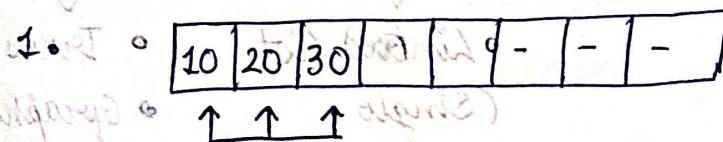


DATA STRUCTURES

- Data structure is a way of organising data in memory so that it can be used efficiently.
- It consists of two parts:
 1. How data is organised in memory.
 2. What all operations can be performed on that data structure.

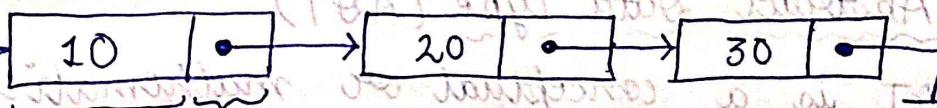


Elements placed one after other

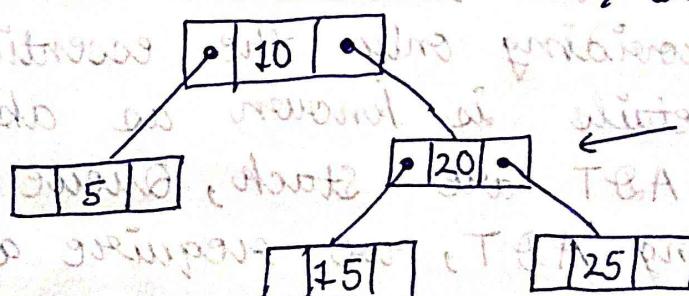
Array (contiguous block of memory)

Advantage: Direct access is available

Linked List (no contiguous block of m/m)



Binary Search Tree (Non-linear D.S.)



children → to insert : start → left → right → end

Classification of ~~data~~ structures

Data structure

Primitive

- int
- float
- character
- boolean

Non-primitive

Simple

- array
- record

Compound

Linear

• Linked List

(Single, Doubly, Circular)

• Trees

• Graphs

Non-linear

Abstract Data Type (ADT)

ADT is a conceptual or mathematical model which specifies different type of operations on that data type. However, it does not specify how these operations will be implemented. It is called abstract because it only tells the operations which ADT can perform without telling how these will be implemented. This process of providing only the essentials while hiding the details is known as abstraction.

Examples of ADT are stack, queue and list.

(For implementing ADT, we require a suitable DS)

In stack, we can perform following operations:

1. Push: To add an element at the top of stack.
2. Pop: To remove an element from top of stack
3. Display: To print the content of stack from top to bottom.

4. Peek : To print the top most element.

These operations are specified in stack ADT which works on LIFO basis. To implement stack, the programmer will be using a suitable data structure such as an array or linked list.

* Implementing Stack using Array

Stack is an ADT which works on LIFO basis. Elements placed at the last is always at the top and the element placed first is at the bottom. Stack is used in no. of applications such as:

1. Converting infix expression to postfix
2. Evaluating the postfix expression
3. Reversing a string
4. Checking the parenthesis
5. Back tracking
6. Storing the return address in function call
7. Memory allocation
8. Compiler design

In stack, we can perform operations like Push, Pop, Display and Peek.

Algorithm to push an element

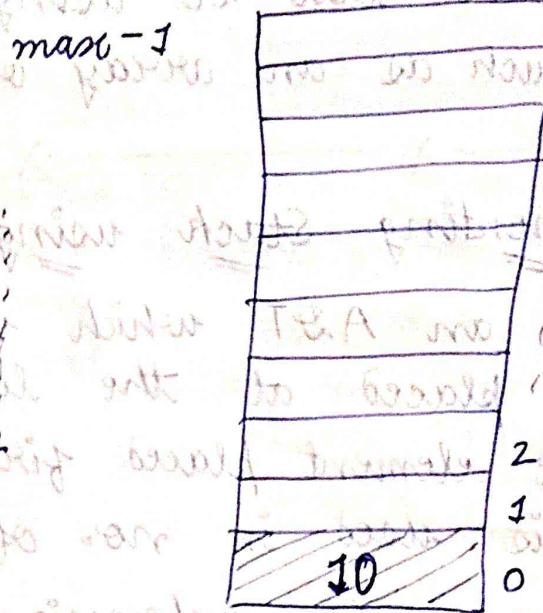
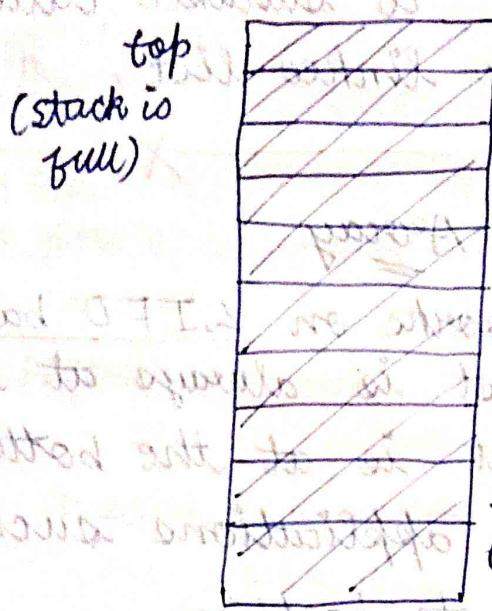
Initially, $\text{top} = -1$ (stack is empty)

- Step 1 : Start
- Step 2 : If $\text{top} = \text{max} - 1$, then display "Stack is full", goto step 6
- Step 3 : Input the element to be pushed

Step 4 : Increment top by 1. ~~int top = -1;~~ ~~int top = 0;~~

Step 5 : Assign element to A[top]

Step 6 : Stop



max-1

top

(on adding
first elem
- end)

Algorithm to Pop an element

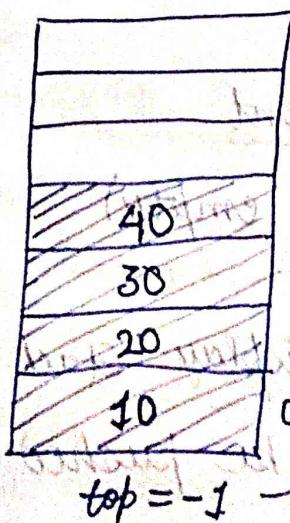
Step 1 : Start

Step 2 : If $\text{top} = -1$, then display "Stack is empty", goto step 5

Step 3 : Display the contents at A[top]

Step 4 : Decrement top by 1

Step 5 : Stop



max-1

top = -1

3 (top+1)
2 top
1 (after popping)

{ After popping,
40 is still the
part of array but
not the part of
stack. 3
0 to top : 8 part

Algorithm to display elements of stack from top to bottom

Step 1: Start

Step 2: If $\text{top} = -1$, then display "Stack is empty", goto step 8

Step 3: Take an integer variable temp.

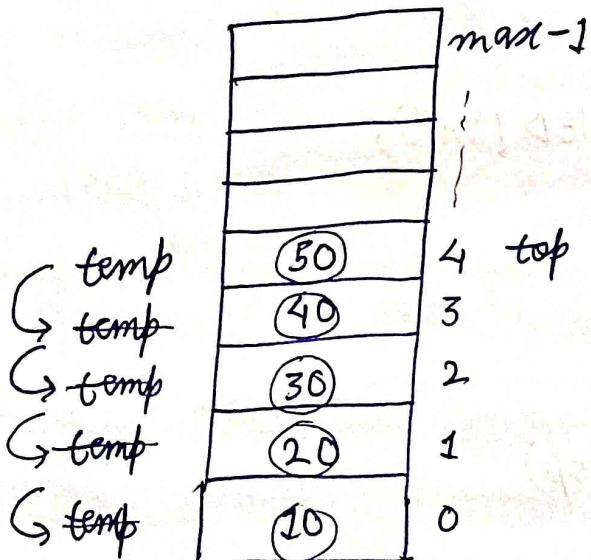
Step 4: Assign top to temp

Step 5: Repeat step 6 and 7 until $\text{temp} = -1$

Step 6: Display contents of $A[\text{temp}]$

Step 7: Decrement temp by 1

Step 8: Stop



$$\text{temp} = -1$$

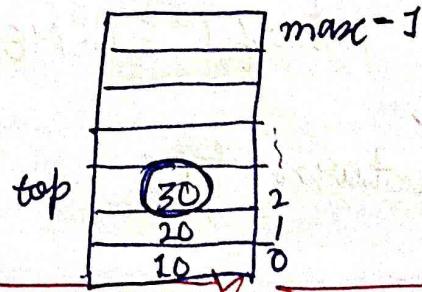
Algorithm to peek

Step 1: Start

Step 2: If $\text{top} = -1$, display "Stack is empty", goto step 4

Step 3: Display $A[\text{top}]$

Step 4: Stop



Menu driven program to implement stack using array

```
# include < stdio.h >
```

```
# define MAX 10 // Maximum size of Stack assigned as 10
```

```
// Push function
```

```
int push (int *stack, int top) {
```

```
    if (top == MAX - 1)
```

```
        printf ("In STACK IS FULL!");
```

```
    else
```

```
        int elem;
```

```
        printf ("Enter the element to be pushed : ");
```

```
        scanf ("%d", &elem);
```

```
        top += 1;
```

```
        stack [top] = elem;
```

```
        printf ("ELEMENT PUSHED!\\n");
```

```
    }
```

```
    return top;
```

```
}
```

```
// Pop function
```

```
int Pop (int *stack, int top)
```

```
    if (top == -1)
```

```
        printf ("In STACK IS EMPTY!");
```

```
    }
```

```
    else
```

```
        printf ("The element at top : %d\\n", stack [top]);
```

```
        top -= 1; // In short "pushes", l = top; if : 5 date
```

```
        printf ("ELEMENT POPPED!\\n");
```

```
    }
```

```
    return top;
```

```
}
```

```
// Display function
void display (int *stack, int top)
{
    if (top == -1)
        printf ("In STACK IS EMPTY!");
    else
    {
        int temp = top;
        printf ("Elements in stack from top to bottom : ");
        while (temp != -1)
        {
            printf ("%d ", stack [temp]);
            temp -= 1;
        }
        printf ("\n");
    }
}
```

// Peek function

```
void Peek (int *stack, int top)
{
    if (top == -1)
        printf ("In STACK IS EMPTY!");
    else
    {
        printf ("Element at the top in the stack is %d\n",
               stack [top]);
    }
}
```

int main()

```
{ int stack [MAX], top = -1, ch;
do
{ printf ("In MENU In 1. Push In 2. Pop In 3. Display In
        4. Peek ");
    printf ("Enter your choice : ");
    scanf ("%d", &ch);
}
```

```

switch(ch)
{
    case 1: top = Push(stack, top);
               break;
    case 2: top = Pop(stack, top);
               break;
    case 3: Display(stack, top);
               break;
    case 4: Peek(stack, top);
               break;
    default: ch = 0; // Exit condition
               break;
}

```

```

y
while(ch);
return 0;

```

y

25/8/2020

* Implementing Queue using Array

Queue is an ADT which works on FIFO (First In First Out). It can be implemented using an array or a linked list data structure. Here, we are implementing queue ADT using an array.

Applications of Queue

1. Processor scheduling (single processor, more than one processes)
2. Device scheduling (single O/P device, more than one user)
3. In implementing BFS (Breadth First Search) for graph traversal algorithm
4. In service provider systems where there are multiple inquiries and less number of handlers.

5. In flight landing and take off systems, as more than one aeroplane compete for single runway at a given time for landing and take off.

Queue operations

1. Enqueue : To add an element at the rear of the queue.
2. Dequeue : To remove an element from front of the queue.
3. Display : To print all elements in queue from front to rear.
4. Peek : To print front element of the queue.

Implementing Queue using an Array

- Assumptions → Static memory allocation
We have an integer array 'a' with size 'max', where 'max' is a constant. Take two integer variables 'front' & 'rear' and initialise them to -1 as initially queue is empty.

• Algorithm to implement Enqueue

Step 1 : Start

Step 2 : If $((\text{rear} + 1) \% \text{max}) = \text{front}$, then display a message, "Queue is full", goto step 7.

Step 3 : Input the element to be inserted.

Step 4 : $\text{rear} = (\text{rear} + 1) \% \text{max}$

Step 5 : Assign the element to $a[\text{rear}]$

Step 6 : If $\text{front} = -1$, it means element inserted is first element, then increment front to 0.

Step 7 : Stop

max = 10

0	1	2	3	4	5	6	7	8	9
10	20	30	40	50	60	70	80	90	100

front

rear

(queue is full)

0	1	2	3	4	5	6	7	8	9
10	20	30	40	50	60	70	80	90	100

rear rear rear front

rear

(queue is full)

If $(\text{rear} + 1) \% \text{max} = \text{front}$,

• For (i), $(9 + 1) \% 10 = 0$

$$10 \% 10 = 0 = \text{RHS}$$

• For (ii), $(2 + 1) \% 10 = 3$

$$3 \% 10 = 3 = \text{RHS}$$

Here, we have used ' $\% \text{max}$ ' in formula so that value of ' $\text{rear} + 1$ ' doesn't exceed ' max '.

e.g. If $\text{rear} = 9$,

$$\text{so, } \text{rear} + 1 = 10 \text{ & } 10 \% 10(\text{max}) = 0$$

So, now, the element will be inserted at index '0'.

• Algorithm to implement Dequeue

Step 1: Start

Step 2: If $\text{front} = -1$, then display a message, "Queue is empty", goto step 6.

(only 1 elem present)

Step 3: If $\text{front} = \text{rear}$, then display the value of $a[\text{front}]$ & reset front & rear to -1 , goto step 6.

Step 4: Display $a[\text{front}]$

Step 5: $\text{front} = (\text{front} + 1) \% \text{max}$

Step 6: Stop

	0	1	2	3	4	5	6	7	8	9	max = 10
a	110	120	130						100		

front

rear

front

{Also called circular queue}

• Algorithm to implement display

Step 1 : Start

Step 2 : If front = -1, then display a message that queue is empty, goto step 8.

Step 3 : Take an integer 'temp' & assign value of 'front' to it.

Step 4 : while (temp != rear), repeat step 5 & 6.

Step 5 : Display a[temp]

Step 6 : temp = (temp + 1) % max

Step 7 : Display a[temp]

Step 8 : Stop

	0	1	2	3	4	5	6	7	8	9	max = 10
	110	120	130			60	70	80	90	100	

temp temp rear

front

temp temp temp temp temp

• Algorithm to implement peek

Step 1 : Start

Step 2 : If front = -1, then display a message that queue is empty, goto step 4.

Step 3 : Display the value of a[front]

Step 4 : Stop

X

X

X

Menu driven program to implement Queue (Circular)
using array

```
# include <stdio.h>
#define MAX 10 // Maximum size of Queue assigned as 10
// Insert function
void Insert(int queue[], int *front, int *rear)
{
    if ((*rear + 1) % MAX == *front)
    {
        printf("QUEUE IS FULL!");
        return;
    }
    else
    {
        int elem;
        printf("Enter the element to be inserted : ");
        scanf("%d", &elem);
        *rear = (*rear + 1) % MAX;
        queue[*rear] = elem;
        if (*front == -1)
        {
            *front += 1;
            printf("ELEMENT INSERTED!");
        }
    }
}

// Serve function
void Serve(int queue[], int *front, int *rear)
{
    if (*front == -1)
    {
        printf("QUEUE IS EMPTY!");
        return;
    }
    else if (*front == *rear)
    {
        printf("Element at front : %d\n", queue[*front]);
        *front = -1;
        *rear = -1;
    }
}
```

printf ("ELEMENT SERVED ! \n");

}

else

{ printf ("Element at front = %d \n", queue[front]);
*front = (*front + 1) % MAX;
printf ("ELEMENT SERVED ! \n");

}

y

// display function

void display (int queue[], int front, ^{int} rear).

{ if (front == -1)

{ printf ("QUEUE IS EMPTY ! \n");

}

else

{ int temp = front;

printf ("Elements in queue from front to
rear : ");

while (temp != rear)

{ printf ("%d ", queue[temp]);

temp = (temp + 1) % MAX;

printf ("%d \n", queue[temp]);

y

y

int main()

{ int queue[MAX], front=-1, rear=-1, ch;

do

{ printf ("1. MENU 2. Insert 3. Serve In
3. Display 4. Exit ");

printf ("Enter your choice : ");

```

scmf("100d", &ch);
switch(ch)
{
    case 1: Insert(queue, &front, &rear);
               break;
    case 2: Some(queue, &front, &rear);
               break;
    case 3: Display(queue, front, rear);
               break;
    case 4: ch=0;
               break;
    default: ch=0; // Exit-condition
               break;
}

```

y

```

while(ch);
return 0;
}

```

}



2/9/2020

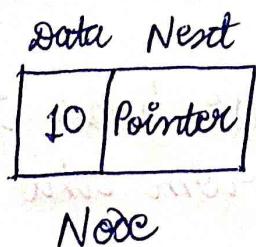
* linked list

It is a linear data structure. In linked list, nodes are not stored at a contiguous location; nodes are linked using pointers.

Types of linked list

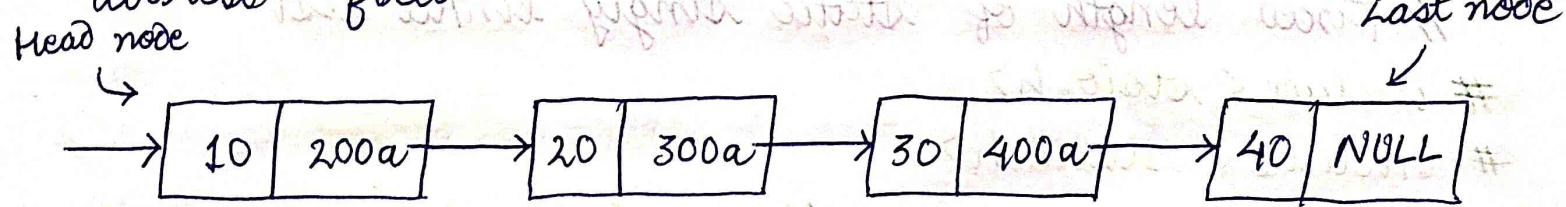
1. Singly linked list
2. Doubly linked list
3. Circular linked list

1. Singly linked list: A node in singly linked list has two fields, i.e., data field & address field (Next). The address field contains the address of next node (using a pointer).



```
struct node
{ int data;
  struct node *Next;
} // Self referential structure
```

A data part stores the data & 'Next' part stores the address of the next node. The last node in the linked list is represented by a NULL in the address field.



It is unidirectional, i.e. we cannot backtrack.

We can access nodes sequentially.

Why Linked List?

Arrays can be used to store linear data of similar types, but arrays have the following limitations:

- The size of the arrays is fixed, so we must know the upper limit of the no. of elements in advance.
- Inserting a new element in an array of elements is expensive in terms of shifting.
- For better memory utilisation (in terms of fragmented memory), linked list is required.

Advantages over arrays

- Dynamic size
- Ease of insertion/deletion
- Heterogeneous data types can be handled.

Drawbacks

- Random access is not allowed. We have to access elements sequentially starting from the first node.
- Extra memory space for a pointer is required with each element of the list.

Example

// Fixed length of static singly linked list

```
#include < stdio.h >
```

```
#include < stdlib.h >
```

typedef struct node // structure of node

```
{ int data;
  struct node *Next; }
```

node type ; Instances

```
int main()
```

```
{ node type A, B, C, D, *ptr = NULL;
```

A.data = 10, B.data = 20, C.data = 30, D.data = 40.
// values are stored in data field of each node

A.Next = &B; // connecting nodes A & B.

B.Next = &C; // connecting nodes B & C

C.Next = &D; // connecting nodes C & D

D.Next = NULL; // Assigning NULL in Next field of last node (D)

ptr = &A; // Assigning address of first node to pointer 'ptr' // ptr = 20a

→ Just continue

```

while (ptr != NULL)
{
    printf ("%d -> ", ptr->data);
    ptr = ptr->Next;
}
return 0;

```

- ✓ 'typedef' has been used to provide alias name for the structure. So, we can write 'nodetype' instead of 'struct node' now.

⇒ Applications of linked list in Computer Science

- Implementation of stack & queue
- Implementation of graphs : Adjacency list representation of graphs.
- Dynamic memory allocation : We use linked list of free blocks
- Maintaining directory of names
- Performing arithmetic operations on long integers
- Manipulation of polynomials by storing constants in the node of linked list representing sparse matrices.

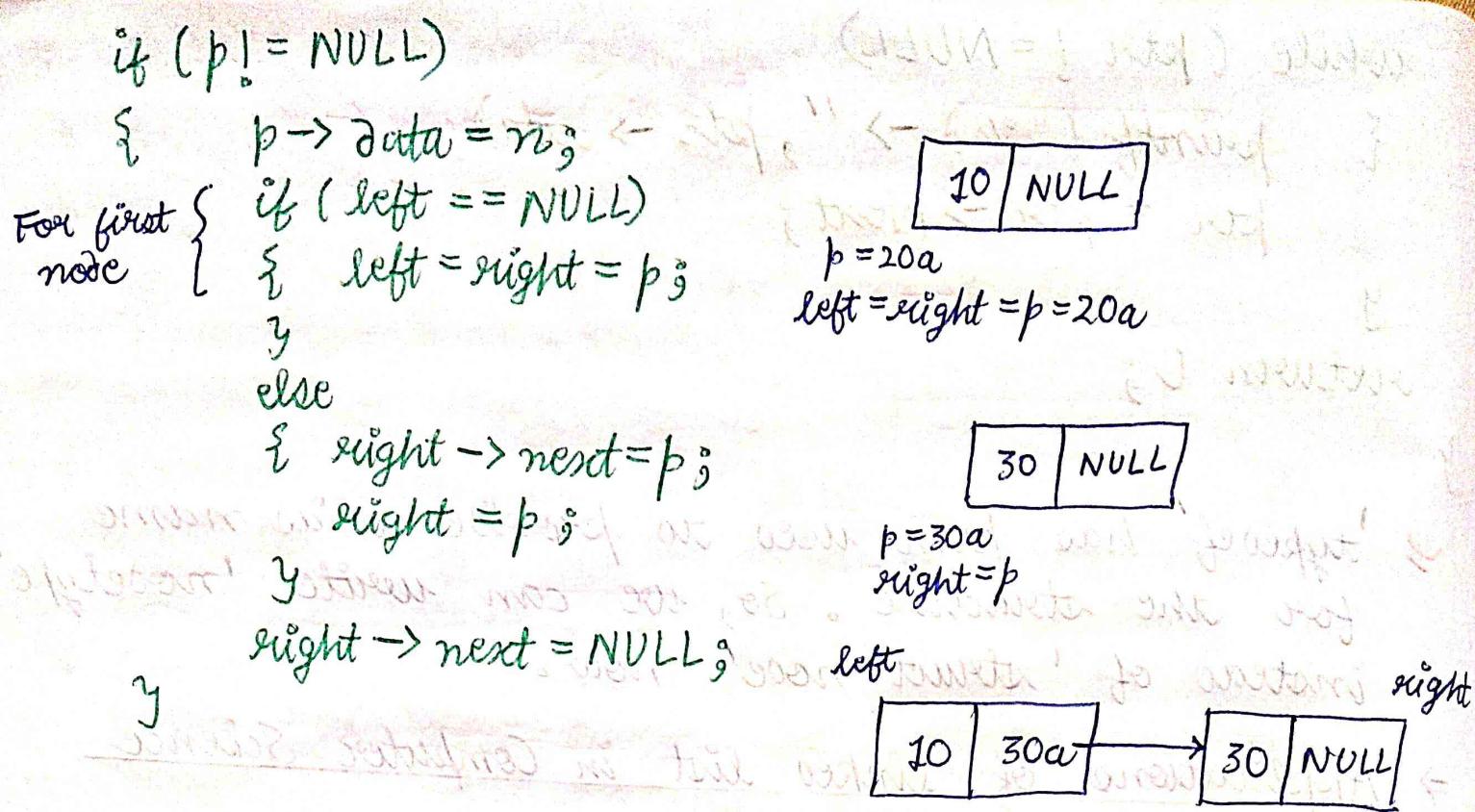
3/9/2020

Singly linked list : Inserting node at right hand side

```

typedef struct node
{
    int data;
    struct node *next;
} nodetype;
nodetype *left = NULL, *right = NULL, *p = NULL;
p = (nodetype *) malloc (sizeof (nodetype));

```



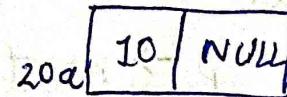
// C program to create a singly linked list by inserting node at right hand side

```
#include < stdio.h >
#include < stdlib.h >
typedef struct node
{
    int data;
    struct node *next;
} nodetype;
int main()
{
    nodetype *left = NULL, *right = NULL, *p = NULL;
    int n, ch;
    do
    {
        printf("Enter value to insert in node:");
        scanf("%d", &n);
        p = (nodetype *)malloc(sizeof(nodetype));
        if (p != NULL)
        {
            p->data = n;
            if (left == NULL)
            {
                left = right = p;
            }
            else
            {
                right->next = p;
                right = p;
            }
        }
    } while (ch != 'n');
}
```

```

else
{
    right -> next = p;
    right = p;
}

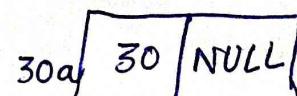
```



```

right -> next = NULL;

```



y

printf("In Do you want to add more node ? In
Press 1 to continue otherwise 0");

```
scanf("%d", &ch);
```

```
} while (ch == 1);
```

printf("In display the list : ");

```
while (left != NULL)
```

```
{ printf("%d", left->data);
```

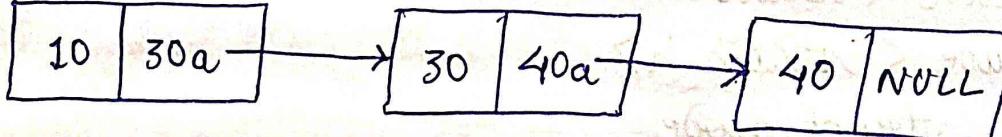
```
left = left->next;
```

y

return 0;

y

left



right

Singly Linked List : Inserting node at left hand side

```
typedef struct node
```

```
{ int data;
```

```
struct node *next;
```

y nodetype

```
nodetype *left = NULL, *p = NULL;
```

```
p = (nodetype *) malloc(sizeof(nodetype));
```

```
if (p != NULL)
```

```
{ p->data = n;
```

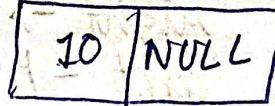
```
if (left == NULL)
```

```
{ left = p;
```

```
y left->next = NULL;
```

else

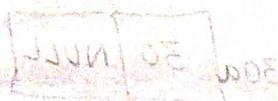
{ $p \rightarrow \text{next} = \text{left};$
 $\text{left} = p;$



$p = 20a$

$\text{left} = p$

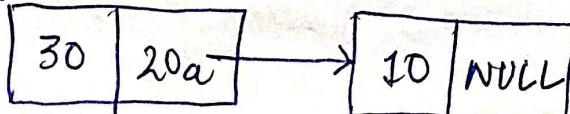
y
y



$b = 30a$

$\text{left} = p$

left



First node
formed

Pointed in
this direction

// C program to create a singly linked list by inserting node at left hand side

```
# include < stdio.h >
```

```
# include < stdlib.h >
```

```
typedef struct node
```

```
{ int data;  
    struct node *next;
```

```
y nodetype;
```

```
int main()
```

```
{ nodetype *left=NULL, *p=NULL;
```

```
int n, ch;
```

```
do
```

```
{ printf("Enter value to insert in node : ");
```

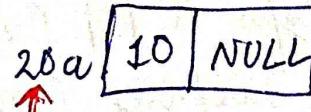
```
scanf("%d", &n);
```

```
p = (nodetype *) malloc(sizeof(nodetype));
```

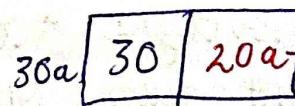
```
if (p != NULL)
```

```
{ p->data = n;
```

```
if (left == NULL)
{   left = p;
    left->next = NULL;
}
```



```
y
else
{   p->next = left;
    left = p;
}
y
```



```
y
printf("In do you want to add more node?
In Press 1 to continue otherwise 0");

```

```
scanf("%d", &ch);
```

```
y while (ch == 1);

```

```
printf("In display the list:");

```

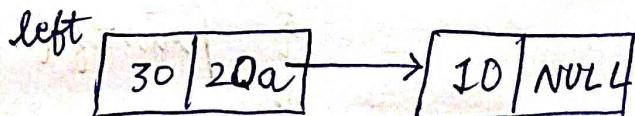
```
while (left != NULL)

```

```
{   printf("%d ", left->data);
    left = left->next;
}
y
```

```
return 0;

```



```
y
```

4/9/2020

* Implementing Stack ADT using Linked list & Single pointer

```
#include <stdio.h>
#include <stdlib.h>
typedef struct node
{   int info;
    struct node *next;
} nodetype;
```

```
nodetype * push(nodetype *, int);
nodetype * pop(nodetype *);
void display(nodetype *);

int main()
{
    nodetype *top = NULL;
    int num, opt;
    char choice;
    do
    {
        printf("Enter opt \n");
        scanf("%d", &opt);
        if (opt == 1)
        {
            scanf("%d", &num);
            top = push(top, num);
        }
        else if (opt == 2)
        {
            if (top != NULL)
            {
                top = pop(top);
            }
        }
        else
        {
            printf("STACK IS EMPTY \n");
        }
        if (opt == 3)
        {
            if (top != NULL)
            {
                display(top);
            }
        }
        else
        {
            printf("STACK IS EMPTY \n");
        }
    } while (choice == 'y');
    printf("INVALID OPTION! \n");
}
```

```
printf ("Do you want to continue? (Y/N) \n");
```

```
scanf ("%c", &choice);
```

```
y while (toupper (choice) == 'Y');
```

```
return 0;
```

```
}
```

```
nodetype * push (nodetype * tp, int nm);
```

```
{ nodetype * p;
```

```
p = (nodetype *) malloc (sizeof (nodetype));
```

```
if (p == NULL)
```

```
{ printf ("NOT ENOUGH MEMORY! \n");
```

```
y
```

```
else
```

```
{ p->info = nm;
```

```
p->next = tp;
```

```
tp = p;
```

```
y
```

```
return tp;
```

```
y
```

```
nodetype * pop (nodetype * tp)
```

```
{ nodetype * temp;
```

```
temp = tp;
```

```
printf ("Element popped is %d ", tp->info);
```

```
tp = tp->next;
```

```
free (temp);
```

```
return tp;
```

```
y
```

```
void display (nodetype * tp)
```

```
{ while (tp != NULL)
```

```
{ printf ("%d ", tp->info);
```

```
tp = tp->next;
```

```
y
```

```
x
```

```
x
```

```
x
```

5/9/2020

* Implementing Queue ADT using Linked List & Single pointer

```
# include< stdio.h>
# include< stdlib.h>
typedef struct node
{
    int data;
    struct node *next;
} nodetype;
nodetype * insert (nodetype * );
nodetype * delete (nodetype * );
void display (nodetype * );
int main()
{
    int ch;
    nodetype *front = NULL, *rear = NULL;
    do
    {
        printf ("1. Insert 2. Delete 3. Display\n
        4. Exit \nEnter your choice : ");
        scanf ("%d", &ch);
        switch(ch)
        {
            case 1: rear = insert (rear);
                if (front == NULL)
                {
                    front = rear; = y
                    break;
                }
            case 2: front = delete (front);
                if (front == NULL)
                {
                    rear = NULL; = y
                    break;
                }
            case 3: display (front);
                break;
        }
    } while (ch != 4);
}
```

front
NULL

rear
NULL

y while (ch <= 3);

return 0;

}

nodetype * insert (nodetype * rear)

{ int n;

nodetype * p = NULL;

p = (nodetype *) malloc(sizeof(nodetype));

if (p != NULL)

{ printf("In Enter value to insert : ");

scanf("%d", &n);

p->data = n;

if (rear == NULL)

{ rear = p;

y

else

{ rear->next = p;

rear = p;

y

rear->next = NULL;

}

return rear;

}

nodetype * delete (nodetype * front)

{ nodetype * p = NULL;

if (front == NULL)

{ printf("Nothing to delete!");

y

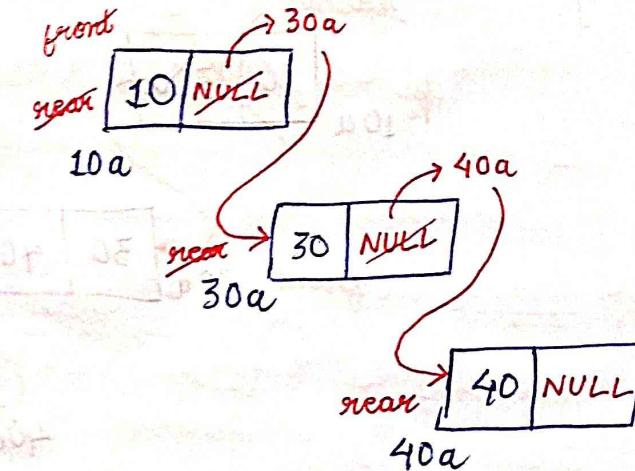
else

{ printf("In Deleted node's data is ->%d", front->data);

p = front;

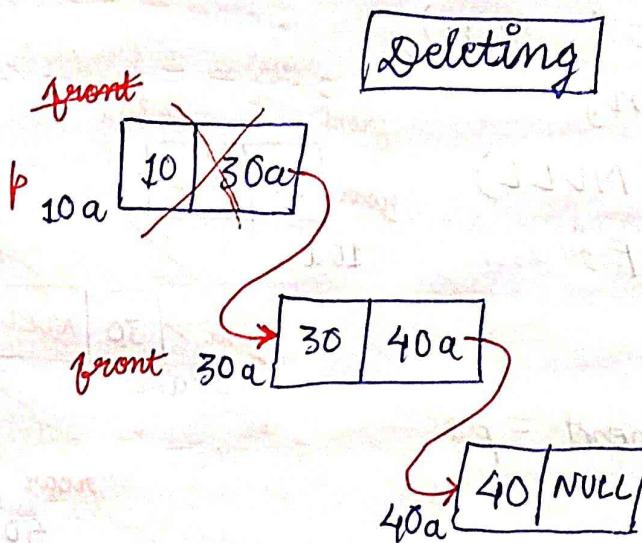
front = front->next;

y free(p);



return front;

```
void display(nodetype *front)
{
    if (front == NULL)
    {
        printf("Nothing to display!");
    }
    else
    {
        while (front != NULL)
        {
            printf("%d ", front->data);
            front = front->next;
        }
    }
}
```



5/9/2020

* Implementing stack ADT using linked list & double pointer

```
#include <stdio.h>
#include <stdlib.h>
typedef struct node
{
    int data;
    struct node *next;
}
```

Now, instead of passing value of 'top', we're passing address of 'top'. ('top' itself is an address of type 'node').

```
void push(nodetype **, int);
void pop(nodetype **);
```

```
int main()
{
    nodetype *top = NULL;
    int num, opt;
    char choice;
    do
    {
        printf("Enter option: \n");
        scanf("%d", &opt);
        if (opt == 1)
        {
            scanf("%d", &num);
            push(&top, num); // call by reference
        }
        else if (opt == 2)
        {
            if (top != NULL)
            {
                pop(&top);
            }
            else
            {
                printf("STACK IS EMPTY \n");
            }
        }
        else if (opt == 3)
        {
            if (top != NULL)
            {
                display(top);
            }
            else
            {
                printf("STACK IS EMPTY \n");
            }
        }
        else
        {
            printf("INVALID OPTION! \n");
        }
        printf("Do you want to continue? (Y/N) \n");
        scanf("%c", &choice);
    } while (choice == 'Y' || choice == 'y');
}
```

y while(toupper(choice) == 'Y');

return 0;

y

```
void push( nodetype **tp, int nm )
{
    nodetype *p;
    p = (nodetype *)malloc(sizeof(nodetype *));
    if (p == NULL)
    {
        printf("NOT ENOUGH MEMORY IN");
        tp = &top;
        *tp = NULL;
    }
    else
    {
        p->data = nm;
        p->next = *tp;
        *tp = p;
    }
}
```

y

void pop(nodetype **tp)

```
{
    nodetype *temp;
    temp = *tp;
    printf("Element popped is %d", *tp->data);
    *tp = *tp->next;
    free(temp);
}
```

y

void display(nodetype *tp)

```
{
    while (tp != NULL)
    {
        printf("%d ", tp->data);
        tp = tp->next;
    }
}
```

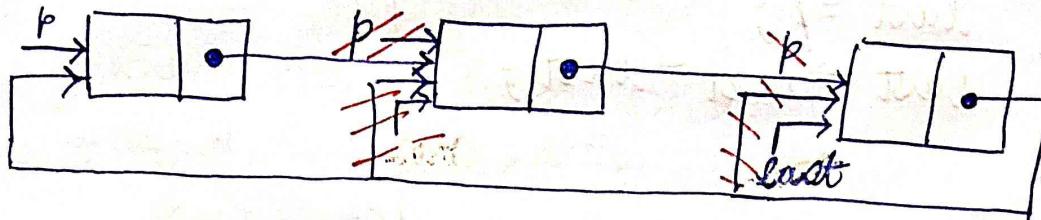
y

X ————— X ————— X —————

6/9/2020

* Circular Linked List

last



$$p \rightarrow \text{next} = \text{last} \rightarrow \text{next}$$

$$\text{last} \rightarrow \text{next} = p$$

Assumptions

Using structure, define a datatype 'nodetype' having first field 'info' of type 'int' and second field 'next' which is an address of type 'nodetype'. Take a 'nodetype' pointer 'last' & initialise it to NULL as initially, linked list is empty.

Algorithm to add an element towards start of circular linked list

Step 1 : Start

Step 2 : Using malloc() function, allocate memory for a 'nodetype' & assign its address to a 'nodetype' pointer 'p'.

Step 3 : if (p == NULL) then print that not enough memory & goto step 7.

Step 4 : Input 'num' & assign it to 'info' field of 'p'.

Step 5 : if (last == NULL) then last = p; last → next = p; goto step 7.

Step 6 : p → next = last → next;

last → next = p;

Step 7 : Stop

7/9/2020

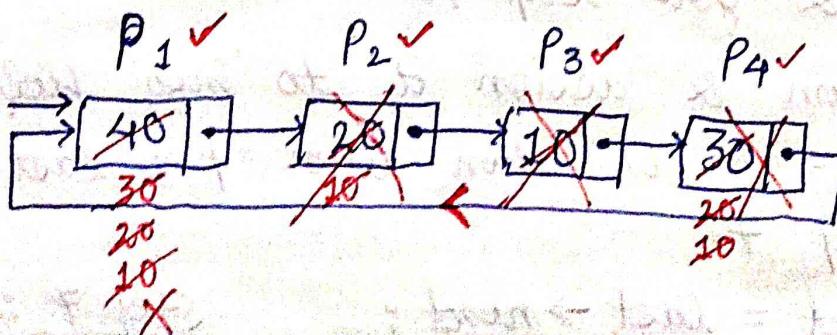
* Circular linked list - Implementation & Application

```
if (last == NULL)
{
    last = p;
    last->next = last;
}
else
{
    p->next = last->next;
    last->next = p;
    last = p;
}
```

Multi-programming (Application)

In multi-programming, there are more than one processes at the state of execution while only one process is getting executed by CPU because there's a single CPU working on more than one processes. So, only one process is getting executed at any given time but there are more than one processes at the state of execution. So, CPU gives fixed time slots to every process.

Assume, there are 4 processes & every process is given a fixed time slot of 10 ns by CPU.



\rightarrow
p
(pointer
to access
circular
linked
list)

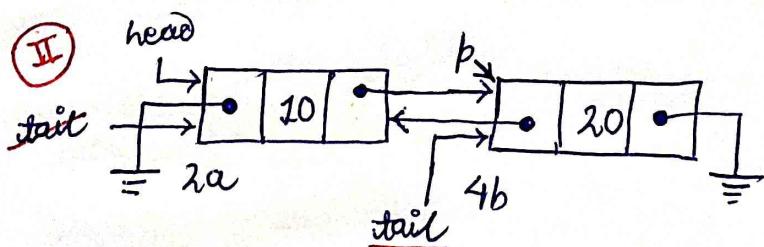
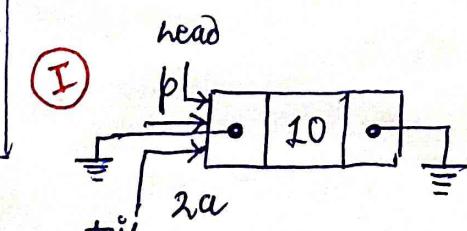
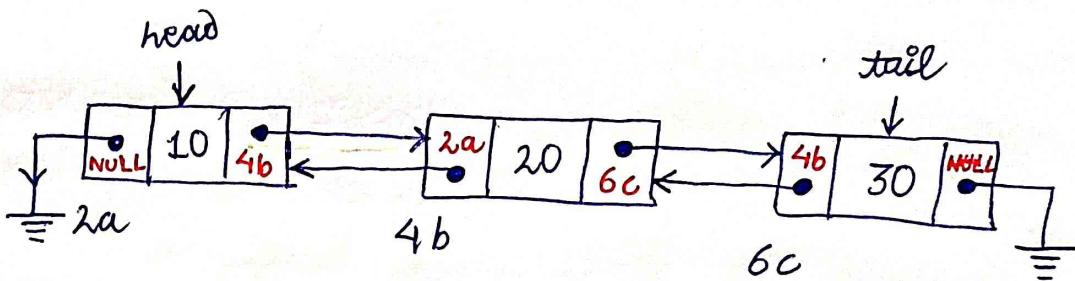
Total time req. = 100 ns

- After 30 ns, P3 is over.
- After 60 ns, P2 is over.
- After 90 ns, P4 is over.
- After 100 ns, P1 is over.

The shortest process got completed first (P_3), i.e. waiting time for shortest process was minimum, then P_2 , P_4 and P_1 get completed sequentially. Thus, process requiring maximum time was completed at last (P_1). Hence, the waiting time has been distributed.

8/9/2020

* Doubly Linked List



(New node is added after the last node)

```

if (tail == NULL)
{
    head = tail = p;
    p->prev = NULL;
}
else
{
    tail->next = p;
    p->prev = tail;
    tail = p;
}
p->next = NULL;

```