Predication

Introduction

- 1. Some branches are difficult to predict even with sophisticated branch predictors
 - Compiler can help avoid some of these branches

Predication

- 1. Branch prediction: Guessing where program is going
 - No penalty if correct (0 instructions)
 - Huge penalty if incorrect (12-stage pipeline with 4 instructions/cycle -> 48)
- 2. Predication: Do work of both directions (taken/not taken)
 - Waste up to 50%
 - Throw away work from wrong path
- 3. Loop: Typically use branch prediction (more loops -> more predictable)
 - Using predication will waste many cycles and makes little sense
- 4. Function call/return: Predict since the branch is always taken
- 5. Large if-then-else: Predict
 - If length of branch is 100 cycles either way, we will waste these 100 cycles compared to the \sim 50 we wasted due to a branch miss in the above example
- 6. Small if-then-else: Predication unless we're certain the branch predictor will be highly accurate
 - If length of branch is 5 cycles either way, we always waste 5 cycles with predication
 - Branch prediction results in 50 wasted cycles in the case of a missed prediction
 - 100% of time we waste 5 cycles with predication -> 5
 - 10% of time we waste 50 cycles with branch prediction -> 5
 - This is the break-even point
 - If size of branch is smaller or branch prediction performs worse, this biases the result in favor of predication

If Conversion

- 1. if(cond) { $x = arr[i]; y++; } else { <math>x = arr[j]; y-; }$
 - x1 = arr[i];
 - x2 = arr[j];
 - y1 = y + 1;
 - y2 = y 1;
 - x = cond ? x1 : x2; -> Now we have two branches
 - y = cond? y1 : y2; -> Now we have two branches
- 2. Need some sort of conditional move instruction in hardware
 - MOV x, x1, cond

Conditional Move

- 1. MIPS
 - MOVZ Rd, Rs, Rt \rightarrow if(Rt == 0) Rd = Rs;
 - MOVN Rd, Rs, Rt \rightarrow if(Rt != 0) Rd = Rs;
- 2. x86
 - CMOVZ, CMOVNZ, CMOVGT, ... -> if(cc) Dst = Src;
 - Based on all of the condition codes

MovZ MovN Quiz

- 1. Consider the following program:
 - BEQZ R1, Else

- ADDI R2, R2, 1
- B End
- Else: ADDI R3, R3, 1
- End:
- 2. How do we rewrite this with conditional moves?
 - ADDI R4, R2, 1
 - ADDI R5, R3, 1
 - MOVN R2, R4, R1
 - MOVZ R3, R5, R1

MovZ MovN Performance

- 1. 80% branch accuracy, 40-instruction penalty on missed branch
 - When the branch is not taken we execute 3 instructions
 - When the branch is taken we execute 2 instructions
 - Assuming the branch is balanced, we execute an average of 2.5 instructions
 - 2.5 + 0.2 * 40 = 10.5 instructions
- 2. If-conversion costs 4 instructions

MovZ MovN Performance Quiz

- 1. Assumptions:
 - Branch is taken 50% of the time
 - 30-instruction misprediction penalty
- 2. If-conversion is better when prediction accuracy is below...
 - 2.5 + Accuracy * 30 = 4
 - Accuracy = (4 2.5) / 30 = 0.05 (5%)
 - If prediction accuracy is below 95%, predication is better

MOVc Summary

- 1. Require compiler support for if-conversion
 - Not fully backwards-compatible (old code won't use these instructions)
- 2. Removes hard-to-predict branches
- 3. More registers needed
 - Keep results from both parts
- 4. More instructions executed
 - Executing both paths
 - MOVc to select actual results
- 5. Full predication: Can change ISA such that all instructions are conditional
 - Requires extensive support in instruction set
 - Every instruction computes a result but may or may not store it

Full Predication Hardware Support

- 1. MOVcc
 - Separate opcode
- 2. Full Predication
 - Add condition bits to every instruction
 - Intel Itanium: (qp) ADD R1, R2, R3
 - qp: Qualifying predicate
 - Instruction has 41 bits, low 6 bits are 1-bit condition registers

Full Predication Example

- 1. Applying full predication to the example program above
 - MP.EQZ P1, P2, R1 (sets qualifying predicate bits)
 - (p2) ADDI R2, R2, 1
 - (p1) ADDI R3, R3, 1
- 2. Don't require additional registers
- 3. No overhead with moving between registers since we're storing directly in the result

Full Predication Quiz

- 1. Consider the following program:
 - BEQZ R2, ELSE
 - ADD R1, R1, 1
 - B DONE
 - ELSE: ADD R1, R1, -1
 - DONE:
- 2. Convert to full predication version
 - MP.EQZ P1, P2, R2 (P1 -> R2 == 0, P2 -> R2 != 0)
 - (P2) ADD R1, R1, 1
 - (P1) ADD R1, R1, -1
- 3. Assuming
 - CPI = 0.5 without mispredictions
 - Misprediction penalty is 10 cycles
- 4. Do if-conversion if BEQZ prediction is < % correct
 - 0.5 * (2 + 3) / 2 + Accuracy * 10 = 3 / 2
 - 1.25 + Accuracy * 10 = 1.5
 - Accuracy = 0.25 / 10 = 1 / 40 = 2.5%
 - Require 97.5% accuracy

Conclusion

- 1. Branch prediction and predication handle control hazards
 - Need to consider how data hazards are handled while maintaining instruction flow in our pipeline