

Yesterday = $\mathcal{Y}\mathcal{L}$

Today = instances of $\mathcal{Y}\mathcal{L}$ in categories
we know / care about

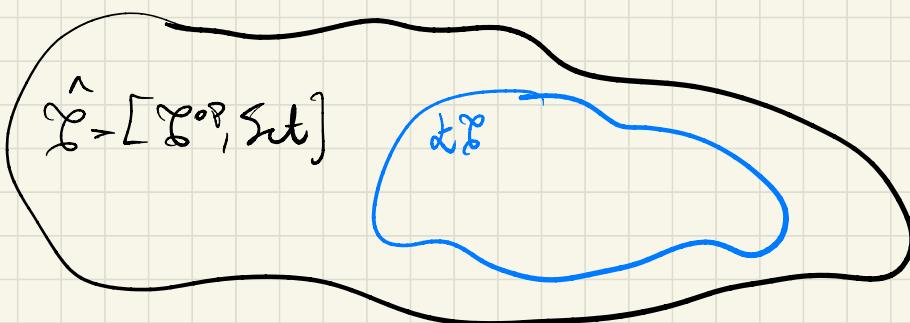
$\mathcal{Y}\mathcal{L} \approx$ Understand some functors well

actually $\mathcal{Y}\mathcal{L} =$ description of a subcategory of $[\mathcal{C}^{\text{op}}, \text{Set}]$

subcategory of representable

$$\cong \text{hom}(-, X) : \mathcal{C}^{\text{op}} \rightarrow \text{Set}$$

subcategory of repr is equivalent to \mathcal{L}



$\mathcal{L}_{\mathcal{C}} : \mathcal{C} \rightarrow [\mathcal{C}^{\text{op}}, \text{Set}]$ is fully faithful

$$\left\{ \mathcal{L}_{\mathcal{C}} X \Rightarrow \mathcal{L}_{\mathcal{C}} Y \right\} \cong \{ X \rightarrow Y \}$$

If you want $X \cong Y$, you can instead:

- i) take arbitrary $A \in \mathcal{C}$
- ii) Prove $\text{hom}(A, X) \cong \text{hom}(A, Y)$ naturally bijective

iii) looks more complicated

but it's an iso of Sets

which may be easier to obtain

than iso of obj in \mathcal{C} .

$$\text{eg } \text{Hom}(A, X) \cong \text{hom}(A, E) \cong$$

\vdots

$$\text{hom}(A, E_n) \cong$$

$$\text{hom}(A, Y) \quad \cup$$

Today, more examples

- i) $\text{Mat}(\mathbb{R})$
- obj are $n \in \mathbb{N}$
 - $n \xrightarrow{\mathbb{R}} m \leftrightarrow n \times m$ matrices
with entries
in \mathbb{R} .

- What are repr functions?
- What are int's $f_m \Rightarrow \text{dm} \cong m \rightarrow m$
- $\{f: A \rightarrow F\} \cong F^A$ det by id_A .

representables @ $k \in \mathbb{N}$

$$h_k: \text{Mat}^{\text{op}} \rightarrow \text{Set}$$

$$n \mapsto \text{Mat}(n, k) = \left\{ \begin{array}{c} \xrightarrow{n} \\ \downarrow k \\ \xrightarrow{m} \end{array} \right\}$$

$$h_k(n \xrightarrow{A} m) = \text{Mat}(m, k) \longrightarrow \text{Mat}(n, k)$$

$$\begin{array}{ccc} \xrightarrow{n} & & \xrightarrow{m} \\ \mathcal{B} & \xrightarrow{k} & \mathcal{B} \\ m & \xrightarrow{\mathcal{B}} & m \end{array} \xrightarrow{\quad} \begin{array}{ccc} \xrightarrow{m} & & \xrightarrow{n} \\ \mathcal{B} & \xrightarrow{k} & \mathcal{B} \\ \xrightarrow{n} \mathcal{A} & \xrightarrow{m} & \xrightarrow{m} \mathcal{A} \end{array} = \mathcal{B} \cdot \mathcal{A}$$

nt

$$\alpha = h_k \rightarrow h_{j_1} \text{ has components:}$$

$$\alpha_n : \text{Mat}(n, k) \xrightarrow{\alpha_n} \text{Mat}(n, j_1)$$

$\left\lfloor \begin{matrix} n \\ k \end{matrix} \right\rfloor \quad \longmapsto \quad \left\lfloor \begin{matrix} n \\ j_1 \end{matrix} \right\rfloor$

a) $h_k \xrightarrow{\delta} h_{k-1}$ by "delete last row"



\mathfrak{B} $h_k(n) \xrightarrow{\delta_n} h_{k-1}(n)$ $\boxed{\delta_n \mathfrak{B}}$

$$h_k(A) \downarrow$$

$$h_k(m) \xrightarrow{\delta_m} h_{k-1}(m)$$

$\boxed{\delta_n \mathfrak{B}} \cdot A$

chop then mult

$\mathfrak{B} \cdot A$

$\boxed{\delta_m(\mathfrak{B} \cdot A)}$

$\equiv *$

mult then chop

* proof by "trust fasc"

b) $\gamma: h_k \rightarrow h_{k+1}$

$\gamma_n: \text{Mat}(n, k) \rightarrow \text{Mat}(n, k+1)$

not natural trans

$h_k \Rightarrow h_{k+1}$

How do we tell apart the natural from the unnatural? YL

YL says $h_k \rightarrow h_j$ natural ones
 arise by mult by some matrix
 (fixed)

YL says $\alpha: h_k \rightarrow h_j$ then

$$A \in \text{Mat}(n, k) : \quad \alpha_n(A) = A - \underbrace{\alpha_k(I_k)}$$

can be computed

$$\alpha_k : \text{Mat}(k, k) \longrightarrow \text{Mat}(k, j)$$

$$\begin{pmatrix} * & * & * \\ * & * & * \\ * & * & * \end{pmatrix} \mapsto \alpha_k(I_k)$$

e.g. δ = "delete last row"

$$\delta_n(A) = A - \underbrace{\delta_k(I_k)}_{\text{must be}} = A \cdot I_{k-1}$$

compute $\cong I_{k-1}$

e.g. τ = swap $2 \leftrightarrow 4$ row

$$\tau_n(A) = A \cdot \underbrace{\tau_k(I_k)}_{\text{L}} \rightarrow \begin{pmatrix} 1 & & & \\ & 2 & & \\ & & 1 & 2 \\ & & -1 & \end{pmatrix}$$

What α ? $\alpha(AB) = A \cdot \alpha B$ *

can only be $\alpha(A) = A \cdot H_A$

H_A fixed, depends on α .

then $\alpha(AB) = (A \cdot B)H_\alpha = A(BH_\alpha) = A(\alpha B)$
assoc!

YEA! These are the only α satisfying *.

$$[(\mathcal{E}/X)^{\text{op}}, \text{Set}]$$

Given $\vdash \frac{A}{X} \rightarrow \left[\mathcal{E}/X \left(\dashv, (u, A) \right) \right] (v, B) = \begin{cases} B & \xrightarrow{u} A \\ X & \downarrow u \end{cases}$

YL: The only way to yield

$\alpha = \vdash (A, u) \rightarrow \vdash (B, v)$ is from α

$$\begin{array}{ccc} A & & B \\ \downarrow u & \longrightarrow & \downarrow v \\ X & & X \end{array}$$

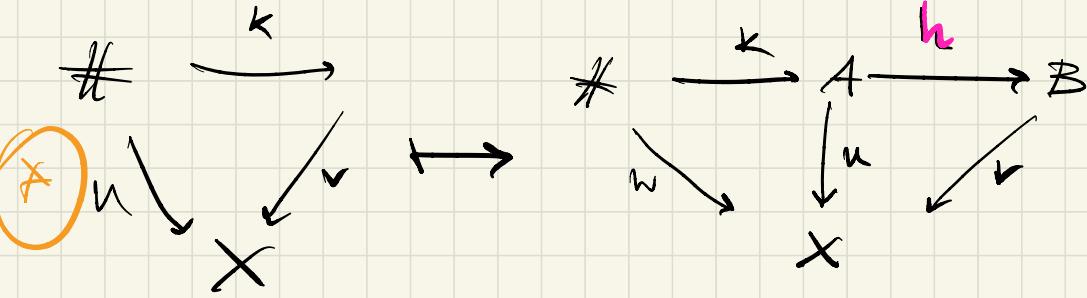
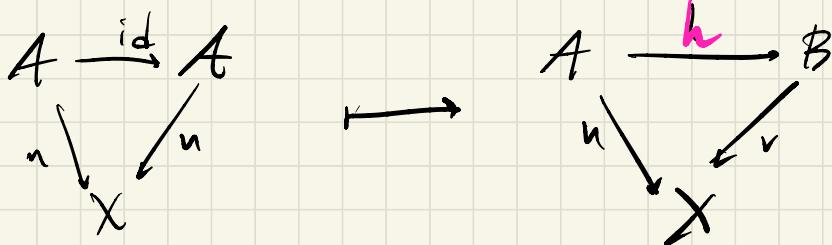
So, $\alpha_{C_W} : \mathcal{E}/X \left((C_W), (A, u) \right) \longrightarrow \mathcal{E}/X \left((C_W), (B, v) \right)$

where to send this?

$$\begin{array}{ccc} \cancel{*} & \xrightarrow{h} & A \\ \cancel{*} & \downarrow u & \nearrow v \\ ? & \xrightarrow{\quad} & \cancel{*} \end{array}$$

YL: Has to be comp by fixed el. Which el?

Look at $\text{id}_X \in \mathcal{E}/X \left((A, u), (A, u) \right)$ and take its image



iii) G group then $BG^{\text{op}} \xrightarrow{\quad} \text{Set}$

To specify an $r = r(\cdot) = A$

$$\forall g \in G \quad = r(g) : A \xrightarrow{\cong} A$$

$g \text{ iso in } BG \rightarrow r(g) \text{ iso in Set}$

satisfying functionality:

$$\left. \begin{array}{l} (r_g \circ r_h)(x) = r_{hg}(x) \\ r_{1_{\text{Id}_G}}(x) = x \end{array} \right\}$$

some data as

group hom

$$(G, \cdot) \longrightarrow (\text{Big } A_1, \circ)$$

So, functor $\widehat{BG^{\text{op}}} \rightarrow \text{Set} \equiv G\text{-Action}$
on A

$$\widehat{BG} \cong G\text{-Sets}$$

\downarrow

[$BG^{\text{op}}, \text{Set}$]

THE

"Representables" are nice G -actions
↑ only one.

$$BG(-, \cdot) : BG^{\text{op}} \rightarrow \text{Set}$$

$$\bullet \mapsto \text{Hom}(\bullet, \cdot) \cong G$$

G is acting on itself

$$G \longrightarrow \text{Big}(G) \quad \leftarrow \text{Grp Hom}$$

$$g \longmapsto (x \mapsto x \cdot g) = m_g$$

hom functor = comp. of morphisms
 = mult of elements

" BG "
 G

m_g is bijective with inverse $y \mapsto y \cdot g^{-1}$.

YL: $\phi: G \longrightarrow \text{Big } G$ is bijective
 $g \longmapsto m_g$

h inj. assume $h \circ g = h \circ h \rightarrow m_g = m_h$

$$\Leftrightarrow \forall x \quad m_g(x) = m_h(x)$$

$$\Leftrightarrow x \cdot g = h \cdot x$$

$$\xrightarrow{x=e} e \cdot g = e \cdot h$$

$$\Rightarrow g = h \quad \checkmark$$

Cayley Thm. $G \leq (\text{Big } G \xrightarrow{\text{Set}}, \circ)$