

# Introduction to Digital Systems

## Part I (4 lectures)

2023/2024

Introduction

Number Systems and Codes

Combinational Logic Design Principles

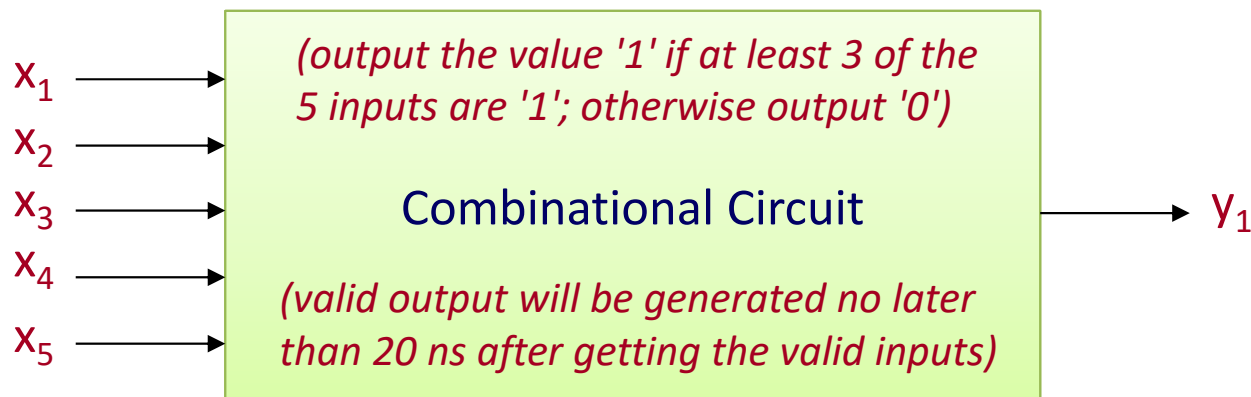
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# Lecture 3 contents

- Combinational circuits
- Boolean algebra
  - axioms
  - theorems
  - duality
  - algebraic simplification of logic functions
  - canonical forms
- Standard representations of logic functions

# Combinational Circuits

- A logic circuit whose outputs depend only on its current inputs is called a **combinational circuit**.  
*Apenas tem dependencia do presente.*
- A combinational circuit is characterized by
  - one or more inputs
  - one or more outputs
  - a functional specification describing each output as a function of the inputs
  - a time specification that includes at least the maximum time it will take the circuit to produce valid output values for an arbitrary set of input values -> **propagation delay**.



# Boolean Algebra

- Formal analysis techniques for digital circuits have their roots in the work of an English mathematician, George Boole.
- In 1854, he invented a two-valued algebraic system, now called **Boolean algebra**.
- In 1938, Bell Labs researcher Claude E. Shannon showed how to adapt Boolean algebra to analyze and describe the behavior of circuits.
- In switching algebra we use a symbolic variable, such as  $x$ , to represent the condition of a logic signal.
- A logic signal is in one of two possible conditions: low or high, off or on, and so on, depending on the technology.
- We say that  $x$  has the value "0" for one of these conditions and "1" for the other.

# Axioms

- The **axioms** (or **postulates**) of a mathematical system are a minimal set of basic definitions that we assume to be true, from which all other information about the system can be derived.
- A variable  $x$  can take on only one of two values:
  - $x = 0$  if  $x \neq 1$
  - $x = 1$  if  $x \neq 0$
- Inversion (definition of NOT):
  - if  $x = 0$  then  $\bar{x} = 1$
  - if  $x = 1$  then  $\bar{x} = 0$
- Definition of the AND and OR operations:
  - $0 \cdot 0 = 0, \quad 1 \cdot 1 = 1, \quad 0 \cdot 1 = 1 \cdot 0 = 0$
  - $1 + 1 = 1, \quad 0 + 0 = 0 \quad 1 + 0 = 0 + 1 = 1$
- These pairs of axioms completely define switching algebra.
- All other facts about the system can be proved using these axioms as a starting point.

# Operator Precedence

- **Operator precedence** is an ordering of logical operators designed to allow the dropping of parentheses in logical expressions.
- The following list gives a hierarchy of precedences for the Boolean operators (from highest to lowest):
  - NOT
  - AND
  - OR

Example:

$$x \cdot \bar{y} + z = (x \cdot (\bar{y})) + z$$

# Single-Variable Theorems

- Switching algebra **theorems** are statements, known to be always true, that permit us to manipulate algebraic expressions to allow simpler analysis or more efficient synthesis of the corresponding circuits.
- Identities:
  - $x + 0 = x$
  - $x \cdot 1 = x$
- Null elements:
  - $x + 1 = 1$
  - $x \cdot 0 = 0$
- Idempotency:
  - $x + x = x$
  - $x \cdot x = x$
- Involution:
  - $\bar{\bar{x}} = x$
- Complements:
  - $x + \bar{x} = 1$
  - $x \cdot \bar{x} = 0$

# Two- and Three-Variable Theorems

- Commutativity:
  - $x + y = y + x$
  - $x \cdot y = y \cdot x$
- Associativity:
  - $(x + y) + z = x + (y + z)$
  - $(x \cdot y) \cdot z = x \cdot (y \cdot z)$
- Distributivity:
  - $x \cdot y + x \cdot z = x \cdot (y + z)$
  - $(x + y) \cdot (x + z) = x + y \cdot z$
- Covering:
  - $x + x \cdot y = x$
  - $x \cdot (x + y) = x$
- Combining:
  - $x \cdot y + x \cdot \bar{y} = x$
  - $(x + y) \cdot (x + \bar{y}) = x$
- Simplification:
  - $x + \bar{x} \cdot y = x + y$
  - $x \cdot (\bar{x} + y) = x \cdot y$
- Consensus:
  - $x \cdot y + \bar{x} \cdot z + y \cdot z = x \cdot y + \bar{x} \cdot z$
  - $(x + y) \cdot (\bar{x} + z) \cdot (y + z) = (x + y) \cdot (\bar{x} + z)$



# Resume of Theorems

- Most theorems in switching algebra are exceedingly simple to prove using a technique called **perfect induction**:
  - prove a theorem by proving that it is true for all possible values (“0” and “1”) of all the variables
- In all of the theorems, **it is possible to replace each variable with an arbitrary logic expression**:
  - $x + x \cdot (a + b \cdot c \cdot \bar{d} \cdot e) = x$
- When realizing the AND(OR) operation, we can connect gate inputs in any order: either one n-input gate or  $(n - 1)$  2-input gates interchangeably, though propagation delay and cost are likely to be higher with multiple 2-input gates:
  - $w \cdot x \cdot y \cdot z = (w \cdot x) \cdot (y \cdot z) = (w \cdot (x \cdot (y \cdot z))) \dots$
- In Boolean algebra, logical addition distributes over logical multiplication:
  - $(x + y) \cdot (x + z) = x + y \cdot z$

The complement of a set "A" is the set of elements not in "A"

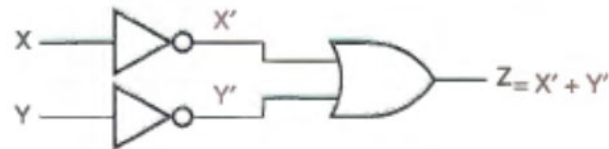
# DeMorgan's Theorems

- An  $n$ -input AND gate whose output is complemented is equivalent to an  $n$ -input OR gate whose inputs are complemented:

- $\overline{x \cdot y} = \bar{x} + \bar{y}$



- $\overline{\prod_{i=0}^{n-1} x_i} = \sum_{i=0}^{n-1} \bar{x}_i$

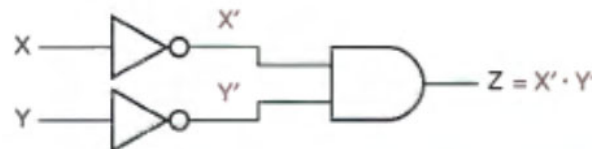


- An  $n$ -input OR gate whose output is complemented is equivalent to an  $n$ -input AND gate whose inputs are complemented:

- $\overline{x + y} = \bar{x} \cdot \bar{y}$



- $\overline{\sum_{i=0}^{n-1} x_i} = \prod_{i=0}^{n-1} \bar{x}_i$



# Generalized DeMorgan's Theorem

- Given any  $n$ -variable fully parenthesized logic expression, its complement can be obtained by swapping  $+$  and  $\cdot$  and complementing all variables:

$$\overline{F(x_0, x_1, \dots, x_{n-1}, +, \cdot)} = F(\overline{x_0}, \overline{x_1}, \dots, \overline{x_{n-1}}, \cdot, +)$$

Example:

$$F(w, x, y, z) = \bar{w} \cdot x + x \cdot y + w \cdot (\bar{x} + \bar{z})$$

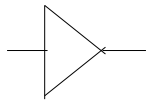
$$\overline{F(w, x, y, z)} = (w + \bar{x}) \cdot (\bar{x} + \bar{y}) \cdot (\bar{w} + x \cdot z)$$

# Principle of Duality

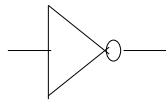
- The **principle of duality** states that any theorem or identity in switching algebra remains true if 0 and 1 are swapped and  $\cdot$  and  $+$  are swapped throughout.
  - Duals of all the axioms are true.
  - Duals of all switching-algebra theorems are true.
- If  $F(x_0, x_1, \dots, x_{n-1}, +, \cdot)$  is a fully parenthesized logic expression involving the variables  $x_0, x_1, \dots, x_{n-1}$ , and the operators  $+$  and  $\cdot$ , then the dual of  $F$ , written  $F^D$  is the same expression with  $+$  and  $\cdot$  swapped:
  - $F^D(x_0, x_1, \dots, x_{n-1}, +, \cdot) = F(x_0, x_1, \dots, x_{n-1}, \cdot, +)$
- The generalized DeMorgan's theorem may now be restated as follows:
  - $\overline{F(x_0, x_1, \dots, x_{n-1})} = F^D(\overline{x_0}, \overline{x_1}, \dots, \overline{x_{n-1}})$

# NAND and NOR Gates

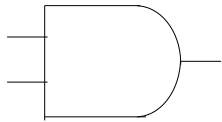
*buffer*



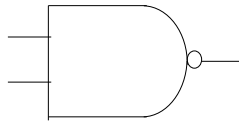
*NOT*



*AND*



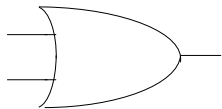
*NAND*



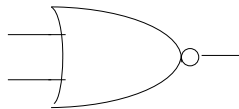
$$\overline{x \cdot y}$$

x	y	x <b>NAND</b> y
0	0	1
0	1	1
1	0	1
1	1	0

*OR*



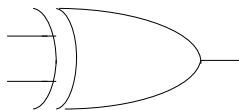
*NOR*



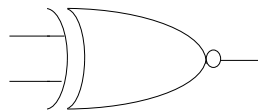
$$\overline{x + y}$$

x	y	x <b>NOR</b> y
0	0	1
0	1	0
1	0	0
1	1	0

*XOR*



*XNOR*



# Functional Completeness

- A **functionally complete set** of Boolean operators is one which can be used to describe the behavior of any digital circuit.
- Examples:
  - {AND, OR, NOT}
  - {AND, NOT}
  - {OR, NOT}
  - {NAND}
  - {NOR}

# NAND and NOR Gates

- To write a Boolean expression only with the operators **NAND**, first put the expression in the **sum-of-products** form and then apply the **involution** theorem ( $\bar{\bar{x}} = x$ ) followed by the **DeMorgan's** theorem ( $\overline{\sum_{i=0}^{n-1} x_i} = \prod_{i=0}^{n-1} \bar{x}_i$ ).
- To write a Boolean expression only with the operators **NOR**, first put the expression in the **product-of-sums** form and then apply the **involution** theorem ( $\bar{\bar{x}} = x$ ) followed by the **DeMorgan's** theorem ( $\overline{\prod_{i=0}^{n-1} x_i} = \sum_{i=0}^{n-1} \bar{x}_i$ ).
- Always assume that complemented versions of input variables are available.

## Examples:

$$x + (y \cdot \bar{z}) = \overline{\overline{x + (y \cdot \bar{z})}} = \overline{\bar{x} \cdot y \cdot \bar{z}}$$

$$x + (y \cdot \bar{z}) = (x + y) \cdot (x + \bar{z}) = \overline{\overline{(x + y) \cdot (x + \bar{z})}} = \overline{\bar{x} + y + \bar{x} + \bar{z}}$$

# Boolean Functions

- A Boolean function  $f(x_0, x_1, \dots, x_{n-1})$  is a match that associates an element of the set  $\{0,1\}$  with each of the  $2^n$  possible combinations that variables can assume.
- There are  $2^{m \times 2^n}$  different Boolean functions that can be implemented in a digital system with  $n$  inputs and  $m$  outputs.



## Examples:

For  $n=1, m=1$ :  $2^{1 \times 2^1} = 4$

x	constant '0'	x	$\bar{x}$	constant '1'
0	0	0	1	1
1	0	1	0	1

For  $n=4, m=3$ :  $2^{3 \times 2^4} = 2^{48} = 281\,474\,976\,710\,656$



# Truth Table

- The most basic representation of a logic function is the **truth table**.
- A truth table simply lists the output of the circuit for every possible input combination.
- Traditionally, the input combinations are arranged in rows in ascending binary counting order, and the corresponding output values are written in a column next to the rows.
- The truth table for an  $n$ -variable logic function has  $2^n$  rows.

Example ( $n=3$ ):

	x	y	z	$f(x,y,z)$
0	0	0	0	0
1	0	0	1	1
2	0	1	0	0
3	0	1	1	0
4	1	0	0	1
5	1	0	1	1
6	1	1	0	1
7	1	1	1	1

# Minterms and Maxterms

- A **literal** is a variable or the complement of a variable. Examples:  $x, y, \bar{x}$ .
- A **product term** is a single literal or a logical product of two or more literals. Examples:  $\bar{z}, x \cdot y, x \cdot \bar{y} \cdot z$ .
- A **sum term** is a single literal or a logical sum of two or more literals. Examples:  $\bar{z}, x + y, x + \bar{y} + z$ .
- A **normal term** is a product or sum term in which no variable appears more than once.
- An  $n$ -variable **minterm** is a normal product term with  $n$  literals. There are  $2^n$  such product terms.
  - A minterm  $m_i$  corresponds to row  $i$  of the truth table.
  - In minterm  $m_i$ , a particular variable appears complemented if the corresponding bit in the binary representation of  $i$  is 0; otherwise, it is uncomplemented.
- An  $n$ -variable **maxterm** is a normal sum term with  $n$  literals. There are  $2^n$  such sum terms.
  - A maxterm  $M_i$  corresponds to row  $i$  of the truth table.
  - In maxterm  $M_i$ , a particular variable appears complemented if the corresponding bit in the binary representation of  $i$  is 1; otherwise, it is uncomplemented.

Example:

	x	y	z	f(x,y,z)
0	0	0	0	0
1	0	0	1	1
2	0	1	0	0
3	0	1	1	0
4	1	0	0	1
5	1	0	1	1
6	1	1	0	1
7	1	1	1	1

$$m_0 = \bar{x} \cdot \bar{y} \cdot \bar{z}$$

$$M_0 = x + y + z$$

$$m_5 = x \cdot \bar{y} \cdot z$$

$$M_5 = \bar{x} + y + \bar{z}$$

$$m_i = \overline{M_i} \quad i = 0, 1, \dots, 2^n - 1$$

# Algebraic Representations

- Any Boolean function can be presented as:
  - a sum of the minterms corresponding to truth-table rows (input combinations) for which the function produces a 1 -> **canonical sum**
  - a product of the maxterms corresponding to truth-table rows (input combinations) for which the function produces a 0 -> **canonical product**

Example:

	x	y	z	f(x,y,z)
0	0	0	0	0
1	0	0	1	1
2	0	1	0	0
3	0	1	1	0
4	1	0	0	1
5	1	0	1	1
6	1	1	0	1
7	1	1	1	1

$$f(x, y, z) = \bar{x} \cdot \bar{y} \cdot z + x \cdot \bar{y} \cdot \bar{z} + x \cdot \bar{y} \cdot z + x \cdot y \cdot \bar{z} + x \cdot y \cdot z$$

$$= \sum_{x,y,z} m(1,4,5,6,7)$$

$$f(x, y, z) = (x + y + z) \cdot (x + \bar{y} + z) \cdot (x + \bar{y} + \bar{z})$$

$$= \prod_{x,y,z} M(0,2,3)$$

# Shannon's Expansion Theorems

- Any Boolean function  $f(x_0, x_1, \dots, x_{n-1})$  can be presented in the following forms:
  - $\overline{x_0} \cdot f(0, x_1, \dots, x_{n-1}) + x_0 \cdot f(1, x_1, \dots, x_{n-1})$
  - $(\overline{x_0} + f(1, x_1, \dots, x_{n-1})) \cdot (x_0 + f(0, x_1, \dots, x_{n-1}))$

Perfect induction:

If  $x_0 = 0$  then:  $f(0, x_1, \dots, x_{n-1}) = 1 \cdot f(0, x_1, \dots, x_{n-1}) + 0 \cdot f(1, x_1, \dots, x_{n-1})$

If  $x_0 = 1$  then:  $f(1, x_1, \dots, x_{n-1}) = 0 \cdot f(0, x_1, \dots, x_{n-1}) + 1 \cdot f(1, x_1, \dots, x_{n-1})$

# Canonical Sum

- Extending to 2 variables:

$$\begin{aligned} f(x_0, x_1, \dots, x_{n-1}) &= \bar{x}_0 \cdot f(0, x_1, \dots, x_{n-1}) + x_0 \cdot f(1, x_1, \dots, x_{n-1}) = \\ &= \bar{x}_0 \cdot \bar{x}_1 \cdot f(0, 0, x_2, \dots, x_{n-1}) + \bar{x}_0 \cdot x_1 \cdot f(0, 1, x_2, \dots, x_{n-1}) + \\ &+ x_0 \cdot \bar{x}_1 \cdot f(1, 0, x_2, \dots, x_{n-1}) + x_0 \cdot x_1 \cdot f(1, 1, x_2, \dots, x_{n-1}) \end{aligned}$$

- Extending to  $n$  variables:

$$f(x_0, x_1, \dots, x_{n-1}) = \sum_{i=0}^{2^n-1} m_i \cdot f_i \quad f_i = f((x_0, x_1, \dots, x_{n-1}) = i)$$

# Canonical Product

- Extend the Shannon expansion theorem  $f(x_0, x_1, \dots, x_{n-1}) = (\overline{x_0} + f(1, x_1, \dots, x_{n-1})) \cdot (x_0 + f(0, x_1, \dots, x_{n-1}))$  to  $n$  variables:

$$- f(x_0, x_1, \dots, x_{n-1}) = \prod_{i=0}^{2^n-1} (f_i + M_i)$$

# 3<sup>rd</sup> and 4<sup>th</sup> Canonical Forms

- 3<sup>rd</sup> canonical form:

$$f(x_0, x_1, \dots, x_{n-1}) = \overline{\overline{f(x_0, x_1, \dots, x_{n-1})}} = \overline{\sum_{i=0}^{2^n-1} f_i \cdot m_i} = \overline{\prod_{i=0}^{2^n-1} f_i \cdot m_i}$$

- 4<sup>th</sup> canonical form:

$$f(x_0, x_1, \dots, x_{n-1}) = \overline{\overline{f(x_0, x_1, \dots, x_{n-1})}} = \overline{\prod_{i=0}^{2^n-1} f_i + M_i} = \overline{\sum_{i=0}^{2^n-1} f_i + M_i}$$

# Canonical Forms

canonical sum  
of products:

$$f(x_0, x_1, \dots, x_{n-1}) = \sum_{i=0}^{2^n-1} m_i \cdot f_i$$

AND-OR

canonical product  
of sums:

$$f(x_0, x_1, \dots, x_{n-1}) = \prod_{i=0}^{2^n-1} (f_i + M_i)$$

OR-AND

3<sup>rd</sup> canonical form:

$$f(x_0, x_1, \dots, x_{n-1}) = \overline{\prod_{i=0}^{2^n-1} f_i \cdot m_i}$$

NAND-NAND

4<sup>th</sup> canonical form :

$$f(x_0, x_1, \dots, x_{n-1}) = \overline{\sum_{i=0}^{2^n-1} f_i + M_i}$$

NOR-NOR



# Canonical Forms (cont.)

**Example:** Derive the canonical forms of function  $f(x,y,z)$ :

$$f(x, y, z) = x \cdot y + \bar{z}$$

	x	y	z	$f(x,y,z)$
0	0	0	0	1
1	0	0	1	0
2	0	1	0	1
3	0	1	1	0
4	1	0	0	1
5	1	0	1	0
6	1	1	0	1
7	1	1	1	1

**1<sup>st</sup>:**  $f(x, y, z) = \sum m(0,2,4,6,7)$

**2<sup>nd</sup>:**  $f(x, y, z) = \prod M(1,3,5)$

**3<sup>rd</sup>:**  $f(x, y, z) = \overline{\prod m(0,2,4,6,7)}$

**4<sup>th</sup>:**  $f(x, y, z) = \overline{\sum M(1,3,5)}$

# Standard Representations of Logic Functions

- A truth table
- Algebraic
- Logic circuit

An algebraic representation frequently includes redundant terms:

$$f(x, y, z) = \bar{x} \cdot \bar{y} \cdot z + x \cdot \bar{y} \cdot \bar{z} + x \cdot \bar{y} \cdot z + x \cdot y \cdot \bar{z} + x \cdot y \cdot z$$

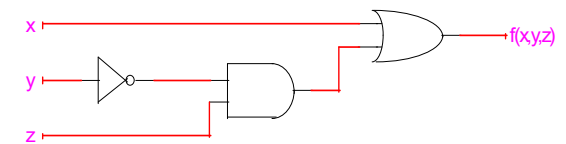
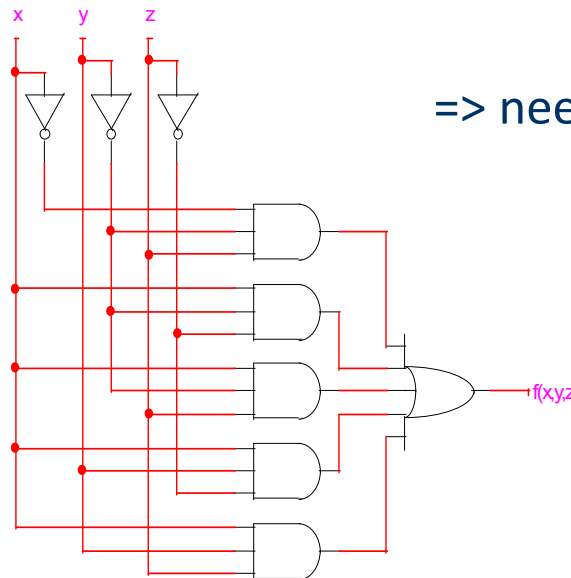
A truth table representation is unique:

x	y	z	f(x,y,z)
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

Logic circuit:

=> need for simplification

$$f(x, y, z) = x + \bar{y} \cdot z$$



# Exercises

- Are the following expressions correct?

$$[x + y \cdot z]^D = x \cdot y + z$$

$$[x + y \cdot z]^D = \bar{x} \cdot (\bar{y} + \bar{z})$$

- A self-dual logic function is a function  $f$  such that  $f = f^D$ . Which of the following functions are self-dual?

$$f_1(x, y, z) = \bar{x} \cdot y + \bar{x} \cdot z + y \cdot z$$

$$f_2(x, y) = \bar{x} \cdot y + x \cdot \bar{y}$$

# Exercises (cont.)

- Express the function  $y$  in the simplest form using only the operator NAND.

$$y = x_1 \cdot (x_2 + \bar{x}_3 \cdot x_4) + x_2$$

$$y = x_1 \cdot x_2 + x_1 \cdot \bar{x}_3 \cdot x_4 + x_2 = x_2 + x_1 \cdot \bar{x}_3 \cdot x_4$$

$$\overline{\overline{x_2 + x_1 \cdot \bar{x}_3 \cdot x_4}} = \overline{\overline{x_2} \cdot \overline{x_1 \cdot \bar{x}_3 \cdot x_4}} = \overline{\overline{x_2} \cdot x_1 \cdot \bar{x}_3 \cdot x_4}$$

- Express the function  $y$  in the simplest form using only the operator NOR.

$$y = x_1 \cdot (x_2 + \bar{x}_3 \cdot x_4) + x_2$$

$$y = (x_1 + x_2) \cdot (x_2 + \bar{x}_3 \cdot x_4 + x_2) = (x_1 + x_2) \cdot (x_2 + \bar{x}_3 \cdot x_4)$$

$$y = (x_1 + x_2) \cdot (x_2 + \bar{x}_3) \cdot (x_2 + x_4)$$

$$\overline{\overline{(x_1 + x_2) \cdot (x_2 + \bar{x}_3) \cdot (x_2 + x_4)}} = \overline{\overline{(x_1 + x_2)} + \overline{\overline{(x_2 + \bar{x}_3)}} + \overline{\overline{(x_2 + x_4)}}}$$

# Exercises (cont.)

- Determine all the canonical forms of the function  $f$ :

$$f(x, y, z) = x \cdot y + \bar{x} \cdot \bar{z} + y \cdot z$$

$$f(x, y, z) = \sum m(0, 2, 3, 6, 7) = \bar{x} \cdot \bar{y} \cdot \bar{z} + \bar{x} \cdot y \cdot \bar{z} + \bar{x} \cdot y \cdot z + x \cdot y \cdot \bar{z} + x \cdot y \cdot z$$

$$f(x, y, z) = \prod M(1, 4, 5) = (x + y + \bar{z}) \cdot (\bar{x} + y + z) \cdot (\bar{x} + y + \bar{z})$$

$$f(x, y, z) = \overline{\prod m(0, 2, 3, 6, 7)} = \overline{(\bar{x} \cdot \bar{y} \cdot \bar{z}) \cdot (\bar{x} \cdot y \cdot \bar{z}) \cdot (\bar{x} \cdot y \cdot z) \cdot (x \cdot y \cdot \bar{z}) \cdot (x \cdot y \cdot z)}$$

$$f(x, y, z) = \overline{\sum M(1, 4, 5)} = \overline{(x + y + \bar{z}) \cdot (\bar{x} + y + z) \cdot (\bar{x} + y + \bar{z})}$$

- Minimize this function.
- Minimize the following functions:

$$f(a, b, c) = \bar{a} \cdot b + \bar{a} \cdot \bar{c} + a \cdot c + a \cdot \bar{b} + b + c$$

$$f(a, b, c) = \bar{a} \cdot \bar{b} \cdot \bar{c} + \bar{a} \cdot b \cdot c + a \cdot b \cdot \bar{c} + a \cdot \bar{b} \cdot c$$