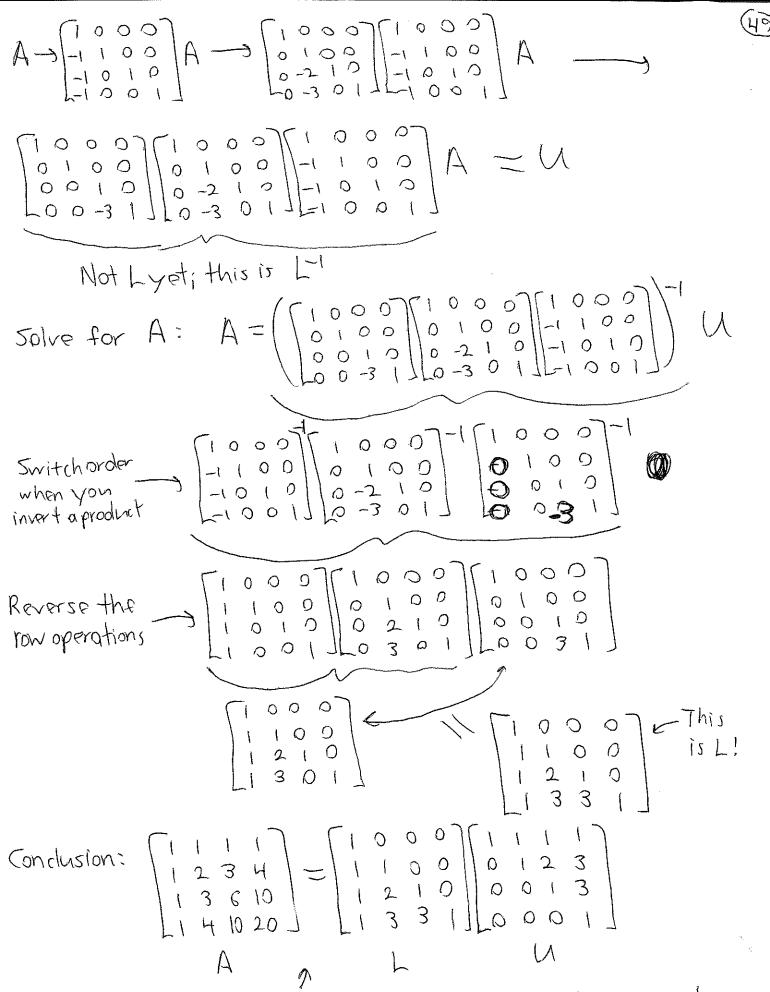
(Reverse motes the row operations to get inverses) $\begin{bmatrix} 1 & 0 & 0 \\ -1/2 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ -1/2 & -1 & 1 \end{bmatrix}$ $50 \quad \begin{bmatrix} 2 & -1 & -1 \\ -1 & 2 & -1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 2 & -1 & -1 \\ 0 & 3/2 & -3/2 \end{bmatrix}$ $\begin{bmatrix} -1/2 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 3/2 & -3/2 \\ -1/2 & -1 & 1 \end{bmatrix} \begin{bmatrix} 0 & 0 & 0 \end{bmatrix}$ Last time: Started LU decomposition Ctriongular Eliminate lower A left variables How? nxn - A = L U matrix - 1 over 1 rupper triangular triangular Same as multiply A by lower triangular elimination matrices not L yet Then: A = [1/2] --- Lm U This is L! Or: multiply by Row 3-2Row4 [0 1 2 3 | Row-3Row3 | 0 1 2 3 | Row 4-3Row4 [0 0 1 3 | 0 0 0 1 3] -1010 L-1001 or multiply by or: Multiply by upper triangular; this is U



Ludecomposition works whenever you can do elimination A -> U without switching rows. How to find the entries of L more efficiently?

Diagonal entries of L: all 1's

Below-diagonal entries: Put Pij in row L, column j if you di

Below-diagonal entries: Put lij in row i, column j if you did
the operation Row i - Row i-lij Row j

Example: $A = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 3 \\ 1 & 9 & 1 \end{bmatrix} \xrightarrow{Row 2 - 1} \xrightarrow{Row 1} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 2 \\ 0 & 3 & 8 \end{bmatrix} \xrightarrow{Row 3 - 1} \xrightarrow{Row 2} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 3 & 8 \end{bmatrix} \xrightarrow{Row 3 - 1} \xrightarrow{Row 2} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 3 & 8 \end{bmatrix} \xrightarrow{Row 3 - 1} \xrightarrow{Row 2} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 3 & 8 \end{bmatrix} \xrightarrow{Row 3 - 1} \xrightarrow{Row 2} \begin{bmatrix} 1 & 1 & 1 \\ 0 & 3 & 8 \end{bmatrix} \xrightarrow{Row 3 - 1} \xrightarrow{Row 3 - 1}$

 $\begin{bmatrix} 0 & 1 & 2 \\ 0 & 0 & 2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$

Check that this is correct: $LU = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 & 2 \\ 1 & 3 & 1 \end{bmatrix} \begin{bmatrix} 0 & 0 & 2 \\ 1 & 4 & 9 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 4 & 9 \end{bmatrix} = A$

Summorize LU decomposition olgorithm:

A lower left triangular

lower triangular

Comes from operation

Row i - Row i - lij Row j

note the

Warning 1: Formula for Lonly works if you do the elimination operations in the right order:

Eliminate Column 1 variables using Row 1 first, Eliminate Column 2 variables using Row 2 next,

Warning 2 = LU decomposition only works if you can go from A to U without switching rows.

Now: What is the use of LU decomposition?

It turns one hard system of equations
$$A \stackrel{>}{\times} = \stackrel{>}{b}$$
 into $\stackrel{\frown}{\bullet}$ $\frac{1}{1}$ $\frac{1}{1}$

 $x_{3}+3x_{4} = 0 \rightarrow x_{3}=-3(2)=-6$ $x_{4} = 2$ $50 \ \dot{x} = \begin{bmatrix} 0 \\ 5 \\ -6 \\ 2 \end{bmatrix}$

Note: The two triangular systems are easy to solve, but we (52) still have to work hard to find L and U.

Solving systems using LU works best when you have a lot of systems with the same A:

 $A\vec{x}_1 = \vec{b}_1, A\vec{x}_2 = \vec{b}_2, A\vec{x}_3 = \vec{b}_{31} - --, A\vec{x}_{1000} = \vec{b}_{1000}$

Even though you have 1000 systems, you only need to find Lu once. Some statistics from the textbook: (if A is nxn)

To find LU, it usually takes about $\frac{2}{3}$ n³ arithmetic operations

To solve a triangular system usually and addition takes about $2n^2$ arithmetic operations

so what if you solve 1000 systems individually?

A
$$\frac{2}{3}$$
 (1) solve triangular solution $\frac{2}{3}$ $\frac{$

Using LU: A \rightarrow U once $\longrightarrow \frac{2}{3}n^3$ Two triangular systems \longrightarrow $4n^2 \rightarrow \frac{2}{3}n^3 + 4000n^2$ operations

If n is big, then n2 is much less than n3, so tolving by Lu is almost 1000 times faster:

$$\frac{2000 \left(\frac{h^{3}}{3} + h^{2}\right)}{\frac{2}{3} h^{3} + 4000 h^{2}} = \frac{2000 h^{3} \left(1 + \frac{3}{h}\right)}{\frac{2}{3} h^{3} \left(1 + \frac{6000}{h}\right)} = 1000 \cdot \frac{1 + \frac{3}{h}}{1 + \frac{6000}{h}}$$
Goes to 1
as $n \to \infty$

Transposes and Permutation

(53

One final matrix operation: transpose

switch rows and columns of a matrix

Mxn matrix A turn rows nxm matrix AT

$$A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} \longrightarrow A^{T} = \begin{bmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{bmatrix}$$

(i,j)-entry of $A^T = (j,i)$ -entry of A

Another viewpoint = "Flip A over the diagonal" to get AT =

$$A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{bmatrix} \longrightarrow A^{T} = \begin{bmatrix} 1 & 4 & 7 \\ 2 & 5 & 8 \\ 3 & 6 & 9 \end{bmatrix}$$

Rules
$$(A+B)^{T} = A^{T} + B^{T}$$
 $(A+B)^{T}_{ij} = (A+B)_{ji}$
 $(AB)^{T} = ??$
 $(A-1)^{T} = ??$
 $= A^{T}_{ij} + B^{T}_{ij} = (A^{T} + B^{T})_{ij}$

(i,j)-entry of $(AB)^T = (j,i)$ -entry of AB= $(jth row of A) \cdot (ith column of B)$ = $(jth column of A^T) \cdot (ith row of B^T)$ - $(ith row of B^T) \cdot (ith row of B^T)$

50 [(AB)T = BTAT) (Switch order, just like for inverse)

$$AA^{-1} = I \longrightarrow (AA^{-1})^{T} = I^{T} \longrightarrow (A^{-1})^{T}A^{T} = I$$

(AT) T lookslike inverse of AT.

$$50:\left[\left(A^{-1}\right)^{T}=\left(A^{T}\right)^{-1}\right]$$

SO = [(A-1)T = (AT)-1] This means that AT is invertible if A is, and At isn't invertible if A isn't.

Example
$$A = \begin{bmatrix} 1 & 3 \\ 1 & 4 \end{bmatrix}$$
 Thresse $A^{-1} = \frac{1}{(1)(4)-(1)(3)} \begin{bmatrix} 4-3 \\ -1 & 1 \end{bmatrix}$

Transpose
$$A^{-1} = \begin{bmatrix} 1 & 1 \\ 3 & 4 \end{bmatrix}$$
Thresse $A^{-1} = \begin{bmatrix} 4 & -1 \\ -3 & 1 \end{bmatrix} = (A^{-1})^{T}$