ICS 271

Fall 2016

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1. $X = \{v_1, v_2, v_3, v_4, v_5, v_6\}$

 $D = \{D_1, D_2, D_3, D_4, D_5, D_6\}$

- $D_1 = \{ \text{desk, easy, dove, else, help, kind, soon, this} \}$
- $D_2 = \{\text{eta, hat, her, him, one}\}$
- $D_3 = \{ at, be, he, it, on \}$
- $D_4 = \{\text{dance, usage, first, loses, fuels, haste, given, sense, sound, think}\}$
- $D_5 = \{\text{dance, usage, first, loses, fuels, haste, given, sense, sound, think}\}$
- $D_6 = \{ \text{desk, easy, dove, else, help, kind, soon, this} \}$

 $C_1 = \{v_2(1) = v_1(2), v_2(2) = v_3(2), v_2(3) = v_4(3), v_5(1) = v_4(4), v_5(3) = v_6(3)\}, v_i(j)$ means the *j*-th letter in word *i*.

$$C_2 = \{v_1 \neq v_6, v_4 \neq v_5\}$$

Cost function: the number of constraints unsatisified. (Range from 0 to 5)

Start: v_1 =this, v_2 =eta, v_3 =at, v_4 =dance, v_5 =haste, v_6 =dove, cost=4

Iter1: v_1 =help, v_2 =eta, v_3 =at, v_4 =dance, v_5 =haste, v_6 =dove, cost=3

Iter2: v_1 =help, v_2 =eta, v_3 =at, v_4 =usage, v_5 =haste, v_6 =dove, cost=2

Iter3: v_1 =help, v_2 =eta, v_3 =at, v_4 =usage, v_5 =given, v_6 =dove, cost=0

2. Assume there's a classes(subjects) list S, the number of elements in S is NoS, s_i is i-th subject in S, i = 1, 2, ..., NoS. We have,

$$S = \{s_i, i = 1, 2, \dots, NoS\}$$

Same as professors(P), classesrooms(C) and timeslots(T). That we have,

$$P = \{p_{\alpha}, \alpha = 1, 2, \dots, NoP\}$$

$$C = \{c_{\beta}, \beta = 1, 2, \dots, NoC\}$$

$$T = \{t_{\gamma}, \gamma = 1, 2, \dots, NoT\}$$

 $X = \{(s_i, p_\alpha, c_\beta, t_\gamma), i = 1, 2, \dots, NoS\}$, it means that subject s_i is taught by professor p_α at classroom c_β on timeslot t_γ .

 $D_j = \{(s_i, p_\alpha, c_\beta, t_\gamma) | i, \alpha, \beta, \gamma \text{ could be any valid value}\}$

C: - Each s_i should only appear once.

- For any two subjects $(s_{i1}, p_{\alpha 1}, c_{\beta 1}, t_{\gamma 1})$ and $(s_{i2}, p_{\alpha 2}, c_{\beta 2}, t_{\gamma 2})$,

if $p_{\alpha 1} = p_{\alpha 2}$, then $t_{\gamma 1} \neq t_{\gamma 2}$, it means a professor could not have two classes at same time;

if $c_{\beta 1} = c_{\beta 2}$, then $t_{\gamma 1} \neq t_{\gamma 2}$, it means one classroom could no have two classes at same time.

3. Assume big rectangle locate in a 2-D coordinate system, range from $[0, X] \times [0, Y]$.

 $X = \{(x_i, y_i, dX_i, dY_i) | i = 1, 2, \dots, n\}, \text{ each } X_i \text{ represents a small rectangle with it's position(bottom)}$

left corner) (x_i, y_i) and size (dX_i, dY_i) .

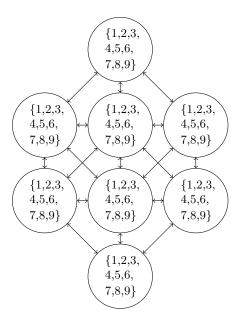
$$D_i = \{(x, y, dX, dY) | dX \leqslant X, dY \leqslant Y\}$$

 C_1 : For each small rectangle (x_i, y_i, dX_i, dY_i) , we have, $x_i \ge 0$ and $x_i + dX \le X$ and $y_i \ge 0$ and $y_i + dY \le Y$

 C_2 : For any of two small rectangles (x_i,y_i,dX_i,dY_i) and (x_j,y_j,dX_j,dY_j) , we have, $x_i+dX_i\leqslant x_j$ or $x_j+dX_j\leqslant x_i$ or $y_i+dY_i\leqslant y_j$ or $y_j+dY_j\leqslant y_i$

4. (a) For any neighboring pair of squares i and j such that $i \neq j$,

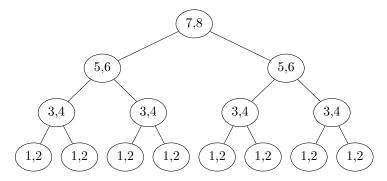
Relation	Domain		
Relation	Domain		
	(1,3), (1,4), (1,5), (1,6), (1,7), (1,8), (1,9),		
	(2,4), (2,5), (2,6), (2,7), (2,8), (2,9),		
	(3,1), (3,5), (3,6), (3,7), (3,8), (3,9),		
	(4,1), (4,2), (4,6), (4,7), (4,8), (4,9),		
$ X_i - X_j \geqslant 2$	(5,1), (5,2), (5,3), (5,7), (5,8), (5,9),		
	(6,1), (6,2), (6,3), (6,4), (6,8), (3,9),		
	(7,1), (7,2), (7,3), (7,4), (7,5), (7,9),		
	(8,1), (8,2), (8,3), (8,4), (8,5), (8,6),		
	(9,1), (9,2), (9,3), (9,4), (9,5), (9,6), (9,7)		



- (b) Yes, it's arc-consistent.
- (c) Yes, it's consistent. One solution is,

	2	
5	8	6
3	1	4
	7	

5. (a) Yes, it's arc-consistent.



(b) Yes, it's consistent. One solution is,

$$X_1 = 8$$
 $X_2 = 6$ $X_3 = 6$
 $X_4 = 4$ $X_5 = 4$ $X_6 = 4$ $X_7 = 4$
 $X_8 = 2$ $X_9 = 2$ $X_{10} = 2$ $X_{11} = 2$
 $X_{12} = 2$ $X_{13} = 2$ $X_{14} = 2$ $X_{15} = 2$

- (c) Do an in-order traversal: 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15
- (d) Since this is a tree the complexity is $O(nd^2)$ when n is the number of variables and d is the domain size.
- 6. Solve problem:

Constraints:

$$O + O = R + 10X_1$$

$$W + W + C_1 = U + 10C_2$$

$$T + T + C_2 = O + 10C_3$$

$$C_3 = F$$

$$all diff(T, W, O, F, U, R)$$

Domain:

$$D_F, D_{C_3} = \{1\}$$

$$D_{C_1}, D_{C_2} = \{0, 1\}$$

$$D_T = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

$$D_W, D_O, D_U, D_R = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

A solution:

A trace to a solution:

Node	State	Action
0	$D_F, D_{C_3} = \{1\}, D_{C_1}, D_{C_2} = \{0, 1\}, D_T = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}, D_W, D_O, D_U, D_R = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$	Initial State
1	$F = 1, D_{C_3} = \{1\}, D_{C_1}, D_{C_2} = \{0, 1\}, D_T = \{2, 3, 4, 5, 6, 7, 8, 9\}, D_W, D_O, D_U, D_R = \{0, 2, 3, 4, 5, 6, 7, 8, 9\}$	Assign $F = 1$
2	$F = 1, C_3 = 1, D_{C_1}, D_{C_2} = \{0, 1\}, D_T = \{5, 6, 7, 8, 9\}, D_W, D_O, D_U, D_R = \{0, 2, 3, 4, 5, 6, 7, 8, 9\}$	Assign $C_3 = 1$
3	$F = 1, C_3 = 1, C_2 = 0, D_{C_1} = \{0, 1\}, D_T = \{5, 6, 7, 8, 9\}, D_W = \{0, 2, 3, 4\}, $ $D_O = \{0, 2, 4, 6, 8\}, D_U, D_R = \{0, 2, 3, 4, 5, 6, 7, 8, 9\}$	Assign $C_2 = 0$
4	$F = 1, C_3 = 1, C_2 = 0, C_1 = 0, D_T = \{5, 6, 7, 8, 9\}, D_W = \{0, 2, 3, 4\},$ $D_O = \{0, 2, 4\}, D_U = \{0, 2, 4, 6, 8\} D_R = \{0, 2, 3, 4, 5, 6, 7, 8, 9\}$	Assign $C_1 = 0$
5	$F = 1, C_3 = 1, C_2 = 0, C_1 = 0, T = 6, O = 2, R = 4, D_W = \{0, 3\},\$ $D_U = \{0, 8\}$	Assign $O = 2$, then $T = 6, R = 4$
6	$F = 1, C_3 = 1, C_2 = 0, C_1 = 0, T = 7, O = 4, R = 8, D_W = \{0, 2, 3\},\ D_U = \{0, 2, 6\}$	Assign $W = 3$ or $W = 0$, no solution in D_U ; return to node 5, assign $O = 4$, then $T = 7, R = 8$
7	$F = 1, C_3 = 1, C_2 = 0, C_1 = 0, T = 7, O = 4, R = 8, D_W = \{3\}, D_U = \{6\}$	Assign $W = 0$ or $W = 2$, then no solution in D_U
8	$F = 1, C_3 = 1, C_2 = 0, C_1 = 0, T = 7, O = 4, R = 8, W = 3, U = 6$	Assign $W = 3$, then $U = 6$. Find Solution!

7. a_1 : arc consistency with domain splitting

 a_2 : variable elimination

 a_3 : stochastic local search

 a_4 : genetic algorithms

- (a) a_1 can determine that there is no solution, if the problem is inconsistent.
- (b) a_2, a_3 can find a solution if one exists.
- (c) a_4 can guarantee to find all solutions.