Assembly Language Programming for ATmega328P

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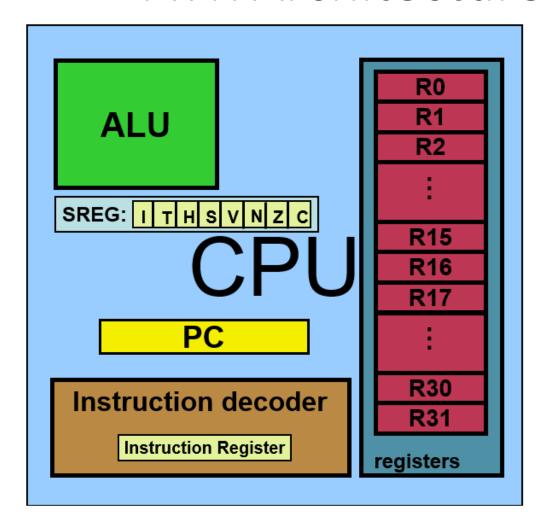




Outline

- Arduino AVR Architecture
- Assembly Language Programming
- Examples of Assembly Programming
- Flow Chart Drawing and Algorithm Development

AVR Architecture of CPU



- Arithmetic Logic Unit (ALU)
- 32 General Purpose Register (GPR) ← each 8-bit
- Program Counter (PC)
- Instruction Decoder
- Status Register (SREG)

AVR Flags/Status Register

- C = Carry Flag, This flag is set whenever there is a carry out from the D7 bit after an arithmetic operation(Addition, subtraction, increment, decrement, etc.). This flag bit is affected after an 8-bit addition or substruction.
- Z = Zero Flag, This flag is affected after an arithmetic or logic operation. If the result is zero then Z = 1, else Z = 0.
- N = Negative Flag, It reflects the result of an arithmetic operation. If the D7 bit of the result is zero, then N = 0 and the result is positive. Else N = 1 and the result is negative.
- V = Overflow Flag,
- S = Sign Flag,
- H = Half Carry Flag, this bit is set if there is a carry from D3 to D4 bit after ADD or SUB instruction.
- T = Bit Copy Storage. Used as a temporary storage for bit. It can be used to copy a bit from a GPR to another GPR.
- I = Global Interrupt Enable

Direct Data Register Method

- To assign input- 0
- To assign output- 1
- DDRB= 0x 05;
- Or DDRB= 0b 00000101; PB0 and PB2 are output pin while rest are input.
- To write High-1
- To write Low-0
- PORTB= 0x01;
- PORTB= 0b00000001; write PB0 as high.

PORTB Register

PB7	PB6	PB5	PB4	PB3	PB2	PB1	PB0
0	0	0	0	0	0	0	1

DDRB Register

PB7	PB6	PB5	PB4	PB3	PB2	PB1	PB0
0	0	0	0	0	1	0	1

Instruction Set Summary

Mnemonics	Operands	Description	Operations	Flags	Clocks
ADD	Rd,Rr	Add two Registers	$Rd \leftarrow Rd + Rr$	Z,C,N,V,H	1
ADC	Rd,Rr	Add with Carry two Registers	$Rd \leftarrow Rd + Rr + C$	Z,C,N,V,H	1
SUB	Rd,Rr	Subtract two Registers	Rd ← Rd - Rr	Z,C,N,V,H	1
SUBI	Rd,K	Subtract Constant from Register	$Rd \leftarrow Rd - K$	Z,C,N,V,H	1
AND	Rd,Rr	Logical AND Registers	Rd ← Rd • Rr	Z,N,V	1
OR	Rd,Rr	Logical OR Registers	$Rd \leftarrow Rd \ V \ Rr$	Z,N,V	1
NEG	Rd	Two's Complement	Rd ← 0x00 - Rd	Z,C,N,V,H	1
INC	Rd	Increment	$Rd \leftarrow Rd + 1$	Z,N,V	1
DEC	Rd	Decrement	Rd ← Rd - 1	Z,N,V	1
MUL	Rd,Rr	Multiply Unsigned	$R1:R0 \leftarrow Rd \times Rr$	Z,C	2

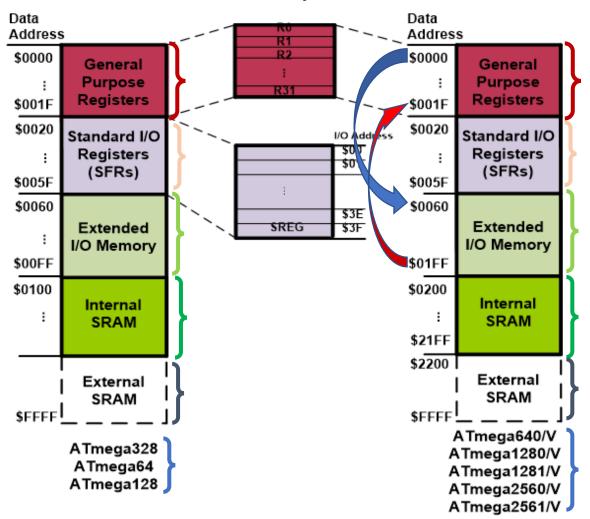
Instruction Set Summary

Mnemonics	Operands	Description	Operations	Flags	Clocks
TST	Rd	Test for Zero or Minus	$Rd \leftarrow Rd \bullet Rd$	Z,N,V	1
CLR	Rd	Clear Register	$Rd \leftarrow Rd \oplus Rd$	Z,N,V	1
MOV	Rd, Rr	Move Between Registers	Rd ← Rr	None	1
LDI	Rd, K	Load Immediate	$Rd \leftarrow K$	None	1
LDS	Rd, K	Load Direct from data space location K	Rd← [K]	None	1
IN	Rd, P	From In Port P/ address to Rd register	$Rd \leftarrow P$	None	1
OUT	P, Rr	From Rr register to Out Port P/ address	P ← Rr	None	1
PUSH	Rr	Push Register on Stack	$STACK \leftarrow Rr$	None	2
POP	Rd	Pop Register from Stack	$Rd \leftarrow STACK$	None	2
NOP		No Operation	None	1	
BREAK		Break	For On-chip Debug Only	None	N/A

Instruction Set Summary

Mnemonics	Operands	Description	Operations	Flags	Clocks
STS	K, Rr	Store Rr to data space location K.	[K] ← Rr	None	1
SBI	A, b	Sets a specified bit in an I/O Register	Sets the b no. bits of the A register	None	1
СВІ	A, b	Resets a specified bit in an I/O Register	Resets the b no bits of the A register	None	1
SBIC	A, b	Skip if Bit in I/O Register is Cleared/ Reset	If b no bit of A register is 0, then skip PC by 2 or 3.	None	1
SBIS A, b Skip if Bit in I/O Register is set		If b no bit of A register is 1, then skip PC by 2 or 3.	None	1	
RJUMP	К	Relative jump	PC← PC+K+1	None	1

Memory Locations



STS 0x60, R0

LDS R31, 0x01FF

Register Contents

SBI DDRB, 3;

DDRB Register

PB7	PB6	PB5	PB4	PB3	PB2	PB1	PB0
0	0	0	0	0	1	0	0

CBI DDRB, 5;

DDRB Register

PB	7	PB6	PB5	PB4	PB3	PB2	PB1	PB0
0		0	0	0	0	0	0	0

SBIC R7, 5; **R7 Register**

8	7	6	5	4	3	2	1
1	1	1	0	1	1	1	1

PC = PC+2 or 3

SBIS R8, 5;

R8 Register

8	7	6	5	4	3	2	1
1	1	1	1	1	1	1	1

PC = PC+2 or 3

Programming Examples

Example: State the contents of R20, R21, and data memory location of 0x120 after executing the following program.

LDI R20, 5; LDI R21, 2; ADD R20,R21; ADD R20,R21; STS 0x120, R20;

Solution:

LDI R20, 5; R20=5

LDI R21, 2; R21=2

ADD R20, R21; R20=5+2=7

ADD R20, R21; R20= 7+2=9

STS 0x120, R20; 0x120 = 9

R20 = 9

R21= 2

0x120 = 9

Example: State the contents of RAM locations of \$212, \$213, \$214, \$215, and \$216 after executing the following program.

LDI R16, 0x99 STS 0x212, R16 LDI R16, 0x85 STS 0x213, R16 LDI R16, 0x3F STS 0x214, R16 LDI R16, 0x63 STS 0x215, R16 LDI R16, 0x12 STS 0x216, R16

Solution:	
0x212	0x99
0x213	0x85
0x214	0x3F
0x215	0x63
0x216	0x12

Difference Between Assembly and C

Assembly Code Example sei; set Global Interrupt Enable **sleep**; enter sleep, waiting for interrupt ; note: will enter sleep before any pending interrupt(s) C Code Example enable interrupt(); /* set Global Interrupt Enable */ sleep(); /* enter sleep, waiting for interrupt */ /* note: will enter sleep before any pending interrupt(s) */

EEPROM Control Register (EECR)

7	6	5	4	3	2	1	0
		EEPN	I(1:0)	EERIE	EEMPE	EEPE	EERE

SBIC EECR, EEPE; If EEPE bit of EECR is 0 then the Program Counter (PC) will skip 2 or 3 steps.

Difference Between Assembly and C

Assembly Code Example

```
EEPROM write:
; Wait for completion of previous write
sbic EECR, EEPE
rjmp EEPROM write
; Set up address (r18:r17) in address register
out EEARH, r18
out EEARL, r17
; Write data (r16) to Data Register
out EEDR,r16
; Write logical one to EEMPE
sbi EECR,EEMPE
; Start eeprom write by setting EEPE
sbi EECR,EEPE
ret
```

C Code Example

```
void EEPROM_write(unsigned int uiAddress,
unsigned char ucData)
/* Wait for completion of previous write */
while(EECR & (1<<EEPE))
/* Set up address and Data Registers */
EEAR = uiAddress:
EEDR = ucData;
/* Write logical one to EEMPE */
EECR |= (1<<EEMPE);
/* Start eeprom write by setting EEPE */
EECR |= (1<<EEPE);
```

The next code examples show assembly and C functions for reading the EEPROM.

The examples assume that interrupts are controlled so that no interrupts will occur during the execution of these functions.

Assembly Code Example

```
EEPROM_read:

; Wait for completion of previous write

sbic EECR,EEPE

rjmp EEPROM_read

; Set up address (r18:r17) in address register

out EEARH, r18

out EEARL, r17

; Start eeprom read by writing EERE

sbi EECR,EERE

; Read data from Data Register

in r16,EEDR

ret
```

C Code Example

```
unsigned char EEPROM_read(unsigned int
uiAddress)
/* Wait for completion of previous write */
while(EECR & (1<<EEPE))
/* Set up address register */
EEAR = uiAddress;
/* Start eeprom read by writing EERE */
EECR |= (1<<EERE);
/* Return data from Data Register */
return EEDR;
```

Bit 0 – IVCE: Interrupt Vector Change Enable
The IVCE bit must be written to logic one to enable change of the IVSEL bit. IVCE is cleared by hardware
four cycles after it is written or when IVSEL is written. Setting the IVCE bit will disable interrupts,

```
Assembly Code Example
Move interrupts:
; Enable change of Interrupt Vectors
ldi r16, (1<<IVCE)
out MCUCR, r16; ;Output to MCUCR: Set
IVCE bit to enable changing interrupt
vectors
; Move interrupts to Boot Flash section
ldi r16, (1<<IVSEL)
out MCUCR, r16 ); Output to MCUCR: Set
IVSEL bit to move interrupts to Boot
Flash section
ret
```

```
void Move_interrupts(void)
{
  /* Enable change of Interrupt Vectors */
  MCUCR = (1<<IVCE);
  /* Move interrupts to Boot Flash section */
  MCUCR = (1<<IVSEL);
}</pre>
```

The assembly code example returns the TCNT1 value in the r17:r16 register pair. If an interrupt occurs between the two instructions accessing the 16-bit register, the interrupt code updates the temporary register by accessing the same or any other of the 16-bit Timer Registers, then the result of the access outside the interrupt will be corrupted. Therefore, when both the main code and the interrupt code update the temporary register, the main code must disable the interrupts during the 16-bit access.

```
Assembly Code Example

; Set TCNT1 to 0x01FF
Idi r17,0x01
Idi r16,0xFF
out TCNT1H,r17
out TCNT1L,r16
; Read TCNT1 into r17:r16
in r16,TCNT1L
in r17,TCNT1H
...
```

```
C Code Example
unsigned int i;
...
/* Set TCNT1 to 0x01FF */
TCNT1 = 0x1FF;
/* Read TCNT1 into i */
i = TCNT1;
...
```

The following code examples show how to do an atomic read of the TCNT1 Register contents.

Reading any of the OCR1A/B or ICR1 Registers can be done by using the same principle

```
Assembly Code Example
                                                        C Code Example
TIM16_ReadTCNT1:
                                                        uunsigned int TIM16 ReadTCNT1( void )
; Save global interrupt flag
  in r18, SREG ; Load the Status Register (SREG) into
                                                        unsigned char sreg;
register r18
                                                        unsigned int i;
; Disable interrupts
                                                        /* Save global interrupt flag */
            ; Clear the global interrupt flag to disable
  cli
                                                        sreg = SREG;
interrupts
                                                        /* Disable interrupts */
; Read TCNT1 into r17:r16
                                                        _CLI();
  in r16, TCNT1L; Load the low byte of TCNT1 into
                                                        /* Read TCNT1 into i */
register r16
                                                        i = TCNT1;
  in r17, TCNT1H; Load the high byte of TCNT1 into
                                                        /* Restore global interrupt flag */
register r17
                                                        SREG = sreg;
                                                        return i;
; Restore global interrupt flag
  out SREG, r18; Restore the original Status Register
value (including the interrupt flag)
            ; Return from the subroutine
 ret
```

The assembly code example returns the TCNT1 value in the r17:r16 register pair.

The following code examples show how to do an atomic write of the TCNT1 Register contents. Writing any of the OCR1A/B or ICR1 Registers can be done by using the same principle

Assembly Code Example

```
; Atomic write to TCNT1
Atomic Write TCNT1:
  ; Save global interrupt flag
  in r18, SREG
  ; Disable interrupts
  cli
  ; Write data to TCNT1
  out TCNT1L, r16; Assuming r16
contains the low byte data
  out TCNT1H, r17; Assuming r17
contains the high byte data
  ; Restore global interrupt flag
  out SREG, r18
  ret
```

C Code Example

```
void TIM16 WriteTCNT1(unsigned int value)
  unsigned char sreg;
  /* Save global interrupt flag */
  sreg = SREG;
  /* Disable interrupts */
  CLI();
  /* Write data to TCNT1 */
  TCNT1 = value;
  /* Restore global interrupt flag */
  SREG = sreg;
```

The assembly code example requires that the r17:r16 register pair contains the value to be written to TCNT1.

Assembly Code Example

```
TIM16_WriteTCNT1:
; Save global interrupt flag
in r18,SREG
; Disable interrupts
cli
; Set TCNT1 to r17:r16
out TCNT1H,r17
out TCNT1L,r16
; Restore global interrupt flag
out SREG,r18
ret
```

C Code Example

```
void TIM16_WriteTCNT1( unsigned int i )
unsigned char sreg;
unsigned int i;
/* Save global interrupt flag */
sreg = SREG;
/* Disable interrupts */
CLI();
/* Set TCNT1 to i */
TCNT1 = i;
/* Restore global interrupt flag */
SREG = sreg;
```

Difference Between Assembly and C

	Assembly Code Example	C Example	Comments
1	ldi r16, (1< <twint) (1<<twsta)="" ="" <br="">(1<<twen) out TWCR, r16</twen) </twint)>	TWCR = (1< <twint) (1<<twsta)="" ="" <br="">(1<<twen)< th=""><th>Send START condition</th></twen)<></twint)>	Send START condition
2	wait1: in r16,TWCR sbrs r16,TWINT rjmp wait1	while (!(TWCR & (1< <twint)));< td=""><td>Wait for TWINT Flag set. This indicates that the START condition has been transmitted</td></twint)));<>	Wait for TWINT Flag set. This indicates that the START condition has been transmitted
	in r16,TWSR andi r16, 0xF8 cpi r16, START brne ERROR	if ((TWSR & 0xF8) != START) ERROR();	Check value of TWI Status Register. Mask prescaler bits. If status different from START go to ERROR
3	ldi r16, SLA_W out TWDR, r16 ldi r16, (1< <twint) (1<<twen)="" out="" r16<="" td="" twcr,="" =""><td>TWDR = SLA_W; TWCR = (1<<twint) (1<<twen);<="" td="" =""><td>Load SLA_W into TWDR Register. Clear TWINT bit in TWCR to start transmission of address</td></twint)></td></twint)>	TWDR = SLA_W; TWCR = (1< <twint) (1<<twen);<="" td="" =""><td>Load SLA_W into TWDR Register. Clear TWINT bit in TWCR to start transmission of address</td></twint)>	Load SLA_W into TWDR Register. Clear TWINT bit in TWCR to start transmission of address
4	wait2: in r16,TWCR sbrs r16,TWINT rjmp wait2	while (!(TWCR & (1< <twint)));< td=""><td>Wait for TWINT Flag set. This indicates that the SLA+W has been transmitted, and ACK/NACK has been received.</td></twint)));<>	Wait for TWINT Flag set. This indicates that the SLA+W has been transmitted, and ACK/NACK has been received.

Difference Between Assembly and C

	<pre>in r16,TWSR andi r16, 0xF8 cpi r16, MT_SLA_ACK brne ERROR</pre>	<pre>if ((TWSR & 0xF8) != MT_SLA_ACK) ERROR();</pre>	Check value of TWI Status Register. Mask prescaler bits. If status different from MT_SLA_ACK go to ERROR
5	ldi r16, DATA out TWDR, r16 ldi r16, (1< <twint) (1<<twen)="" out="" r16<="" td="" twcr,="" =""><td>TWDR = DATA; TWCR = (1<<twint) (1<<twen);<="" td=""><td>Load DATA into TWDR Register. Clear TWINT bit in TWCR to start transmission of data</td></twint)></td></twint)>	TWDR = DATA; TWCR = (1< <twint) (1<<twen);<="" td=""><td>Load DATA into TWDR Register. Clear TWINT bit in TWCR to start transmission of data</td></twint)>	Load DATA into TWDR Register. Clear TWINT bit in TWCR to start transmission of data
6	wait3: in r16,TWCR sbrs r16,TWINT rjmp wait3	while (!(TWCR & (1< <twint)));< td=""><td>Wait for TWINT Flag set. This indicates that the DATA has been transmitted, and ACK/NACK has been received.</td></twint)));<>	Wait for TWINT Flag set. This indicates that the DATA has been transmitted, and ACK/NACK has been received.
7	<pre>in r16,TWSR andi r16, 0xF8 cpi r16, MT_DATA_ACK brne ERROR</pre>	<pre>if ((TWSR & 0xF8) != MT_DATA_ACK) ERROR();</pre>	Check value of TWI Status Register. Mask prescaler bits. If status different from MT_DATA_ACK go to ERROR
•	<pre>ldi r16, (1<<twint) (1<<twen)="" td="" ="" <=""><td>TWCR = (1<<twint) (1<<twen)="" ="" <br="">(1<<twsto);< td=""><td>Transmit STOP condition</td></twsto);<></twint)></td></twint)></pre>	TWCR = (1< <twint) (1<<twen)="" ="" <br="">(1<<twsto);< td=""><td>Transmit STOP condition</td></twsto);<></twint)>	Transmit STOP condition

Flowchart

The first design of flowchart goes back to 1945 which was designed by John Von Neumann. Unlike an algorithm, Flowchart uses different symbols to design a solution to a problem. It is another commonly used programming tool. By looking at a Flowchart, one can understand the operations and sequence of operations performed in a system.

A flowchart is often considered a blueprint of a design used for solving a specific problem. A flowchart is a diagrammatic representation of an algorithm. Flowcharts are very helpful in writing programs and explaining programs to others.

Though, flowcharts are useful in efficient coding, debugging, and analysis of a program, drawing flowcharts for very complicated in case of complex programs is often ignored. Flowcharts use special shapes to represent different types of actions or steps in a process. Lines and arrows show the sequence of the steps and the relationships among them. These are known as flowchart symbols. So, flowchart symbols are specific shapes used to create a visual representation of a program.

Advantages of Flowchart

- 1. A flowchart is an excellent way of communicating the logic of a program.
- 2. Easy and efficient to analyze the problem using a flowchart.
- 3. During the program development cycle, the flowchart plays the role of a blueprint, which makes the program development process easier.
- 4. After the successful development of a program, it needs continuous timely maintenance during the course of its operation. The flowchart makes a program or system maintenance easier.
- 5. It is easy to convert the flowchart into any programming language code.

Commonly used Operators in Flowchart

While creating flowchart, the following mathematical operators are used.

Operator	Meaning	Example	Remarks
+	Addition	A + B	Values in A and B will be added
-	Subtraction	A - B	Values of B will be subtracted from that of A
*	Multiplication	A * B	Values of B will be multiplied by that of A
/	Division	A/B	Values of A will be divided by that of B
٨	Power	A^3	Values of A will be raised by a power of 3
%	Remainder	A%B	Values of A will be divided by that of A to get a remainder (if any)

Commonly used Operators in Flowchart

While creating flowchart, the following relational operators are used.

Operator	Meaning	Example	Remarks
<	Less than	A < B	Value in A is less than that of B
<=	Less than or equal to	A <= B	Value in A is less than or equal to that of B
= or ==	Equal to	A = B or A == B	Values of A and B are equal
# or !=	Not Equal to	A # B, A != B	Values of A and B are not equal
>	Greater than	A > B	Value in A is greater than that of B
>=	Greater than or equal to	A >= B	Value in A is greater than or equal to that of B

Commonly used Operators in Flowchart

While creating flowchart, the following logical operators are used.

Operator	Meaning	Example	Remarks
AND	Logical AND	A^B	Values in A and B will be ANDed
OR	Logical OR	AVB	Values in A and B will be ORed
NOT	Logical NOT	!A	Values in A will be inverted

Commonly used Symbols in Flowchart

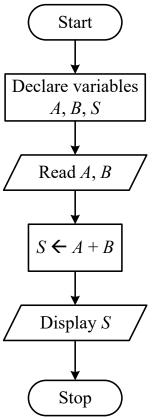
Symbol Namo	Symbol	function
Symbol Name	Symbol	tunction
Oval		Used to represent start and end of flowchart
Parallelogram		Used for input and output operation
Rectangle		Processing: Used for arithmetic operations and data-manipulations
Diamond	\Diamond	Decision making. Used to represent the operation in which there are two/three alternatives, true and false etc
Arrows	$\leftarrow^{\uparrow}_{\downarrow}$	Flow line Used to indicate the flow of logic by connecting symbols
Circle	0	Page Connector
		Off Page Connector
		Predefined Process /Function Used to represent a group of statements performing one processing task.
		Preprocessor
	 	Comments

While creating a flowchart, the following symbols are used.

- Rectangle Shape Represents a process
- Oval or Pill Shape Represents the start or end
- Diamond Shape Represents a decision
- Parallelogram Represents input/output
- Circle Represents reference/page connector
- Arrow Represents a connection between shapes

Examples of Flowchart

Draw a flowchart along with its algorithm to add two numbers entered by the user.

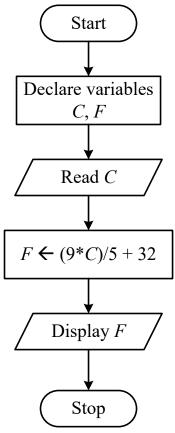


An algorithm to find the sum of two numbers:

- > Step 1: Start
- > Step 2: Declare three variables
- ➤ Step 3: Read numbers and store them in the declared variables
- Step 4: Sum these two numbers and assign them to the declared variable
- Step 5: Display the sum
- > Step 6: Stop

Examples of Flowchart

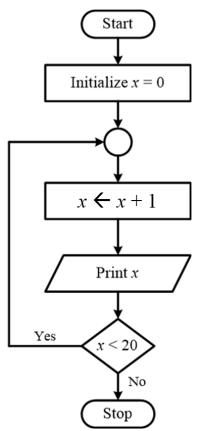
Draw a flowchart along with its algorithm to convert temperature from Celsius to Fahrenheit.



- Step 1: Start
- Step 2: Declare two variables
- ➤ Step 3: Read the temperature value in Celsius scale and store it in one of the declared variables
- > Step 4: Apply the formula of temperature conversion i.e., F = (9*C)/5 + 32 to the convert Celsius value in Fahrenheit scale and store it in the declared variable
- > Step 5: Display the temperature in Fahrenheit scale
- > Step 6: Stop

Examples of Flowchart

Draw a flowchart along with its algorithm to print numbers from 1 to 20.



- > Step 1: Start
- > Step 2: Declare and initialize a variable with a 0 value
- ➤ Step 3: Increment the variable value by 1 and store it in the same declared variable
- > Step 4: Print the value stored in the variable
- Step 5: If the variable is < 20 then go to step 3, otherwise go to the next step</p>
- > Step 6: Stop

Thank You