

Section: K18FG

Report On SIMULATION BASED PROCESS SCHEDULER

For Course Operating Systems

(CSE 316)

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Problem Description

Problem states that consider n number of processes and design a process management scheduler, which assign a processor to each process on the basics given conditions. Problem say's take arrival time of the process and the required burst time of that process as input from the user at run time and calculate the completion time, turnaround time and waiting time of all the process and also calculates the average turnaround time and waiting time of the processes. While calculating the completion time of the process, we have to follow some set of rules.

- The scheduler schedules the processes by interrupting the processor after every 3 units of time and does consider the completion of the process in this iteration. (First iteration)
- The schedulers then checks for the number of processes waiting for the processor and allots the processor to the process but interrupting the processor after every 6 units of time and considers the completion of the process in this iteration. (Second iteration)
- The scheduler after the second iteration checks for the number of processes waiting for the processor and now provides the processor to the process with the least time requirement to go in the terminated state.

Basically we have to implement 2 algorithms for design such type scheduler,

- 1. Round-Robin
- 2. Shortest job first

Round-Robin: Round-Robin algorithm having a fixed time quantum to execute each process in this type of scheduling approach every process will execute at a certain time and no process will go to starvation.

Shortest job First: It is a non-primitive scheduling algorithm where it works on the basis of arrival time and burst time it assigned CPU to that very process which required less time to fully execute and then move to next process in this process large process may be go at starvation.

Test Cases

Test Case: 1

Sample Input:

Process	Arrival time	Burst time
P ₁	0	18
P ₂	2	23
P ₃	4	13
P ₄	13	10

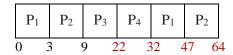
Expected Output:

Process	Arrival time	Burst time	Completion time	Turnaround time	Waiting time
P ₁	0	18	47	47	29
P ₂	2	23	64	62	39
P ₃	4	13	22	18	5
P ₄	13	10	32	19	9

Average Turnaround time: 36.5 Average Waiting time: 20.5

<u>Explanation</u>: The time quantum of 1 iteration is 3 unit of time, the time quantum of 2 iteration is 6 unit of time and after 2 iteration provide CPU of that process which required less time to go terminated state.

Gantt chart of Process Completion:



Test Case: 2

Sample Input:

Process	Arrival time	Burst time
P_1	0	20
P ₂	3	36
P ₃	13	19

P ₄	20	42
P ₅	36	69
P ₆	41	35

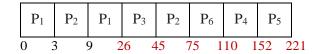
Expected Output:

Process	Arrival time	Burst time	Completion time	Turnaround time	Waiting time
P ₁	0	20	26	26	6
P ₂	3	36	75	72	36
P ₃	13	19	45	32	13
P ₄	20	42	152	132	90
P ₅	36	69	221	185	116
P ₆	41	35	110	69	34

Average Turnaround time: 86 Average Waiting time: 49.1667

<u>Explanation</u>: The time quantum of 1 iteration is 3 unit of time, the time quantum of 2 iteration is 6 unit of time and after 2 iteration provide CPU of that process which required less time to go terminated state.

Gantt chart of Process Completion:



Test Case: 3Sample Input:

	1	
Process	Arrival time	Burst time
P ₁	0	6
P ₂	3	16
P ₃	5	2
P ₄	9	8
P ₅	10	7

P ₆	22	5
P ₇	14	13
P ₈	26	4
P ₉	17	19
P ₁₀	28	3

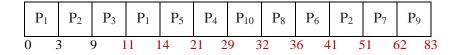
Expected Output:

Process	Arrival time	Burst time	Completion time	Turnaround time	Waiting time
P ₁	0	6	14	14	8
P ₂	3	16	51	48	32
P ₃	5	2	11	6	4
P ₄	9	8	29	20	12
P ₅	10	7	21	11	4
P ₆	22	5	41	19	14
P ₇	14	13	64	50	37
P ₈	26	4	36	10	6
P ₉	17	19	83	66	47
P ₁₀	28	3	32	4	1

Average Turnaround time: 24.8 Average Waiting time: 16.5

Explanation: The time quantum of 1 iteration is 3 unit of time, the time quantum of 2 iteration is 6 unit of time and after 2 iteration provide CPU of that process which required less time to go terminated state.

Gantt chart of Process Completion:



Algorithm

```
Step -1: Declare array arival_time[], busrt_time[], completion_time[],
waiting_time[], trunaround_time[] and int variable size.
Step -2: Take input number of process in size.
Step -3: Repeat for int I = 0,1,2... Size-1
      Take arrival time and burst time of the process as input.
Step -4: Remaining process = size , time_quantam_i1=3, time_quantam_i2=6;
      current time=0 and Repeat step -5 to step - 9 while remaining process != 0
Step -5: if(remanning_bt[Process_no]<time_quantam_i1 && remanning_bt
       [Process_no]>=0&&current_time<9)
           current_time+=remanning_bt[Process_no];
           remanning bt[Process no]=0
           indicator = 1
           time_quantam_itration++
           remmaning process--
Step -6: else if(remanning_bt[Process_no]>0&&current_time<9)
           if(time_quantam_itration==1)
           remanning_bt[Process_no]-=time_quantam_i1
           time_quantam_itration++
           current_time+=time_quantam_i1
           else if(time quantam itration==2)
           remanning_bt[Process_no]-=time_quantam_i2
           current_time+=time_quantam_i2
Step -7: else if(remanning_bt[Process_no]<9 && remanning_bt
       [Process no]>=3&&current time<9)
           current_time+=remanning_bt[Process_no];
           remanning bt[Process no]=0;
```

```
remmaning_process--, indicator = 1
Step -8: else if(remanning_bt[Process_no]>3&&current_time<9)
           remanning_bt[Process_no]-=time_quantam_i2
           current_time+=time_quantam_i2
Step -9: if(remanning bt[Process no]==0 && indicator==1)
           Remening_proc--
           completion time[Process no]=current time
           trunaround_time[Process_no]=completion_time[Process_no]-
            arival_time[Process_no]
           waiting_time[Process_no]=trunaround_time[Process_no]-
           bust_time[Process_no]
           indicator = 0
Step -10: End of step 4 loop, set:- Process_no=0
         Repeat Step -11 to Step - 12 while remaining process != 0
Step -11: Pick a job whose burst time is lesser among all of them.
Step -12: Execute that job and set
           completion_time[Process_no]=current_time
           trunaround_time[Process_no]=completion_time[Process_no]-
            arival_time[Process_no]
           waiting_time[Process_no]=trunaround_time[Process_no]-
           bust_time[Process_no]
Step -13: End of step 10 loop.
Step -14: Repeat for int I = 0,1,2... Size-1
        A) Calculate average waiting time and average burst time.
        B) Print all the information.
```

Time Complexity

Step -1 and Step -2 having time complexity constant time complexity: O(1).

Step -3 having time complexity no of process : O(n).

Step -4 to Step-9 having time complexity two iteration : O(m).

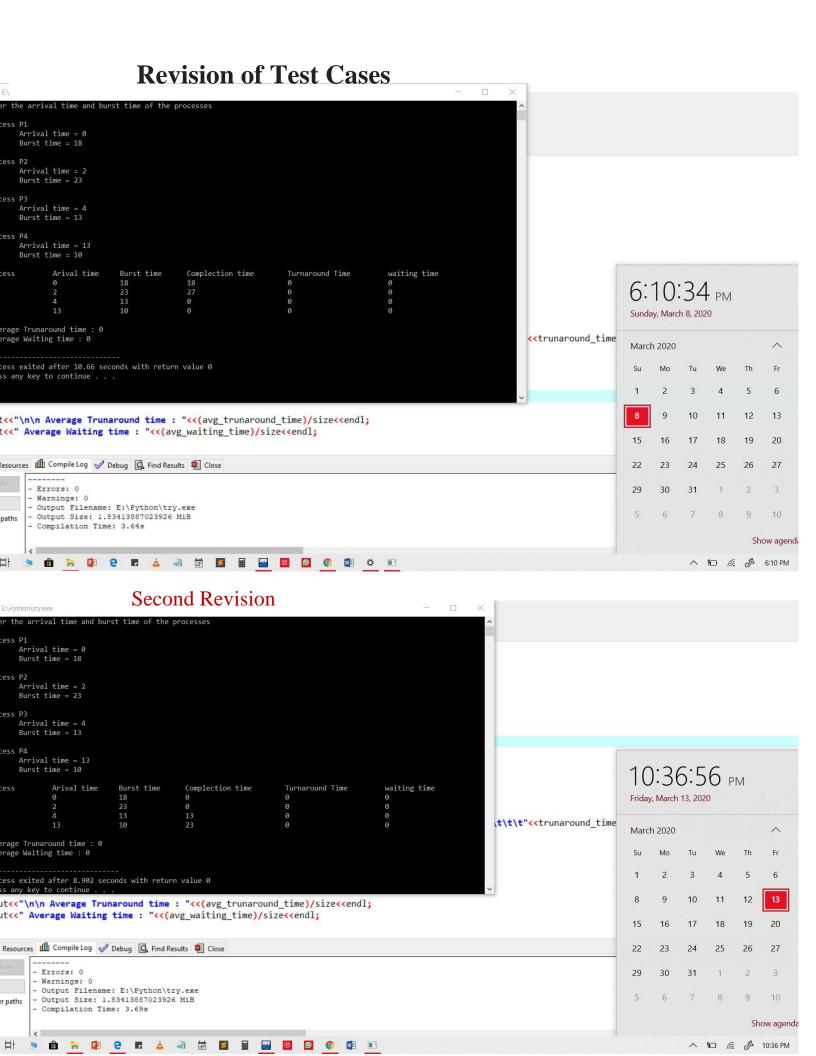
Step -10 to Step 13 having time complexity : $O(n^2)$.

Step -11 having time complexity no of process : O(n).

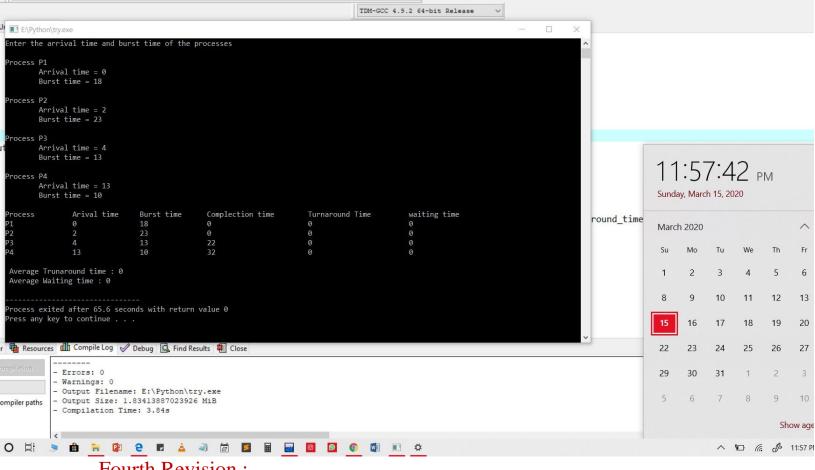
Overall Time Complexity is : $O(n^2)$.

Boundary Conditions

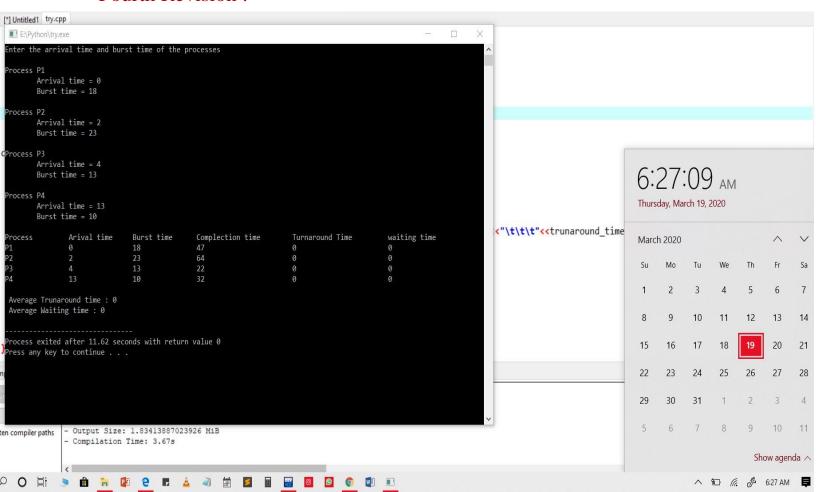
- 1. The first iteration should be follow Round Robin scheduling algorithm with time quantum of 3.
 - Arrival time must be greater and equal to 0
 - If any process having burst time less than 3, then set its completion time is burst time and remove from ready queue.
- 2. Second iteration also should be follow the Round Robin scheduling algorithm but having time quantum of 6.
 - Arrival time must be greater and equal to 0
 - And also check the ready queue if any process waiting already then execute first that very process.
 - If any process having burst time less than 9, then set its completion time is burst time and remove from ready queue.
- 3. After the second iteration follow then Shortest Job First scheduling algorithm.
 - Current time quantum must be greater or equals to 9
 - And also check the ready queue if any process waiting already then
 execute first that very process according to which have less time
 required to fully execute.
 - While picking up any shortest job to execute must be check that its arrival time equals of greater than current execution time.



Third Revision:



Fourth Revision:



Fifth Revision: Successful

