



# **TsuKing: Coordinating DNS Resolvers and Queries into Potent DDoS Amplifiers**

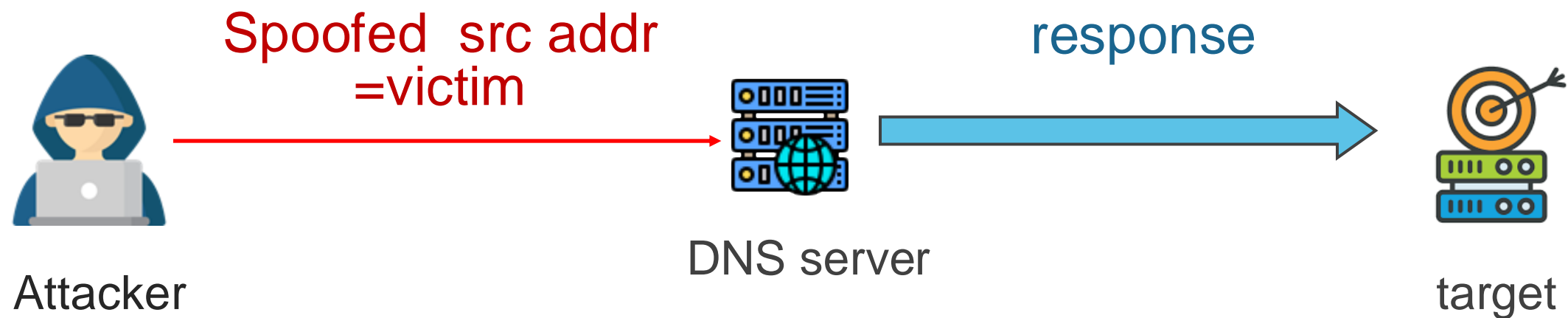
Speaker: Haixin Duan

Slides Contributors: Wei Xu & Xiang Li & Chaoyi Lu

Tsinghua University, Dec. 2023

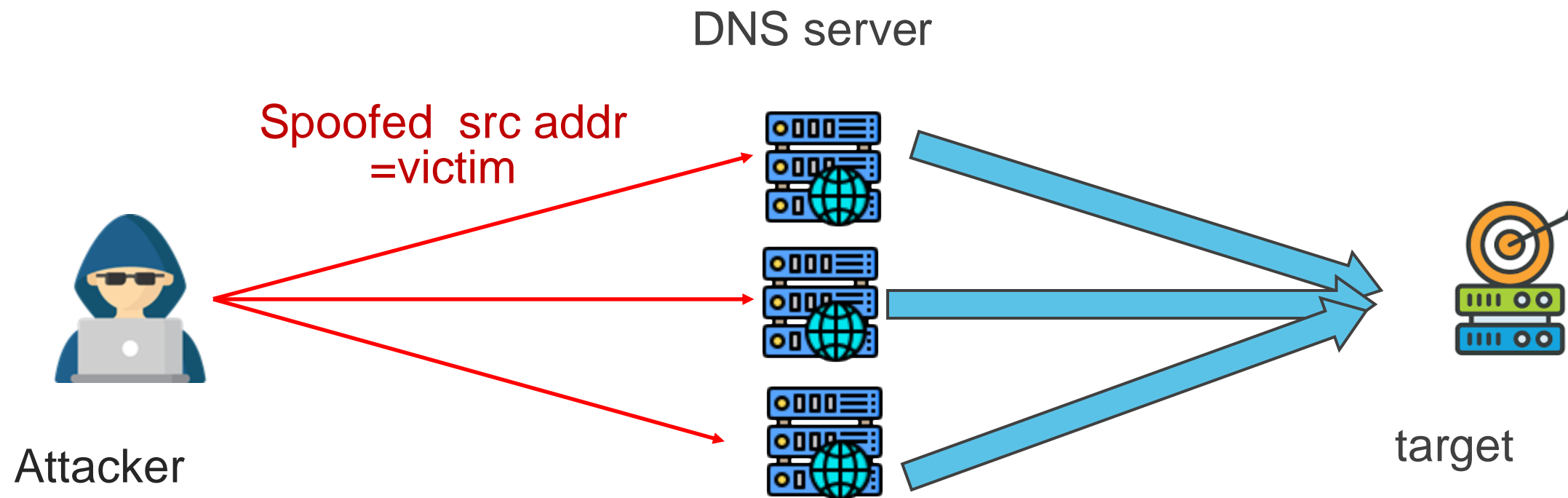


# Reflection and Amplification DDOS by DNS



- Maximize the amplification potential of a **single DNS server**.
- IP Address Spoofing

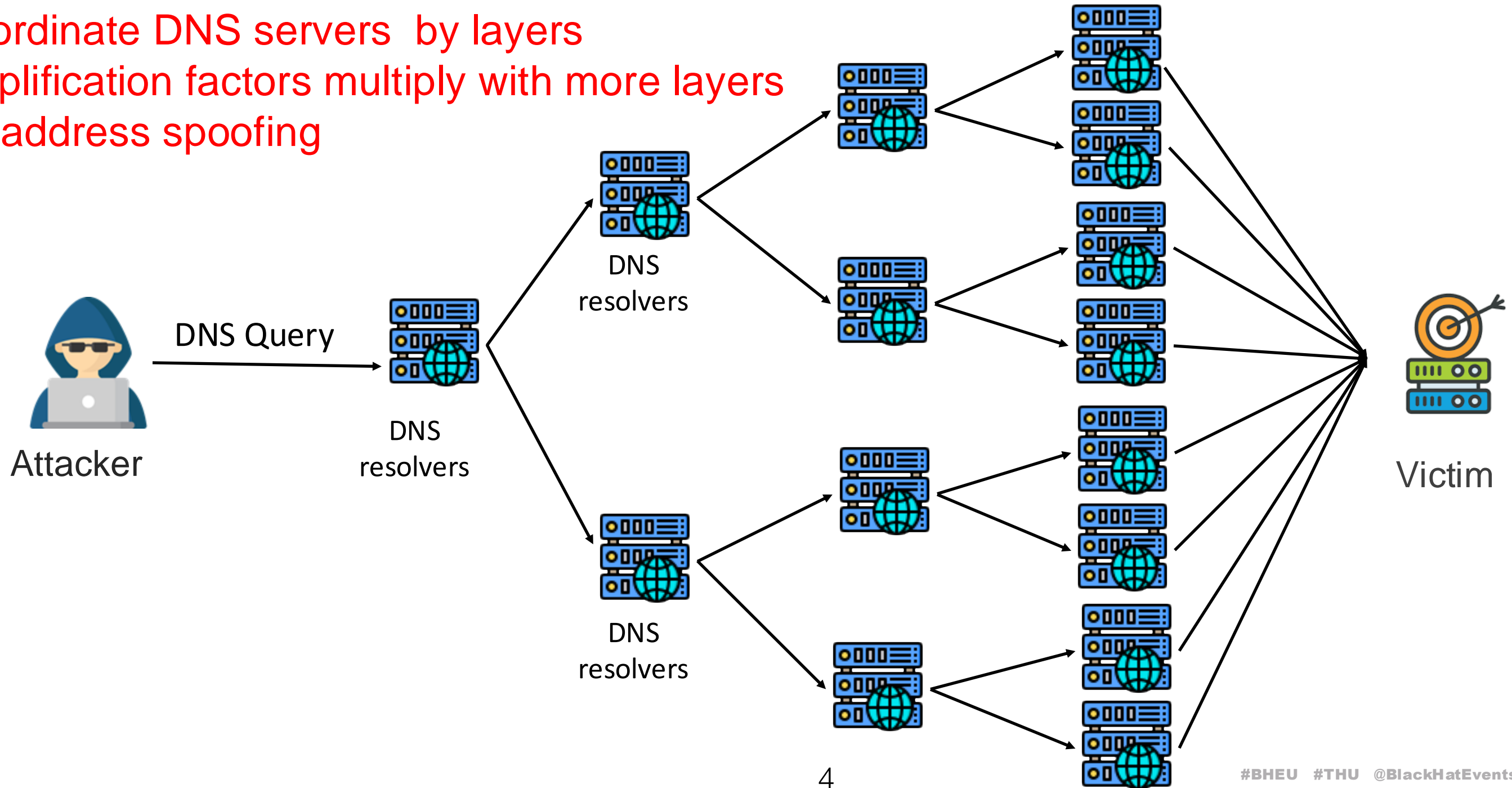
# Current DOS attack by DNS



Send more queries to more servers, **in parallel**  
Amplification factor unchanged

# Tsuking Attack is different

- Coordinate DNS servers by layers
- Amplification factors multiply with more layers
- No address spoofing





# TsuKing: Tsunami + King

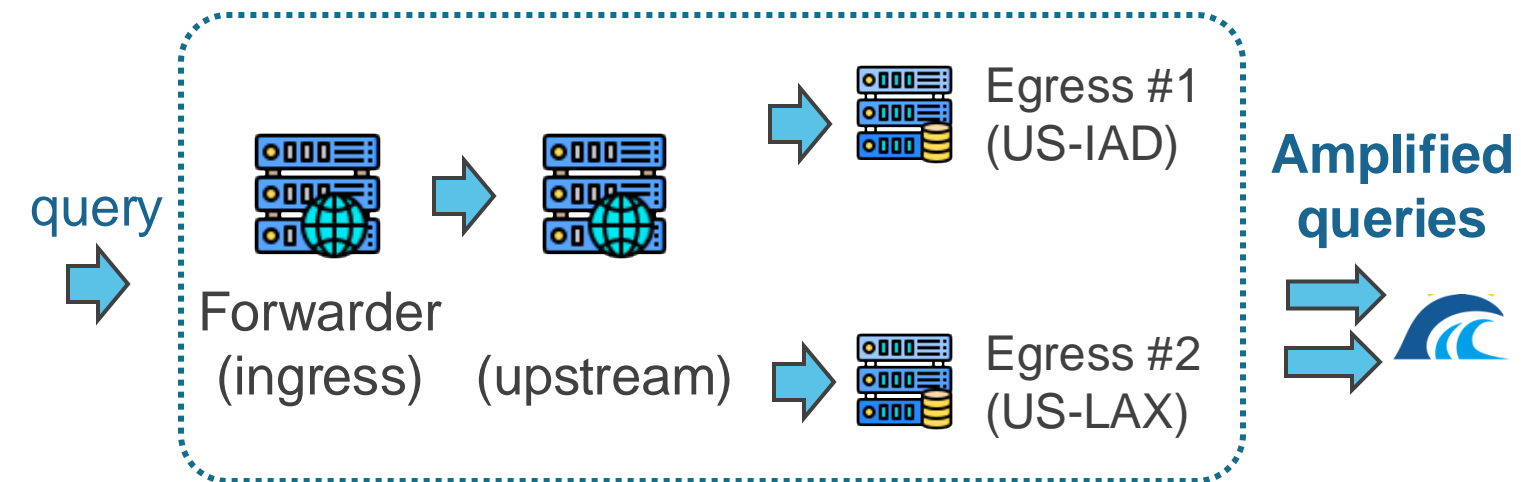


# TsuKing: Tsunami + King



*(Traffic amplification ability)*

- ❖ cause: bogus implementation & complex resolving infrastructure

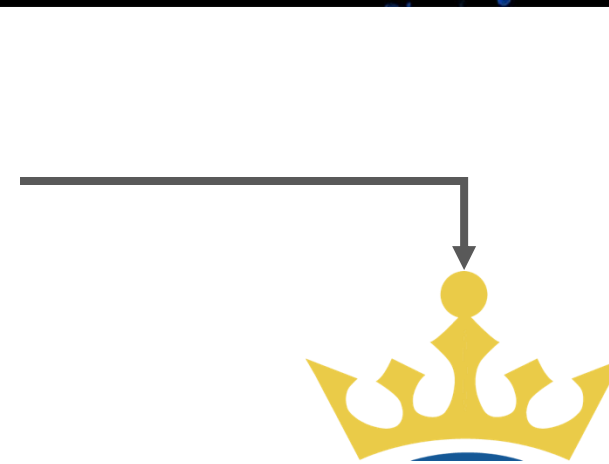


# TsuKing: Tsunami + King [1]



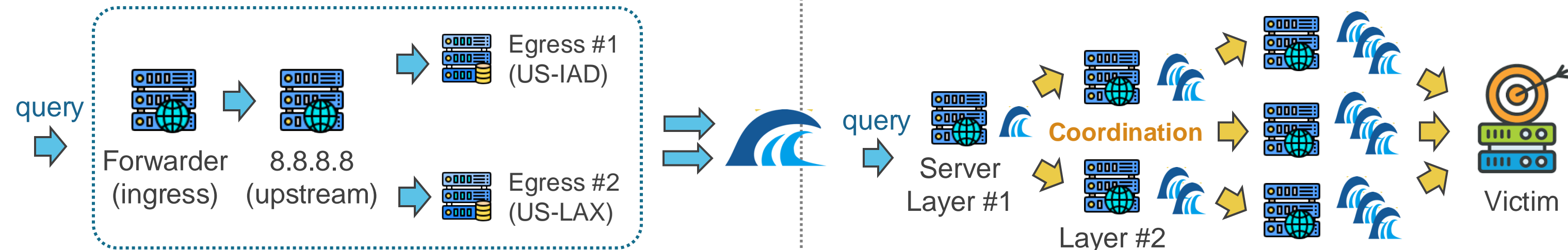
*(Traffic amplification ability)*

- ❖ Cause: DNS implementation choices & complex service infrastructure



*(Server coordination ability)*

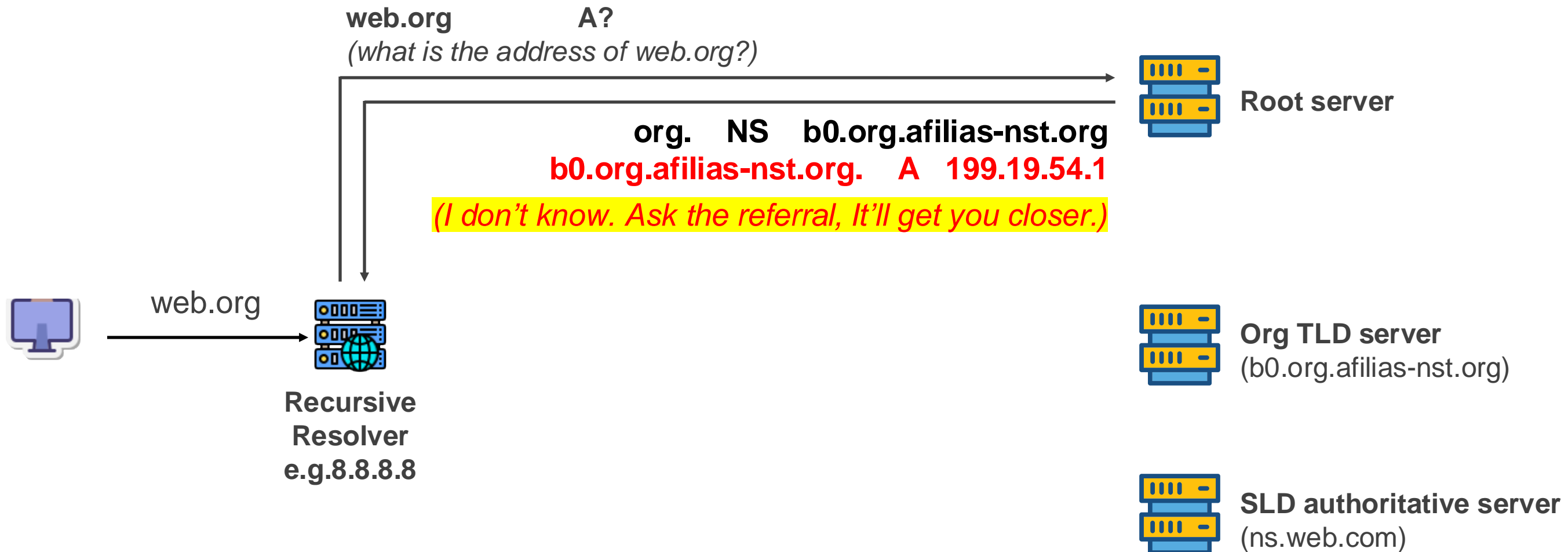
- ❖ Coordinates DNS server systems -> **3,000+X amplification factor (king of DoS)**



[1] **King**: estimating latency between arbitrary internet end hosts, ACM CCR 2002

# DNS resolution guided by *referrals*

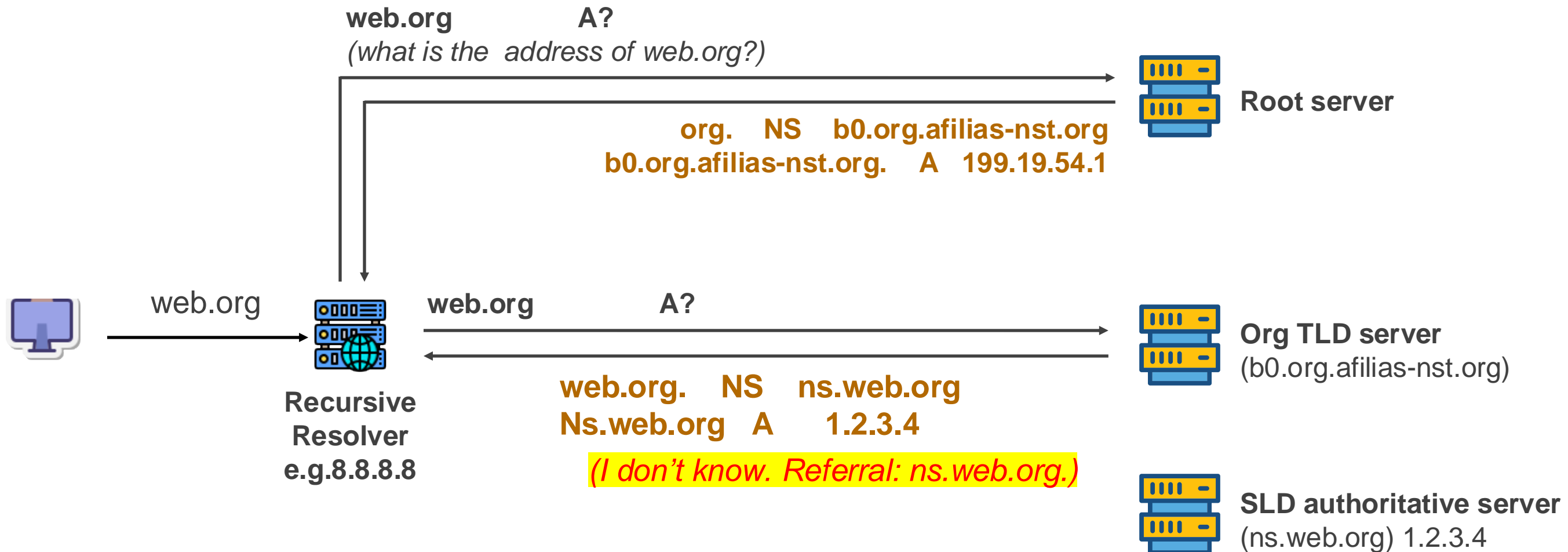
❖ Referrals *tell recursive resolvers who to ask next*





# DNS resolution guided by *referrals*

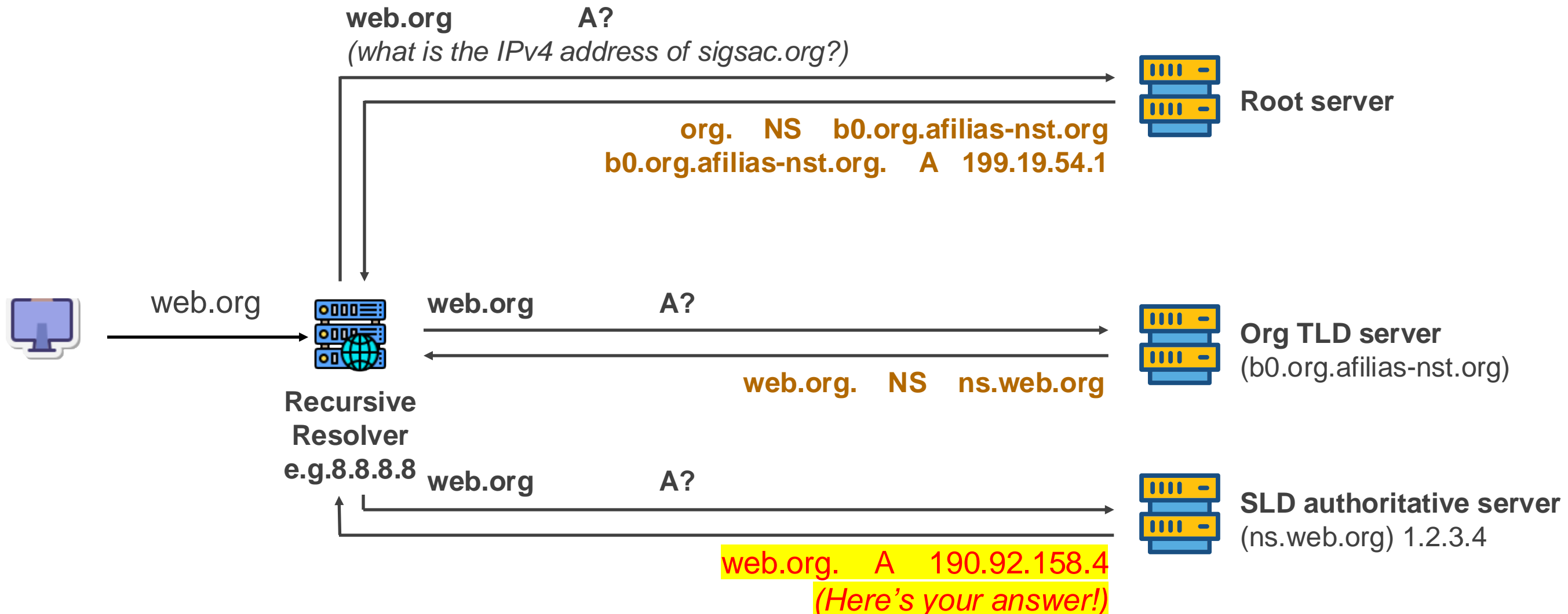
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# DNS resolution guided by *referrals*

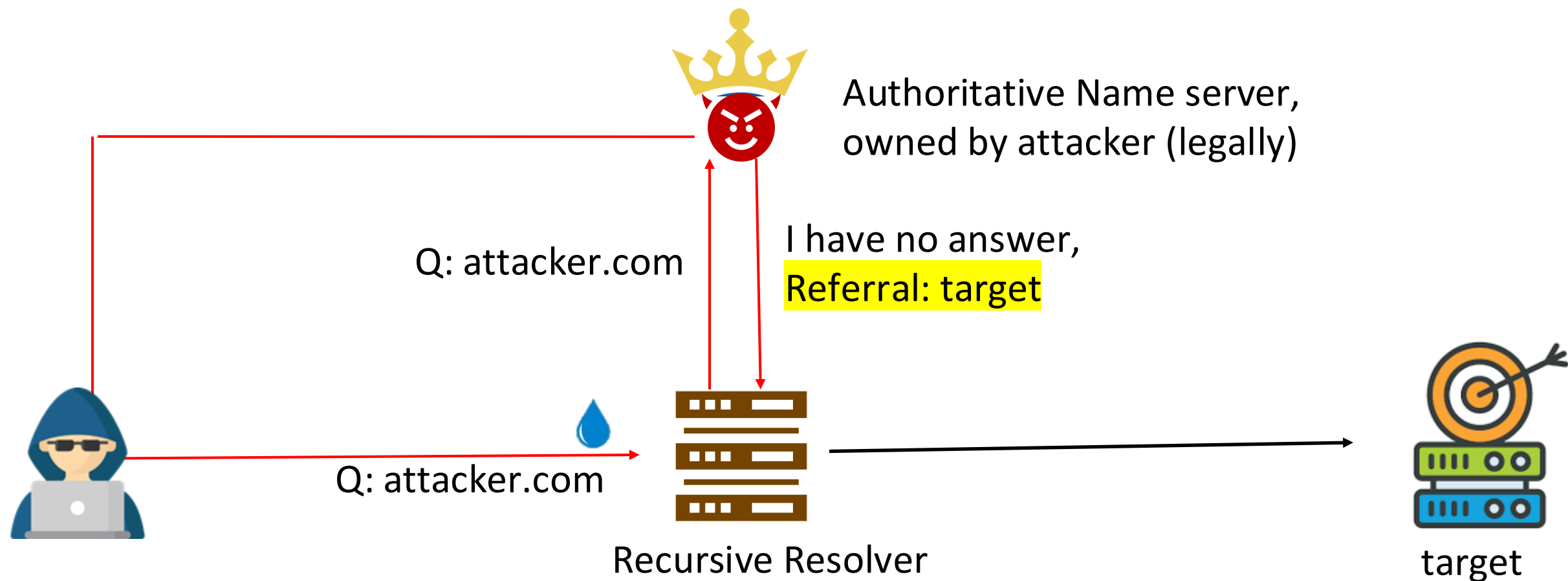
## ❖ Recursive DNS resolution guided by *referrals*

❖ Referrals *tell recursive resolvers who to ask next*

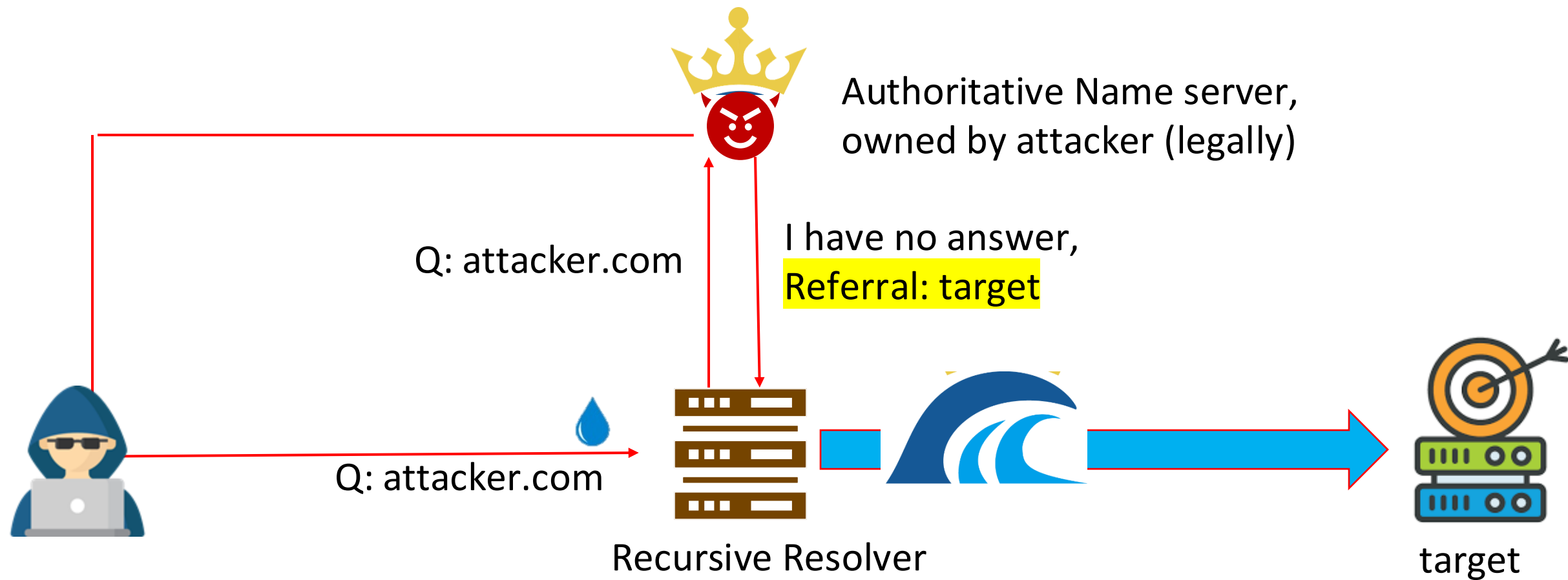


# Threat model of TsuKing

- ❖ Attacker sends DNS query to recursive resolver for his own domain name



# Threat model of TsuKing



But where does the tsunami come from?

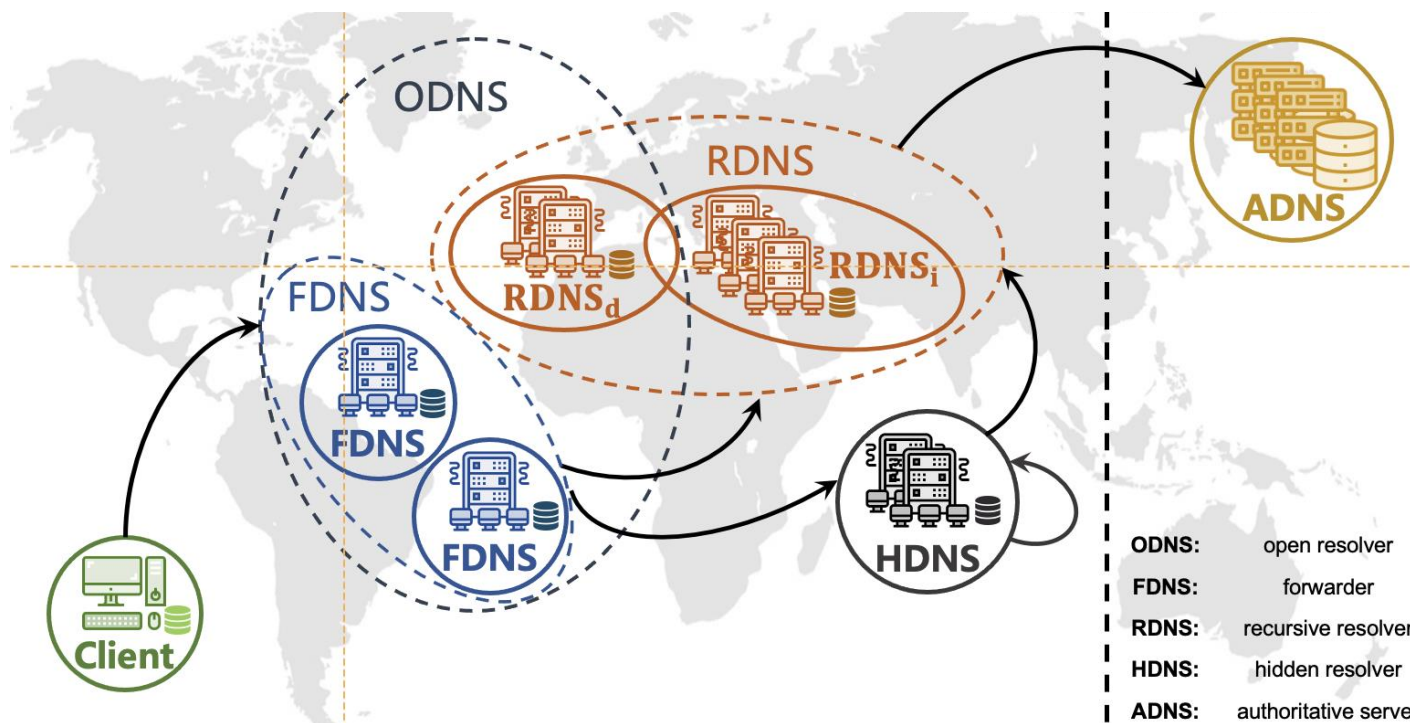


# DNS as a complex infrastructure

## ❖ Multiple *types* and *layers* of DNS servers

- ❖ DNS forwarders → pass queries to upstream (e.g., another forwarder)
- ❖ Large public DNS services → complexes of load balancers, caches, egress servers, etc.

## The complex DNS infrastructure



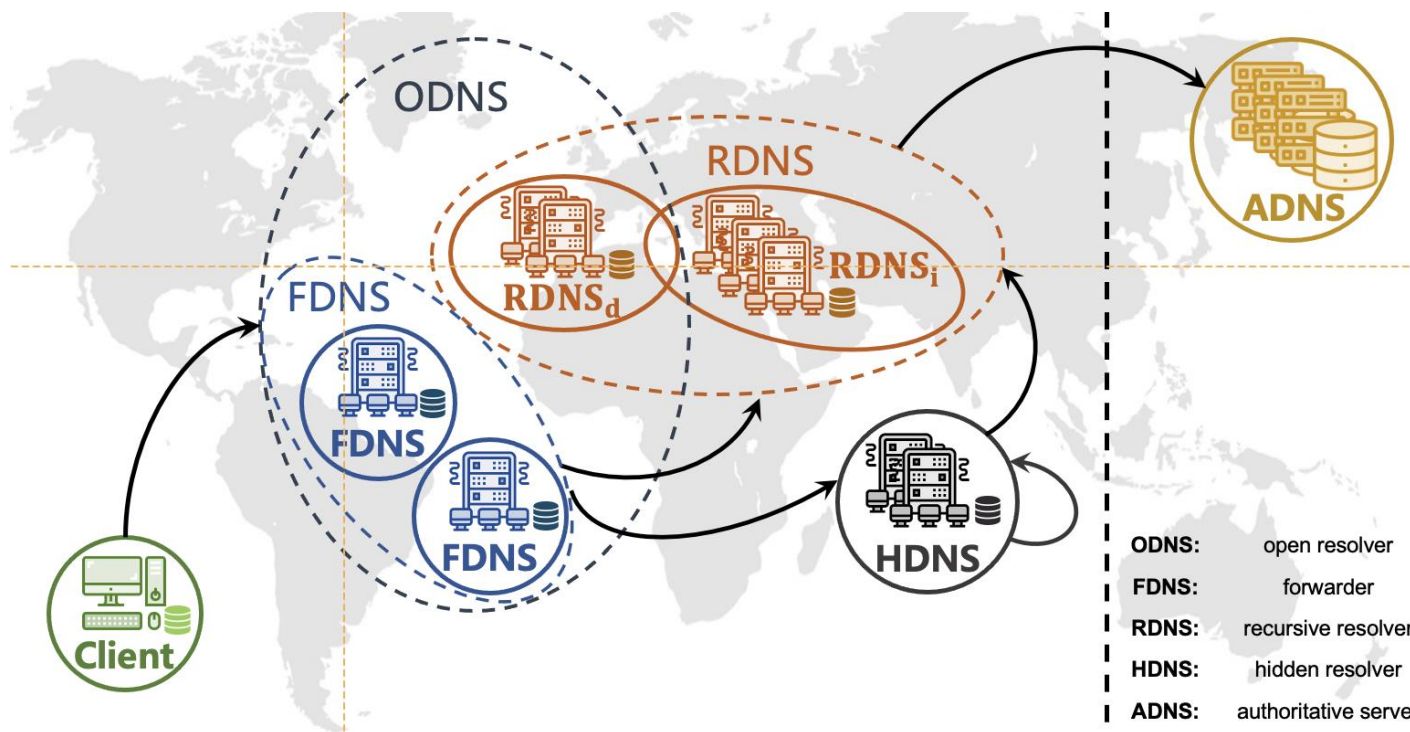
Schomp, et al. On Measuring the Client-side DNS Infrastructure, IMC 2013

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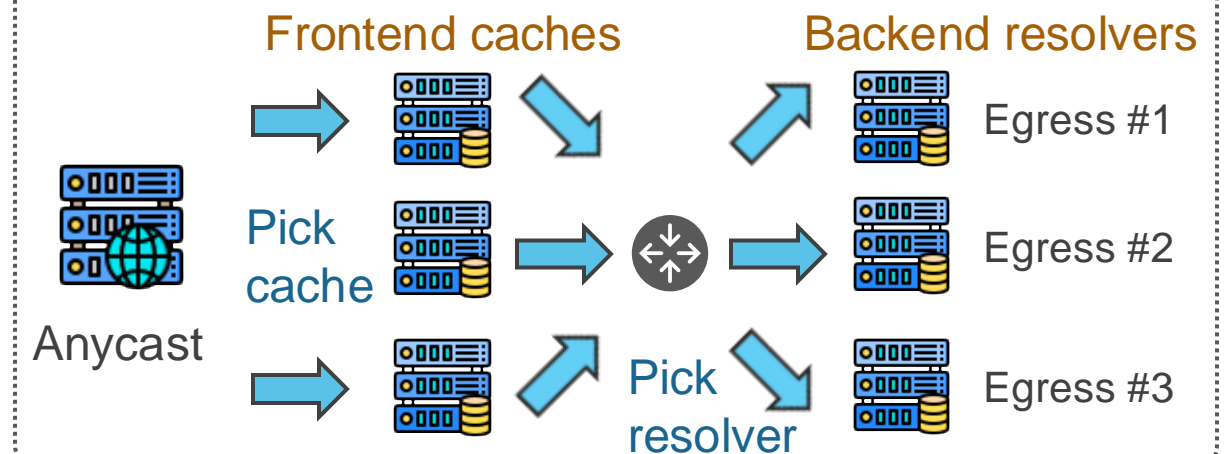
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## The complex DNS infrastructure



## Large public DNS service

(e.g., Google Public DNS)



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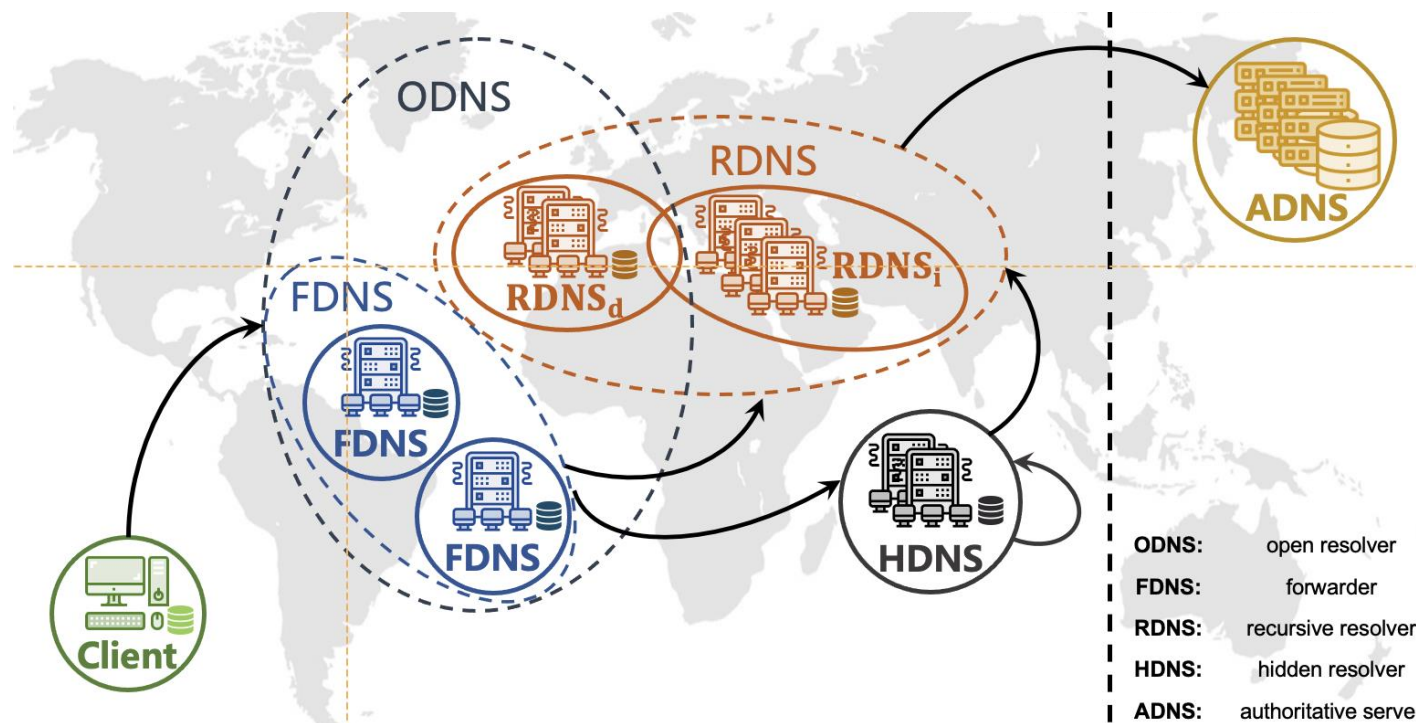


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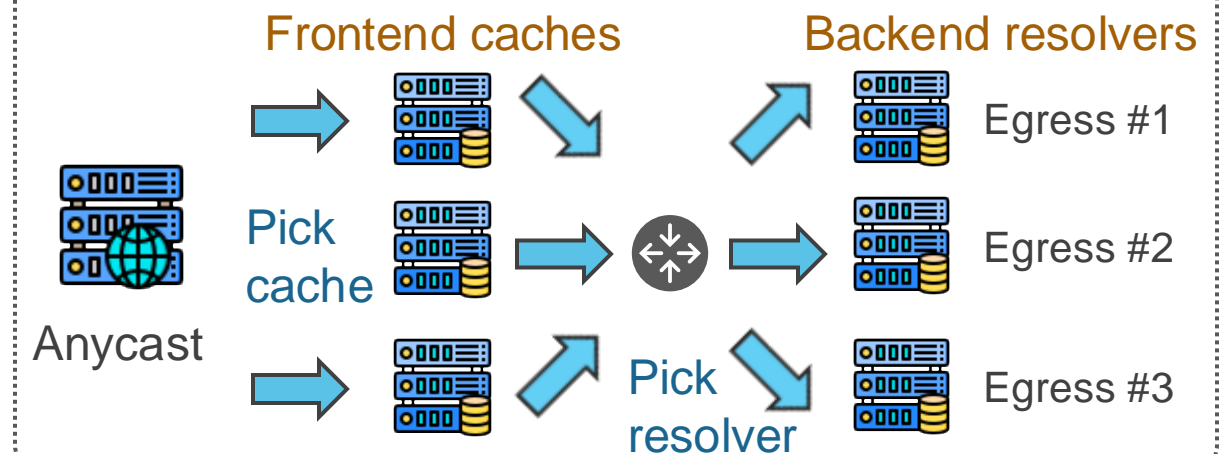
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## Large public DNS service

(e.g., Google Public DNS)



**2.27 Million**  
Open DNS servers

\* Data from Censys,  
Oct 2023

#BHEU #THU @BlackHatEvents

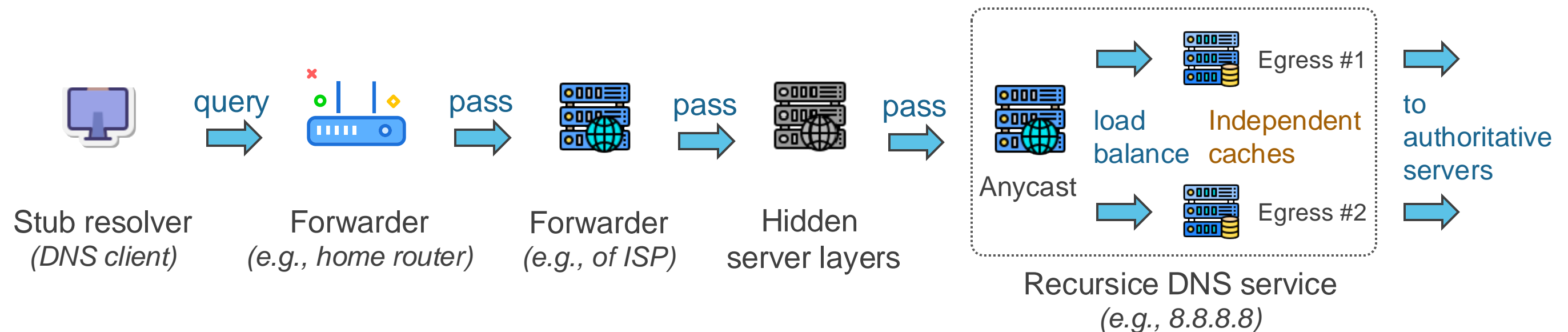
# A typical domain name resolution path

- ❖ Multiple *types* and *layers* of DNS servers

- ❖ DNS forwarders → pass queries to upstream (e.g., another forwarder)

- ❖ Large public DNS services → complexes of load balancers, caches, egress servers, etc.

- ❖ A *typical* DNS resolution path now looks like this



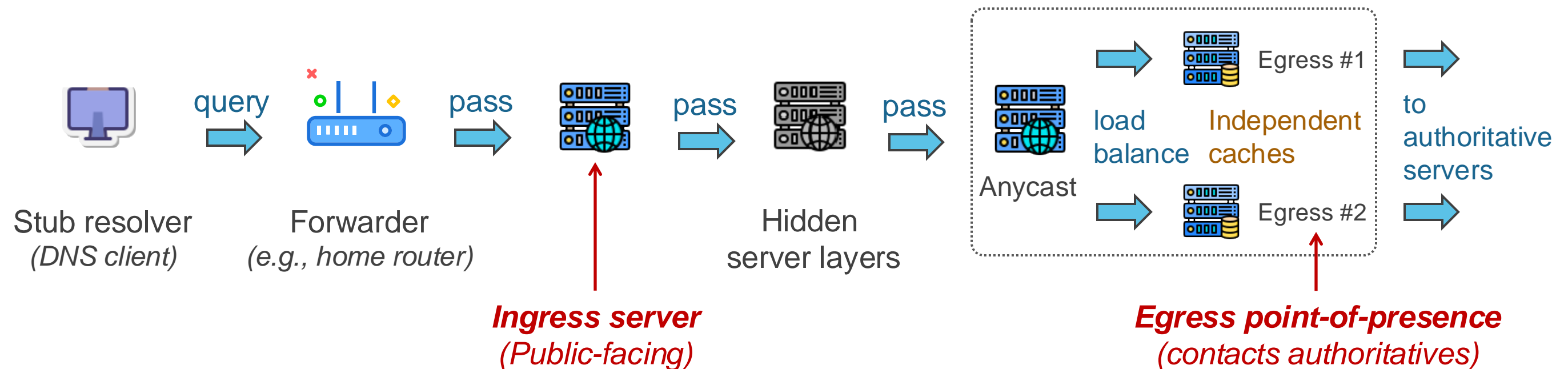


# DNS as a complex infrastructure

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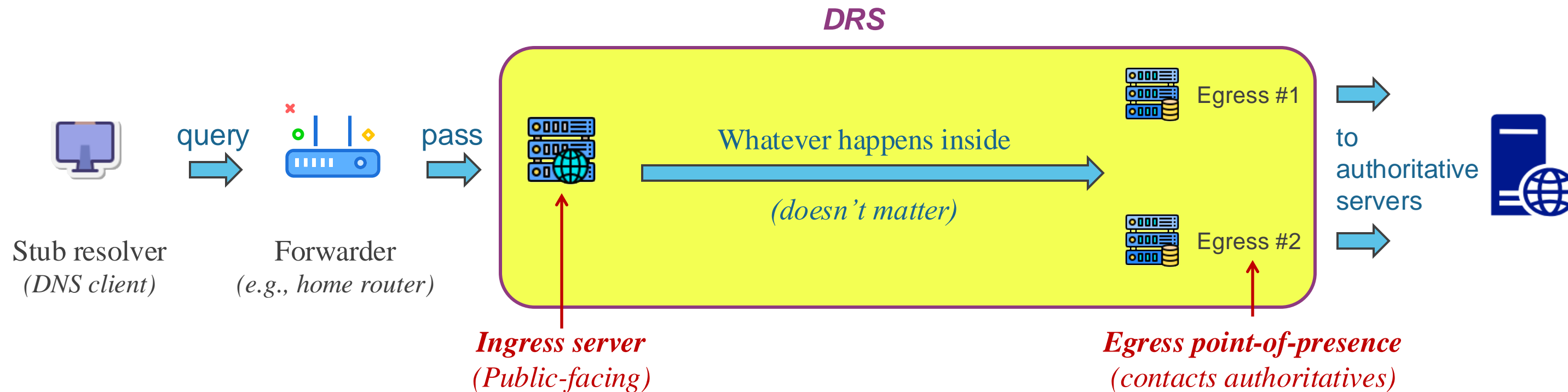


# Definition of DNS Resolver System(DRS)

## ❖ DNS resolver system (DRS)

❖ A public-facing DNS server, together with everything between it and authoritative servers

## ❖ Black box inside



# DNS as a complex infrastructure

**OK, I get it.  
DNS resolver is a complex system.**

But where does the power, or amplification, come from?

# Amplification ability: DNS retries

- ❖ DNS query could fail for variety of reasons
  - ❖ Packet lost, server fail, routing problems
- ❖ So upon failure, please *retry* for a few more times
  - ❖ Adopted by mainstream DNS software
  - ❖ ***THE amplification potential exploited by our attack***

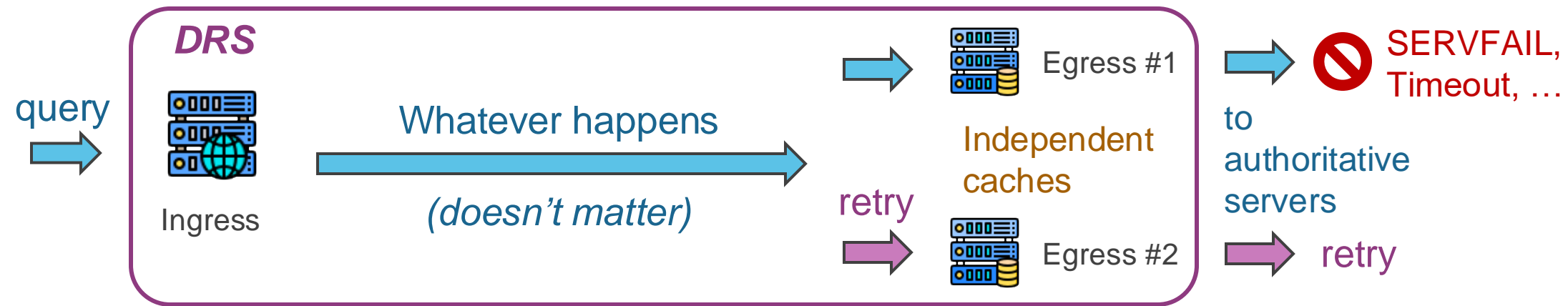
DNS software	# of retries
BIND9	13
Unbound	9
Knot	3



# Amplification ability: DNS retries

❖ For a DRS, retries may exit from *different egresses*

- ❖ Egress servers don't share cache
- ❖ Prevents *query aggregation* and *cache hits*



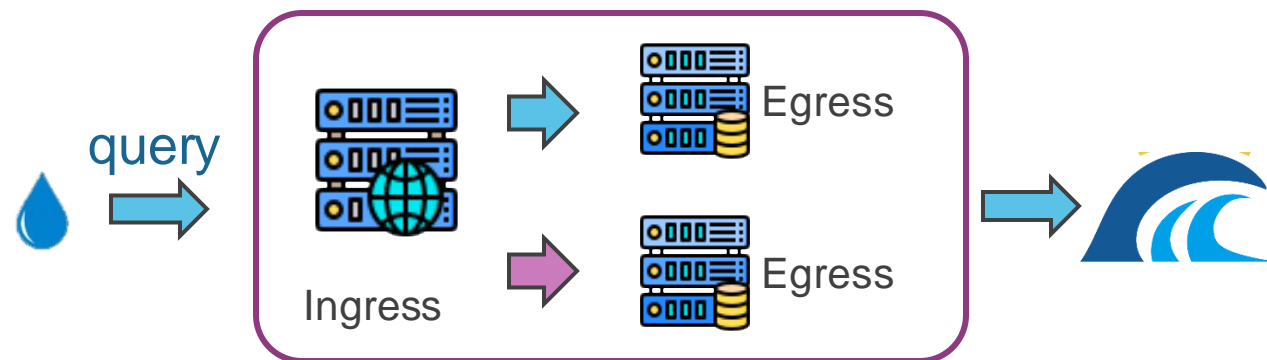
# Amplification ability: DNS retries

**Wait... You exploit retries?**

That's not even enough  
to cause ripples!

# Attack variant I: DNS-Retry

- ❖ Some bogus DRS implementations that retry aggressively
- ❖ In **1.3M DRS**, 2.4% (>30,000) retry more than 100 times
- ❖ 529 DRSES retry more than 1,000 times
- ❖ Max # of retries by one DRS: **117,541**



Amplification by one DRS only is big enough

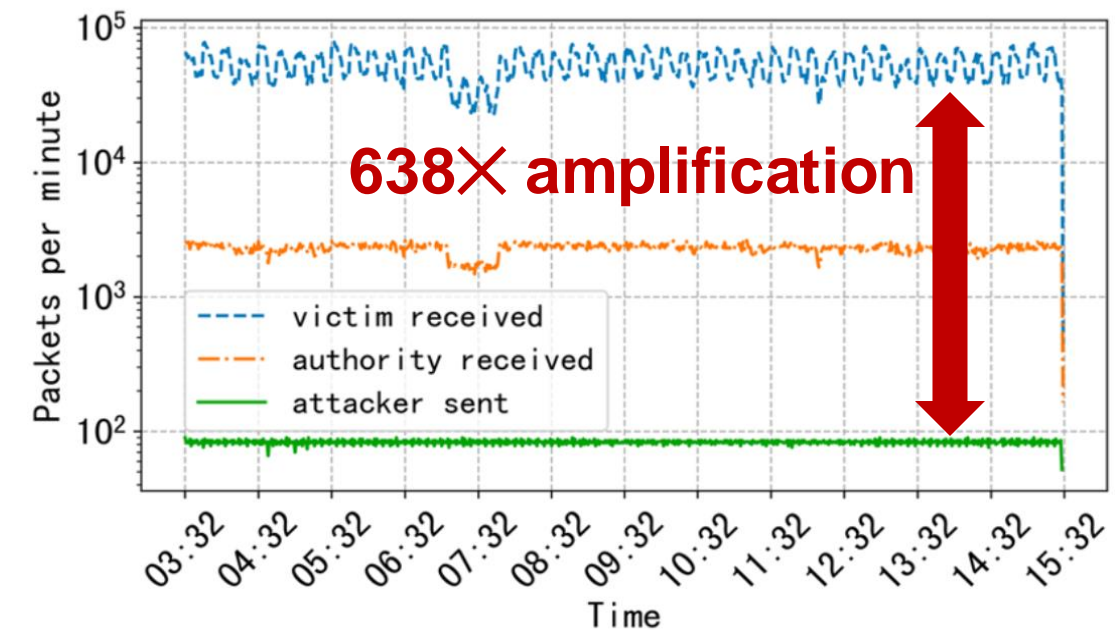
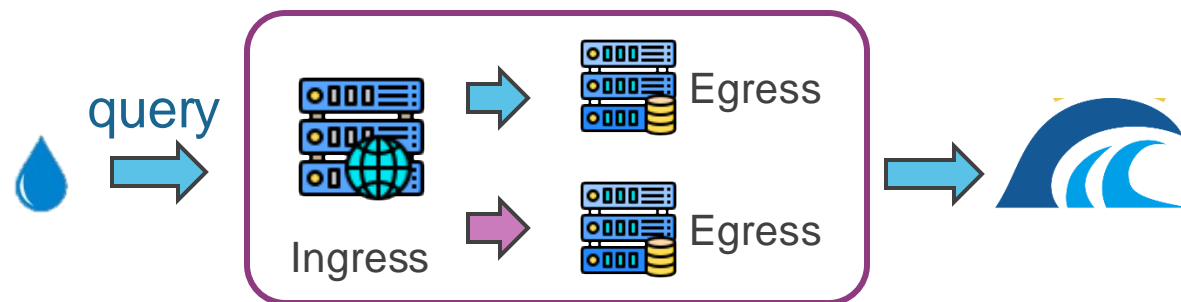
In 1.3M open DNS Resolver System(DRS)

# of retries	# of open DRSES	% of tested
> 2	925,500	69.8%
> 10	407,581	30.7%
> 100	31,660	2.4%
> 1,000	529	0.04%

# DNS-Retry Evaluation

## ❖ Evaluation in controlled environment

- ❖ Select 10 DRSeS that retry aggressively
- ❖ Attacker sends 1.3 pkt/s → **Victim receives 882 pkt/s**





**Alright, but lots of them are not  
aggressive at all.  
Only modest retries...**



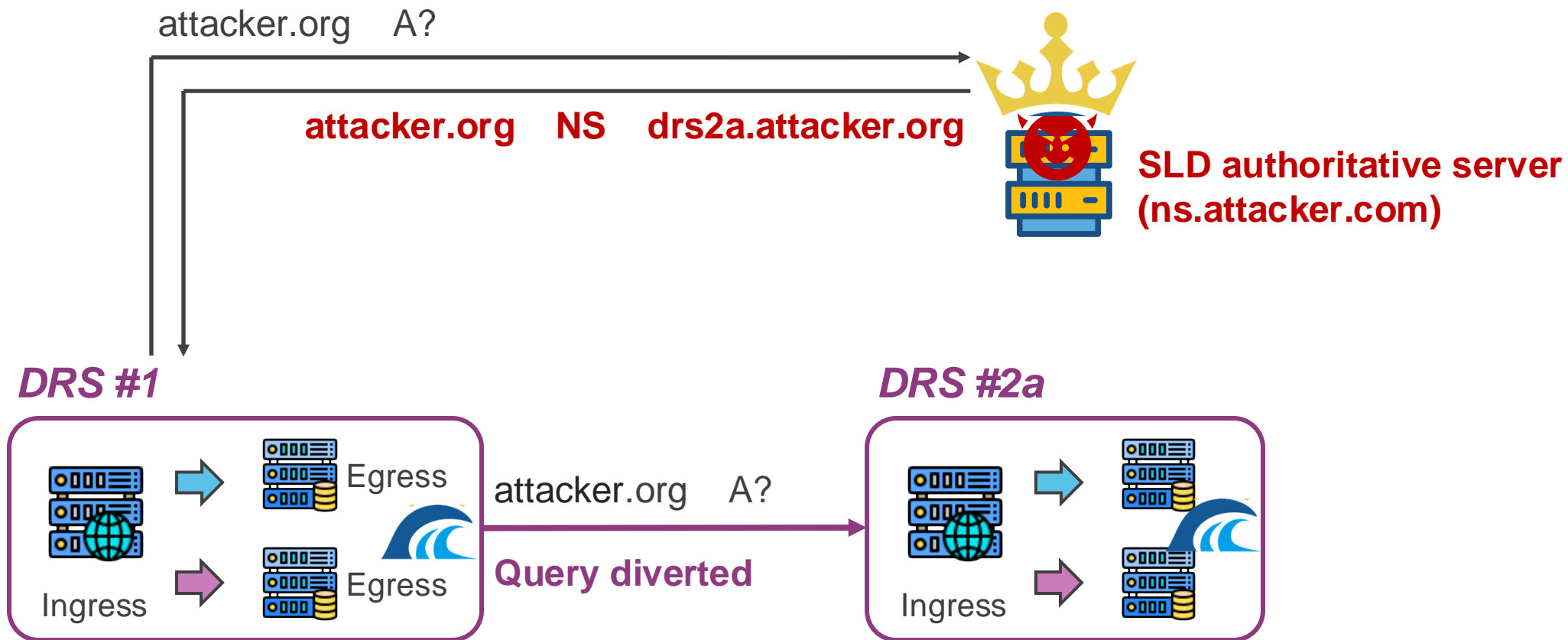


Let's *chain* these ripples into bigger waves!



# Attack variant II: DNS-Chain

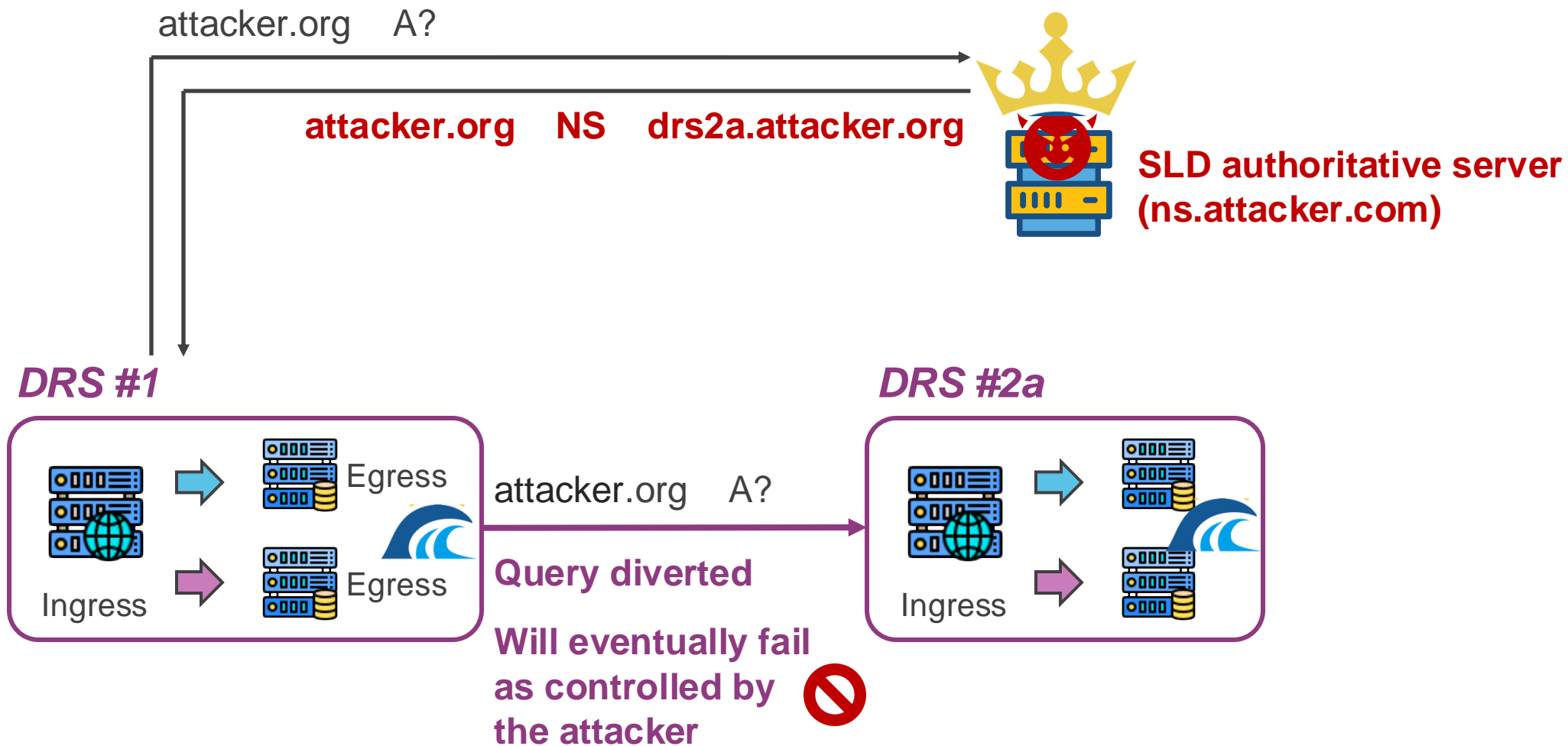
## ❖ Recursive DNS resolution guided by *evil referrals*





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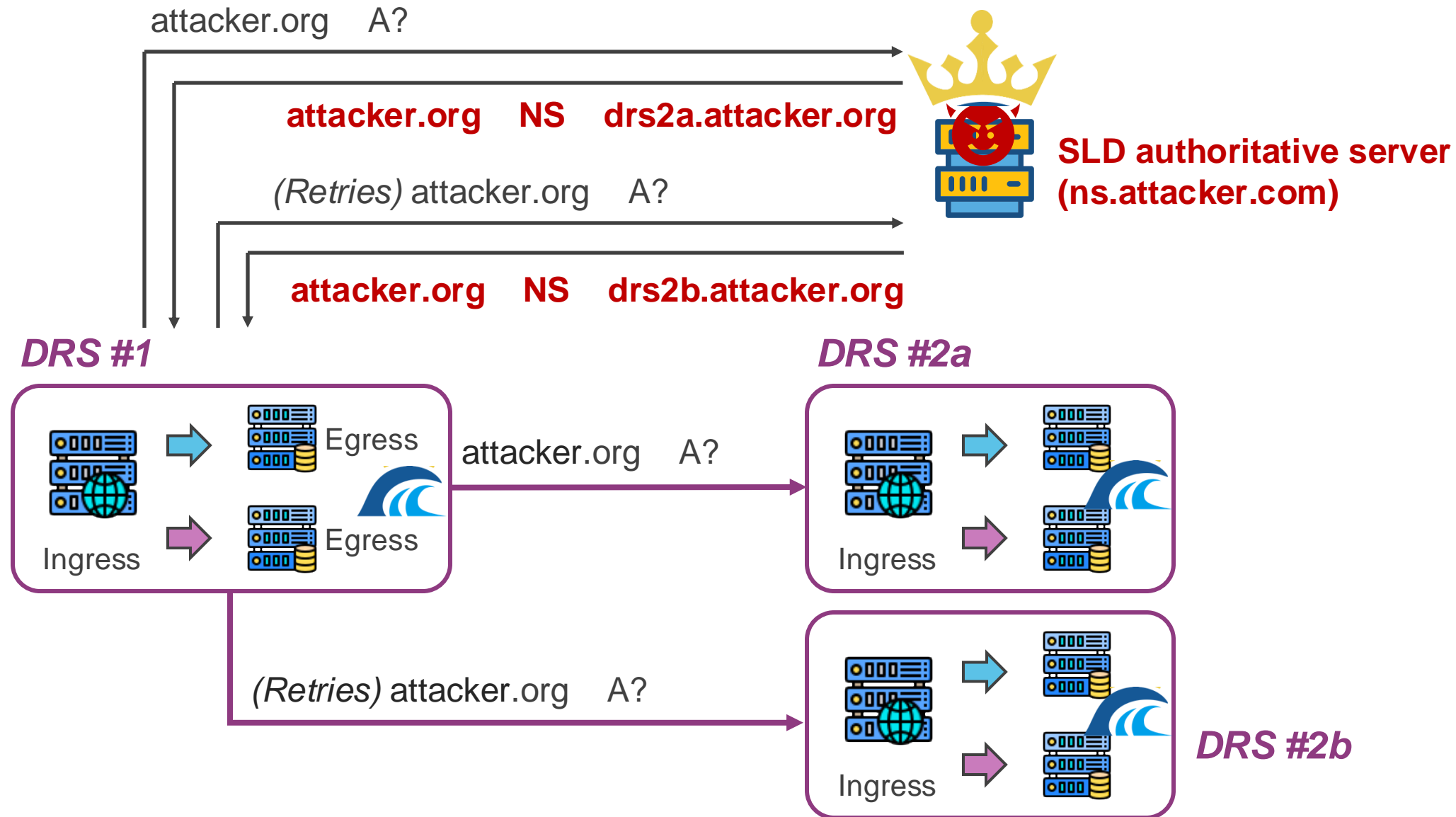
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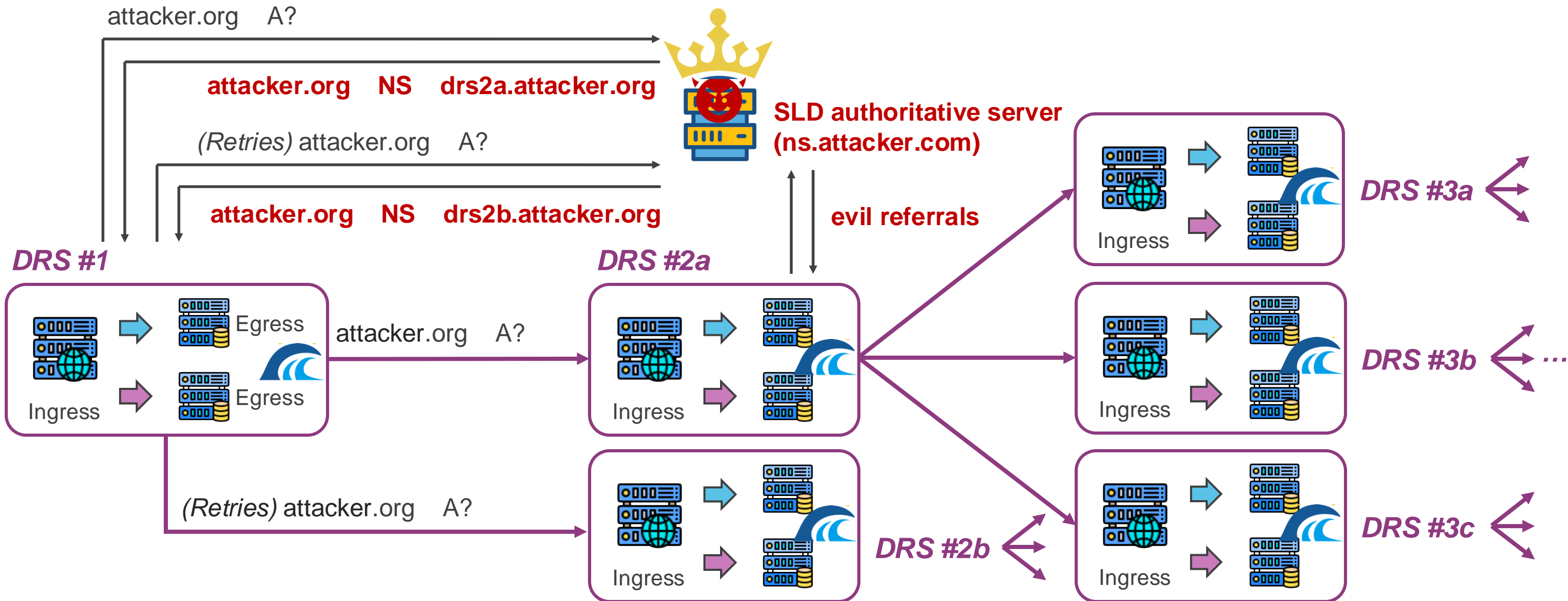
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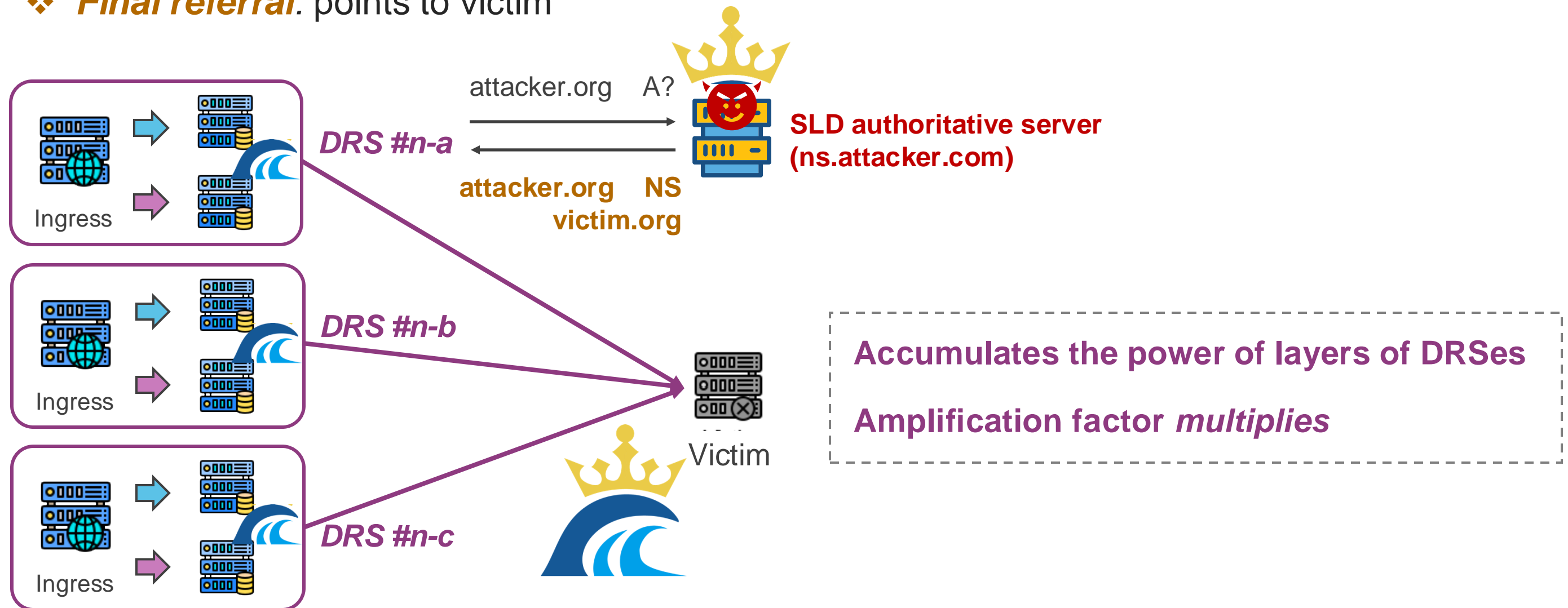
## ❖ Recursive DNS resolution guided by *evil referrals*



# Attack variant II: DNS-Chain

❖ Recursive DNS resolution guided by *evil referrals*

❖ *Final referral*: points to victim





**Seems plausible,  
but can many DRSeS be used?**

What are the conditions of successful attacks?



# Conditions of successful attacks

## ❖ DRS *not honoring cleared RD bit* in DNS header

❖ RD (recursion desired) =0: *do not perform recursion, find answers locally in cache*

❖ Usually *cleared by egress*, as authoritative servers cannot perform recursion

❖ DRS honors RD → *chain cannot continue*

❖ **27.2% of 1.3M tested DRSES do not honor**

Transaction ID	QR	Opcode	<b>RD</b>	Flags	Z	RCODE
QDCOUNT	ANCOUNT					
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## ❖ DRS not deployed with negative caching [RFC 2308]

❖ Negative caching records DNS failures → *effectively eliminates retries*

❖ **43% of 1.3M tested DRSES do not deploy**

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❖ **43% of 1.3M tested DRSES do not deploy**

## ❖ DRS has multiple egresses: *the more, the better*

❖ **52% of 1.3M tested DRSES have over 10 egresses**

# Evaluation of DNS-Chain

## ❖ Evaluation in controlled environment

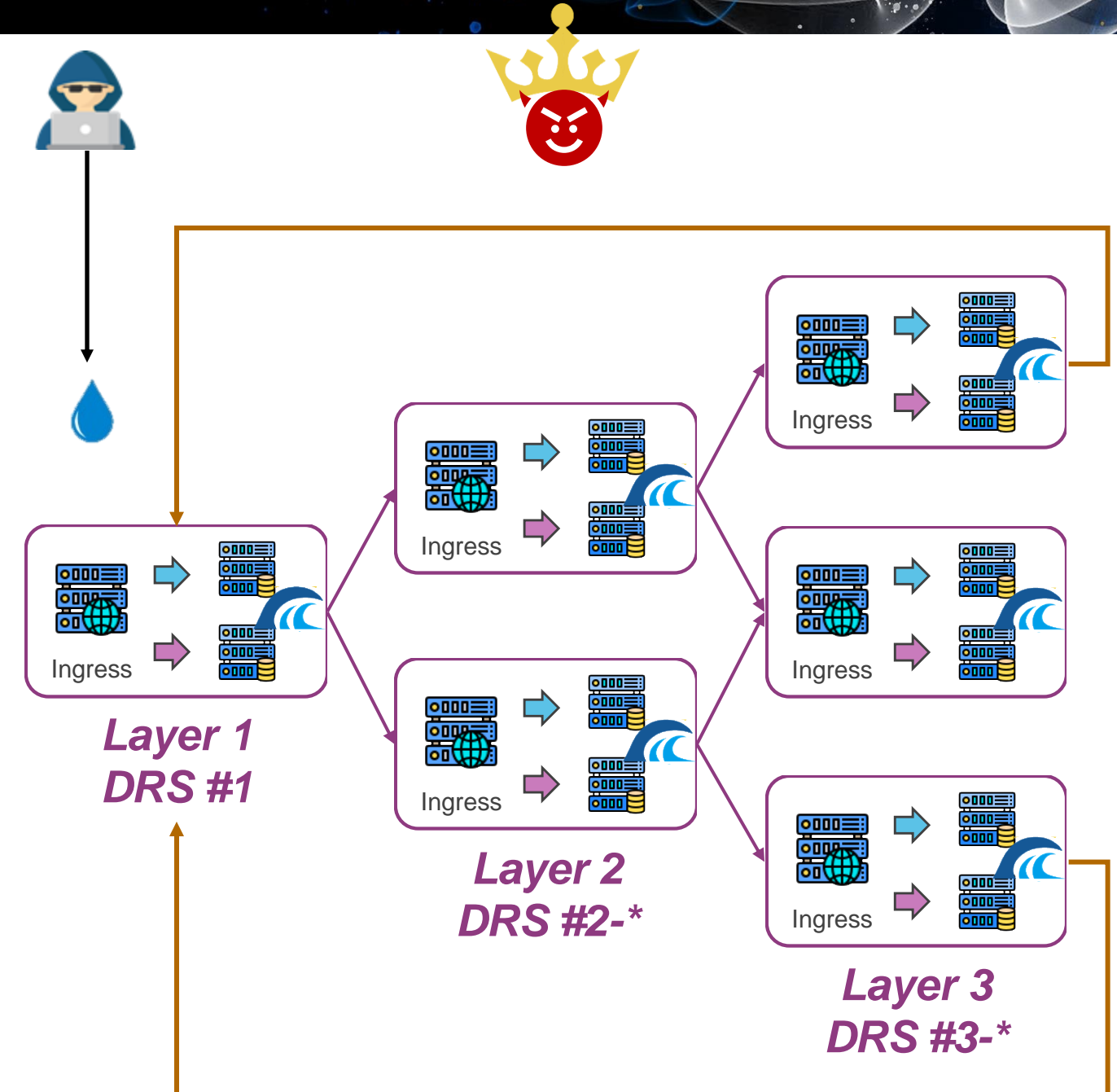
- ❖ We select from exploitable DRSES and coordinate them into *layers*

Setting	# of DRSES coordinated in each layer							Amp. factor
	Layer 1	Layer 2	Layer 3	Layer 4	Layer 5	Layer 6	Layer 7	
# 1	1	4	8	-	-	-	-	288
# 2	1	4	8	16	32	-	-	591
# 3	1	4	8	16	32	64	128	3,702



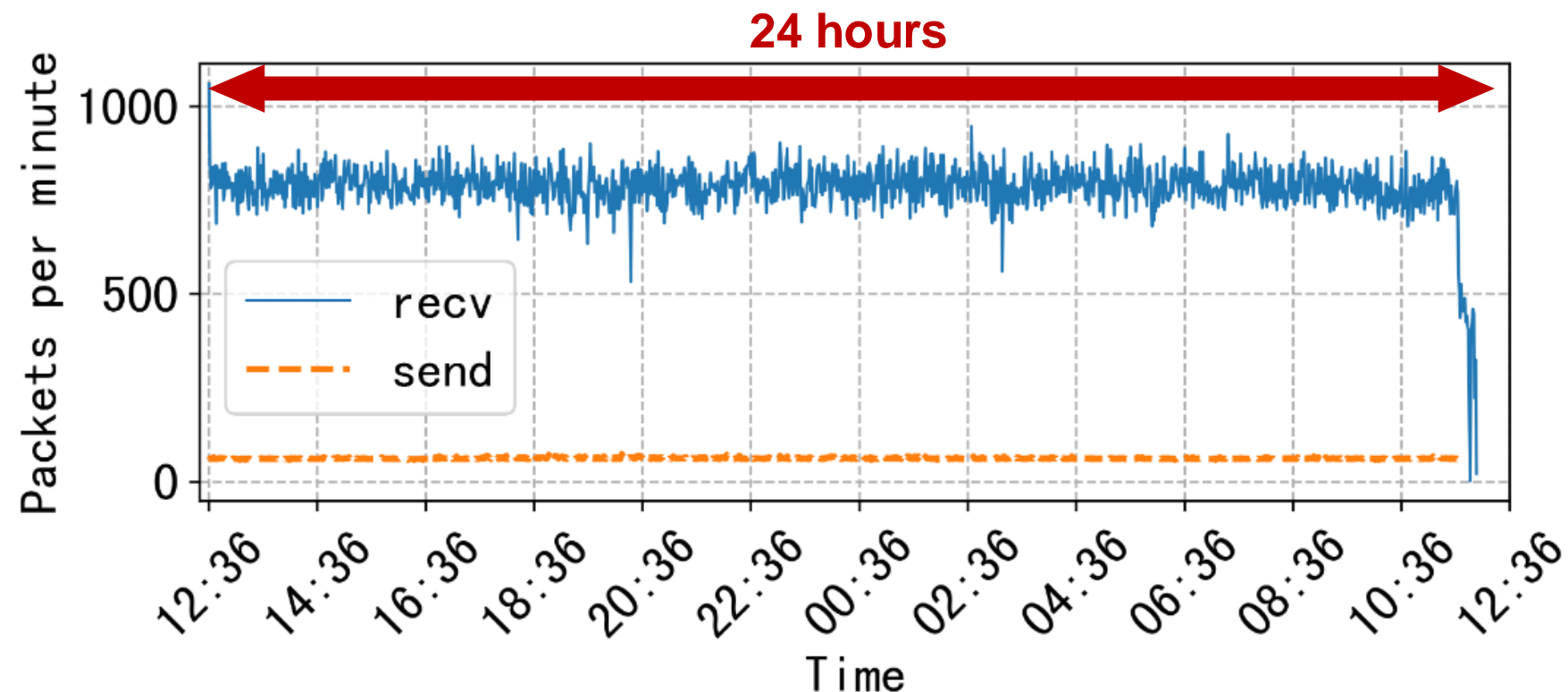
# Attack variant III: DNS-Loop

- ❖ Modified from DNSChain, creating a *loop* of retry queries
  - ❖ *Final referral*: points back to DRS #1
- ❖ The victim and goal change now
  - ❖ *ALL DRSes in the loop* become victims
  - ❖ Goal is to exhaust their resources
  - ❖ *Increasing amplification factor is a non-goal*
- ❖ Attackers may also
  - ❖ Inject new rounds of retries to the loop
  - ❖ Simply by querying DRS #1



# DNS-Loop Evaluation

- ❖ Evaluation in controlled environment - can the loop last?
- ❖ Coordinates 7 layers of DRSeS in the real network
- ❖ layer #0 is our server, with *rate limit at 1 pkt/s (due to ethical considerations)*
- ❖ Send **only one** DNS query Layer 0, to trigger the loop
- ❖ **Loop lasts for 24 hours until deliberate stop**





## What can we do to prevent this attack?

Correct bogus implementations such that attack conditions cannot be fulfilled.

# Mitigations

## ❖ Avoid aggressive retries

- ❖ A **modest number of retries** should suffice, as adopted by mainstream software

## ❖ Follow DNS specifications

- ❖ **Honor the DNS flags:** if RD tells not to perform recursion, just don't

## ❖ Deploy additional mechanisms that add protection

- ❖ **Negative caching:** good to reduce retries
- ❖ **Egress and cache management:** reduce independence between egress servers



# Acknowledgement

## DNS Software Vendors



## DNS service providers



# Questions?



**Paper website: <https://tsuking.net>**

## **Contributors of the slides:**

- ❖ Wei Xu ([xu-w21@mails.tsinghua.edu.cn](mailto:xu-w21@mails.tsinghua.edu.cn))
- ❖ Xiang Li ([x-l19@mails.tsinghua.edu.cn](mailto:x-l19@mails.tsinghua.edu.cn))
- ❖ Chaoyi Lu ([luchaoyi@tsinghua.edu.cn](mailto:luchaoyi@tsinghua.edu.cn))
- ❖ Haixin Duan ([duanhx@tsinghua.edu.cn](mailto:duanhx@tsinghua.edu.cn))

