MA 225 Problem Set 11

exercises 1. Write three blueprints for a proof that If blah blah and so on, then A is finite.: each blue print should use a different approach. Consider $\phi: A \to \mathbb{N}_k$

...

Since if $\phi(a) = \phi(b)$ means a = b, ϕ is injective.

...

Since $\forall g \in \mathbb{N}_k, \exists a \in A : \phi(a) = g$, A is surjective. So A is finite.

Consider $\phi: A \to \mathbb{N}_k$.

...

So there exists a $\phi^{-1}: \mathbb{N}_k \to A$. So there exists a bijection between A and N_k , A is finite.

Consider the set A with the rule blah blah blah

...

So $A = \emptyset$, so A is finite.

2. Write two blueprints for a proof that If blah blah and so on, then A is infinite..

Let A be finite.

Blah blah blah...

So $x \in A$. By the original definition, $x \notin A$. So A is infinite by contradiction.

Consider $x \in B$ where B is an infinite set.

Since $x \in A$, $B \subseteq A$. So A is infinite.

3. Why does problem 11 make the assumption that S has at least two elements? Because if it has less than 2, then there are no numbers, c which a < c < b as there is no b.

proofs Write a complete proof for each of the following statements.

1. Long ago, we gave a bit of a handwavey argument that the product of two finite sets was finite. Let A be a finite set (say, with cardinality n) and B be a finite set (say, with cardinality m). Give a concrete proof that $A \times B$ is finite by explicitly constructing a bijection between $A \times B$ and some \mathbb{N}_k .

Claim 1. $A \times B$ is finite if A and B are finite.

Proof. $A \approx \mathbb{N}_n$ and $B \approx \mathbb{N}_m$. By proof 3, $A \times B \approx N_n \times N_k$. Since there are n * k objects in $\mathbb{N}_n \times \mathbb{N}_k$, $\mathbb{N}_n \times \mathbb{N}_k \approx \mathbb{N}_j$. So $A \times B \approx \mathbb{N}_j$ (Transitivity) and therefore is finite.

- 2. If A_1, \ldots, A_p are each finite sets, then $\underset{k=1}{\overset{p}{\times}} A_k = \{(x_1, \ldots, x_p) | x_i \in A_i\}$ is finite.
 - Claim 2. If A_1, \ldots, A_p are each finite sets, then $\underset{k=1}{\overset{p}{\times}} A_k = \{(x_1, \ldots, x_p) | x_i \in A_i\}$ is finite.

Proof. Base Case(p = 1): A_1 is finite (by def)

Inductive hypothesis(p = r): $\underset{k=1}{\overset{p}{\times}} A_k$ is finite.

Inductive step(r+1): We'll show \times A_k is finite: Since A_{r+1} and $\underset{k=1}{\overset{r}{\swarrow}} A_k$ are finite, $\underset{k=1}{\overset{r}{\swarrow}} A_k \times A_{r+1}$ is finite (proof 1). So by PMI, $\underset{k=1}{\overset{r}{\swarrow}} A_k =$ $\{(x_1,\ldots,x_n)|x_i\in A_i\}$ is finite. 3. Write out in detail the proof that if $A \approx X$ and $B \approx Y$, then $A \times B \approx X \times Y$. Claim 3. If $A \approx X$ and $B \approx Y$, then $A \times B \approx X \times Y$. *Proof.* Consider the bijections $\psi: A \to X$ and $\phi: B \to Y$. Let $\psi \times \phi : A \times B \to X \times Y$ which $(a, b) \mapsto (\psi(a), \phi(b))$ where $a \in A$ and $b \in B$. Injectivity: Let $\psi \times \phi((a_1,b_1)) = \psi \times \phi((a_2,b_2))$. In particular, $\psi(a_1) = \psi(a_2)$, so $a_1 = a_2$ (injectivity of ψ). Also $\phi(b_1) = \phi(b_2)$, so $b_1 = b_2$. Therefore $\psi \times \phi$ is injective. Surjectivity: Let $(x,y) \in X \times Y$ where $x \in X$ and $y \in Y$. So $\exists a \in A$ such that $\psi(a) = x$ and $\exists b \in B$ such that $\phi(b) = y$. So $\psi \times \phi((x,y)) = (a,b)$, so $\psi \times \phi$ is bijective. Therefore $A \times B \approx X \times Y$. 4. $(\star\star)$ It's a fact of arithmetic that for any natural numbers $b, m, n, b^{m+n} = b^m b^n$. State, and prove, the set equivalence which reduces to this fact if the three sets involved happen to be finite. Proof. 5. $(\star\star)$ It's a fact of arithmetic that for any natural numbers $b, m, n, b^{mn} = (b^m)^n$. State, and prove, the set equivalence which reduces to this fact if the three sets involved happen to be finite. 6. Let A be an infinite set and R an equivalence relation on A with the property that each equivalence class $[a]_R$ is finite. Show that $A_{/R}$ is infinite. **Claim 4.** If A is an infinite set and R an equivalence relation on A with the property that each equivalence class $[a]_R$ is finite, then $A_{/R}$ is finite. *Proof.* Assume $A_{/R}$ is finite. So $\bigcup_{x \in A_{/R}} x = A$. Since $\bigcup_{x \in A_{/R}} x$ is finite (Proof 9 Problem set 10) and A is not. So by contradiction, $A_{/R}$ cannot be finite, and therefore must be infinite. 7. Show that the union of any collection of finitely many denumerable sets is denumerable; that is, for any $n \in \mathbb{N}$, if A_1, \ldots, A_n are each denumerable, then $\bigcup A_k$ is denumerable. (*Hint.* n is a natural number.) Claim 5. The union of any collection of finitely many denumerable sets is denumerable; that is, for any $n \in \mathbb{N}$, if A_1, \ldots, A_n are each denumerable, then $\bigcup_{i=1}^n A_k$ is denumerable.

Proof. Base case(n=1): $\bigcup_{k=1}^{n} A_k = A_1$ (denumerable by definition)

Inductive Hypothesis(r = n): $\bigcup_{k=1}^{r} A_k$ is denumerable.

Inductive Step:

$$\bigcup_{k=1}^{r+1} A_k \qquad \qquad \bigcup_{k=1}^r A_k \cup A_r$$

So $\bigcup_{k=1}^r A_k \approx \mathbb{N}$ (inductive hypothesis). Since adding a denumerable number keeps a set

denumerable,
$$\bigcup_{k=1}^r A_k \cup A_r$$
 is denumerable. By PMI, $\bigcup_{k=1}^n A_k$ is denumerable. \square

8. Let A be a countable set and R an equivalence relation on A. Show that $A_{/R}$ is countable. Claim 6. $A_{/R}$ is countable.

Proof. Assume $A_{/R}$ is not countable and A is countable. So $\bigcup_{a \in A} [a]_R = A$. However we earlier established A is countable, and $\bigcup_{a \in A} [a]_R$ isn't, this is a contradiction and so $A_{/R}$ is countable.

- 9. Show that the set of all lines in the plane with rational slope and rational intercept is denumerable. (*Hint*. Write each line in slope-intercept form.)
- 10. (\star) No powerset is denumerable.

Claim 7. No powerset is denumerable.

Proof. Let A be a set.

Case A is finite: Then $card(2^A) = 2^{card(A)}$. Since this is finite, 2^A is not denumerable.

Case A is denumerable: Assume 2^A is denumerable. So $2^A = A_1, A_2, \cdots$ for infinite sets. Let $B = b_1, b_2, \ldots$ be a subset of A. Consider A_n term. Change the nth element to k in A_n and define $b_n = k$. Since $b_n \notin 2^A$ (since it differs by definition from every element in 2^A), this is a contradiction, as by definition $b_n \in 2^A$. So 2^A isn't denumerable in this case.

Case A is uncountable: Since 2^A is not smaller than A, 2^A is not denumerable. So no powerset is denumerable.

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- 11. $(\star\star)$ Let $S \subseteq \mathbb{R}$ be a set with the property: given any $a,b \in S$ with a < b, there is some $c \in S$ with a < c < b. Assume S has at least two elements.
 - (a) Show that S is infinite.

Claim 8. S is infinite.

Proof. Assume S is finite. Then $S = s_1, ...s_n$ where $n \in \mathbb{N}$ and $s_1 < s_2 < ... < s_n$. Consider $c = \frac{s_1 + s_2}{2}$. Since $s_1, s_2 \in S$ and $s_1 < c < s_2$. So $c \in S$. However we claimed to have a complete list, and since there were no elements between s_1 and s_2 , this is a contradiction. So S is infinite.

(b) Show that between any two distinct elements of S, there are infinitely many elements of S.

Claim 9. Between any two distinct elements of S, there are infinitely many elements of S.

Proof. Let $a,b \in A$. Consider the set C where $a,b \in C$ where $C=a,c_2...b$ and $a < c_2 < ... < b$. So C is infinite (previous proof). Since all elements of C are < a and > b, there exists an infinite amount of elements between any two distinct points in A.