

A CONSTRUCTION WITH YONEDA STRUCTURES

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ABSTRACT. In **Cat**, we call an object μ -presentable (for μ a regular cardinal) when $\text{hom}(a, -)$ commutes with μ -filtered colimits.

- A category \mathcal{C} is μ -accessible when it has a small subcategory $\mathcal{G} \subseteq \mathcal{C}$ whose objects are all μ -presentable and such that every object of \mathcal{C} is a μ -filtered colimit of objects in \mathcal{G} (such a \mathcal{G} is called a *generator* for \mathcal{C}).
- A category \mathcal{C} is called *locally μ -presentable* if it is μ -accessible and co-complete (equivalently, μ -accessible and complete).

I want to reproduce these two definitions inside a 2-category \mathcal{K} with a Yoneda structure. To do it, I take the definitions in **Cat** and try to rephrase them formally so that they become straightforwardly referred to the 0-,1- and 2-cells in \mathcal{K} .

Claim : There is a monad $T_{\mathcal{C},\mu}$ on $[\mathcal{C}, \mathbf{Set}]$ such that $T_{\mathcal{C}}(F) \cong F$ iff F commutes with μ -filtered colimits.

The unit of this monad induces by restriction a 2-cell

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} & \xrightarrow{y} & [\mathcal{C}, \mathbf{Set}] \\ \downarrow y & \alpha \swarrow & \nearrow T_{\mathcal{C},\mu} \\ [\mathcal{C}, \mathbf{Set}] & & \end{array}$$

Now if \mathcal{C} has a generator \mathcal{G} , composing with $j : \mathcal{G} \rightarrow \mathcal{C}$ one has $\alpha * j : Tyj(a) \cong yj(a) = \text{hom}(a, -)$, since all objects \mathcal{G} are μ -presentable.

The fact that every object in \mathcal{C} is a colimit of objects in \mathcal{G} can be expressed by saying that j is dense: $\text{Lan}_j j \cong 1$. So, it seems that

Claim : \mathcal{C} is μ -accessible iff

- it has a dense generator $j : \mathcal{G} \rightarrow \mathcal{C}$,
- such that if $\alpha := \eta * y$ is the restriction of the unit of $T_{\mathcal{C},\mu}$ to representables, the whiskering $\alpha * j$ is invertible.

Is it possible to find T ? If yes, we're really near to define accessible and locally presentable objects in \mathcal{K} with a YS: in fact, the definition of locally presentable is way easier: since \mathcal{C} is locally presentable if it is a localization of a category of presheaves, it seems enough to say that \mathcal{C} is locally presentable if it has a generator j such that $\mathcal{G}(j, 1)$ has a left adjoint and it is fully faithful.

I would like to know

- if this circle of ideas makes any sense at all; has anybody tried to do “the formal category theory of locally presentable and accessible objects”?
- if (and how) the existence of T can be verified, and when it exists in $\mathcal{K} \neq \mathbf{Cat}$ what is its definition; I expect it to exist as a consequence of a specific property of \mathcal{K} (or maybe even from a new axiom?)

- if the necessity to work with the *contravariant* Yoneda in

$$\begin{array}{ccc}
 \mathcal{C}^{\text{op}} & \xrightarrow{y} & [\mathcal{C}, \mathbf{Set}] \\
 y \downarrow & \alpha \nearrow & \\
 [\mathcal{C}, \mathbf{Set}] & \xleftarrow{T_{e,\mu}} &
 \end{array}$$

poses a problem: are we dealing with a...right Yoneda structure (ie with right extensions and liftings)? If yes, can the theorems about these dualized Yoneda structure be dualized in the straightforward way?
