible techniques for vp_x [...82] is then made explicit in our framework:

T = {inheritance, generic type, overriding, overloading, strategy pattern, template pattern, ...}

It includes any traditional technique that promotes a form of reusability for code assets[..83] 84.

A vp with variants is not a by-product of such [...⁸⁵] traditional technique [17, 69]. Therefore, to capture them properly, we need also to tag them in reusable code assets in some way. In this context,

Definition 3. Tagging a variation point or variant means to map the variation point or variant concept to its concrete variable asset, or variable asset element, in reusable code assets.

This mapping is needed to keep the association between the captured $vp[..^{86}]$, or variant, concept and its respective variable code asset.

For example ^a, from Listing 1, the superclass <code>AbstractLine2D</code> is a <code>vp</code> and we abstract/capture it as <code>vp_AbsLine2D</code>. Whereas, tagging has in purpose to ensure their association, namely (<code>AbstractLine2D</code>, <code>vp_AbsLine2D</code>). Similarly, its subclasses <code>StraightLine2D</code>, <code>LineSegment2D</code>, and <code>Ray2D</code> are the <code>vp</code>'s variants, which we abstract as <code>v_Line2D</code>, <code>v_Segment2D</code>, and <code>v_Ray2D</code>, respectively. Their respective [... ltags are, (<code>StraightLine2D</code>, <code>v_Line2D</code>, (<code>LineSegment2D</code>, <code>v_Ray2D</code>, and (<code>Ray2D</code>, <code>v_Ray2D</code>).

4.1.1. Tagging properties

We provide the following nomenclature for describing the tagging properties of vp-s with variants. Let the set VP of variation points in [..87] reusable code assets be:

$$VP = \{vp_a, vp_b, vp_c, ..., vp_n\}, \text{ where } a, b, c, n \in \mathbb{N}.$$
 (1)

The set V of all realized variants for the $vp_x \in VP$ is then:

$$V = \{v_{x1}, v_{x2}, v_{x3}, ..., v_{xn}\}, \text{ where } x, n \in \mathbb{N}$$
 (2)

The set CA of all [...88] reusable code assets, with variability or [...89] variable themselves, [...90] lis:

$$CA = \{ca_m, ca_n, ca_n, \ldots, ca_n\}, \text{ where } m, n, o, z \in \mathbb{N}.$$
 (3)

The [..91] variable elements in a reusable code asset, for example, for $ca_x \in CA$, which [..92] represent some vp-s or their variants, [..93] are:

$$ca_{x} = \{ea_{x1}, ea_{x2}, ea_{x3}, ..., ea_{xn}\}, \text{ where } x, n \in \mathbb{N}.$$

$$(4)$$

According to Definition 3, the abstractions of each vp and variant from sets (1) and (2) are associated with their concrete implementation in reusable code assets, for example, with a varying file, class, or method, which we named

^aIn the following, all illustrative examples will be given into similar blue boxes.

^bremoved: associations are:

⁸²removed: with its variants

⁸³removed: in a given variability-rich system

⁸⁴Although this set includes any of such techniques from object-oriented, functional programming, or design patterns, we exemplify and validate the framework with a limited set of them. In total, we use up to ten techniques, which are common and are used in our considered variability-rich systems (given in Section 6).

⁸⁵ removed: a

⁸⁶removed: or variant

⁸⁷ removed: code assets of a targeted system

⁸⁸removed: code assets of the system

⁸⁹ removed: being

 $^{^{90}}$ removed: which have an association to some vp-s or their variants,

⁹¹ removed: set of

⁹² removed: have an association to

⁹³ removed: is

as in sets (3) and (4). Analysing some of the traditional variability implementation techniques, we [...⁹⁴]determine four possible associations for vp-s and [...⁹⁵] two for variants. They also represent all possible associations between a vp[...⁹⁶], or a variant and their respective [...⁹⁷] variable code assets, except for one association, which is explained in the following. All these associations are described in the following and can be grouped into three kinds, named as *single location*, *implicit*, or *spread*.

Single location. This kind of association represents [...⁹⁸] a vp or variant [...⁹⁹] that abstracts a single location in reusable code assets. [...¹⁰⁰] The three possible cases are:

• (vp_x, ea_{xn}) , where n represents a specific element of the ca_x in (4). This is the case when vp_x represents only one element of an asset.

For instance, (vp_AbsLine2D, AbstractLine2D), represented by a class.

• (vp_x, ca_x) , when the reusable code asset ca_x itself is variable.

If $vp_AbsLine2D$ is optional vp and together with its three variants is contained in a file, then the whole file is variable. Thus, $(vp_AbsLine2D, fileName)$.

• (v_{xy}, ca_x) or (v_{xy}, ea_{xn}) , where y in v_{xy} represents a specific variant of V in (2) and n represents a specific variable asset of CA in (3) or a variable element of ca_x in (4). This is the case when the variant represents only one varying asset or one varying element of it.

(v_Line2D, StraightLine2D), (v_Segment2D, LineSegment2D), or (v_Ray2D, Ray2D), where each of these variants is implemented by a single class in reusable code assets.

Implicit. This kind of association represents the case when a vp [..¹⁰¹] is implicit and cannot easily be tagged in reusable code assets. Therefore,

• (vp_x, \emptyset) , when the vp_x is implicit, hard to be tagged.

For example, this is the case when an *if-else* statement is used [72].

We do not include in the framework [..102] the case when a [..103] variant is unimplemented, (v_{xy}, \emptyset) , because, for example, an abstract class may have infinite subclasses as variants and abstracting them would be impossible and unnecessary.

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94removed: discerned five
95removed: five
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97removed: code assets.
98removed: the case when
99removed: is tagged to
100removed: The four
101removed: or variant
102removed: two cases: (i)
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¹⁰³removed: vp is introduced but has no variants [72], as we consider that it is part of the core assets [80], and (ii) the case when a

Table 1
Logical relations of variation points, and of variants in a variation point

Logical Relation	Description
Mandatory	The [¹⁰⁹] <i>vp</i> or variant is part of each software product
Optional	The 1.110 vp or variant can be part of the software product or not
Multi Coexisting (Or)	One or more than one of the variants in a $[^{111}]vp$ can be part of the software product
Alternative (Xor)	Only one of the alternative variants in a $[^{112}]vp$ can be part of the software product
Mutual exclusion	When, during configuration, the selection of [¹¹³] vp or variant requires the exclusion of another [¹¹⁴] vp or variant, and vice versa
Requires	When the selection of a $[^{115}]vp$ or variant requires the selection of another $[^{116}]vp$ or variant

Spread. This kind of association represents the case when the same vp or variant is [..¹⁰⁴] present in several locations in [..¹⁰⁵] reusable code asset. Therefore,

- $(vp_x, \{ea_{x1}, ea_{x2}, ..., ea_{xn}\})$, when the vp_x is found in more than one place in reusable code assets. The vp_x can be associated also with any code asset from the set CA in (3).
- $(v_{xy}, \{ea_{x1}, ea_{x2}, ..., ea_{xn}\})$, when the variant is found in more than one place in reusable code assets. The v_{xy} can be associated also with any code asset from the set CA in (3).

```
[..^a] For example, this is the case when [..^b] parameterization is used as an implementation technique.

are moved: This

bremoved: a vp or variant is implemented by more than one core asset or their elements. For example, the (vp_distance, { distance(p), distance(x, y)}) is a vp that encircle both overriden methods, therefore it needs to be tagged to both of them.
```

Moreover, these spread cases have a slightly different meaning from the crosscutting case of features, for example, when preprocessor directives are used. Commonly, preprocessors are used to annotate the lines of code that implement the functionality of a domain feature, which may crosscut the current design of code. Whereas, we consider that vp-s and variants align with the design of code.

In some existing approaches vp-s or variants are labeled depending on their resolution nature (e.g., choice, substitution [24]), or depending on their location in assets (e.g., kernel-param-vp, kernel-abstract-vp [31]). For now, we keep a more inclusive label, simply as a vp and variant concept, which can reflect in the future their resolution nature, location, or evolution.

4.1.2. Capturing characteristic properties

[...¹⁰⁶] Depending on the used variability implementation technique, the nature of a code asset element that represents a vp or one of its variants varies. Their variety is presented by the characteristic properties of vp-s with variants [76]. [..¹⁰⁷] As part of our framework, we introduce here a formal definition of the four properties that need to be captured, [..¹⁰⁸] as they are essential for variability management, such as for checking, resolving, and evolving the variability. They are the *logical relations* between variants in a vp, or between vp-s, their *binding time*, *evolution*, and *granularity*.

Logical relation. The logical relations between variants in a [..¹¹⁷] variation point, vp_x , that are commonly faced in practice are similar to the relations between features in an FM (cf. Section 2.1). In addition, vp_x with variants can

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104removed: tagged to
105removed: code assets
106removed: In addition to the tagging properties, depending
107removed: Therefore, as
108removed: given that
117removed: vpx variation point
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Table 2Binding times of *vp*-s with variants (adapted from [16, 20])

Binding time	Values	Description
Static binding (S)	(S) compilation / link (S) build / assembly (S) programming time (S/D) configuration (S/D) deploy and redeploy	The variability is resolved early during the development cycle, that is, the decision for a variant in a variation point is made early/statically.
Dynamic binding (D)	(D) runtime (start-up) (D) pure runtime (operational mode)	The variability is resolved later during the development cycle, that is, the decision for a variant in a variation point is made as late as possible/dynamically.

have dependencies [..¹¹⁸][18], which are also similar to the dependencies between features in an FM, such as *mutual* exclusion, and requires. All these possible relations are shown in Table 1.

A single technique, $t_x \in \mathbb{T}$, can offer at least one of these logical relations between variants. For example, inheritance can be used for implementing the *alternative* variants, overriding for implementing the *or* variants, or aggregation for implementing the *optional* variant(s) of a vp.

Concretely, [.. a] in Listing 1 we capture the *alternative* logical relation between variants v_Line2D, v_Segment2D, and v_Ray2D [.. b] (cf. Section 4.1), which is realized by inheritance.

Therefore, we define the set of possible logical relations (\mathbb{LG}) of vp-s and of variants in a vp as:

 $\mathbb{LG} = \{\text{mandatory, optional, alternative, or, mutual exclusion, requires}\}$

Binding time. Within [..¹¹⁹] a domain, each vp is associated with a binding time, which is the time when the variability is decided or the vp is resolved with its variants during a product derivation. Then, [..¹²⁰] vp-s may require being resolved at different development phases. Generally, depending from the chosen variability implementation technique, a vp can be resolved early during the development cycle (e.g., statically, when the decision for a variant is made at compile time), or later during the development cycle (e.g., dynamically, at runtime) [..¹²¹] [15, 3].

At the level of a whole SPL there are distinguished *binding units*, which are important for identifying which vp-s should be bound together. Specifically, when several vp-s participate in implementing some major functionality in a system then they may require to be resolved together as a unit [20, Ch. 4],[47, 46].

In some domains, a vp may require more than one binding time (*i.e.*, multiple binding times), for example when a feature requires a static and/or a dynamic binding [65]. In this case, it represents a real challenge for existing techniques, as they offer only a single binding time. To accomplish a flexible binding time, different techniques are usually required to be combined. However, these cases seem to be very specific and rare. [..¹²²] Therefore in this work, we consider that a vp with its variants is realized by a single variability implementation technique that offers a single binding time [..¹²³].

The binding time here is the time when the variability should be resolved and should not be confused with the time when it will be introduced, which is differentiated by Syahnberg *et al.* [72]. As for the common kinds of binding times,

aremoved: after the three

^bremoved: in Listing 1 are tagged, we capture their *alternative* logical relation,

¹¹⁸removed: [19]

¹¹⁹removed: the code assets of a system

¹²⁰removed: different

¹²¹removed: [15, 3?]

 $^{^{122}}$ removed: On the other hand, there are distinguished *binding units*, which are important for identifying which vp-s should be bound together. Specifically, when several vp-s participate in implementing some major functionality in a system then they may require to be resolved together as a unit [20, Ch. 4],[47, 46].

¹²³ removed: and is bound independently

A Framework for Managing the Imperfect Modularity of Variability Implementations

a taxonomy is available [20, 16][..¹²⁴], which is also shown in Table 2.

Thus, the set of possible static and dynamic binding times (\mathbb{BT}) of a vp_{v} is:

 $\mathbb{BT} = \{\text{compilation}, \text{assembly}, \text{programming}, \text{configuration}, \text{deploy}, \text{startup}, \text{runtime}, \text{pure runtime}\}$

For example, in Listing 1 we capture that the vp_AbsLine2D is bound during the *runtime* to one of its three alternative variants, for instance, to v_Segment2D.

Evolution. Depending on whether the specified variability in the FM is meant to be evolved with new features, [..125] a specific vp [..126] can evolve with new variants in the future [..127] or not. Thus, depending on the used technique, [..128] and on the constraints they create over the evolution of the variability they implement, vp-s can be open or closed to evolution.

Therefore, the evolution set (IEW) is:

```
IEV = \{open, closed\}
```

For example, we capture a *vp* as closed when it is implemented as an *enum type* in Java, and open when it is implemented simply as an *abstract class*.

In [..a]JavaGeom PL, the abstract class $vp_AbstractLine2D$ in Listing 1 is a vp that [..b]is captured as an open $vp_AbsLine2D$, meaning that its variants can evolve in the future.

```
<sup>a</sup>removed: the JavaGeom system
```

Granularity. A vp or variant in reusable code assets can have a specific granularity depending on the size of variability and the used technique [6, p.59][41]. Specifically, a vp or variant, as abstractions of variability, may represent (i) a coarse-grained element that is going to vary, such as a file, a package, a class, or an interface, (ii) a medium-grained element, such as a method, or a field inside a class, or (iii) a fine-grained element, such as an expression, a statement, or a block of code. A summary of them is provided in Table 3 129 . Thus, when traditional techniques are used, the following granularity set (\mathbb{GR}) contains the elements that actually show the granularity levels at which a variability implementation technique may be used to implement a vp or variant.

GIR = {file, package, interface, class, method, field, statement, generic type, argument, ...} {file, package, interface, class, method,

For example, in Listing 1 we should be able to capture the variant v_Line2D in class level, which is realized by the class StraightLine2D in lines 21-24; or, the variant v_distance in method level, which is realized by distance(...) method in lines 6 - 9 or 11 - 18.

4.2. Modeling the variability of reusable code assets

[..132] Each captured vp with its variants [..133] of reusable code assets are realized by a single technique, [..134] therefore they are somehow related.

bremoved: we capture

¹²⁴ removed: While some approaches have a broader view on binding times, for instance, advocating for different binding times at runtime [16?], we use the seven most common static and dynamic binding times in our model, as

¹²⁵removed: new *vp*-s may emerge or

¹²⁶ removed: may

¹²⁷ removed: .

¹²⁸removed: the new or existing

¹²⁹A [..¹³⁰] more complete table, [..¹³¹] including the granularity of language constructs coming from the emerging techniques, such as feature modules as coarse-grained elements, is available elsewhere [76].

¹³²removed: Given that each

¹³³ removed: is

¹³⁴ removed: they are then

Table 3Granularity of *vp*-s and of their variants

Granularity	Values	Description
Coarse-grained	Component, framework with plug-ins as variants, file, package, class, interface, etc.	The specified variability has an effect on the coarsest grained elements of the implementation structure.
Medium-grained	Method, field inside a class, etc.	The specified variability has an effect on the medium grained elements of the implementation structure.
Fine-grained	Expression, statement, block of code within a method, generic type, argument, etc.	The specified variability has an effect on the finest grained elements of the implementation structure.

For example, in Listing 1 the three variants, v_Line2D, v_Segment2D, and v_Ray2D, are related by inheritance with the vp_AbsLine2D, which technique offers an *alternative* [.. a] relation between variants, a *runtime* binding of variants, and [.. b] open evolution of the vp itself.

^aremoved: relationship ^bremoved: an

Therefore, we model their [..135] relations by associating each vp with its realized variants and the used implementation technique. The associated technique $t_x \in \mathbb{T}$ of the $vp_x \in VP$, which relation we write as (vp_x, t_x) , describes three main properties of the vp_x : the logical relation for its variants (lg_x) , the binding time (bt_x) , and the evolution (ev_y) :

$$t_x = \{lg_x, bt_x, ev_x, ...\}, \text{ where } lg_x \in ILG, bt_x \in IBT, \text{ and } ev_x \in IEW$$

The three dots in t_x indicate that the other properties can be added, such as the granularity level. Based on the above example, we can establish

In this way, we aim to address the second step of our framework (cf. Figure 5) for modeling the implemented variability while keeping the [..¹³⁶] consistency among the vp-s with variants, their implementation techniques, and the elements of code asset that they abstract.

4.2.1. Types of variation points

[..137] When modeling the implemented variability, we observed that the following five types of vp-s can be distinguished. [..138]

```
X = \{vp, vp \text{ unimplemented}, vp \text{ technical}, vp \text{ nested}, vp \text{ optional}\}
```

This is not an exhaustive list of them, other types of vp-s may also be possible, but this set was proved to be sufficient during the framework application (cf. Section 6). They are given in the set X below.

$$X = \{vp, vp_unimplemented, vp_technical, vp_nested, vp_optional\}$$

A formalization for each of them is given in the following.

Ordinary. An ordinary variation point (vp) is an introduced and implemented vp by a specific technique, including all its variants, which are also implemented. Therefore, when the vp_x is an ordinary vp we model [...¹³⁹] the variability

¹³⁵ removed: relation

¹³⁶ removed: association

¹³⁷ removed: For

¹³⁸removed: They are given in the set X below.

¹³⁹ removed: this variability in code assets

in a reusable code asset as a set of its variants and the characteristic properties derived from the vp_x 's implementation technique [...¹⁴⁰] t_x . This leads to the following definition (*cf.* Equations (1) and (2)):

$$vp_{x} = \{V, t_{x}\} = \{\{v_{x1}, v_{x2}, v_{x3}, ..., v_{xn}\}, t_{x}\}, \text{ where } n \in \mathbb{N}$$
 (5)

 $[..^{141}]^{142}[..^{143}]$ We will represent it graphically as:

Such ordinary vp with its variants [...¹⁴⁴] from Listing 1 [...¹⁴⁵] is:

Unimplemented. An unimplemented variation point (vp) is an introduced vp but without predefined variants, meaning that its variants are unknown during the domain engineering phase. When the vp_x is an unimplemented vp, we model it as:

$$vp_{x} = \{\{\emptyset\}, t_{x}\}$$

$$(vp_{x}, t_{x})$$

$$vp_{x} = \{\{\emptyset\}, t_{x}\}, \text{ respectively}$$

respectively,



It means that vp_x is introduced by the technique t_x , whereas its variants will be implemented later.

Technical. A technical variation point (vp) is introduced and implemented only for supporting internally the implementation of another vp which has a direct mapping to a domain feature. The mere difference with the other vp-s is that it does not have a direct mapping to a domain feature, thus the name technical. Moreover, a technical vp is in principle an ordinary vp and needs to be modeled at the implementation level as it needs, for example, to be resolved during product derivation. For the sake of generalization, we consider that a technical vp can also be an unimplemented vp. When a technical variation point vp_v is within a variant v_{x3} of an ordinary vp_x , it is modeled as below.

$$\begin{aligned} \mathsf{vp_x} &= \{\{v_{x1}, v_{x2}, v_{x3}, ..., v_{xn}\}, t_x\} \\ \mathsf{vp_y} &= \{\{v_{y1}, v_{y2}, v_{y3}, ..., v_{yn}\}, t_y\} \\ &\quad \text{where } \mathsf{n} \in \mathbb{N} \text{ and } \mathsf{vp_y} \subset v_{x3} \end{aligned}$$

¹⁴⁰ removed:,

¹⁴¹removed: Based on this, we propose the following graphical representation of it

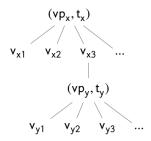
^{142 143}

¹⁴³removed: :

¹⁴⁴ removed:,

¹⁴⁵ removed: .

respectively,



For example, let us consider that $vp_distance$ is the abstraction of the method level variability in Listing 1. It [...a] is a technical vp of $vp_AbsLine2D$, within the [...b] varint v_Line2D , as it is not modeled in the FM. It has two alternative variants, $v_distPoint2D$ (cf. lines 6 - 9) and $v_distDouble$ (cf. lines 11 - 18), which are realized using the technique of overloading. This variability is then modeled as follows.

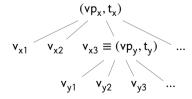
aremoved: can be considered as

^bremoved: variant

Nested. A nested variation point (vp) represents a variable part in a reusable code asset which becomes the common part for some other variants. On the contrary to a technical vp, when the vp_y is a nested variation point of vp_x then a variable part (e.g., the variant v_{x3}) becomes the common part for the variability of vp_y (e.g., for $v_{y1}, v_{y2},$ and v_{y3}). This is modeled as follows.

$$\begin{split} \mathsf{vp_x} &= \{\{v_{x1}, v_{x2}, v_{x3}, ..., v_{xn}\}, t_x\} = \{\{v_{x1}, v_{x2}, \mathsf{vp_y}, ..., v_{xn}\}, t_x\} = \\ &= \{\{v_{x1}, v_{x2}, \{\{v_{y1}, v_{y2}, v_{y3}, ..., v_{yn}\}, t_y\}, ..., v_{xn}\}, t_x\} \\ &\qquad \qquad \text{where } \mathsf{n} \in \mathbb{N} \text{ and } v_{x3} \equiv \mathsf{vp_y} \end{split}$$

respectively,

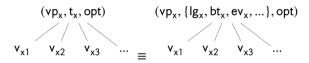


For example, v_Line2D can be a *vp* for its (un)implemented variants.

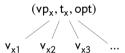
Optional. An optional variation point (vp) [..¹⁴⁶] is defined when the vp [..¹⁴⁷] itself is optional, but not its variants. However, when it is included or excluded in a [..¹⁴⁸] product, so are its variants. [..¹⁴⁹]

[..150] When vp_x is optional then [..151] the tuple (vp_x, t_x) [..152] becomes a triple (vp_x, t_x) , optional variants in a vp and the optionality of the vp itself[..154]. An optional vp can be ordinary, unimplemented, and can have other nested or technical vp-s. This is modeled as:

$$(vp_x, opt) = \{\{v_{x1}, v_{x2}, v_{x3}, ..., v_{xn}\}, t_x\}, \text{ where } n \in \mathbb{N}$$



respectively,



[..155] Also, to make a vp as optional it may not require a technique [..156] (i.e., simply the vp [..157] is not included in [..158] a final product). Therefore, we should document / model it.

4.2.2. Fragmented variability modeling

[..159] Since each vp [..160] is implemented by a specific technique, it can be related or not with the other vp-s. In which case the whole modeled variability at the implementation level will have a forest-like structure[..161]

¹⁴⁶ removed: represents a

 $^{^{147}}$ removed: in code assets that together with its all variants is optional. This means that, whenever an optional vp

¹⁴⁸ removed: software

¹⁴⁹removed: Moreover, variants may have any of the logical relations among themselves.

¹⁵⁰ removed: In our framework, when

¹⁵¹removed: we represent the

¹⁵²removed: tuple as a triple, that is,

¹⁵³removed: the

 $^{^{154}}$ removed: , which is different from the optional logical relation among variants in a vp

¹⁵⁵removed: To

¹⁵⁶ removed: , that is,

¹⁵⁷removed: with its variants

¹⁵⁸removed: the final software product

¹⁵⁹removed: In addition to the logical relation among variants in a vp, different vp-s

¹⁶⁰ removed: and their variants may also be related, for instance, the nested, technical, and/or ordinary vp-s with their variants. This implies that

¹⁶¹removed: , where the *vp*-s and their variants that are related will belong to the same tree, whereas those that are not related will belong to different trees.

```
((vp_AbsLine2D, {alternative, runtime, open}), AbstractLine2D)

(v_Line2D, StraightLine2D) (v_Segment2D, LineSegment2D) (v_Ray2D, Ray2D)

((vp_distance, {alternative, compile, open}), distance())

(v_distPoint2D, distance(p)) (v_distDouble, distance(x,y))
```

Figure 6: Documentation of the implemented variability in a TVM for the JavaGeom PL (cf. Listing 1). The italic names stand for tags of vp-s and variants to the variable asset elements

[..162]. Therefore, instead of modeling the whole implemented variability at once and in one place, we propose to model it in a fragmented way. [..163] What we consider as a fragment aims at being flexible. Specifically, it is a [..164] reusable code asset that can be a package, a file, or a class. In other words, a fragment can be any unit that has its inner variability and it is worth being modeled locally or separately.

For this reason, we [..¹⁶⁵] designed special models, named as *technical variability models* (TVM). [..¹⁶⁶] They contain the abstractions of vp-s with variants, [..¹⁶⁷] their tags to the variable asset elements, and [..¹⁶⁸] describe the realized variability of specific reusable code assets. [..¹⁶⁹] Specifically, a TVM is the power set of all vp-s in [..¹⁷⁰] a reusable code asset. For example, let us suppose that the vp-s in (1) capture the variability of the code asset ca_x in (3) and (4), then the TVM, is:

$$TVM_{x} = Pow(VP_{x}) = \{\{\{v_{a1}, v_{a2}, ..., v_{an}\}, t_{a}\}, \{\{v_{b1}, v_{b2}, ..., v_{bn}\}, t_{b}\}, ..., \{\{v_{z1}, v_{z2}, ..., v_{zn}\}, t_{z}\}\}$$
(6)

where a. b. n. x. $z \in \mathbb{N}$.

As a result, the whole existing variability of code assets [..171] in [..172] CA (in (3)) is modeled in technical variability models (TVMs)[..173], which are created and maintained locally and closer to the code assets. It is important to note that, for simplicity, the tags of variants to the variable asset elements in the TVM_x in [..174](6) are not shown, while they are illustrated in Figure 6. This is because, as specified in Section 4.1.1, any of the variants in TVM_x can have one or more tags to the reusable code assets in [..175](3) and/or [..176](4).

¹⁶²removed: In our framework

 $^{^{163}}$ removed: That is, each fragment may include at least one tree with a vp and its variants. While there is a problem of reconciliation between evolving code and variabilities with separate models, we choose this solution for three reasons: (i) to follow the expected forest-like structure of variability in code assets, (ii) to take into account the large amount of implemented variability in real systems, which can be even larger than the domain variability (capturing it in a single model may hamper scalability), and (iii) to be able to reason locally over a fragment, for example, to check its consistency. Moreover, instead of a traditional global variability model for a variability-rich system, a modular variability model opens the path toward the development of software ecosystems and product lines of product lines [43].

¹⁶⁴removed: code asset with variability, which

¹⁶⁵removed: propose to model the variability of each fragment in a separate model

¹⁶⁶ removed: It will

¹⁶⁷removed: tags to their respective code

¹⁶⁸removed: their characteristic properties within some given

¹⁶⁹ removed: Then

¹⁷⁰removed: some given code assets

¹⁷¹removed: of a system, that is, of CA

¹⁷²removed: Equation (3),

¹⁷³removed: .

¹⁷⁴removed: Equation (6)

¹⁷⁵removed: Equation (3)

¹⁷⁶removed: Equation (4)

In a more illustrative way, in Figure 6 is shown the tvm_line that model the implemented variability in Listing 1. Specifically, it shows the two vp-s, the ordinary vp_AbsLine2D and the technical vp_distance, with their variants and tags to the variable asset elements for the excerpt of JavaGeom [...a]PL. Similarly, the remained variability in JavaGeom PL can be modeled within the same or in other TVMs.

aremoved: system

[..¹⁷⁷] All these TVMs together constitute the *main technical variability model* (MTVM)[..¹⁷⁸]. Therefore, the MTVM [..¹⁷⁹] is defined as the set of all TVMs, given in the following.

$$\mathsf{MTVM} = \{\mathsf{TVM}_{\mathsf{m}}, \mathsf{TVM}_{\mathsf{n}}, \mathsf{TVM}_{\mathsf{o}}, ..., \mathsf{TVM}_{\mathsf{r}}\} \tag{7}$$

[..¹⁸⁰] This MTVM contains the modeled variability of all code assets of [..¹⁸¹] an SPL. Unlike the organization of features in an FM as a tree structure, in MTVM the *vp*-s with variants [..¹⁸²] reside in a forest-like structure. [..¹⁸³] Moreover, the original meaning of a *vp* (given in Section 2.2.2[..¹⁸⁴]) is extended:

Definition 4. A variation point is a location in [...¹⁸⁵] reusable code assets that represents a node for collecting, attaching to the core implementation, and configuring a set of variable units (variable elements of code assets) that are related in functionality. In particular, it is the place at which the variation occurs [36], and represents the used technique to realize the variability.

4.3. Tracing the variability of reusable code assets

Variability traceability is usually organized around two types of trace links: *realization*, and *use* [5, 63]. The first [..¹⁸⁶] link is used in domain engineering for relating the specified variability at the domain level with the artifacts that realize it, such as in code level. [..¹⁸⁷] The trace link *use* relates an artifact during the application engineering process to its respective core assets that are developed during the domain engineering process [63]. It is used to show from which core assets an artifact is elicited during product derivation. Usually, the *realization* trace link is implied by "variability traceability" and in the same meaning we use it also here, for the third step of our framework.

4.3.1. Realization trace links

Let us suppose that $f_x \in FM$ is a [...¹⁸⁸] variable feature at the specification level, then

$$FM = \{f_1, f_2, f_3, ..., f_n\}, \text{ where } n \in \mathbb{N}$$
 (8)

is the set of features in [..¹⁸⁹] the feature model (FM). [..¹⁹⁰] Whereas, a single compound feature f_x with its [..¹⁹¹] variant features and their logical relation (*i.e.*, logic) [..¹⁹²] in an illustrative way [..¹⁹³] is given as in the following. The logic can have any of the values from Table 1 in Section 4.1.2 [..¹⁹⁴] in the feature model.

¹⁷⁷removed: A different number of TVMs can be used in a given system, while all of them together model the variability of its all code assets. Conceptually, all TVMs

¹⁷⁸removed:, which is specific to the given system. As part of our framework

¹⁷⁹ removed: can be

¹⁸⁰removed: Thus, the

¹⁸¹removed: a given system. Moreover, unlike

¹⁸²removed: in the MTVM

¹⁸³removed: Based on these two steps of the framework, we noticed that

¹⁸⁴removed: can be further refined

¹⁸⁵removed: code assets of a system

¹⁸⁶removed: type of links

¹⁸⁷removed: Similar usage of this type of link can be found for example between application and platform features [?]. The use type of links

¹⁸⁸removed: domain feature of a given system

¹⁸⁹removed: its

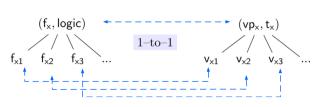
¹⁹⁰removed: In a similar way with the illustration of a *vp* with variants during the framework's first two steps, a

¹⁹¹ removed: varying

¹⁹² removed: can be given

¹⁹³ removed: as follows

¹⁹⁴removed: .



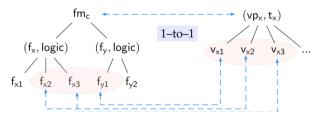
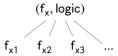


Figure 7: 1-to-1 mapping: ideal mapping

Figure 8: 1-to-1 mapping: a special case of this mapping



For mapping features [..195] to vp-s and variants [..196] in TVMs, we use a single bidirectional type of trace links implementedBy (\longleftrightarrow) or implements (\longleftrightarrow), which presents the *variability realization* trace link at the implementation level. Although the nature of trace links is bidirectional, a single bidirectional type is needed to distinguish which is the source artifact and which is the target artifact of the trace link (according to the trace link definition given in Section 2.3). Namely, implementedBy and implements make possible this differentiation. Otherwise, a link such as $f_x \leftrightarrow vp_x$ would become hard to read as the source and target artifacts are not distinguished.

In general, this mapping of variability from the domain level to the implementation level is [..¹⁹⁷]n-to-m (cf. step ③ in Figure 5). For example, if a domain feature is implemented by three vp-s, then three trace links are needed to describe this mapping. Specifically, from the analysed diversity of vp-s with variants [76], traces of variability may consist of (1) 1-to-1 mapping, (2) 1-to-m mapping, and/or (3) n-to-1 mapping.

1-to-1 Mapping. This is the ideal mapping where each specified feature is realized by a single [...¹⁹⁸] variable unit (i.e., variable element in code assets) that is represented by a vp or variant concept. Such mapping looks as in Figure 7.

In this case, we say that a compound feature f_x and its [..¹⁹⁹] variant features f_{x1} , f_{x2} , and f_{x3} are implemented ideally by a single variation point vp_x with its variants v_{x1} , v_{x2} , and v_{x3} , or conversely. Thus, their mapping is ideal or 1-to-1. Therefore, for feature f_x we write:

$$(f_x \mapsto vp_x), or implementedBy(f_x, vp_x), respectively implements(vp_x, f_x)$$
 (9)

For example, from Figures 1 and 3 in JavaGeom PL, (StraightCurve2D \longmapsto vp_AbsLine2D). That is, the feature StraightCurve2D is implemented by the single variation point vp_AbsLine2D, thus mapped to it.

Similarly, based on the three other features illustrated in Figure 7, a [..200] variant feature f_{xn} [..201] can be implemented by a single variant v_{xn} of v_{px} , that is,

$$(f_{xn} \mapsto v_{xn}), or \text{ implementedBy}(f_{xn}, v_{xn}), \text{ respectively implements}(v_{xn}, f_{xn})$$
 (10)

¹⁹⁵ removed: of an FM

 $^{^{196}}$ removed: of TVMs

¹⁹⁷removed: expected to be

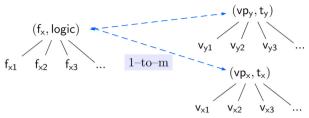
¹⁹⁸removed: code asset (element), which

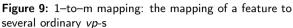
¹⁹⁹removed: varying

²⁰⁰removed: varying

²⁰¹removed: is implemented ideally

A Framework for Managing the Imperfect Modularity of Variability Implementations





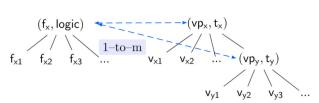


Figure 10: 1–to–m mapping: the mapping of a feature to several nested vp-s

For example, from Figures 1 and 3 in JavaGeom PL, (LineSegment2D \longmapsto v_Segment2D). That is, the [..a] feature LineSegment2D is implemented by the single [..b] variant v_Segment2D, thus mapped to it.

^aremoved: LineSegment2D feature
^bremoved: v_Segment2D variant

[..202] The vp_x and v_{xn} can be from different TVMs. Therefore, we may need to specify, for example, that vp_x , $v_{xn} \in TVM_x$. Moreover, the vp-s with variants in a TVM can be from different code assets.

[..203] We distinguish a special case in the ideal mapping when a vp with variants implement some features that do not correspond to the same hierarchy in the FM. In an illustrative way, this looks as in Figure 8. According to [20], this happens when the variation point has a broader scope, by relating the functionalities that belong in different subsystems. But, this can be resolved by refactoring the organization of features in the FM or [..204] the implementation of vp-s with variants.

1–to–m Mapping. In this mapping, each domain feature is realized by more than one vp or variant. In Figures 9 and 10 are illustrated the cases when a feature is mapped to two vp-s, being ordinary or nested. In this case, a feature f_x is implemented by several vp-s, which can be from the same subdomain of code [...²⁰⁵] asset or not, then

$$(f_x \longmapsto \{vp_x, vp_y, vp_z, ...\}), \ \textit{or} \ \texttt{implementedBy} \ (f_x, \{vp_x, vp_y, vp_z, ...\}) \tag{11}$$

In our observations, this case happens depending on the type of vp-s (cf. Section 4.2.1). Concretely, when vp_x is an ordinary variation point and vp_y its technical variation point, then their mapping looks as in the following:

For example, from Figures 1 and 3, (StraightCurve2D \longmapsto {vp_AbsLine2D, vp_distance}). That is, the [..a] feature StraightCurve2D is implemented by two variation points, vp_AbsLine2D, and vp_distance, thus mapped to them.

aremoved: StraightCurve2D feature

n-to-1 *Mapping*. In this multiple mapping, a vp_x with its variants implement partially several variable features, which can be from the same or different parts of the tree hierarchy in the FM. They are illustrated in Figure 11 and are traced as $[...^{206}]^{207}$

$$(\{f_x, f_y, f_z, ...\} \longmapsto \mathsf{vp}_x), \ \textit{or} \ \mathsf{implements} \ (\mathsf{vp}_x, \{f_x, f_y, f_z, ...\}) \tag{12}$$

208

²⁰²removed: In general, the

²⁰⁶removed:

207 208

²⁰³removed: In the ideal mapping we

²⁰⁴removed: by refactoring

²⁰⁵removed: assets

 $^{^{208}}$ removed: This mapping is relevant to point out the feature interactions or feature crosscutting to vp-s and variants in code assets of a variability-rich system, when combinations of features lead to the resolution of a specific vp-s.

A Framework for Managing the Imperfect Modularity of Variability Implementations

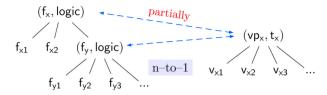


Figure 11: n-to-1 mapping: when several features are partially implemented by a single variation point

The overall mapping between features and vp-s with variants (see step 3 in Figure 5) is expected to be a partial mapping. This is because, some features in FM can be abstract features, by not requiring an implementation [79], or they can be deferred to be implemented later, for example, when a variation point is unimplemented.

4.3.2. Establishing trace links

An automated approach for establishing the trace links may be desirable. It depends on the multiplicity of mapping between features and vp-s with variants. Specifically, when the mapping is ideal then the trace links can be automated by using the same names for features and their respective vp-s or variants. For instance, in the JavaGeom [...²⁰⁹]PL we mostly used the same names for [...²¹⁰] variants and features, except that we added the prefix v_{\perp} for a variant and similarly the prefix vp_{\perp} for a variation point, making features like Ray2D and variants like v_{\perp} Ray2D. Whereas, in the case of multiple mapping (e.g., 1-to-m), the vp-s or variants that implement the same feature can be grouped under the same name with the feature's name that they implement. But, this requires extra work during the variability modeling in TVMs.

In our three step traceability approach (cf. Figure 5), establishing the trace links (for step ③) requires both variability models at the specification and implementation levels (derived in steps ① and ②), that is, the FM and TVMs, respectively. Then the established trace links can be kept separate and orthogonal to the FM and the TVMs. As a result, we consider that the proposed framework addresses the first challenge (cf. $Challenge_1$ in Section 3) for modeling and tracing the imperfectly modular variability of [...211] reusable code assets.

5. Framework implementation

Aiming at a prototyped tool support to help in addressing *Challenge*₃ (*cf.*Section 3), we implemented the proposed framework as a textual *domain specific language* (DSL), a DSL being a small language that is narrowly focused on a particular problem domain [82, 83, 27], We chose Scala to implement the framework as it is well-known for its first-class support for designing expressive DSLs [58, 30, Ch. 6]. In addition, the interoperability between Java and Scala enables one to use our DSL in almost all Java-based variability-rich systems.

In the following is provided the syntax of our *variability modeling DSL* (VM-DSL). It can be used to implement the framework as another internal DSL in other general-purpose languages, or as an external DSL. It consists of a set of language constructs for variability modeling and tracing, thus realizing the defined framework in Section 4. In addition to the variability modeling of code assets, VM-DSL supports the conversion of TVMs and trace links into propositional logic. This enables the implementation of many scenarios of variability management, such as consistency checking between variability models at the specification and implementation levels [13, 52, 75]. We will detail this usage of our framework in Section 7.

The prototype implementation of our DSL is publicly available at https://github.com/ternava/variability-cchecking/. Its detailed description and usage guidance are given at https://ternava.github.io/vm-dsl/, including detailed applications to four SPL case studies.

5.1. Variation points and variants support

In Listing 2 is shown the VM-DSL syntax for defining (i.e.,capturing) vp-s with variants and modeling the

²⁰⁹removed: system

²¹⁰removed: *vp*-s, variants,

²¹¹removed: a targeted system's

```
⟨vps-list⟩
                ::= \langle vp \rangle | \langle vp \rangle \langle EOL \rangle \langle vps-list \rangle
\langle variants-list \rangle ::= \langle variant \rangle | \langle variant \rangle \langle EOL \rangle \langle variants-list \rangle
                 ::= \langle vp\text{-}name \rangle '=' \langle vp\text{-}tvpe \rangle '(' 'asset' '(' \langle tag \rangle ')' ')'
\langle vp \rangle
                ::= \(\forall variant-name \rangle \) '=' 'Variant' '(' 'asset' '(' \(\lambda tag \rangle ')' ')'
⟨variant⟩
                  | \(\forall variant-name \) := string name
⟨vp-name⟩
                 ::= 'VP' | 'oVP' | 'tVP' | 'nVP'
⟨vp-tvpe⟩
\langle tag \rangle
                 ::= \( class-name \) \( \langle method-name \) \( \langle field-name \)
⟨fragment⟩ ::= 'fragment' '(' ⟨tvm-name⟩ ')' '{'
                       ⟨vp-name⟩ 'is' ⟨logical-relation⟩ 'with_variants' '(''{\'ariants}')')'
                       'use' \(\lambda technique\rangle\) 'with_binding' \(\lambda binding-time\rangle\) 'and_evolution' \(\lambda volution\rangle\) \(\lambda EOL\rangle\)
⟨tvm-name⟩ ::= ⟨class-name⟩ | ⟨file-name⟩ | ⟨package-name⟩ | string name
⟨logical-relation⟩ ::= 'MND' | 'OPT' | 'ALT' | 'MUL'
⟨variants⟩ ::= ⟨variant-name⟩ | ⟨variant-name⟩ ', '⟨variants⟩
\(\lambda technique\rangle ::= 'Inheritance' | 'Overloading' | 'Strategy_Pattern' | 'Template_Pattern' | etc
⟨binding-time⟩ ::= 'Compile' | 'StartUp' | 'Runtime' | etc
⟨evolution⟩ ::= 'Open' | 'Close'
```

Listing 2: Syntax for defining variation points and variants, and modeling the implemented variability in terms of them

```
implemented variabilities into a technical variability model (TVM). [..<sup>212</sup> ][..<sup>213</sup> ][..<sup>214</sup> ][..<sup>215</sup> ] [..<sup>216</sup> ] The TVM consists of two parts:
```

- The captured abstractions of vp-s, $\langle vps$ - $list \rangle$, with variants, $\langle variants$ - $list \rangle$, that are associated with the variable elements in code assets, $\langle tag \rangle$.
- Then, the variabilities are modeled within a $\langle fragment \rangle$ module.

```
[..219][..220][..221]
```

5.2. Variability modeling domain-specific language (VM-DSL)

²²¹removed: Section 3), we developed a prototyped tool to support the three steps of our framework. Different mapping studies [???] and surveys on industrial practices [? 34] show that, from the variability modeling tooled approaches, textual variability modeling languages have an increased attention. The main stated reasons are against the scalability issues and the difficulties in comprehending a large amount of variability when competing graphical tooled approaches are used. We thus propose the following textual variability modeling language in our context of code assets.

In the first part, for tagging the variabilities (i.e., $\langle tag \rangle$) the VM-DSL relies on the reflection capabilities in Scala.

[..²²²]
[..²²³]
[..²²⁴]

[..225] This helps to keep the relationship and strong consistency between the defined concepts of vp-s [..226] and variants (i.e.,a $\langle vp$ -name) [..227] and $\langle variant$ -name)) with their concrete implementation in reusable code assets. Then, the VM-DSL supports the five types of vp-s [..228] defined in our framework[..229][..230]. Specifically, an ordinary vp is defined as VP, an optional vp is defined as oVP, a technical vp is defined as an ordinary vp (i.e., as VP) but without variants[..231]

[..232], whereas a nested vp [..233] is defined as nVP.

[..235] Whereas, in the second part[..236], variability is modeled within a fragment construct by specifying the different properties of a vp[..237] and its $\langle variants \rangle$, such as the $\langle logical-relation \rangle$, [..238] $\langle technique \rangle$, [..239] $\langle binding-time \rangle$, and $\langle evolution \rangle$ [..240]

Technically, we defined [...²⁴¹] their values as extensions of their respective sealed classes ²⁴² in Scala, thus well defining their possible values. In addition, the VM-DSL is made expressive by using implicit conversions in Scala. This Scala feature allows the DSL to convert one object to another and [...²⁴³] give the user the impression that methods are dynamically added to an existing class. We used it especially to implement the fragment [...²⁴⁴] module, in order to express in a more natural language the modeling of variability, as in Listing 2. As a result, the language

²²⁸removed: given

²²⁹removed: (

²³²removed: Technically, all captured *vp*-s

²³⁵removed: To meet this demand

²³⁹removed: ⟨logical-relation⟩,

²²²removed: We implemented the proposed language as a textual *domain-specific language* (DSL). A DSL is a small language that is narrowly focused on a particular problem domain [82, 83, 27]. Technically, we chose the Scala language to implement the framework as it is well-known for its first-class support for designing expressive DSLs [58, 30, Ch. 6]. In addition, the interoperability between Java and Scala enables one to use our DSL in almost all Java-based variability-rich systems, such as JavaGeom.

²²³removed: In Listings 2 and 3 is provided the syntax of our *variability modeling domain-specific language* (VM-DSL). Its design and syntax is guided by the following three main demands, which aim to support the three framework steps.

²²⁴removed: One should be able to abstract the implemented variability of code assets of a given system in terms of vp-s with variants while keeping a tag to their respective code assets.

 $^{^{225}}$ removed: To meet this demand, the first part of syntax in Listing 2 reveals the VM-DSL support for abstracting and tagging, that is, capturing, 226 removed: with their variants. Specifically, the $\langle vp \rangle$ indicates that each vp in code assets is captured by giving a name, a type, and a tag to code assets through setting the values for

²²⁷removed: $\langle vp\text{-}type \rangle$, and $\langle tag \rangle$, respectively. While the $\langle vp\text{-}name \rangle$ can have any string values, the $\langle vp\text{-}type \rangle$ can have one of the four predefined types of a vp

²³⁰removed: Section 4.2.1), namely VP, oVP, tVP, or nVP. It should be noted that

 $^{^{231}}$ removed: . Whereas, the $\langle tag \rangle$ is set to the path of the respective code assets, which can be a $\langle class-name \rangle$, a $\langle method-name \rangle$, or a $\langle field-name \rangle$. Similarly, the $\langle variant \rangle$ indicates that each variant in code assets is captured by giving a name, using a language construct, and giving a tag to code assets through setting the value for $\langle variant-name \rangle$, using Variant construct, and setting the value for $\langle tag \rangle$, respectively.

 $^{^{233}}$ removed: and variants are kept into their list structures, $\langle vps\text{-}list \rangle$ and $\langle variants\text{-}list \rangle$, respectively. To set the $\langle tag \rangle$ value our VM-DSL relies on the reflection capabilities in Scala. This helps to keep the relationship and strong consistency between the defined concepts of vp-s and variants, that is, between a $\langle vp\text{-}name \rangle$ or a $\langle variant\text{-}name \rangle$, as well as their concrete implementation in code assets.

²³⁴removed: One should be able to model the variability of a given system by capturing four characteristic properties of each vp with variants, namely their used implementation technique, logical relation, binding time, and evolution. Then, depending on the needs, a dedicated variability model can be set up for different parts of the system.

²³⁶removed: of Listing 2 is given the language syntax for modeling the variability within the fragment construct where the four characteristic ²³⁷removed: and its variants are specified, namely the ⟨technique⟩,

 $^{^{238}}$ removed: $\langle binding\text{-}time \rangle$, and $\langle evolution \rangle$. The fragment construct can contain more than one vp with its variants and requires a $\langle tvm\text{-}name \rangle$. The values for the

²⁴⁰removed: are set to one of their 5, 4, 3, and 2 predefined values, respectively. Their predefined values are enumerated in Listing 2. Still, they can be extended in the future with other values, for example, to cover a larger set of techniques and binding times. All together, the syntax in Listing 2 is used to set up a single variability model, named as a technical variability model (TVM) in our framework. The filename of each variability model and the ⟨tvm-name⟩ of the fragment construct are designed to distinguish different variability models within a given system.

²⁴¹removed: the \(\lambda technique\rangle\), \(\lambda logical-relation\rangle\), \(\lambda binding-time\rangle\), and \(\lambda evolution\rangle\)

²⁴²A sealed class cannot have any new subclasses added except the ones that are defined in the same file [58]

²⁴³removed: gives

²⁴⁴removed: construct

```
⟨trace-links⟩ ::= 'traces' '{' ⟨links⟩ '}'
\langle links \rangle ::= \langle vp-name \rangle 'implements' \langle feature-name \rangle \mid \langle variant-name \rangle 'implements' \langle feature-name \rangle
⟨vp-name⟩ | ⟨variant-name⟩ | ⟨feature-name⟩ := string name
```

Listing 3: Syntax for defining trace links

supports fragmented modeling as described in the previous section, meaning that several [...²⁴⁵][...²⁴⁶]TVMs can be defined[..²⁴⁷]. [..248]

5.2. Tracing support

In addition to language constructs for variability modeling, the VM-DSL provides support for establishing the trace links, as defined in the third step of the framework. [..²⁴⁹] It provides a traces construct with an implements construct for establishing the trace links between the specified features and the modeled variability in TVMs. It supports the creation of 1-to-m trace links, where a single feature [...²⁵⁰] $\langle feature-name \rangle$ [...²⁵¹] can be mapped to several vp-s [...²⁵²] (*vp-name*) or variants (*variant-name*). [..²⁵³] Obviously, 1-to-1 mapping [..²⁵⁴] is also supported. In our current prototyped implementation, the n-to-1 mapping is not supported, but this limitation has a low impact for two reasons: (i) variability management [..255] mainly considers vp-s with variants at the implementation level [..256] as a refinement of features from the domain level [35][...²⁵⁷], (ii) the amount of variability at the implementation level is expected to be larger than the variability at the domain level [63, Ch. 4]. This means that a feature [...²⁵⁸] can have a mapping to several vp-s and/or variants, thus VM-DSL's support for 1-to-m mapping, while the inverse case is quite rare.

Technically, these traces are internally kept in a map structure as Map [feature; vp] and Map [feature; variant], where the feature is the key whereas the vp or variant [...²⁵⁹] are its value. Moreover, a key in the map structure can have several values.

```
[...260][...261]
   [...262 ][...263 ][...264 ]
 <sup>245</sup>removed: variability models (
 <sup>246</sup>removed: technical variability models -TVMs)
 <sup>247</sup>removed: for a given system
 <sup>248</sup>removed: One should be able to trace the modeled variability of a given system's code assets to its domain features.
 <sup>249</sup>removed: The language syntax is given in Listing 3.
 <sup>250</sup>removed: , that is,
 251 removed:,
 <sup>252</sup>removed: or variants, that is,
 <sup>253</sup>removed: This type of trace link also includes the support for the
 <sup>254</sup>removed: .
 <sup>255</sup>removed: approaches mainly consider that
 256 removed: are
 257 removed: and
 <sup>258</sup>removed: , in general, is expected to
 <sup>260</sup>removed: The provided syntax of our variability modeling DSL (VM-DSL) in Listings 2 and 3 can be used to implement the framework as
```

another internal DSL in other general-purpose languages or as an external DSL.

²⁶²removed: The prototype implementation of our DSL is publicly available at

²⁶¹removed: In addition to the variability modeling of code assets, VM-DSL supports the conversion of TVMs and trace links into propositional logic. This enables the implementation of many scenarios of variability management, such as consistency checking between variability models at the specification and implementation levels [13, 52, 75]. We will detail this usage of our framework in Section 7.

²⁶³removed: . Its description and usage guidance are also given at

²⁶⁴removed: , including detailed applications to four variability-rich systems.

```
/* tvm_line.java */
    package tvm // all TVMs are kept within a package
2
3
    import dsl.
4
    import scala.reflect.runtime.universe._
5
    val vp_AbsLine2D = VP(asset(typeOf[AbstractLine2D].typeSymbol))
    val v_Line2D
                        = Variant(asset(typeOf[StraightLine2D].typeSymbol))
                       = Variant(asset(typeOf[LineSegment2D].typeSymbol))
    val v_Segment2D
                        = Variant(asset(typeOf[Ray2D].typeSymbol))
    val v_Ray2D
10
11
    import fragment._
12
    fragment("geom2D.line") {
13
             vp_AbsLine2D is ALT with_variants (v_Line2D, v_Segment2D, v_Ray2D) use
14
             INHERITANCE with_binding RUN_TIME and_evolution OPEN
15
16
```

Listing 4: The captured and modeled variability for Listing 1

5.3. VM-DSL usage example

Following the given syntax in Listings 2 and 3, Listings 4 and 5 show, respectively, [...²⁶⁵] an usage example of the implemented VM-DSL. Specifically, Listing 4 [..²⁶⁶] exhibits some modeling of the variability of code assets of [..²⁶⁷] JavaGeom PL in Listing 1. [..²⁶⁸] Listing 5 shows how to express the trace links [..²⁶⁹] between the variability of code assets in [..²⁷⁰] Listing 1 and features in Figure 1.

[..271] At first, the internal DSL has to be imported [..272] linto the current scope, such as package, file, or class, where variability needs to be separately documented into a TVM (line 4). In line 5, the usage of Scala reflection is made available. Then, an [..273] ordinary vp (line 7) with its three variants (lines 8-10) are defined/captured. They are used to model the variability within the fragment construct (lines 12-16). [..274]

[..275] Whereas, in Listing 5 each vp and variant [..276] is traced to their respective features in Figure 1 that they implement, by using the traces construct. In Listings 4 and 5, the comments in lines 1 and 2 indicate that each TVM can be kept into a separate file, including the trace links, and in a dedicated package within a given [..277] SPL.

By providing the VM-DSL prototype as an implementation of the proposed framework, we consider that the *Challenge*₃ is partially addressed. [..²⁷⁸] In the following, we report [..²⁷⁹] on the experimental evaluation of the proposed VM-DSL.

6. Applications

The provided framework is expected to be applicable to realistic [...²⁸⁰] product line domains where traditional techniques are used. Therefore, to show its applicability, we decided to apply its tool support, the VM-DSL, [...²⁸¹] to

```
265 removed: show a
   <sup>266</sup>removed: shows the modeled
   <sup>267</sup>removed: the JavaGeom system given
   <sup>268</sup>removed: Then,
   <sup>269</sup>removed: are established between the modeled
   <sup>270</sup>removed: Listing 4 and domain
   <sup>271</sup>removed: From Listing 4, the VM-DSL
   272 removed: at first
   273 removed: ordinary
  <sup>274</sup>removed: As mentioned in Demand ??, the relationship of captured vp-s and variants to their respective code assets is kept by using Scala
reflection mechanisms. This makes possible to only capture vp-s and variants that have a concrete realization incode assets, so that inconsistent vp-s
and variants without a right tagging to code assets are disclosed at compile time.
   <sup>275</sup>removed: From Listing 5, each captured
   <sup>276</sup>removed: in Listing 4
   <sup>277</sup>removed: system
   <sup>278</sup>removed: To address the rest of it
   <sup>279</sup>removed: in the following
   <sup>280</sup>removed: variability-rich systems
   <sup>281</sup>removed: in several relevant variability-rich systems, as subject systems,
```

```
/* tracelinks.java*/
import dsl.traces._

\PAR

traces {
    vp_AbsLine2D implements StraightCurve2D
    v_Line2D implements Line2D
    v_Segment2D implements Segment2D
    v_Ray2D implements Ray2D
}
```

Listing 5: Traced variability between Listing 1 and Figure 1 through the captured defined vp-s with variants in Listing 4

model and trace [..²⁸²] the implemented variability in several relevant case studies.

6.1. Research questions

We organized the evaluation by first defining four main research questions [..²⁸³] that further address $Challenge_3$, on the applicability of our [..²⁸⁴] proposed framework (cf. Section 3). [..²⁸⁵] They are given below[..²⁸⁶], including their rationale.

- **RQ₁:** What variability implementation techniques are encountered within [...²⁸⁷] an SPL that uses traditional techniques? This question is intended to answer [..²⁸⁸] whether several techniques are easily used in combination. To this end, we analyse how often each technique is used and [..²⁸⁹] what combination of techniques are often used within four SPL case study.
- **RQ₂:** What are the encountered types of *vp*-s within [..²⁹⁰] an SPL? The purpose of this question is to highlight the types (*ordinary*, *technical*, *optional*, *nested* ²⁹¹) of *vp*-s that are often found within [..²⁹²] an SPL, the diversity of the characteristic properties of these *vp*-s, and the number of TVMs needed to model the variability [..²⁹³] within an SPL. To this end, we [..²⁹⁴] analyse the modeled variability of reusable code assets of four [..²⁹⁵] case studies.
- **RQ₃:** What is the encountered multiplicity of variability trace links within [..²⁹⁶] an SPL? The purpose is to show what is the encountered multiplicity of trace links when traditional variability implementation techniques are used. Therefore, we analyse the encountered multiplicity of trace links [..²⁹⁷] in one of the four [..²⁹⁸] case studies.
- **RQ₄:** What is the overhead for modeling the variability by using the VM-DSL? The purpose is to evaluate the design of the VM-DSL by investigating how many additional *lines of code* (LoC) are needed to model the variability within [..²⁹⁹] an SPL and what is their report with the overall LoC of reusable code assets. To this

```
<sup>282</sup>removed: their implemented variability
<sup>283</sup>removed: , which purpose is to further address
<sup>284</sup>removed: tooled
<sup>285</sup>removed: The four research questions
<sup>286</sup>removed: .
<sup>287</sup>removed: a variability-rich system
<sup>288</sup>removed: the encountered combination of used variability implementation techniques in variability-rich systems
<sup>289</sup>removed: in what combination in four subject systems
<sup>290</sup>removed: a variability-rich system
<sup>291</sup>As defined in Section 4.2.1
<sup>292</sup>removed: a variability-rich system
<sup>293</sup>removed: of a system
<sup>294</sup>removed: modeled and analysed the variability of
<sup>295</sup>removed: subject systems
<sup>296</sup>removed: a variability-rich system
<sup>297</sup>removed: between domain features and vp-s with variants at the code assets
<sup>298</sup>removed: subject systems
<sup>299</sup>removed: a variability-rich system
```

end, we define a measure that helps to evaluate this overhead, then we calculate it for four [..³⁰⁰] case studies for observation.

6.2. Analysed subject systems

We expected that the application of our tooled framework in different [..³⁰¹] SPLs of different domains leads to varying collected data, especially for questions RQ_1 to RQ_3 . As capturing variability implementations is a time consuming process, we applied VM-DSL on several [..³⁰²] case studies, but we could not build very large experiments. [..³⁰³] We first selected three small [..³⁰⁵][..³⁰⁶] case studies that were easily available, have an extensive usage in the SPL engineering community with simple and well-defined domains, so that we can rely on well-established feature models to relate the variability implementations with. We complement these [..³⁰⁷] cases with a fourth [..³⁰⁸] case study being a real feature-rich API[..³⁰⁹][..³¹⁰], which broadens the scope of our validation. The list of [..³¹¹] the fourth case study is as follows:

- **Graph** PL is a family of graph applications, which is proposed as a case study for evaluating the product line methodologies [50].
- **Arcade Game Maker** PL is a family of three arcade games: brickles, pong, and bowling. It is introduced by [..³¹²] the SEI ³¹³ as a case study for experimenting and learning the development of software product lines [2].
- Microwave Oven PL is a family of microwaves, which is proposed to show the usage of UML for developing a set of software-intensive systems from requirements to detailed design [32].
- The fourth [..315] case study, **JavaGeom** PL, introduced in Section 1 is a large real open-source feature-rich API for creating geometric shapes in Java [48].

Before [..³¹⁶] applying the VM-DSL, we had to enhance [..³¹⁷] the four considered systems to get complete ones. [..³¹⁸] For the first three case studies, we had the feature models at hand, and we implemented the SPL in Scala as the language supports the targeted [..³²⁰] implementation techniques, and it was easier to apply the VM-DSL within the same development environment.

Obviously, implementing ourselves the [..321] SPLs that are going to use the VM-DSL is creating a bias on our experiments, especially on questions RQ_1 and RQ_2 . To reduce it, we tried as much as possible to apply good design practices in the implementation, not especially looking for adding several techniques, but rather picking up the right technique to get a proper design of the systems, as developers of a [..322] feature-rich system would have done. In

```
<sup>300</sup>removed: subject systems and make observations
  <sup>301</sup>removed: variability-rich systems
  302 removed: subject systems
  303 removed: For instance, we could not choose big systems using preprocessor directives, such as Linux, as they did not use object-oriented
traditional techniques as we defined them in the scope of our work, and using our DSL on top of preprocessors to model their variability sounded
inappropriate. Therefore, we considered several different object-oriented pre-existing systems, such as Java-based examples in SPL2go 304 and
mostly those used in the SPL engineering community. From them, we
  <sup>305</sup>removed: subject systems (with less than 1K LoC,
  306removed: Table 4)
  <sup>307</sup>removed: subjects
  <sup>308</sup>removed: subject system
  <sup>309</sup>removed: (with over 35K LoC,
  310 removed: Table 4),
  <sup>311</sup>removed: four subject systems, with their LoC, used programming language, and the number of domain features, is summarized in Table 4.
Specifically,
  312removed: SEI
  313https://www.sei.cmu.edu/[..314]
  315 removed: subject system,
  <sup>316</sup>removed: modeling the implemented variability of these subject systems, through
  317 removed: each of them
  318 removed: For the first three subject systems we had their feature models at hand. Therefore, we needed their implementations. Actually,
different implementations of these systems exist. From those that we are aware of, for instance, Graph system is available in SPL2go 319, but it is
built as an SPL using FeatureIDE, while in our scope are variability-rich systems that are not organized as an SPL. Therefore, we implemented these
subject systems using Scala, as it
```

³²⁰removed: variability implementation techniques

³²¹removed: variability-rich systems ³²²removed: variability-rich

Table 4
The implementation details of the four used subject systems

Subject system	Language	LoC	Files	Features	Implemented Features
Graph [50]	Scala	323	3	19	16
Arcade Game Maker [2]	Scala	651	17	27	26
Microwave Oven [32]	Scala	914	19	26	23
JavaGeom [48]	Java	35,456	142	941	941

particular, whenever it was possible, we used those techniques that were specified in the design of each given [...323] case study. For instance, the Arcade Game Maker PL [2] has a detailed design that makes explicit the suggested techniques to implement each vp with variants. Similarly, a detailed design is available also for two other [...324] case studies, Graph PL [50] and Microwave Oven PL [32]. As for JavaGeom PL, we had to analyse its domain [...325] as the Java source code was available, but no feature model was existing. JavaGeom is made of 142 Java source files (excluding tests) and 35 456 lines of code. Considering that the implementation of its 3D geometric shapes is under implementation, we only took its 122 files of 2D geometric shapes that constitute 92% of the code. We then analysed manually the textual description of features in its documentation and their implementation in code. We extracted 941 concrete features (cf. Table 4), which are rather close to the implementation. A summary of the size and implementation of these four systems, based on their lines of code, files, and number of features, is given in Table 4. The domains of these [...326] case studies can further be understood by their feature models, which are made available online 327 .

[..328] Next, we used the VM-DSL to model the variability of [..329] code assets in all these case studies in a forward approach. For the [..330] three case studies we implemented, we model it during their domain implementation or just after their whole implementation. For the JavaGeom [..331] PL case, as we analysed its variability in advance and recorded in a spreadsheet all its νp -s with variants, we used the VM-DSL to model only some of its variability implementations 332 . Finally, we analysed the modeled variability of all [..333] case studies regarding the four given research questions in Section 6.1. We discuss these results in the following subsections.

6.3. Encountered implementation techniques

In order to answer [..334] research question RQ_1 [..335], we analysed the [..336] variability modeled with the VM-DSL in each case study and counted the occurrences of the variability implementation techniques. [..337] Results are shown in Figure 12. Specifically, in [..338] Graph PL, four techniques are used, namely strategy pattern, template pattern, parameters, and include package. Although the *include package* may not be considered as an implementation technique, we counted it here to show one way how an *optional vp*, included in our framework, can be realized. In such a case, an implemented vp within a package will be included in the final product only if the package that contains it is included too. In the Arcade Game Maker and Microwave Oven [..339] PLs, only inheritance is used as a single implementation technique. As mentioned in Section 6.2, we only used inheritance as it is specified in their detailed design as the most appropriate technique to implement vp-s with variants in their domains. Presumably, this can be

```
323 removed: subject system
```

³²⁴ removed: subjects, Graph

³²⁵ removed: and build its feature model,

³²⁶removed: subject systems

³²⁷ https://github.com/ternava/Variability-CChecking/wiki

³²⁸removed: After we prepared the feature model and code assets for each subject system

³²⁹removed: their code assets

³³⁰removed: first three subject systems, we modeled their variability

³³¹ removed: subject

³³² Details are available here: https://github.com/ternava/Variability-CChecking/tree/master/src/SPL_Analysed_Examples

³³³removed: subject systems

³³⁴ removed: the first research question(

³³⁵ removed:)

³³⁶removed: resulted variability models from the VM-DSL usage in each subject system

³³⁷ removed: The results

³³⁸removed: the Graph subject

³³⁹removed: subjects

A Framework for Managing the Imperfect Modularity of Variability Implementations

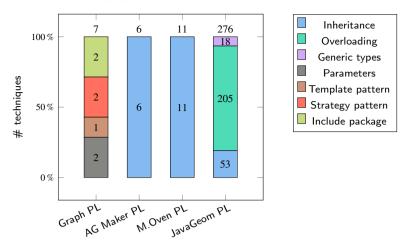


Figure 12: The used techniques, their number, and their combination, in each subject system

explained with their domain of variability and [..³⁴⁰] small size. In JavaGeom [..³⁴¹] PL up to three techniques are used, namely inheritance, overloading, and generic types. In our analysis we omit the technique of overriding as it is used constantly to realize the specialization in the inheritance hierarchy.

The summarized results are given in Figure 12. They show that in maximum [..342] up to four techniques are used in combination. Then, within a [..343] case study a specific technique is used at least once within an SPL, such as inheritance in three case studies, Arcade Game Maker PL, Microwave Oven PL, and JavaGeom PL. In Arcade Game Maker PL it is used 6 times or 100%, in Microwave Oven PL it is used 11 times or 100%, whereas in JavaGeom PL it is used 53 times or 19%. In total, seven techniques are used in the four [..344] case studies. As a partial answer to the first research question RQ_1 , this shows that a combined set of different techniques [..345] are easily used to implement the variability [..346] within an SPL. These data especially enforce our overall motivation that, in many variability-rich systems where traditional techniques are used, it is common that several of them are used in combination. In addition to the reported experiment, a combined usage of traditional techniques is also observed by a recent tooled approach, symfinder [78]. It provides an automated approach for identifying vp-s with variants at class and method level for Java-based variability-rich systems. Its evaluation on eight very large systems, such as AWT library in the JDK 8, Apache Maven 3, and JFreeChart, shows that in each system a combined set of techniques are used. Therefore, the necessity to capture and model the diverse vp-s with variants of such code assets shows a strong need for the support we provide in our framework.

6.4. Modeled variability

In each [..347] case study we also analysed the [..348] variability in each file within a package. Whenever we found that a variability implementation technique is used at class or method level we abstracted it in terms of vp-s with its variants, including their logical relation, binding time, and evolution properties. The resulting data are given in [..349] Table 4.

6.4.1. Encountered characteristic properties of variation points

```
340 removed: a few number of features.
341 removed: subject
342 removed: of
343 removed: subject system
344 removed: subject systems
345 removed: is
346 removed: of a variability-rich system
347 removed: subject system
348 removed: modeled
349 removed: Table 5
```

Table 5
Number of technical variability models (TVMs) and the captured *vp*-s with variants in each subject system

[³⁵⁵]Case study	TVMs	<i>vp</i> -s	logical relation	binding time	evolution	granularity	variant s
Graph PL	3	7	1 mandatory 6 optional	7 runtime	7 open 1 method	6 class	12
Arcade Game Maker PL	3	6	6 mandatory	6 runtime	6 open	6 class	17
Microwave Oven PL	8	11	11 mandatory	11 runtime	11 open	11 class	20
JavaGeom PL	11	276	276 mandatory	53 runtime 223 compile time	53 open 223 close	53 class 205 method 18 type	665

[..350] We show in Table 5 the resulting number of vp-s and variants in each considered [..351] SPL. For instance, the Graph [..352] PL has 7 vp-s with 12 variants, the smallest number, whereas the JavaGeom [..353] PL has 276 vp-s with 665 variants, being the largest number. These data show to some extent the ability of our tooled framework to model a considerable number of vp-s with variants at class and method level for [..354] an SPL.

This table also gives details of the characteristic properties for the modeled vp-s within each [..356] PL. From the resulting 300 vp-s in four [..357] case studies (cf. Table 5), 294 vp-s are modeled as mandatory, [..358] and the remained 6 vp-s [..359] as optional, all being found in the Graph [..360] PL. Each of them is realized into separate files and can be incrementally added or not to the core of its code assets. This is not the case for the vp-s in Microwave Oven PL and Arcade Game Maker PL, as the implementation of vp-s crosscuts files. Regarding the JavaGeom PL, it may have optional vp-s too, but we classified all of them as mandatory vp-s as knowing whether a vp-s is optional required deeper domain knowledge. Then, the 294 mandatory vp-s [..361] [..362] are modeled as follows:

- 71 vp-s have alternative variants, with runtime binding, open evolution, and class granularity;
- 18 vp-s have alternative variants, with compile time binding, closed evolution, and type granularity; and
- 205 vp-s have alternative variants, with compile time binding, closed evolution, and method granularity.

We observed that in JavaGeom PL 41% of the medium-grained variabilities happen at the *constructor level* of a class, the rest is at the *method level*. Then, in all [.. 363] case studies, the vp-s have between 1 and 9 variants. A close observation showed that most of them have only 2 variants, for example, around 54% of vp-s in JavaGeom PL have only 2 variants.

These data show that vp-s [...³⁶⁴] within an SPL, or among them, have indeed diverse characteristic properties. Specifically, even with a small set of [...³⁶⁵] case studies, we observed that each vp has its own logical relation between variants, their binding time, ability to be evolved in the future, and a granularity. Moreover, we also observed that these properties look in direct proportion with the used techniques, that is, the larger the set of used techniques the larger is expected to be the diversity of vp-s. This answers partially the second research question (RQ_2).

```
350 removed: In Table 5 are shown
351 removed: subject system
352 removed: subject
353 removed: subject
354 removed: a variability-rich system
356 removed: subject system
357 removed: subject systems
358 removed: that is, 1 vp in Graph, 6 vp-s in Arcade Game Maker, 11 vp-s in Microwave Oven, and 276 vp-s in JavaGeom. Whereas,
359 removed: are
360 removed: subject.
361 removed: (
362 removed: values in the column logical relation in Table 5)
363 removed: subject systems,
364 removed: of a given variability-rich system
365 removed: subject systems
```

A Framework for Managing the Imperfect Modularity of Variability Implementations

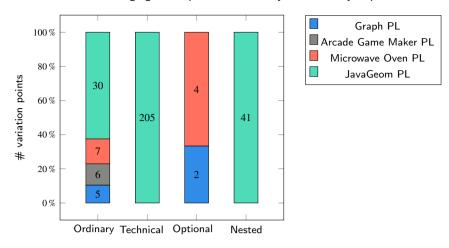


Figure 13: The used techniques, their number, and their combination, in four subject systems

6.4.2. Encountered types of variation points

Considering that each technique is abstracted by a vp, the number of used techniques denotes the total number of vp-s in a [...³⁶⁶] case study. Thus, from Figure 12, the modeled variability in Graph [...³⁶⁷] PL has 7 vp-s, in Arcade Game Maker PL has 6 vp-s, in Microwave Oven PL has 11 vp-s, and in JavaGeom PL has 276 vp-s. Then, each vp is modeled using one of the five types of vp-s, ordinary, unimplemented, technical, nested, or optional. Their used types are summarized in Figure 13, revealing several findings.

First, except in Arcade Game Maker PL, more than one type of *vp*-s exist in each [..³⁶⁸] case study. For instance, in JavaGeom PL three types of *vp*-s are used to model its variabilities, namely *ordinary*, *technical*, and *nested vp*-s. In total, we encountered four types of *vp*-s, *ordinary*, *technical*, *optional*, and *nested* (*cf.* Figure 13). The *unimplemented vp*-s may exist but they were implicit. For example, it was hard to tell which class is introduced in code assets in order to be extended with new variants in the future. These results show that in total there are far more *technical* and *nested vp*-s (82%) than *ordinary* and *optional vp*-s (18%). This is best highlighted from the largest [..³⁶⁹] case study of JavaGeom PL, where there are around 89% more *technical* and *nested vp*-s than *ordinary vp*-s. Furthermore, this can be important for the resolution of variability and indicates that there is a multiple mapping between domain features and *vp*-s in the implementation.

These observations also partially answer the second research question (RQ_2) , showing that different types of [...³⁷⁰] vp are needed to model the variability of code assets in our [..³⁷¹] case studies. The presence of *technical* and *nested vp*-s indicate that vp-s are [..³⁷²] refinement of domain features, a characteristic that matters for establishing the variability trace links.

6.4.3. Number of TVMs

In general, we [..373] used more than one TVM from our framework to model the [..374] variability in each product line. Table 5 shows the resulted number of TVMs with their vp-s and variants. For example, we modeled the variability of the Graph [..375] PL in 3 TVMs, which contain in total 7 vp-s with 12 variants, whereas JavaGeom [..376] PL took 11 TVMs containing 276 vp-s with 665 variants.

It is important to mention that all TVMs in JavaGeom PL are at the package level, as we intended to minimize

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the dependencies of vp-s across different TVMs, whereas in three other [..³⁷⁷] case studies they are at the file level. Although the granularity of TVMs seems to be related [..³⁷⁸] with the size of the [..³⁷⁹] product line, a combination of TVMs at the package, file, or class level would have been also possible. For instance, Figure 14 shows the eight TVMs at [..³⁸⁰] file level for the Microwave Oven [..³⁸¹] PL.

Technically, we stored all modeled variability of a [..³⁸²]PL directly into a map structure representing the power_set, that is, the *main* TVM (MTVM). This indicates that, in addition to all individual TVMs, the VM-DSL makes available all modeled variability of [..³⁸³]an SPL into a single model. Still, from the beginning, someone may choose to model all variability of [..³⁸⁴]some given code assets into a single model.

This answers the remained part of the second research question (RQ_2) . Specifically, these results show that more than one TVM can be used, at package, file, or class level, to model the variability of [..³⁸⁵] an SPL. This fragmented modeling has been shown to have notable advantages, for example, on variability consistency checking [75].

6.5. Encountered multiplicity of trace links

In order to answer the third research question (RQ_3) , we analysed the encountered multiplicity of trace links between domain features and vp-s with variants of [...³⁸⁶] a given PL.

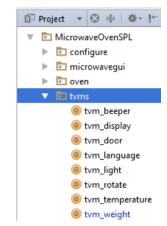


Figure 14: The eight TVMs into separate files for Microwave Oven PL

To trace variability, the variability models at both abstraction levels are required, that is, the FM and the TVMs. In our approach, both of them should be in some way within the same variability management tooling so that the links can be established. In our case, we developed an extension to the VM-DSL so that a domain FM can be entered in our tool. We then conducted this part of the experiment with the [..³⁸⁷] case study of the Microwave Oven PL. We decided for it as, in comparison to JavaGeom PL, we had knowledge for each of its *vp*-s with variants. Then, in comparison to Graph PL and Arcade Game Maker PL, it has the largest number of *vp*-s with variants (*cf.* Table 5).

After we prepared features of Microwave Oven PL into a textual format ³⁸⁸, we used the part of VM-DSL given in Listings 3 and 5 for tracing them to its modeled *vp*-s with variants. The results ³⁸⁹ showed that 17 *vp*-s and variants of Microwave Oven PL (*cf.* Table 5) map to their respective 17 features (*cf.* Table 4), that is, they have 1–to–1 mapping. Whereas, 2 *vp*-s with their 2 variants (*cf.* Table 5), each, are mapped to only 2 features (*cf.* Table 4), that is, they have 1–to–3 mapping. For instance, Microwave Oven PL has a compound feature WeightSensor with two *alternative* features AnalogWeight and DigitalWeight, and also a feature Light. The compound feature WeightSensor is implemented by a single variation point [..³⁹⁰] vp_WeightSensor and each of its alternative features is implemented by a single variant, v_AnalogWeight and v_DigitalWeight. Thus, each feature is traced to one *vp* or variant, therefore their mapping is 1–to–1. As for the feature Light, it is implemented by vp_Light and its two variants v_LightOn and v_LightOf. In this case, a single feature is traced to one *vp* and two variants, therefore their mapping is 1–to–3.

As we found the two forms of mapping in this single [...391] case study, which forms are also supported by the

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³⁸⁸ Its features into textual format are available here: https://github.com/ternava/Variability-CChecking/blob/master/out/production/Variability-CChecking/SPLexamples/fm_prop_microwaveoven_spl.txt

³⁸⁹The resulted mapping is available here: https://github.com/ternava/Variability-CChecking/blob/master/src/SPL_Analysed_Examples/MicrowaveOvenSPL/FeatureVPtraces.scala.

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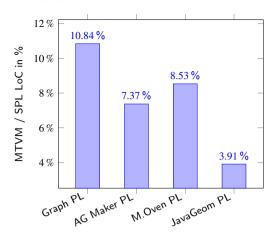


Figure 15: The MTVM's LoC to system's LoC ratio in four subject systems

VM-DSL, we consider that these results can bring an answer to the third research question (RQ_3). They indicate that, even within a small [...³⁹²] product line, when traditional variability implementation techniques are used, the mapping of domain features to vp-s with variants of code assets is imperfect, that is, the 1-to-m multiplicity of trace links may be encountered.

6.6. MTVM to System ratio

In order to evaluate the overhead for modeling the variability by using the VM-DSL, we rely on the fact that the usage of the VM-DSL contributes to the overall lines of code in [..393] an SPL. Therefore, we expect to answer the last research question (RQ_4) [..394] by studying the MTVM to [..395] SPL ratio, defined below.

Definition 5. The MTVM to SPL ratio is the percentage of lines of code in all technical variability models (TVMs) toward the lines of code for realizing the functionality in code assets of [...³⁹⁶] an SPL.

From Listing 2 one can observe that for capturing a vp or variant only a single line of code is needed, that is, only $\langle vp \rangle$ or $\langle variant \rangle$. For modeling the variability another line of code is also needed, using $\langle fragment \rangle$, or two lines for a nested vp. It is important to clarify that, for example, lines 14-15 in Listing 4 represent a single line of code in VM-DSL and are calculated so, but we have separated them into 2 or 3 lines for the sake of readability. Then, each new TVM requires always five additional lines of code (cf. Listing 4, lines 2, 4, 5, 12 [...³⁹⁷] and 13).

For each [..³⁹⁸] SPL, the calculated MTVM to [..³⁹⁹] SPL ratio is given in Figure 15. It shows that on average 7.66% of LoC are used to model their variabilities [..⁴⁰⁰] in TVMs and the rest [..⁴⁰¹] are used for realizing their respective [..⁴⁰²] SPL functionality. It can be observed that, for instance, Graph PL has the fewest lines of code (*cf.* Table 4, 323 LoC) but needs around 11% LoC for modeling [..⁴⁰³] their variability. Then, Microwave Oven [..⁴⁰⁴] PL has more lines of code (914 LoC) than Arcade Game Maker PL (651 LoC)[..⁴⁰⁵], but it uses more LoC for variability modeling (8.53% compared to 7.37%, respectively).

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This indicates that the overhead for using the VM-DSL, calculated by the MTVM to [..406] SPL ratio, does not depend directly on the lines of functional code [..407] within an SPL. It actually is proportional to the number of vp-s, variants, and TVMs, thus answering the last research question (RQ_4). As such, this ratio is a basic indicator for measuring the overhead for modeling the variability by the VM-DSL. In the future, it could be extended by taking into account other elements, for example, the required time to identify for modeling vp-s with variants of some given reusable code assets and to produce their LoC.

6.7. VM-DSL limitations

As a means to further address $Challenge_3$, in this section we showed how our proposed VM-DSL addresses the four research questions (cf. [...⁴⁰⁸][...⁴⁰⁹]Section 6.1) by applying it in four [...⁴¹⁰]case studies. The answers to these questions show the degree at which the VM-DSL, as a proposed prototype implementation of our framework, can be used to capture, model, and trace the imperfectly modular variability of a given [...⁴¹¹]SPL. Still, we observed some limitations.

A limitation of our VM-DSL, but not of the framework itself, is that we could not apply it for tracing the variability at the finest granularity level (e.g., at the expression level), as using [..412] reflection in Scala for tagging the variability is not possible at that level. A solution would be to refactor this finest-grained variability, although using [..413] reflection is not mandatory. Variability could also be captured by using a form of annotations [34]. However, these annotations would still require a manual process, as well as the development of another tool for parsing the annotated code and *automatically* extracting the variability information. In addition, annotations would stay amalgamated with code. On the other hand, [..414] reflection helps us to maintain the strong consistency between the vp-s with variants and their implementation in code assets. They help to keep the variability information separated from the code assets such that the TVMs can be part of the code asset; not amalgamated with its code, or separated into their own files.

In addition, we acknowledge the weak consistency (*i.e.*, corresponding relationship) between the implementation technique itself and the modeled logical relation, binding time, and the evolution properties of a *vp* with its variants (*cf.* Listing 4, lines 14-15). For instance, the technique of inheritance is identified and modeled by the person using the VM-DSL, instead of being automatically detected by a tool, such as *symfinder* [78]. Inconsistencies might happen in this part of the modeled variability. However such inconsistencies would happen also in the case when annotations are used, as they are also given manually.

7. Framework usage

To complement the previous validation and directly tackle the $Challenge_2$ (cf. Section 3), we demonstrate here the usage of our framework for variability management. We aim at showing that our proposed approach for modeling some implemented variability and establishing trace links can be exploited to check the consistency of available models of variability. To do so, we rely on one of our previous work [75], which was only focusing on consistency checking, and put it into perspective to show how its requirements are met by our [...⁴¹⁵] framework.

7.1. Requirements for consistency checking

Within an SPL, it is a well-known desirable property to check whether the specified [..⁴¹⁶] variabilities (*e.g.*, [..⁴¹⁷] a feature model as in Figure 1) and implemented [..⁴¹⁸] variabilities (*e.g.*, [..⁴¹⁹] in reusable code assets as in Listing 1) are consistent [52]. To drive our analysis, we first define a consistency rule that represents what must be satisfied by the variability models at specification and implementation levels.

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Definition 6. [Consistency rule] Within an SPL[..⁴²⁰], where the specified domain variability and the implemented variability convey the same functionality, they also should represent the same set of software products.

Taking into account the mappings between a feature model and the modeled implemented variabilities, we identify and aim to address the following three consistency requirements.

 $Reqmt_1$: One should be able to check the consistency of [..421] variability when the mapping between features and vp-s with variants is 1-to-n, including [..422] and 1-to-1 mapping.

Reqmt₂: In case that some variability is implemented by using an improper technique, an inconsistency may appear, for example, when an *alternative* logical relation between features is implemented by a *vp* that has an *Or* logical relation between its variants. Thus, one should be able to check the consistency of features logical relations in an FM with regard to the used variability implementation techniques.

Reqmt₃: To make consistency checking practical, early detection of inconsistencies is also needed, to reduce costs of correction of erroneous implementations, but also as some variability may be deferred to be implemented later.

We expect that these requirements will be supported by our framework, and its VM-DSL implementation, by relying on trace links, different *vp*-s types, and TVMs to build a consistency checking method.

7.2. Design and implementation of the consistency checking method [1,423]

The proposed consistency checking method takes as input a feature model (FM), representing the domain variability, as well as the trace links (TLs), and a TVM or [..⁴²⁴] subset of them, [..⁴²⁵] taken from our framework. The method enables early detection of variability inconsistencies by relying on three main checking steps, namely *slicing*, *substitution*, and *assertion*. Each of these steps is based on propositional logic⁴²⁶.

- **Slicing.** During this step, a single TVM, or a subset of them from the MTVM [..⁴²⁷] in an SPL, is selected and used to slice the FM. The purpose of the slicing operation [1] is to select for checking only some of the desired variability, thus making possible an early consistency checking that only focuses on the already implemented features.
- **Substitution.** After the slicing step, the consistency rule states that domain features convey the same functionality with their traced *vp*-s with variants at the implementation level. Technically, this implies that abstractions of features should be substitutable with the abstractions of their traced *vp*-s or variants. The purpose is to check whether features are mapped correctly to their corresponding *vp*-s and variants.
- **Assertion.** Finally, if features can be substituted with their traced *vp*-s or variants, then an inconsistency is detected by asserting if the number of feature configurations is equal with the number of *vp*-s with variants configurations.

⁴²⁰removed: or variability-rich system

⁴²¹removed: variabilities

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⁴²³removed: For checking the consistency of the entire variability of a variability-rich system, that is, between all its features and *vp*-s with variants under their 1–to–m mapping, it is required that (*i*) the whole specified variability in an FM is implemented and modeled in terms of *vp*-s with variants in an MTVM, and (*ii*) all their trace links are established. This restricts checking to a complete system, which itself is likely to be represented by large variability models [?], harder to check, but also harder to trace and fix inconsistencies after all of them are shown at once [?]. In addition, it is common that some variability implementation is deferred, for example, during the application engineering phase. Therefore, some partial checking is then highly desirable. Besides, even for illustrative variability-rich systems with a small set of features, the propositional formula of variability models becomes quite large. Moreover, in realistic variability-rich systems, checking for inconsistencies only within a single FM still expose scalability issues [?]. To overcome these problems, we propose a *consistency checking method* for detecting the variability inconsistencies earlier during the development process.

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⁴²⁶In the following, the considered feature models are only propositional ones.

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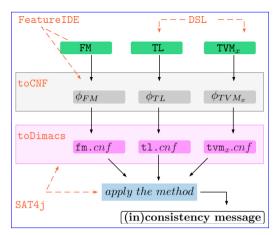


Figure 16: Prototype toolchain for consistency checking. Inputs: feature model (FM), trace links (TL), and technical variability model (TVM_x). Their respective propositional formula: ϕ_{FM} , ϕ_{TL} , and ϕ_{TVM_x} . Their respective DIMACS format: fm.cnf, tl.cnf, and $tvm_x.cnf$.

The consistency checking method has been implemented using the Scala language 428 . As our VM-DSL is also realized in Scala (*cf.* Section 5), we could use the TVMs and trace links directly. The resulting toolchain is depicted in Figure 16 with the three inputs, FM, TLs, and [...⁴²⁹] TVMs (or a subset of them). While the FM is converted to a propositional formula using FeatureIDE [51], we implemented the conversion for the TLs and the TVMs. Next, their propositional logic formulas are converted to a *conjunctive normal form* (CNF), namely ϕ_{FM} , ϕ_{TL} , and ϕ_{TVM_x} , respectively. Then, their CNF formulas are encoded into DIMACS CNF format, namely fm.cnf, tl.cnf, and $tvm_x.cnf$, respectively, so that SAT solving techniques can be used for consistency checking, using the Sat4j library. More details on the method and on the performance of its implementation are provided in a previous work [75].

The implemented method has also been applied to the Microwave Oven [.. 430][.. 431][.. 432]PL case study. This PL has a feature model with 23 implemented features (cf. Table 4), 8 TVMs with 11 vp-s and 20 variants in total (cf. Table 5). It also contains the 1-to-m trace links established during Section 6.5. By selecting a single or a subset of these TVMs at a time, the consistency checking is made possible early in the development process, that is, as soon as the specified variability is implemented. But, we also expected that the consistency checking time for the resulting sliced formulas should be smaller. As for comparison, we recorded the execution time for some TVMs in this [.. 433]PL. The measurements are performed on a PC with 2.50 GHz and 4 GB RAM, on Windows 10. For instance, we selected for checking one TVM with 7 vp-s and variants. The execution time was almost instantaneous. Specifically, the execution time for the sliced FM was 2 ms, for the TVM itself was 7 ms, and the time for checking their consistency was 5 ms. Then, for comparison, we measured the execution time for checking the validity of only the whole FM of Microwave Oven [.. 434]PL, which has 26 features and 720 configurations. It took around 0.13 seconds. As we expected, the execution times for checking the consistency of partial variability compared to the validity of only the FM is smaller.

Figure 17 shows the output of an inconsistency detection. According to our *consistency rule*, the implemented method will report any (in)consistency between FM and TVMs as a difference between their numbers of configurations. In this example, the number of configurations for the FM is 5 and for the [...⁴³⁵]TVM(s) is 31, which indicates that

⁴²⁸Its implementation is publicly available at https://github.com/ternava/variability-cchecking.

⁴²⁹ removed: TVM

⁴³⁰removed: subject system. We report only on this subject system as, in contrast to Graph and Arcade Game Maker subjects, we encountered multiple trace links between its domain features and captured *vp*-s with variants (

⁴³¹removed: Section 6.5), which is one of the requirements for our consistency checking method (

⁴³²removed:). Then, in comparison with the JavaGeom, we used its externally created feature model [32], while for JavaGeom we had to create it directly from the code assets. Thus, to avoid any further bias and data manipulation, we decided to evaluate our method with the Microwave Oven system. This variability-rich system

⁴³³removed: subject system

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tvm language
          import module._
        module("Language scala") {
Language i MUL with variants (English, French, Italian, Spanish, German) use
INHERITANCE with binding RUN_TIME and_evolution OPEN
Run
     TVM moven toCNF
         Nr of Configurations for the TVM(s) is: 31.0
        1 10 13 -14 -15 -16 -17 0
1 10 -13 14 -15 -16 -17 0
1 10 -13 -14 15 -16 -17 0
         1 10 -13 -14 -15 16 -17 0
         1 10 -13 -14 -15 -16 17 0
+[]
         Nr of Configurations for the sliced FM is 5.0
         1 10 13 -14 -15 -16 -17 101 102 -103 -104 -105 -106 0
=
         1 10 -13 14 -15 -16 -17 101 -102 103 -104 -105 -106 0
         1 10 -13 -14 -15 -16 17 101 -102 -103 104 -105 -106 0
100
         1 10 -13 -14 -15 16 -17 101 -102 -103 -104 105 -106 0
×
         1 10 -13 -14 15 -16 -17 101 -102 -103 -104 -105 106 0
         Nr of Configurations for the sliced FM and TVM(s) is 5
        (Models are Inconsistent with each other!,5.0,31.0,5.0
         Process finished with exit code 0
```

Figure 17: Detected inconsistency between the FM and a TVM in the Microwave Oven variability-rich system

the FM and the selected [..436] TVM(s) variabilities are inconsistent. [..437]

7.3. Evaluation

We first evaluate the consistency checking method according to the three given requirements, $Reqmt_1$ to $Reqmt_3$. Implementing the method, we successfully could check the variability when the mapping between features and vp-s with variants is 1–to–1 and 1–to–m, thus fulfilling the first requirement ($Reqmt_1$). The second requirement, $Reqmt_2$, implies checking whether the properties of features, such as their logical relation, correspond to the logical relations of vp-s with variants. Within the method, this checking process happens during the substitution and assertion steps. Then, for all modeled TVMs in the [...⁴³⁸] case study, we were able to select one TVM at a time or a subset of them, using the slicing step. This makes possible to check earlier the implemented variability, meeting requirement $Reqmt_3$. We successfully checked the consistency of [...⁴³⁹] variability under four logical relations of features, respectively vp-s with variants, namely mandatory, optional, or, and alternative ones. In this case we partially fulfilled the second requirement, as we have not checked the consistency regarding other properties, for instance, binding time. The main reason was that these other properties were missing in the FM.

Moreover, the proposed method denotes the usage of our designed TVMs and trace links given in Section 4.2 and Section 4.3, respectively. Its evaluation with the Microwave Oven PL shows that, it handles the case of checking consistency when (i) variability is implemented by different variability implementation techniques and (ii) vp-s with variants can have an 1-to-1 and 1-to-m mapping to the specified features in an FM.

Overall, this consistency checking method and its prototype implementation demonstrate a typical usage of our framework for variability management, thus addressing the second challenge (cf. Challenge₂ in Section 3). Further, the fulfilled requirements (Reqmt₁, Reqmt₁, and Reqmt₃) and the successful application to one of our [...⁴⁴⁰] case studies give us confidence in the good applicability of our framework and its [...⁴⁴¹] implemented VM-DSL, as a concrete solution for modeling and tracing the variability of reusable code assets.

⁴³⁶ removed: TVMs

⁴³⁷removed: The reported numerical format of an inconsistency helps to show only that an inconsistency exists between the FM and the checked TVMs of a system. In this prototyped usage of our framework, the focus of our method is only to detect inconsistencies. Helping developers in addressing this inconsistency is a known non-trivial problem [52], a desirable enhancement would thus be to point to the cause of the detected inconsistency and to recommend a solution for it. For instance, it could be done by explicitly showing which features are inconsistent to which *vp*-s with variants, or whether the inconsistency has happened because the trace links may not be well-established. Nothing in the framework hampers the realization of this extension.

⁴³⁸removed: subject system

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8. Threats to validity

In addition to the VM-DSL limitations given in Section 6.7, we discuss here several validity threats in our validation.

8.1. Internal Validity

8.2. Construct Validity

A first threat to validity is related to [..442] the selected case studies. We used three academic [..443] case studies that we implemented ourselves. As we discussed in Section 6.2 [..444] this threatens the validity of the [..445] research questions that aim at characterizing the encountered implementation techniques. [..446] We tried to implement [..447] these three systems only by following good practices in object-oriented design[..448], but there could be other forms of implementation that would handle the considered variability while being not detected by our framework. Still, the proposed framework supports most of the traditional techniques already surveyed for variability implementations [29, 28, 60, 72, 76], which are themselves very common, such as inheritance and overloading.

8.2. Internal Validity

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[..^{449}] [..^{450}] A second threat is related to the [..^{451}]
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[..452] Jdifficulty we faced during the application of the VM-DSL to [..453] the case studies, especially on recognizing whether all used techniques are variability related. For instance, it was sometimes hard to distinguish whether an inheritance hierarchy implements some [..454] product line variability (i.e., which [..455] Jdifferentiates products in an SPL[..456]) or some internal functionality that should not be modeled or traced. [..457] In our experiments, we decided to classify all used techniques as variability related techniques. However, in [..458] realistic SPL domains, this may not be always true [..459] [55]. For instance, there are approaches that consider that all C preprocessor directives in [..460] reusable code assets are used for implementing some PL variability, whereas some others disagree [84]. Specifically, Zhang et al. [84] state that "from our experience most #ifdef blocks (e.g., 87.6% in the Danfoss SPL) are actually not variability related, but for other purposes such as include guards or macro substitution". Although this is not a concern that is particular to our framework or its realization, it shows a limitation to the generalization of the approach. Recent works [..461] [78, 55] on automatic detection of some object-oriented techniques rely on the density of variability implementations within a set of classes to determine areas of interest in terms of variability, but this technique is not precise enough and still needs some human expertise to map features to variability implementations [..462]

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448 removed: and, whenever was possible, by using the specified techniques in their original design, namely in [50, 2, 32]. But,

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⁴⁴⁹removed: There are several approaches that reveal that domain features crosscut in code assets of variability-rich systems, for instance, in the case of the Arcade Game Maker subject system [?]. In our framework, we consider that a feature may have a trace to a *vp* or variant, which have a direct tag to code assets. It then may seem that features crosscutting in code assets are omitted, forming a threat on *ignoring* some facts. Actually, we acknowledged it at the very beginning and proposed to trace features to *vp*-s and variants, which represent the actual implemented variability of code assets of a system, instead of tracing features directly to code assets. Then, we consider that features to *vp*-s and variants can have n-to-1 multiplicity of trace links, that is, they may crosscut but to the *vp*-s and variants.

⁴⁵⁰removed: For evaluating the consistency checking method, we used a subject system that has only ordinary and optional *vp*-s. It can be viewed as a threat

⁴⁵¹removed: *incompleteness* of the considered types of *vp*-s for evaluating the method, such as nested and technical *vp*-s. Considering them and improving the method accordingly is inline with our future work.

[55].

8.2. Conclusion Validity

..463

8.2. External Validity

While we believe the different experiments show the feasibility of the propositions for modeling and tracing the considered techniques, the number and diversity of [...464] studied cases do not enable us to conclude on a large applicability of the proposed framework and implementation.

On the covered techniques in the [..⁴⁶⁵] case studies, their number [..⁴⁶⁶] with the three academic [..⁴⁶⁷] case studies were four, but [..⁴⁶⁸] with the fourth real [..⁴⁶⁹] case study, JavaGeom, - a realistic feature-rich system, the total number of used techniques became seven. It cannot be determined whether more real [..⁴⁷⁰] case studies would enable the validation over a larger set of techniques. Nevertheless, as we discussed for the [..⁴⁷¹] internal threat, the covered techniques by the framework are the most common ones. We also believe that such validation can successfully be extended to other variability-rich systems that use other techniques, for example, design patterns, as some recent works show that they are present in many object-oriented variability-rich systems [78].

As for the framework usage, we only analysed one scenario with consistency checking, but we believe it is a very relevant one to also [..472] the feasibility of the usage. Obviously more extensions would bring a better validation on this aspect.

9. Related work

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9.1. Variability management approaches

Variability traceability. Several approaches on variability management provide detailed relationship or mapping models between the variability of core assets at different abstraction levels. Most of these models came in form of a general mapping model [13, 12], metadata model [59], basic traceability model [54], a representation independent variability meta-model [68, 67], a conceptual model for traceability [14], or a framework for variability modeling and management in all abstraction levels [70, 20, Ch.9]. A common aspect of these approaches is a formal specification of the relations or associations among the specified variability (*i.e.*, features in an FM), the *vp*-s with variants, as variability abstractions in different core assets, and core assets themselves. However, in all these approaches the first steps on capturing and modeling the variability are not considered. In this work, we provide a detailed framework, which has in purpose to handle the management of variability at the implementation level. Specifically, we distinguish the variability model at specification and implementation levels in a similar way to Becker [13] and Metzger *et al.* [52], and we keep the variability information separated from the code assets as by Berg *et al.* [14] and Pohl *et al.* [63]. However, we go further by detailing the earlier steps of traceability, such as the capture and modeling of variability in terms of *vp*-s and variants, which is the main difference with, or a kind of complement of, other existing approaches.

⁴⁶³removed: A threat to conclusion validity is the *lack of expert evaluation* of our tooled framework by software architects in their variability-rich systems. Such an evaluation would help us to have a stronger interpretation of results and to further identify the possible improvements of our tooled framework, for instance, by reporting the time spent and difficulties to model the variability of a whole real variability-rich system by our VM-DSL. Considering its importance, evaluating our tooled framework in real systems and by their experts is also part of our future work.

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⁴⁷³removed: In the context of product lines and product families, several approaches related to our framework exist on variability management, including variability modeling and variability traceability of code assets. On the other side, there are alternative approaches to our framework. Further, several textual variability modeling languages related to our VM-DSL have been proposed. In the following we discuss the related and alternative approaches to our framework, as well as related works to our DSL implementation.

Besides these approaches, there are other proposals that specifically address the traceability of variability. Mostly, they address the general issues of variability traceability management in SPL engineering, such as establishing, using, and/or managing the trace links, without much focus on how the variability in different abstraction levels, or types of core assets, is realized or identified, and modeled. For instance, AMPLE [5] is a general framework based on a reference meta-model for tracing the variability of different types of artifacts by several types of trace links. It discusses the main traceability dimensions in SPL engineering, specifically when model-driven engineering and aspect-orientation are applied. Quite similarly, XTraQue [37] and the approach by Bayer and Widen [11] target tracing the variability in different types of artifacts. In XTraQue the main idea is to translate the product line documents, as the UML documents, into XML documents and to trace their variability by using several types of trace links.

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Variability modeling. [..⁴⁷⁵] As for the types of *vp*-s and variants, some other works exist[..⁴⁷⁶], for instance, by Jacobson *et al.* [36], Schmid *et al.* [68], Sinnema *et al.* [70, 69] and Capilla *et al.* [20]. [..⁴⁷⁷] Although they are not variability traceability approaches, we used their notion of *vp*-s and *variants* and several of their findings [..⁴⁷⁸] of the diversity of *vp*-s. Yet, we consider that the [..⁴⁷⁹] range of *vp*-s [..⁴⁸⁰] comes from the used variability implementation techniques themselves. In addition, for modeling the variability in a fragmented way we get inspired by the variability-aware module system [43] where each module defines its own local variability model. However, in this approach, a [..⁴⁸¹].c file is considered as a module with inner variability, which has only *compile time* binding. In [..⁴⁸²] contrary, what is considered as a fragment in our framework aims at being flexible, such as a package, file, or class. Then, possibly, fragments are kept separated from the code assets. Besides, we consider that variability is implemented by traditional techniques, not including the preprocessor directives.

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Variability management tools. Furthermore, there are [..⁴⁸⁴] commercial tools for variability management, such as [..⁴⁸⁵] pure::variants [73][..⁴⁸⁶] . pure::variants provides a set of integrated tools to support each phase of the SPL development process, including the implementation phase. Their variability management approach is not much related to the used programming paradigm. Thus, it can be used to manage the variability implemented by preprocessors in C/C++, or implemented by other functional or object-oriented techniques. In the last case, its captured way of variability relies on annotations, a form of macro-like calls. In contrast, instead of annotating the domain features [..⁴⁸⁸] in the lines of code, we use a textual DSL to model vp-s with variants[..⁴⁸⁹], that is, the structure of code, through the capture of the used implementation technique inside a package, file, or class. [..⁴⁹⁰]

When preprocessors in C are used [49], the used directives are commonly recognized as *vp*-s and variants which are more explicit compared to the case when traditional techniques are used. If we consider preprocessor directives, they also can offer all kinds of logical relations between variants in a *vp*. Although compared to the case

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when several traditional techniques are used, they offer a single binding time for vp-s. Still, as preprocessors are a form of annotations with variability information between variants, it could be interesting to apply our framework for modeling and tracing the implemented variability when only this technique is used.

Finally, while the main purpose in pure::variants is variability resolution, we provide a slightly more general approach by not becoming specific on how the vp-s and trace links are going to be used. For instance, one could decide to use them for consistency checking (cf. Section 7), or variability resolution.

9.2. Variability traceability approaches

Besides these variability management approaches, there are other proposals that specifically address the traceability of variability. Mostly, they address the general issues of variability traceability management in SPL engineering, such as establishing, using, and/or managing the trace links, without much focus on how the variability in different abstraction levels, or types of core assets, is realized or identified, and modeled. For instance, AMPLE [5] is a general framework based on a reference meta-model for tracing the variability of different types of artifacts by several types of trace links. It discusses the main traceability dimensions in SPL engineering, specifically when model-driven engineering and aspect-orientation are applied. Quite similarly, XTraQue [37] and the approach by Bayer and Widen [11] target tracing the variability in different types of artifacts. In XTraQue the main idea is to translate the product line documents, as the UML documents, into XML documents and to trace their variability by using several types of trace links.

Although they consider that the variability in core assets is already made explicit in some way and need just to be traced, we consider that these works are complementary with our contribution. Specifically, while they deal more with the abstract level of variability, we provide details on how to capture and model the variability of code assets.

9.3. Orthogonal approaches to variability traceability

When traditional techniques are used for implementing the variability in an SPL[..⁴⁹¹], the code is not shaped in terms of features [36, 17, 72, 20]. In these techniques, the main concerns being separated are objects and/or functions, depending on the used programming paradigm. Thus, a feature can be implemented as a file, class, method, plugin, package, a combination of them or by any part of them. Therefore, as is shown in our framework, the trace relation is n-to-m between the specified features to the vp-s with variants at the implementation level. But, there are approaches [..⁴⁹²] which shape the reusable code assets in terms of features [6, 53, 66, 8], or the specified features hold a direct representation in code [34, 44], that is, their mapping is 1-to-1. We consider that these are alternative approaches to the variability traceability itself, and to our framework, as the traceability of features is straightforward in this case.

9.4. Textual variability modeling approaches

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10. Conclusion and future work

Summary. Variability management is a complicated activity during the engineering phases of software product lines or variability-rich systems. In many settings, variability is specified in terms of features, and they are implemented by using a combined set of traditional techniques, such as inheritance, overloading, or design patterns. In such [...⁴⁹⁴] case, they do not have an explicit mapping at the implementation level, which leads to a form of imperfectly modular variability and strongly hinders the variability management at the code level.

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⁴⁹³removed: In addition to the graphical-based variability modeling approaches, there are several textual variability modeling languages. Specifically, a recent study provides a summary and a comparison, regarding their provided support, of up to fourteen textual variability modeling languages [?]. In almost all of them, the main language concept is feature, whereas the main concepts in our VM-DSL are *vp*-s and variants. Besides, from their mentioned aspects that are relevant when designing new languages, we provide support for *binding time*, but not only - through capturing characteristic properties of *vp*-s with variants, *analyzability* - through supporting the conversion of TVMs and trace links to propositional logic, *expressiveness* - through using the futures of Scala for realizing expressive domain-specific languages, and *extensibility* - through providing the language syntax. Moreover, we provide a textual variability modeling language that aims to capture and model the implemented variability of variability-rich systems, which may not be necessarily organized as a software product line.

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In this work, we provided a three step framework for managing this imperfectly modular variability of [...⁴⁹⁵] reusable code assets. It supports (i) the capture of the implemented variability in terms of variation points (vp-s) with variants, (ii) the fragmented modeling of these vp-s and variants [...⁴⁹⁶] with their tags to specific code assets, [...⁴⁹⁷] in technical variability models (TVMs), and (iii) establishing the n-to-m trace links between [...⁴⁹⁸] features and their respective vp-s or variants. In particular, during the first step different types of vp-s with variants are captured, including their diverse characteristic properties, namely their logical relation, binding time, evolution property, and granularity. These types and properties are used to model the variability of code assets [...⁴⁹⁹] within TVMs, which form a forest-like structure of variability at the implementation level. This enables a form of fragmented variability modeling which, together with the established trace links, can be used for different purposes, such as for resolving the variability, evolving, checking its consistency, or simply [...⁵⁰⁰] understanding the addressed variability in reusable code assets.

We implemented the framework as an internal [..⁵⁰¹] [..⁵⁰²] domain specific language (DSL) in Scala, and applied it to four different [..⁵⁰³] SPL case studies. The results show that in realistic systems, a combined set of traditional techniques can indeed be used to implement the variability. For instance, in the four used [..⁵⁰⁴] case studies, we captured and modeled the implemented variability with up to seven common variability implementation techniques, namely inheritance, overloading, generic types, parameters, template pattern, strategy pattern, and include package. Different types of *vp*-s with their diverse characteristic properties and variants were also encountered. Furthermore, a different number of TVMs were successfully used to model the variability. In addition, the 1–to–m variability trace links were evident even in an academic [..⁵⁰⁵] case study. While being on a small set of [..⁵⁰⁶] case studies, these results show the feasibility of our framework and its potential application in realistic variability-rich systems with a small [..⁵⁰⁷] DSL.

Additionally, we have proposed a method for checking the consistency between the specified and implemented variabilities. The method was also implemented and applied to [..⁵⁰⁸] a case study. Its application shows that the captured variability of code assets and their established traced links to the domain features can be successfully exploited to check their consistency. This shows the potential usage of our contribution as a framework for variability management, which was our broader objective.

Future Work. In future work, we first plan to complement the framework to support the modeling and reasoning over dependencies between vp-s with variants, which dependencies are different from the cross-tree constraints between features in feature models at the specification level. In addition, as we model variability in several TVMs in the framework, we will also have to support vp-s with variants in one TVM having dependencies to vp-s with variants in other TVMs. In the literature, there are already some frameworks that handle dependencies of vp-s [19, 68], which can serve as a starting point for this extension. Further, we will evaluate the scalability of our [...⁵⁰⁹] framework on a larger set of [...⁵¹⁰] case studies, and with an application by other users. In particular, by applying it to a broader range of [...⁵¹¹] case studies, we expect to validate our framework and the DSL on a larger set of variability implementation techniques, going beyond the seven used techniques in this contribution.

Ongoing work that tackles the automatic capturing and modeling of [...⁵¹²] variation points in similar kinds of

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code structure [...⁵¹³][78, 55] could also be considered to be coupled to the proposed contribution, so to reduce the burden of [...⁵¹⁴] variability capturing and modeling [..⁵¹⁵] only on relevant or intricate variations points. In addition, we are working on the interoperability of our tool with other variability management tools, such as pure::variants, by using the variability exchange language (VEL) [81], as a coming standard for exchanging variability information between variability management tools. By this, we will aim at demonstrating the usage of our framework for resolving variability during product derivation.

With these extensions and more evaluations, we hope that new insights can emerge, for example, to highlight patterns of variability in the code structure and thus to better help practitioners in managing variability and code reuse.

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